

# *Infinite Automata 2025/26*

## Exercise Sheet 1

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**Exercise 1.1.** Given multisets  $A$  and  $B$  each comprising of at least  $n^2$  positive integers not greater than  $n$ , show that there exist subsets  $A' \subseteq A$  and  $B' \subseteq B$  such that  $A', B' \neq \emptyset$  and

$$\sum_{a \in A'} a = \sum_{b \in B'} b.$$

**Exercise 1.2.** (Strengthening of Exercise 1.1.) Given multisets  $A$  and  $B$  each comprising of at least  $2n$  positive integers not greater than  $n$ , show that there exist subsets  $A' \subseteq A$  and  $B' \subseteq B$  such that  $A', B' \neq \emptyset$  and

$$\sum_{a \in A'} a = \sum_{b \in B'} b.$$

**Exercise 1.3.** (Generalisation of Exercise 1.2.) Given a multiset  $A$  of  $n$  positive integers not greater than  $m$  and a multiset  $B$  of  $m$  positive integers not greater than  $n$ , show that there exist subsets  $A' \subseteq A$  and  $B' \subseteq B$  such that  $A', B' \neq \emptyset$  and

$$\sum_{a \in A'} a = \sum_{b \in B'} b.$$

**Exercise 1.4.** Show that the reachability problem for pushdown automata is decidable in polynomial time. We shall assume that the automaton starts in a distinguished initial state and with an empty stack and we ask whether it can reach a distinguished final state also with an empty stack.

*Hint 1. Compute an equivalent context-free grammar.*

*Hint 2. For each pair of states  $(p, q)$ , let  $X_{p,q}$  be a nonterminal deriving the words which can be accepted starting from state  $p$  with an empty stack and ending in state  $q$  with an empty stack.*

**Exercise 1.5.** Prove that reachability problem for two-stack pushdown automata is undecidable.

*Hint. Reduce from the reachability problem (the halting problem) for Turing machines with one tape.*