

Preservation theorems for transducer outputs

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Combinatorial virtues feat. the Fibonacci word

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Morphic: $x_{\text{fib}} = \sigma_{\text{fib}}(x_{\text{fib}})$, where $\sigma_{\text{fib}}(0) = 01$, $\sigma_{\text{fib}}(1) = 0$

Recurrent: every factor occurs infinitely often

Furthermore, it is linearly, and hence uniformly recurrent



(linearly) bounded gaps between occurrences of u

Well-defined factor frequencies

For every u , $\lim_{N \rightarrow \infty} \frac{1}{N} |\{i : i < N, x[i, \infty) \in u\Sigma^\omega\}|$ is well-defined

Suppose a word $x \in \Sigma^\omega$ is

Morphic

(Linearly, Uniformly) Recurrent

Well-defined factor frequencies

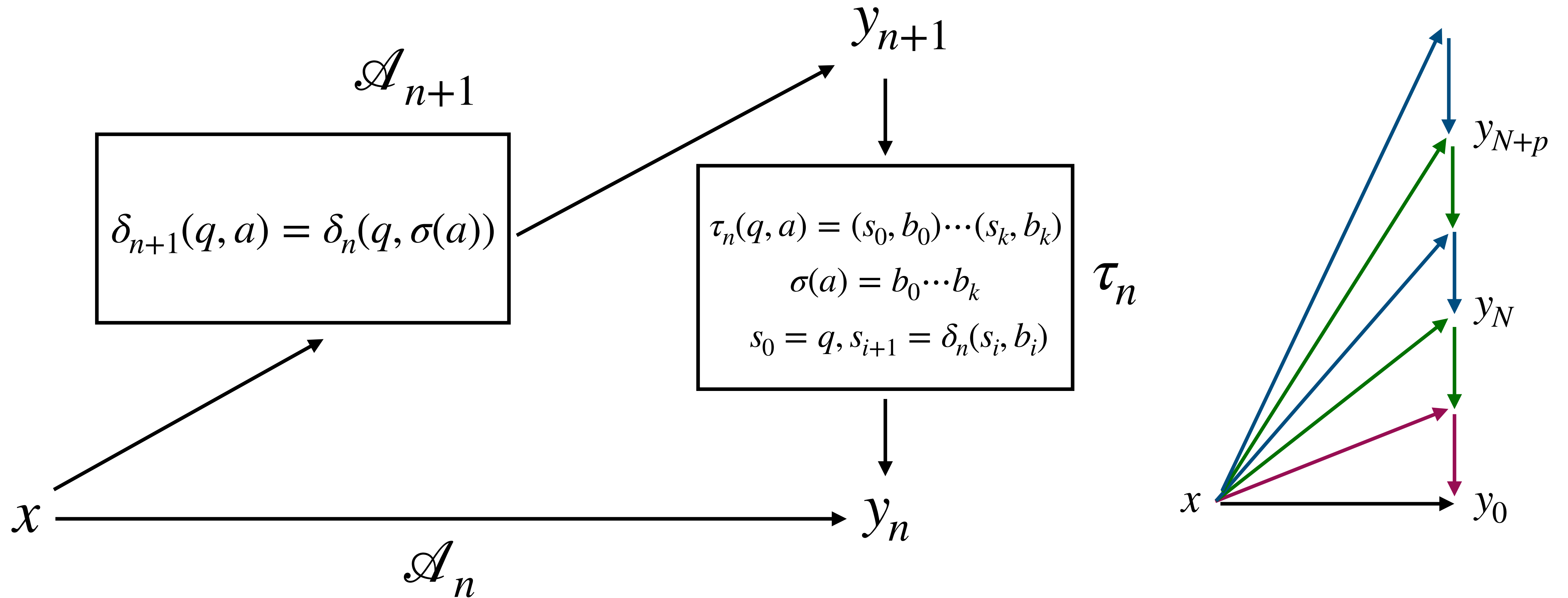
We feed x to an automaton $\mathcal{A} = (Q, q_{\text{init}}, \Sigma, \delta)$ to get run $y \in (Q \times \Sigma)^\omega$

Are these properties preserved?

