

# Complexity Theory - Tutorial

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## Part 1

1. Decide whether the following statements are true or false and justify your decision.
  - (a) For every class  $C$  and  $A \in C$ ; if  $A$  is  $C$ -complete under Karp reductions, then  $C \subset P^A$
  - (b) If  $L$  is NP-complete and  $L'$  is coNP complete, then  $L \cap L'$  is  $NP \cap coNP$  complete.
  - (c) If  $PH=PSPACE$ , then  $PH$  collapses to some level  $\Sigma_k$
  - (d) There is an undecidable language in  $P/poly$
  - (e) If  $P=NP$ , then  $NP=coNP$ .
  - (f) If  $\#L$  is  $\#P$ -complete, then  $L$  is NP-complete (you may assume  $P \neq NP$ ).
  - (g) Show that every circuit with only  $\wedge$  and  $\vee$  gates can be replaced by a circuit containing only *majority* gates.
2.  $DOUBLE-SAT = \{\phi : \phi \text{ is a Boolean formula with 2 satisfying assignments}\}$ . Show that  $DOUBLE-SAT$  is NP-complete.
3. Show that  $UPATH \in RL$ . You may use the following fact: if graph  $G$  has a path from  $s$  to  $t$ , then a random walk of length  $8 \cdot |V(G)| \cdot |E(G)|$  visits  $t$  with probability  $\geq \frac{1}{2}$
4. Show that  $SPACE(f(n)) \subset RSPACE(f(n)) \subset NSPACE(f(n))$

## Part 2

5. We showed that there is an undecidable language in  $P/poly$ . Show that there are decidable languages in  $P/poly$  that are not in  $P$ . Use the following waypoints:
  - (a) show that there is a decidable language  $L$  that is not in  $EXP$

- (b) define  $L' := \{1^m : m \in L\}$  and show that  $L'$  is decidable and member of P/poly
- (c) show that  $L'$  is not in P
6. Suppose that we have a poly-time algorithm  $A$  such that for  $\phi \in \text{USAT}$ ,  $A(\phi) = 1$  and for  $\phi \notin \text{SAT}$ ,  $A(\phi) = 0$ . Show that then  $\text{NP} = \text{RP}$ . Afterwards, show that  $\text{AM}[2] = \text{BPP}$ .
7. Show that regular languages are in  $\text{NC}^1$ . (Hint: note that the final state can be reached from the initial one in  $n$  steps if there is some intermediate state  $q$  that can be reached from the initial one in  $\frac{n}{2}$  steps and the final state can be reached from  $q$  also in  $\frac{n}{2}$  steps. Repeat the same idea to the leaves.)

**NOTE:** Some of the problems and solutions are taken from

- <http://soltys.cs.csuci.edu/blog/wp-content/oldpage/cu-f07/chp5.pdf>
- <http://cse.iitkgp.ac.in/~abhij/course/theory/CC/Spring04/soln5.pdf>