

Types for Information Flow Control: Labeling Granularity and Semantic Models (Technical appendix)

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1 Part I: Details of what's in the paper

1.1 Fine-grained IFC enforcement (FG)

1.1.1 FG type system

Syntax, types, constraints:

Expressions	e	::=	$x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, x.e) \mid \text{new } e \mid !e \mid e := e$
Labels	ℓ, pc	::=	$\perp \mid \top \mid l \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
(Labeled) Types	τ	::=	A^ℓ
Unlabeled types	A	::=	$\mathbf{b} \mid \text{unit} \mid \tau \xrightarrow{\ell_e} \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \tau$

Type system: $\boxed{\Gamma \vdash_{pc} e : \tau}$

$$\begin{array}{c}
\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var} \qquad \frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2}{\Gamma \vdash_{pc} \lambda x.e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp} \text{FG-lam} \\
\\
\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau_1 \quad \mathcal{L} \vdash \tau_2 \searrow \ell \quad \mathcal{L} \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Gamma \vdash_{pc} e_1 e_2 : \tau_2} \text{FG-app} \\
\\
\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \quad \Gamma \vdash_{pc} e_2 : \tau_2}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp} \text{FG-prod} \\
\\
\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e) : \tau_1} \text{FG-fst} \qquad \frac{\Gamma \vdash_{pc} e : \tau_1}{\Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp} \text{FG-inl} \\
\\
\frac{\Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau} \text{FG-case} \\
\\
\frac{\Gamma \vdash_{pc'} e : \tau' \quad \mathcal{L} \vdash pc \sqsubseteq pc' \quad \mathcal{L} \vdash \tau' <: \tau}{\Gamma \vdash_{pc} e : \tau} \text{FG-sub} \qquad \frac{\Gamma \vdash_{pc} e : \tau \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^\perp} \text{FG-ref} \\
\\
\frac{\Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e : \tau'} \text{FG-deref} \\
\\
\frac{\Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau \quad \mathcal{L} \vdash \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_1 := e_2 : \text{unit}} \text{FG-assign} \qquad \frac{}{\Gamma \vdash_{pc} () : \text{unit}^\perp} \text{FG-unitI}
\end{array}$$

Figure 1: Type system for FG

Lemma 1.1 (Reflexivity of subtyping). *The following hold:*

1. For all τ : $\mathcal{L} \vdash \tau <: \tau$
2. For all A : $\mathcal{L} \vdash A <: A$

$$\begin{array}{c}
\frac{\mathcal{L} \vdash \ell \sqsubseteq \ell' \quad \mathcal{L} \vdash A <: A'}{\mathcal{L} \vdash A^\ell <: A'^{\ell'}} \text{FGsub-label} \qquad \frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base} \\
\\
\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref} \qquad \frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod} \\
\\
\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum} \\
\\
\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2 \quad \mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell'_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow} \qquad \frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}
\end{array}$$

Figure 2: FG subtyping.

Proof. Proof by simultaneous induction on τ and A .

Proof of statement (1)

Let $\tau = A^\ell$. Then, we have:

$$\frac{\frac{}{\mathcal{L} \vdash A <: A} \text{IH(2)} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash A^\ell <: A'^{\ell'}} \text{FGsub-label}$$

Proof of statement (2)

We proceed by cases on A .

1. $A = \mathbf{b}$:

$$\frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

2. $A = \text{ref } \tau$:

$$\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

3. $A = \tau_1 \times \tau_2$:

$$\frac{\frac{}{\mathcal{L} \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\mathcal{L} \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau_1 \times \tau_2}$$

4. $A = \tau_1 + \tau_2$:

$$\frac{\frac{}{\mathcal{L} \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\mathcal{L} \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau_1 + \tau_2}$$

5. $A = \tau_1 \xrightarrow{\ell_e} \tau_2$:

$$\frac{\frac{}{\mathcal{L} \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\mathcal{L} \vdash \tau_2 <: \tau_2} \text{IH(2) on } \tau_2 \quad \frac{}{\mathcal{L} \vdash \ell_e \sqsubseteq \ell_e}}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau_1 \xrightarrow{\ell_e} \tau_2}}$$