The run of a substitutive word is morphic

$$x \in \Sigma^\omega, x = \sigma(x)$$

$$(Q, q_{\text{init}}, \Sigma, \delta_0) = \mathcal{A}_0$$



$$y_{N+p} = y_N, y_0 = \tau_0 \cdots \tau_{N-1}(y_N), y_N = \tau_N \cdots \tau_{N+p-1}(y_N + p)$$

Suppose a word $x \in \Sigma^\omega$ is

Morphic ✓

(Linearly, Uniformly) Recurrent

Well-defined factor frequencies

Does the run $y \in (Q \times \Sigma)^\omega$ of \mathcal{A} on x have these properties?

Theorem (Semënov et al., Pritykin).

If x is uniformly (resp. linearly) recurrent, then

y has a suffix that is uniformly (resp. linearly) recurrent

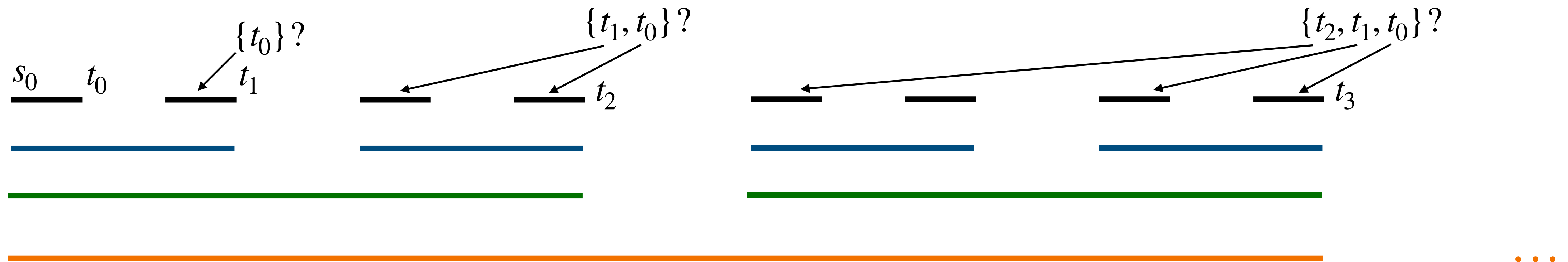
We can further prove:

- The corresponding statement holds for recurrent words
- Furthermore, we can effectively identify the suffix

A flavour of the Semënov argument

Assume that \mathcal{A} is a permutation automaton, i.e.,
each letter induces a permutation of Q (the transitions are reversible)

If x is recurrent, so is the run y : any occurrence of v in y has another to its right
Suppose v is the run on u_0 starting in s_0 and ending in t_0



There are finitely many states, this process saturates in at most $|Q| + 1$ iterations

We can also argue that uniform and linear recurrence are preserved



The other special case: no reversibility

Let \mathcal{A} be a reset automaton, i.e.,
the transition function induced by each letter is either the identity or constant
do nothing reset to a state

