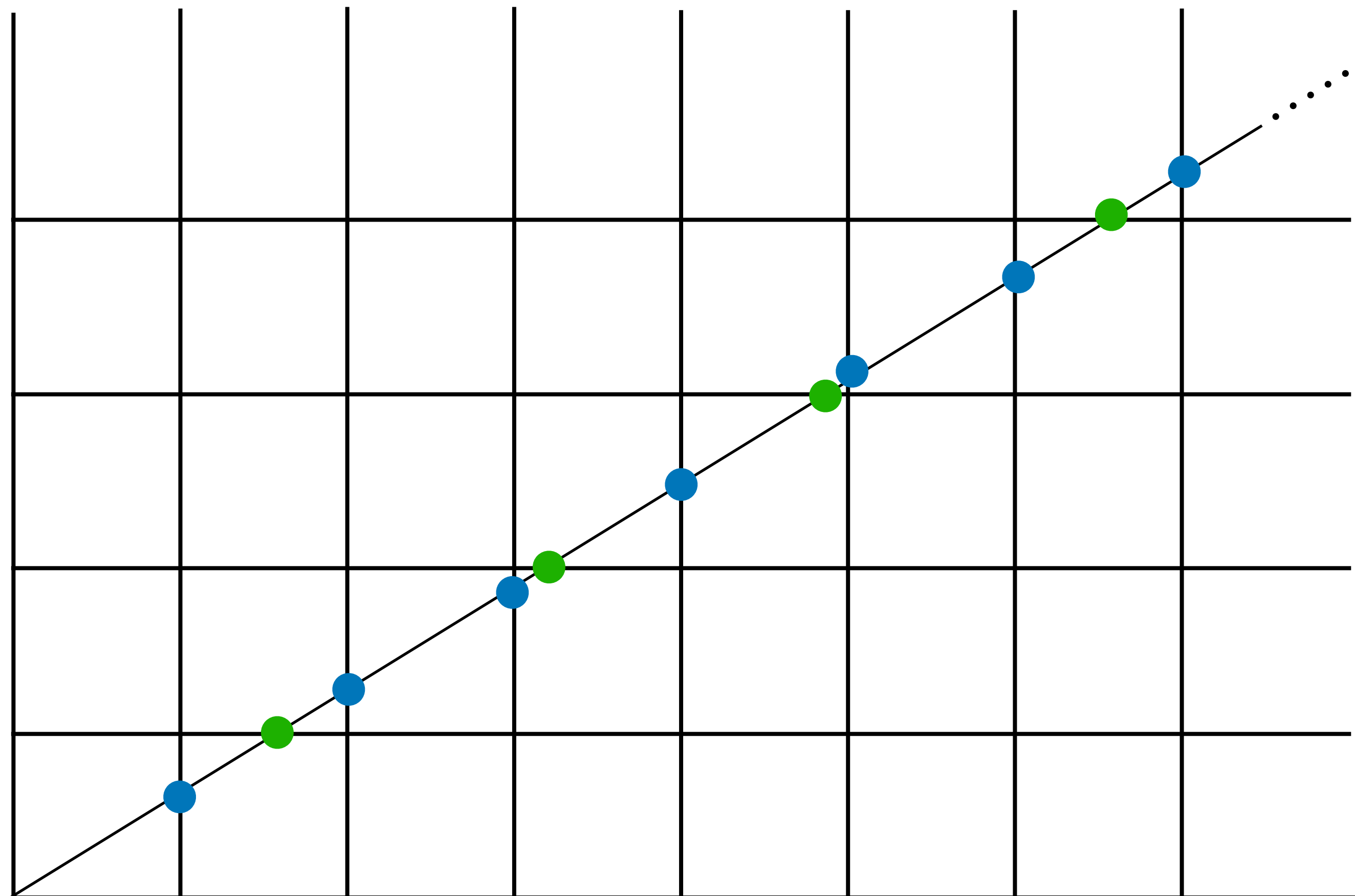


Automata on S-adic words

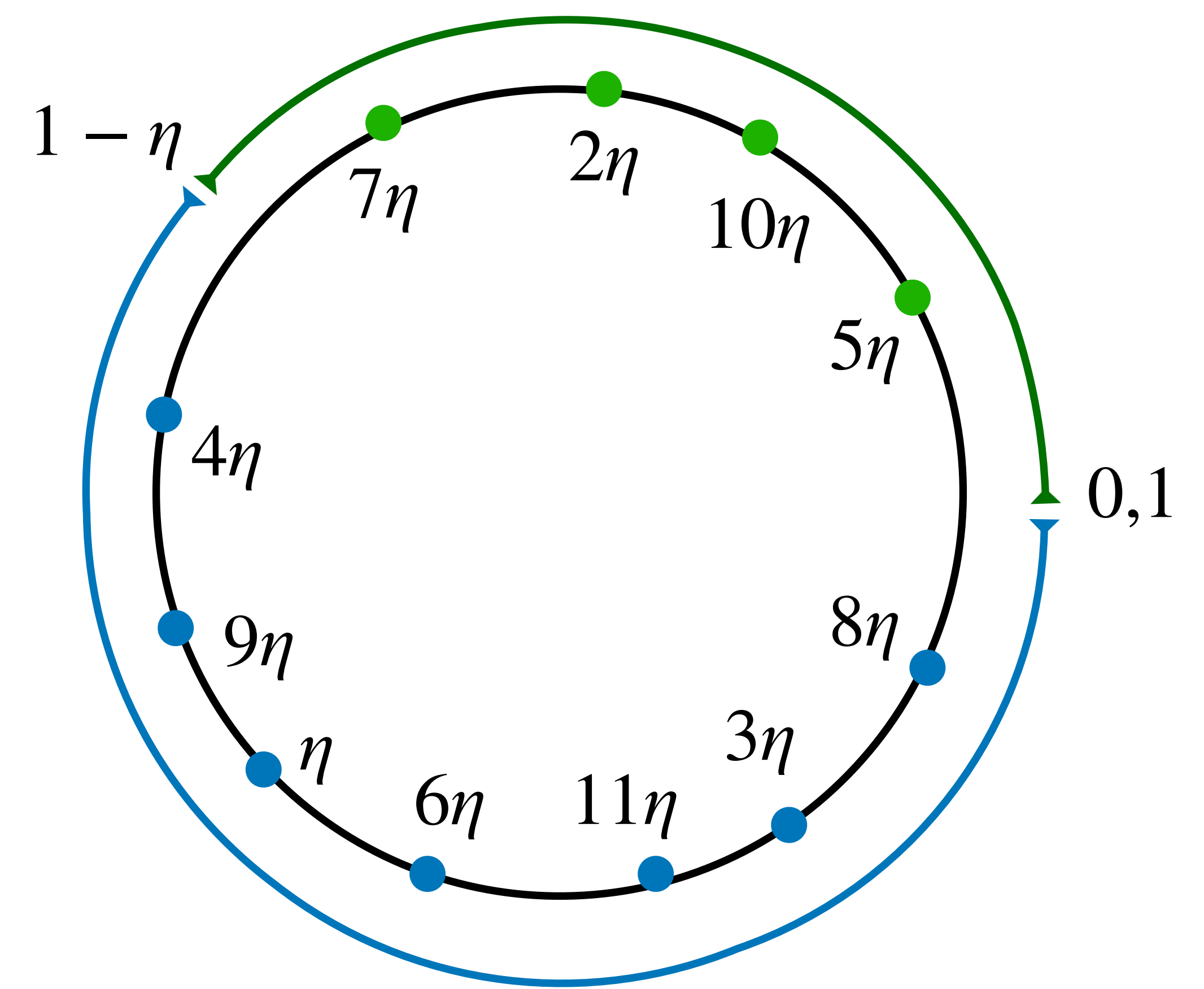
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Some families of symbolic dynamical systems, such as **Sturmian words**, have several equivalent characterisations and enjoy neat combinatorial properties that reflect **self-similarity**

Definition by example: the Fibonacci word



Cutting sequence, obtained by drawing a line of slope η , and recording the order of crossing grid lines
Here $\eta = 1/\phi$, where ϕ is the golden ratio



●●●●●●●●●● pattern obtained as the itinerary of a rotation by η around a circle of circumference 1

$$\mathbf{x}(k) = \lfloor (k+2) \cdot \eta \rfloor - \lfloor (k+1) \cdot \eta \rfloor$$

Formula for k -th letter of Sturmian word $\mathbf{x} \in \{0,1\}^\omega$

Any Sturmian word has exactly $n+1$ length- n factors, e.g., ●●●, ●●●, ●●●, ●●●

The Problem

Given an ω -regular language L , effectively characterise the set of Sturmian words in L

Key Fact