6. $A = \text{unit}$:

$$\frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}}$$

□

1.1.2 FG semantics

Judgement: $(H, e) \Downarrow_i (H', v)$

The semantics are described in Figure 3

$$\begin{array}{c} \frac{(H, e_1) \Downarrow_i (H', \lambda x. e_i) \quad (H', e_2) \Downarrow_j (H'', v_2) \quad (H'', e_i[v_2/x]) \Downarrow_k (H''', v_3)}{(H, e_1 e_2) \Downarrow_{i+j+k+1} (H''', v_3)} \text{fg-app} \\ \\ \frac{(H, e_1) \Downarrow_i (H', v_1) \quad (H', e_2) \Downarrow_j (H'', v_2)}{(H, (e_1, e_2)) \Downarrow_{i+j+1} (H'', (v_1, v_2))} \text{fg-prod} \quad \frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{fst}(e)) \Downarrow_{i+1} (H', v_1)} \text{fg-fst} \\ \\ \frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{snd}(e)) \Downarrow_{i+1} (H', v_2)} \text{fg-snd} \quad \frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inl}(e)) \Downarrow_{i+1} (H', \text{inl}(v))} \text{fg-inl} \\ \\ \frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inr}(e)) \Downarrow_{i+1} (H', \text{inr}(v))} \text{fg-inr} \quad \frac{(H, e) \Downarrow_i (H', \text{inl } v) \quad (H', e_1[v/x]) \Downarrow_j (H'', v_1)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_1)} \text{fg-case1} \\ \\ \frac{(H, e) \Downarrow_i (H', \text{inr } v) \quad (H', e_2[v/x]) \Downarrow_j (H'', v_2)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_2)} \text{fg-case2} \\ \\ \frac{(H, e) \Downarrow_i (H', v) \quad a \notin \text{dom}(H)}{(H, \text{new}(e)) \Downarrow_{i+1} (H'[a \mapsto v], a)} \text{fg-ref} \quad \frac{(H, e) \Downarrow_i (H', a)}{(H, !e) \Downarrow_{i+1} (H', H(a))} \text{fg-deref} \\ \\ \frac{(H, e_1) \Downarrow_i (H', a) \quad (H', e_2) \Downarrow_j (H'', v)}{(H, e_1 := e_2) \Downarrow_{i+j+1} (H''[a \mapsto v], ())} \text{fg-assign} \quad \frac{e \in \{x, \lambda y. -\}}{(H, e) \Downarrow_0 (H, e)} \text{fg-val} \end{array}$$

Figure 3: FG semantics

1.1.3 Logical relation for FG

$W : ((\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \leftrightarrow \text{Loc}))$

Definition 1.2 (θ_2 extends θ_1). $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

Definition 1.3 (W_2 extends W_1). $W_1 \sqsubseteq W_2 \triangleq$

1. $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2. $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

Definition 1.4 (Binary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^A &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V^A &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^A &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A\} \\
[\tau_1 + \tau_2]_V^A &\triangleq \{(W, n, \mathbf{inl} \ v, \mathbf{inl} \ v') \mid (W, n, v, v') \in [\tau_1]_V^A\} \cup \\
&\quad \{(W, n, \mathbf{inr} \ v, \mathbf{inr} \ v') \mid (W, n, v, v') \in [\tau_2]_V^A\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\
&\quad ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e}) \wedge \\
&\quad \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})\} \\
[\mathbf{ref} \ \tau]_V^A &\triangleq \{(W, n, a_1, a_2) \mid \\
&\quad (a_1, a_2) \in W.\hat{\beta} \wedge W.\theta_1(a_1) = W.\theta_2(a_2) = \tau\}
\end{aligned}$$

$$[\mathbf{A}^{\ell'}]_V^A \triangleq \begin{cases} \{(W, n, v_1, v_2) \mid (W, n, v_1, v_2) \in [\mathbf{A}]_V^A\} & \ell' \sqsubseteq \mathbf{A} \\ \{(W, n, v_1, v_2) \mid \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\mathbf{A}]_V\} & \ell' \not\sqsubseteq \mathbf{A} \end{cases}$$