Let x be recurrent and let the first reset letter occur at position N

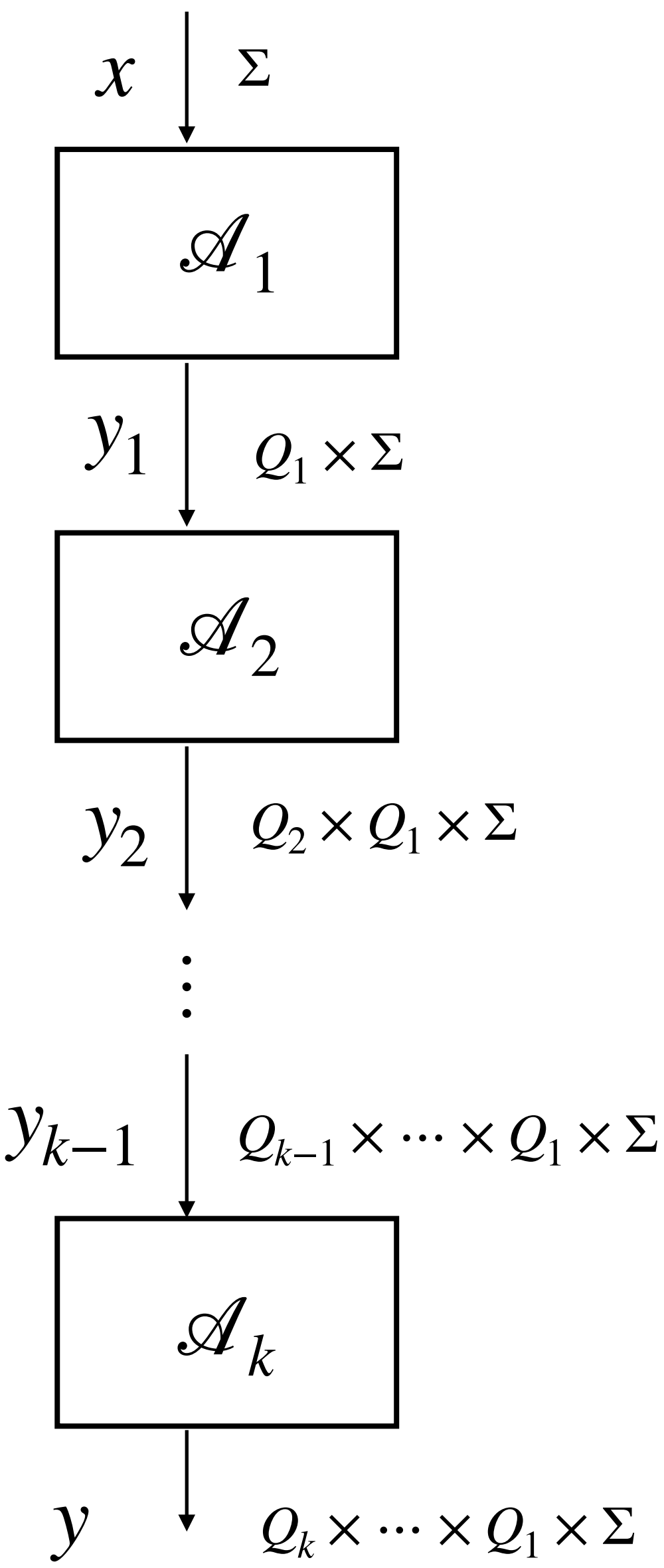
The run y of \mathcal{A} on x is recurrent starting at position $N + 1$