Any Sturmian word \mathbf{x} is the image of a Sturmian word \mathbf{x}' under a Sturmian substitution

Sturmian substitutions $S = \{\lambda_0, \lambda_1\}$, where $\lambda_0(0) = 0$, $\lambda_0(1) = 01$, $\lambda_1(0) = 10$, $\lambda_1(1) = 1$

Consequence: A Sturmian word \mathbf{x} is directed by a sequence $\mathbf{s} = s_0 s_1 \dots \in S^\omega$, i.e., there are Sturmian words $\mathbf{x}^{(0)} = \mathbf{x}$, $\mathbf{x}^{(1)}$, $\mathbf{x}^{(2)}$, ... satisfying $\mathbf{x}^{(n)} = s_n(\mathbf{x}^{(n+1)})$

This allows to further deduce $\mathbf{x} = \lim s_0 s_1 \dots s_n(0)$

Sturmian word \mathbf{x}

●●●●●●●●●● generated from \mathbf{s} , has sophisticated combinatorial properties

Directive sequence \mathbf{s}

●●●●●●●●●● obtained from slope, is subject to simpler constraints

Theorem. Given ω -regular $L \subseteq \{0,1\}^\omega$, we can compute ω -regular $L' \subseteq S^\omega$ such that a Sturmian word $\mathbf{x} \in L$ if and only if its directive sequence $\mathbf{s} \in L'$