Definition 1.5 (Binary expression relation).

$$\begin{aligned}
[\tau]_E^A &\triangleq \{(W, n, e_1, e_2) \mid \\
&\quad \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge \\
&\quad (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\
&\quad \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_V^A\}
\end{aligned}$$

Definition 1.6 (Unary value relation).

$$\begin{aligned}
[\mathbf{b}]_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V &\triangleq \{(\theta, m, v \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in [\tau_1]_V \wedge (\theta, m, v_2) \in [\tau_2]_V\} \\
[\tau_1 + \tau_2]_V &\triangleq \{(\theta, m, \mathbf{inl} \ v) \mid (\theta, m, v) \in [\tau_1]_V\} \cup \{(\theta, m, \mathbf{inr} \ v) \mid (\theta, m, v) \in [\tau_2]_V\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies \\
&\quad (\theta', j, e[v/x]) \in [\tau_2]_E^{\ell_e}\} \\
[\mathbf{ref} \ \tau]_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \tau\}
\end{aligned}$$

$$[\mathbf{A}^{\ell'}]_V \triangleq [\mathbf{A}]_V$$

Definition 1.7 (Unary expression relation).

$$\begin{aligned}
[\tau]_E^{pc} &\triangleq \{(\theta, n, e) \mid \forall H. (n, H) \triangleright \theta \wedge \forall j < n. (H, e) \Downarrow_j (H', v') \implies \\
&\quad \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in [\tau]_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)\}
\end{aligned}$$

Definition 1.8 (Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n-1, H(a)) \in \lfloor \theta(a) \rfloor_V$$

Definition 1.9 (Binary heap well formedness).

$$\begin{aligned} (n, H_1, H_2) \triangleright^A W \triangleq & \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ & (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ & \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ & (W, n-1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ & \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

Definition 1.10 (Binary substitution). $\gamma : \text{Var} \mapsto (\text{Val}, \text{Val})$

Definition 1.11 (Unary substitution). $\delta : \text{Var} \mapsto \text{Val}$

Definition 1.12 (Unary interpretation of Γ).

$$\lfloor \Gamma \rfloor_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in \lfloor \Gamma(x) \rfloor_V\}$$

Definition 1.13 (Binary interpretation of Γ).

$$\lfloor \Gamma \rfloor_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A\}$$

1.1.4 Soundness proof for FG

Lemma 1.14 (Binary value relation subsumes unary value relation). $\forall W, v_1, v_2, \mathcal{A}, n.$

The following holds:

1. $\forall \mathcal{A}.$

$$(W, n, v_1, v_2) \in \lceil \mathcal{A} \rceil_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \lfloor \mathcal{A} \rfloor_V$$

2. $\forall \tau.$

$$(W, n, v_1, v_2) \in \lceil \tau \rceil_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \lfloor \tau \rfloor_V$$

Proof. Proof by simultaneous induction on \mathcal{A} and τ

Proof of statement (1)

We analyze the various cases of \mathcal{A} in the last step:

1. Case **b**, unit:

From Definition 1.6

2. Case $\tau_1 \times \tau_2$:

$$\underline{\text{Given:}} (W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in \lceil \tau_1 \times \tau_2 \rceil_V^A$$

To prove:

$$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in \lfloor \tau_1 \times \tau_2 \rfloor_V \quad (\text{P01})$$

and

$$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in \lfloor \tau_1 \times \tau_2 \rfloor_V \quad (\text{P02})$$

From Definition 1.4 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A \quad (\text{P1})$$

IH1a: $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$ and

IH1b: $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a: $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$ and

IH2b: $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some m we need to prove

$$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$$

Similarly from (P02) we know that given some m we need to prove

$$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$$

We instantiate IH1a and IH2a with the given m from (P01) to get

$$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V \text{ and } (W.\theta_1, m, v_{i2}) \in [\tau_2]_V$$