The linear and uniform recurrence of y , as applicable, can readily be argued ■

Preservation of recurrence holds for cascades of special automata

permutation automata, reset automata



$$Q = Q_k \times \dots \times Q_1$$

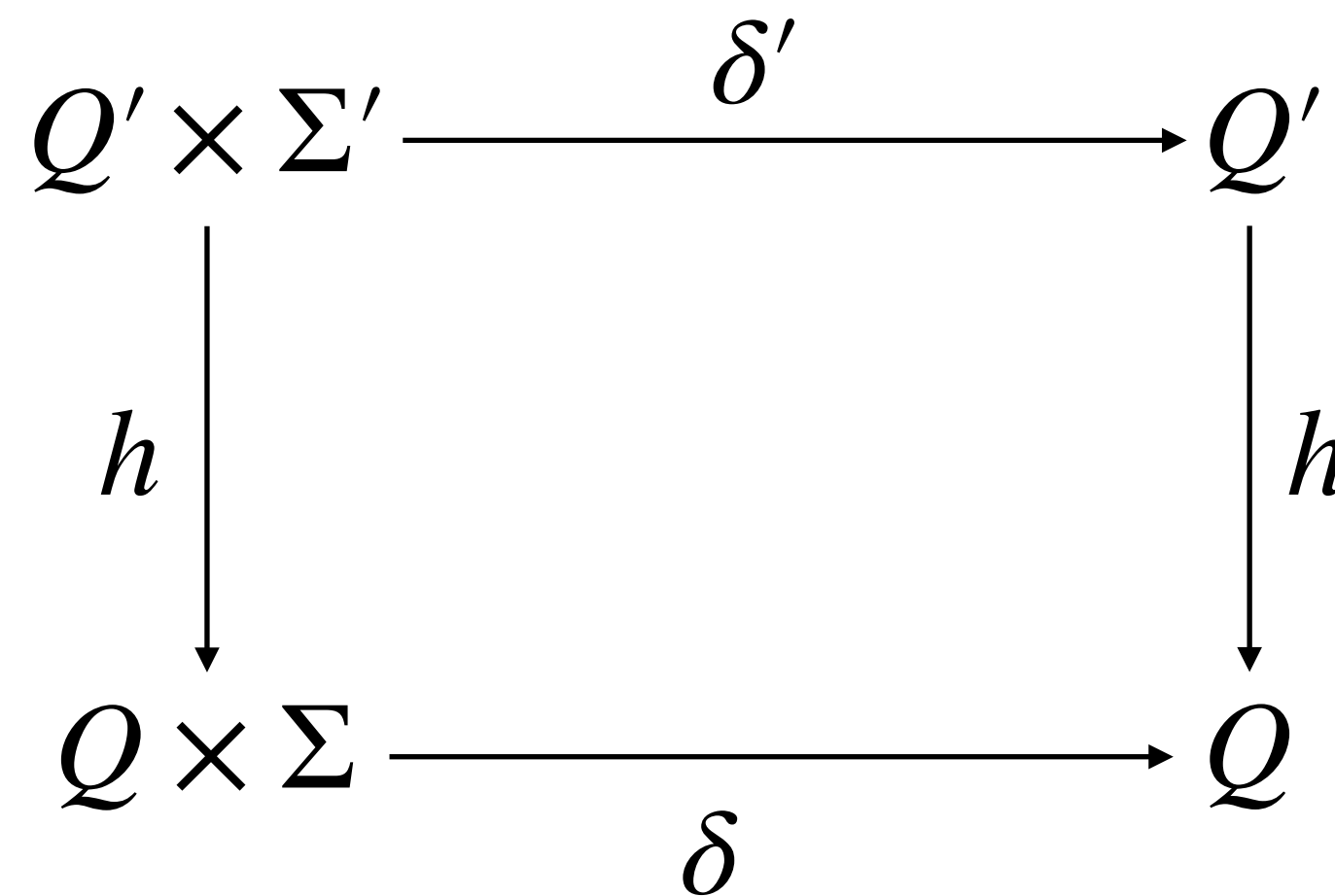
$$\begin{aligned} &\delta((q^k, \dots, q^1), a) \\ &= \\ &\left(\delta_k(q^k, (q^{k-1}, \dots, q^1, a)), \right. \\ &\quad \dots, \\ &\quad \delta_j(q^j, (q^{j-1}, \dots, q^1, a)), \\ &\quad \dots, \\ &\quad \left. \delta_1(q^1, a) \right) \end{aligned}$$

Preservation of recurrence holds for cascades of special automata

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Theorem (Krohn-Rhodes).

Every automaton can be simulated by a cascade of special automata



commutative diagram of primed automaton simulating unprimed automaton through morphism h

Preservation of recurrence holds for all automata



Suppose a word $x \in \Sigma^\omega$ has some of the following properties

Morphic ✓

(Linearly, Uniformly) Recurrent ✓

Well-defined factor frequencies

Yes for reset, and hence counter-free automata¹

Does the run $y \in (Q \times \Sigma)^\omega$ of \mathcal{A} on x have these properties?

¹ A corollary of the Krohn-Rhodes theorem is that any counter-free automaton can be simulated by a cascade of resets.

The other special case: no reversibility

Let \mathcal{A} be a reset automaton, i.e.,
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The linear and uniform recurrence of y , as applicable, can readily be argued ■

Suppose a word $x \in \Sigma^\omega$ has some of the following properties

Morphic ✓

(Linearly, Uniformly) Recurrent ✓

Well-defined factor frequencies

Yes for reset, and hence counter-free automata¹

Not necessarily for permutation automata

Yes, if Boshernitzan's condition holds

E.g., primitive morphic words, linearly recurrent words, Sturmian words, almost all IETs, a large class of 1D-toric words, Arnoux-Rauzy words with bounded partial quotients, ...

Does the run $y \in (Q \times \Sigma)^\omega$ of \mathcal{A} on x have these properties?

Primitive morphic \longrightarrow Linearly recurrent \longrightarrow Boshernitzan \longrightarrow Factor frequencies

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The first 13 letters of the Fibonacci word

0 : 8 occurrences, frequency 0.615

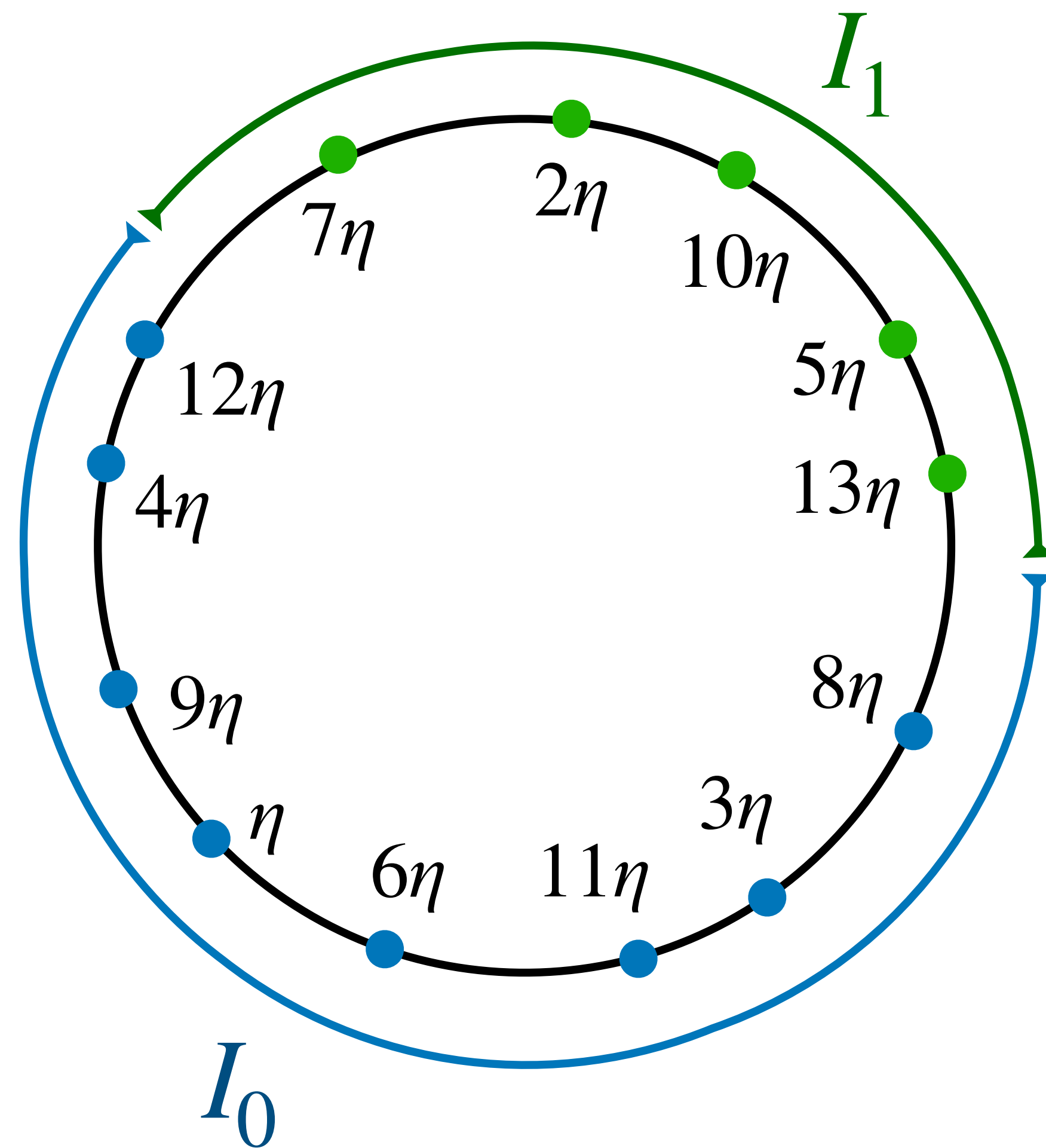
1 : 5 occurrences, frequency 0.385

In the limit,

0 occurs with frequency $1/\phi \approx 0.618$

1 occurs with frequency $1/\phi^2 \approx 0.382$

●●●●●... word obtained
 as the coding of a rotation by $\eta = 1/\phi$
 around a circle of circumference 1



$$I_{au} = I_a \cap (I_u - \eta)$$

frequency of factor u
 =
 length of arc I_u

Trace generated by **Topological Dynamical System**
Fibonacci word Rotation

Observation attributed to **Causing Region**
factor u arc I_u

The punchline, when ergodic theory applies, is:

**temporal frequency of observation
equals
spatial measure of causing region**

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The "default" topological dynamical system

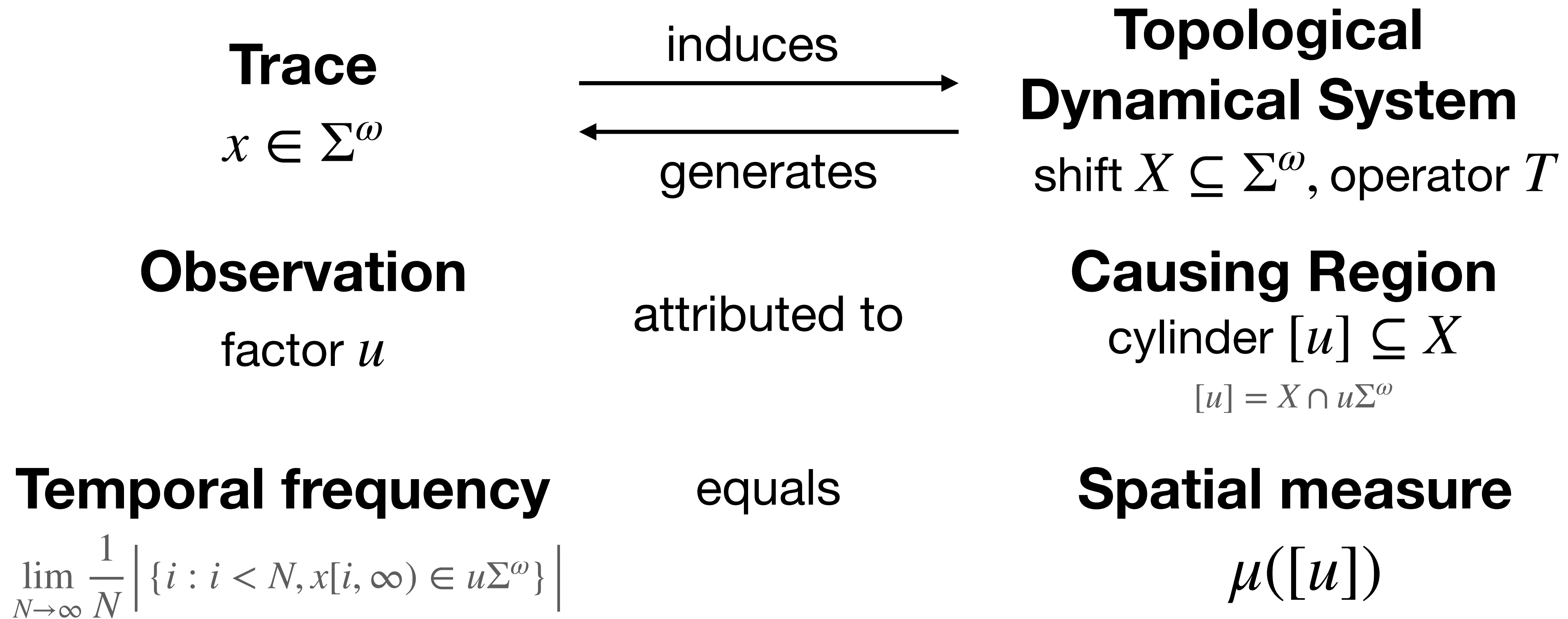
trace $x \in \Sigma^\omega$

endow Σ^ω with the product topology

shift operator $T : \Sigma^\omega \rightarrow \Sigma^\omega$, $T(a_0a_1a_2\cdots) = a_1a_2\cdots$

Shift (space) $X \subseteq \Sigma^\omega$: closure of orbit (x, Tx, T^2x, \dots)

When and how does this realise the ergodic dream?



We want sensible, invariant measures, i.e., $\mu(L) = \mu(T^{-1}L)$

We consider shifts X that have a unique invariant measure
In such shifts, all words have well-defined factor frequencies
(Oxtoby's theorem)

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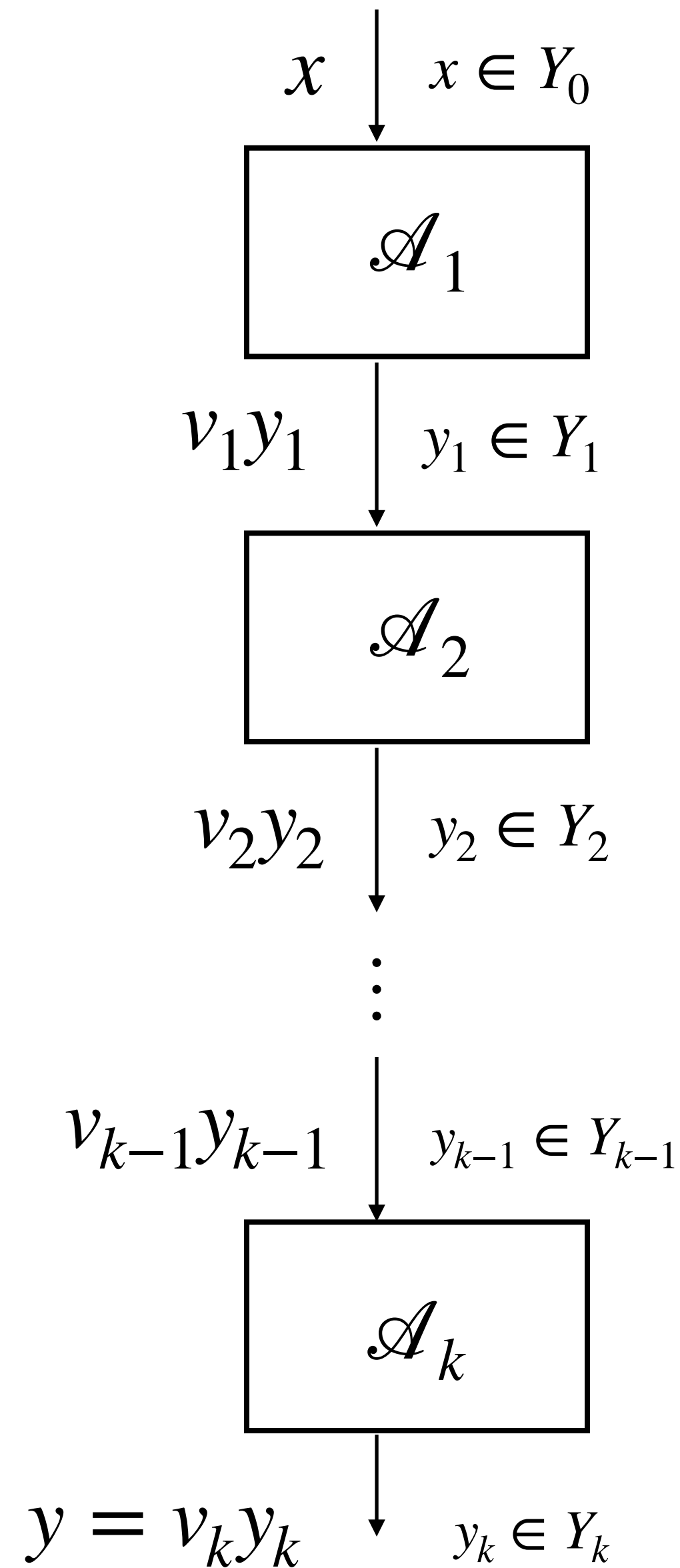
Primitive morphic \longrightarrow Linearly recurrent \longrightarrow Boshernitzan \longrightarrow Factor frequencies

Boshernitzan's condition

If a minimal shift X has an invariant measure μ such that

$$\liminf_{n \rightarrow \infty} \left(n \cdot \min_{\{u: |u|=n\}} \mu([u]) \right) > 0$$

then μ is the unique invariant measure of X



Lemma

If Y_i satisfies Boshernitzan's condition
then so does Y_{i+1}

Corollary

If Y_0 satisfies Boshernitzan's condition
then x, y have well-defined factor frequencies

Bonus

If x is primitive morphic with computable factor frequencies
then so is y_k

Suppose a word $x \in \Sigma^\omega$ has some of the following properties

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Does the run $y \in (Q \times \Sigma)^\omega$ of \mathcal{A} on x have these properties?

Thank You!

Extra technical slides

Consider $x \in X$ and permutation automaton \mathcal{A}

The run y is obtained from the skew product $G \rtimes X$

The group G is the transition group of \mathcal{A} , the morphism $\varphi : \Sigma \rightarrow G$

The dynamical transition of $G \rtimes X$ is given by T

$$T \circ (g, ax) = (g \cdot \varphi(a), x)$$

Let v be the run on u starting in q , let C be the coset that maps q_{init} to q

We hope that the frequency of v is the measure of $(C, [u])$

But is $G \rtimes X$ even minimal?

Does its relevant minimal component have a unique invariant measure?

But is $G \rtimes X$ even minimal?

Key construct: cobounding map α from X to set \mathcal{C} of cosets of $H \leq G$

By definition, α is required to be continuous

X is partitioned into finitely many (clopen) cylinders on which α is constant

The argument uses the minimality and compactness of X

Define $Y_\alpha = \{(g, x) : g \in \alpha(x)\} \subseteq G \rtimes X$

The point: if H is a minimal subgroup for which a cobounding map exists then Y_α is minimal

Note, however, that there are $|G|/|H|$ choices of Y_α , we make the one aligning with the run of \mathcal{A}

Does Y_α have a unique invariant measure?

Let H be the minimal subgroup for which cobounding α exists

Let \mathcal{C} be the set of right cosets of H

Recall $Y_\alpha = \{(g, x) : g \in \alpha(x)\}$. Does Y_α have a unique invariant measure?

Consider $Z_\alpha = \{(C, x) : \alpha(x) = C\} \subseteq \mathcal{C} \rtimes X$

Z_α is isomorphic to X , and in particular has a unique invariant measure

To define the invariant measure on Z_α , recall that X is partitioned into cylinders on which α is constant

To determine whether a cylinder exists in Z_α at all, we can use Semënov's theorem

An invariant measure for $Y_\alpha : \mu_Y(g, L) = \frac{1}{|H|} \mu_Z(Hg, L)$

Intuitively, this distributes mass equally to all elements of the coset

Could this be *the* invariant measure for Y ?

An invariant measure for $Y_\alpha : \mu_Y(g, [u]) = \frac{1}{|H|} \mu_Z(Hg, [u]) = \frac{1}{|H|} \mu_X([u])$

For u long enough

Intuitively, this distributes mass equally to all elements of the coset

Could this be *the* invariant measure for Y ?

Suppose μ_X is the unique invariant measure by virtue of satisfying

$$\liminf_{n \rightarrow \infty} \min_{|u|=n} n\mu_X([u]) > 0,$$

the Boshernitzan condition which is sufficient for unique ergodicity

Then, by inspection, μ_Y also satisfies the Boshernitzan condition and is indeed the unique invariant measure of Y_α

Let \mathcal{A} be a permutation automaton and let x be such that its shift X has a unique invariant measure that satisfies the Boshernitzan condition

The run y has a shift $Y \cong Y_\alpha$ which also has a unique invariant measure that satisfies the Boshernitzan condition, and hence y has well-defined factor frequencies

We can also show that the Boshernitzan condition is preserved by reset automata
modulo a finite prefix of the run

In the Boshernitzan case, we can then apply the Krohn-Rhodes theorem and establish that the run inherits the property of well-defined factor frequencies