Then from Definition 1.6, we get

$$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$$

Similarly we instantiate IH1b and IH2b with the given m from (P02) to get

$$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V \text{ and } (W.\theta_2, m, v_{j2}) \in [\tau_2]_V$$

Then from Definition 1.6, we get

$$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$$

3. Case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v_1 = \text{inl}(v_{i1})$ and $v_2 = \text{inl}(v_{j1})$

Given: $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V \quad (\text{S01})$$

and

$$\forall m. (W.\theta_2, m, \text{inl}(v_{i2})) \in [\tau_1 + \tau_2]_V \quad (\text{S02})$$

From Definition 1.4 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \quad (\text{S0})$$

IH1: $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$ and

IH2: $\forall m_2. (W.\theta_2, m_2, v_{j1}) \in [\tau_1]_V$

From (S01) we know that given some m and we are required to prove:

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

Also from (S02) we know that given some m and we are required to prove:

$$(W.\theta_2, m, \text{inl}(v_{i2})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH1 with m from (S01) to get

$$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$$

Therefore from Definition 1.6, we get

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH2 with m from (S02) to get

$$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$$

Therefore from Definition 1.6, we get

$$(W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$$

(b) $v_1 = \text{inr}(v_{i2})$ and $v_2 = \text{inr}(v_{j2})$

Symmetric case as (a)

4. Case $\tau_1 \xrightarrow{\ell_e} \tau_2$:

$$\text{Given: } (W, n, \lambda x.e_1, \lambda x.e_2) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$$

This means from Definition 1.4 we know that

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A &\implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in [\tau_1]_V &\implies (\theta_l, i, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e}) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V &\implies (\theta_l, k, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e}) \quad (\text{L0}) \end{aligned}$$

To prove:

(a) $\forall m. (W.\theta_1, m, \lambda x.e_1) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$:

This means from Definition 1.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies (\theta', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e}$$

This further means that we have some θ', j and v s.t

$$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in [\tau_1]_V$$

$$\text{And we need to prove: } (\theta', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e}$$

Instantiating θ_l, i and v_c in the second conjunct of L0 with θ', j and v respectively and since we know that $W.\theta_1 \sqsubseteq \theta'$ and $(\theta', j, v) \in [\tau_1]_V$

$$\text{Therefore we get } (\theta', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e}$$

(b) $\forall m. (W.\theta_2, m, \lambda x.e_2) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$:

Similar reasoning with e_2

5. Case ref τ :

From Definition 1.4 and 1.6

Proof of statement (2)

Let $\tau = A^\ell$

2 cases arise:

1. $\ell \sqsubseteq \mathcal{A}$:

From IH (statement(1))

2. $\ell \not\sqsubseteq \mathcal{A}$:

Directly from Definition 1.4

□

Lemma 1.15 (Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m'.$

1. $\forall A. (\theta, m, v) \in [A]_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in [A]_V$
2. $\forall \tau. (\theta, m, v) \in [\tau]_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in [\tau]_V$

Proof. Proof by simultaneous induction on A and τ

Proof of statement (1)

We analyze the various cases of A in the last step:

1. case **b, unit**:

Directly from Definition 1.6

2. case $\tau_1 \times \tau_2$:

Given: $(\theta, m, (v_1, v_2)) \in [\tau_1 \times \tau_2]_V$

To prove: $(\theta', m', (v_1, v_2)) \in [\tau_1 \times \tau_2]_V$

This means from Definition 1.6 we know that

$(\theta, m, v_1) \in [\tau_1]_V \wedge (\theta, m, v_2) \in [\tau_2]_V$

IH1 : $(\theta', m', v_1) \in [\tau_1]_V$

IH2 : $(\theta', m', v_2) \in [\tau_2]_V$

We get the desired from IH1, IH2 and Definition 1.6

3. case $\tau_1 + \tau_2$:

2 cases arise:

- (a) $v = \text{inl}(v_1)$:

Given: $(\theta, m, (\text{inl } v_1)) \in [\tau_1 + \tau_2]_V$

To prove: $(\theta', m', \text{inl } v_1) \in [\tau_1 + \tau_2]_V$

This means from Definition 1.6 we know that

$(\theta, m, v_1) \in [\tau_1]_V$

IH : $(\theta', m', v_1) \in [\tau_1]_V$

Therefore from IH and Definition 1.6 we get the desired

- (b) $v = \text{inr}(v_2)$

Symmetric case

4. case $\tau_1 \xrightarrow{\ell_e} \tau_2$:

Given: $(\theta, m, (\lambda x. e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

To prove: $(\theta', m', (\lambda x. e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

This means from Definition 1.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in [\tau_1]_V \implies (\theta'', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e} \quad (1)$$

Similarly from Definition 1.6 we know that we are required to prove

$$\forall \theta''' . \theta' \sqsubseteq \theta''' \wedge \forall k < m' . \forall v_1 . (\theta''', k, v_1) \in [\tau_1]_V \implies (\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$$

This means that given some θ''', k and v_1 such that $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in [\tau_1]_V$

And we are required to prove $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating Equation 57 with θ''', k and v_1 and since we know that $\theta' \sqsubseteq \theta'''$ and $\theta \sqsubseteq \theta'$ therefore we have $\theta \sqsubseteq \theta'''$. Also, we know that $k < m' < m$ and $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

5. case ref τ :

From Definition 1.6 and Definition 1.2

Proof of statement (2)

Let $\tau = A^\ell$

Since $[A^\ell]_V = [A]_V$, therefore from IH (statement 1) □

Lemma 1.16 (Monotonicity binary). *The following holds:*

$\forall W, W', v_1, v_2, A, n, n'$.

1. $\forall A$. $(W, n, v_1, v_2) \in [A]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [A]_V^A$
2. $\forall \tau$. $(W, n, v_1, v_2) \in [\tau]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [\tau]_V^A$

Proof. Proof by simultaneous induction on A and τ

Proof of statement (1)

We analyze the different cases of A in the last step:

1. Case **b**, unit:

From Definition 1.4

2. Case $\tau_1 \times \tau_2$:

Given: $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove: $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

From Definition 1.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$

IH1 : $(W', n', v_{i1}, v_{j1}) \in [\tau_1]_V^A$

IH2 : $(W', n', v_{i2}, v_{j2}) \in [\tau_2]_V^A$

From IH1, IH2 and Definition 1.4 we get the desired.

3. Case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v_1 = \text{inl } v_{i1}$ and $v_2 = \text{inl } v_{i2}$:

Given: $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

To prove: $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

From Definition 1.4 we know that we are given

$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

IH : $(W', n', v_{i1}, v_{i2}) \in [\tau_1]_V^A$

Therefore from Definition 1.4 we get

$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_V^A$

(b) $v_1 = \text{inr}(v_{12})$ and $v_2 = \text{inr}(v_{22})$:

Symmetric case

4. Case $\tau_1 \xrightarrow{\ell_e} \tau_2$:

Given: $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

To prove: $(\theta', n', (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

This means from Definition 1.4 we know that the following holds

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A)$
(BM-A0)

$\forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e})$ (BM-A1)

$\forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})$ (BM-A2)

Similarly from Definition 1.4 we know that we are required to prove

(a) $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$:

This means that we are given some $W'' \sqsupseteq W', k < n'$ and v'_1, v'_2 s.t

$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

And we are required to prove: $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with W'', k and v'_1, v'_2 we get

$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

(b) $\forall \theta'_l \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e})$:

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_1, k$ and v'_c s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we are required to prove: $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A1 with θ'_l, k and v'_c we get

$(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

(c) $\forall \theta'_l \sqsupseteq W'.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e})$:

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_2, k$ and v'_c s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we are required to prove: $(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A2 with θ'_l, k and v'_c we get

$(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

5. Case ref τ :

From Definition 1.4 and Definition 1.3

Proof of statement (2)

Let $\tau = \mathbf{A}^\ell$

2 cases arise:

1. $\ell \sqsubseteq \mathcal{A}$:

From IH (statement 1)

2. $\ell \not\sqsubseteq \mathcal{A}$:

From Lemma 1.15 and Definition 1.4

□

Lemma 1.17 (Unary monotonicity for Γ). $\forall \theta, \theta', \delta, \Gamma, n, n'$.
 $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in [\Gamma]_V$

Proof. Given: $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$

To prove: $(\theta', n', \delta) \in [\Gamma]_V$

From Definition 1.12 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V$$

And again from Definition 1.12 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$$

• $\text{dom}(\Gamma) \subseteq \text{dom}(\delta)$:

Given

• $\forall x \in \text{dom}(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$:

Since we know that $\forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V$ (given)

Therefore from Lemma 1.15 we get

$$\forall x \in \text{dom}(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$$

□

Lemma 1.18 (Binary monotonicity for Γ). $\forall W, W', \delta, \Gamma, n, n'$.

$$(W, n, \gamma) \in [\Gamma]_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in [\Gamma]_V$$

Proof. Given: $(W, n, \gamma) \in [\Gamma]_V \wedge n' < n \wedge W \sqsubseteq W'$

To prove: $(W', n', \gamma) \in [\Gamma]_V$

From Definition 1.13 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And again from Definition 1.12 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

- $dom(\Gamma) \subseteq dom(\gamma)$:

Given

- $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$:

Since we know that $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$ (given)

Therefore from Lemma 1.16 we get

$$\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

□

Lemma 1.19 (Unary monotonicity for H). $\forall \theta, H, n, n'$.

$$(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$$

Proof. Given: $(n, H) \triangleright \theta \wedge n' < n$

To prove: $(n', H) \triangleright \theta$

From Definition 1.8 it is given that

$$dom(\theta) \subseteq dom(H) \wedge \forall a \in dom(\theta). (\theta, n - 1, H(a)) \in [\theta(a)]_V$$

And again from Definition 1.12 we are required to prove that

$$dom(\theta) \subseteq dom(H) \wedge \forall a \in dom(\theta). (\theta, n' - 1, H(a)) \in [\theta'(a)]_V$$

- $dom(\theta) \subseteq dom(H)$:

Given

- $\forall a \in dom(\theta). (\theta, n' - 1, H(a)) \in [\theta'(a)]_V$:

Since we know that $\forall a \in dom(\theta). (\theta, n - 1, H(a)) \in [\theta(a)]_V$ (given)

Therefore from Lemma 1.15 we get

$$\forall a \in dom(\theta). (\theta, n' - 1, H(a)) \in [\theta'(a)]_V$$

□

Lemma 1.20 (Binary monotonicity for heaps). $\forall W, H_1, H_2, n, n'$.

$$(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$$

Proof. Given: $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$

To prove: $(n', H_1, H_2) \triangleright W$

From Definition 1.9 it is given that

$$dom(W.\theta_1) \subseteq dom(H_1) \wedge dom(W.\theta_2) \subseteq dom(H_2) \wedge$$

$$(W.\hat{\beta}) \subseteq (dom(W.\theta_1) \times dom(W.\theta_2)) \wedge$$

$$\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge$$

$$(W, n - 1, H_1(a_1), H_2(a_2)) \in [W.\theta_1(a_1)]_V^A \wedge$$

$$\forall i \in \{1, 2\}. \forall m. \forall a_i \in dom(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in [W.\theta_i(a_i)]_V$$

And again from Definition 1.9 we are required to prove:

- $dom(W.\theta_1) \subseteq dom(H_1) \wedge dom(W.\theta_2) \subseteq dom(H_2)$:

Given

- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$:
Given
- $\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2) \text{ and } (W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A)$:
 $\forall (a_1, a_2) \in (W.\hat{\beta})$.
 - $(W.\theta_1(a_1) = W.\theta_2(a_2))$: Given
 - $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$:
Given and from Lemma 1.16
- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$:
Given

□

Theorem 1.21 (Fundamental theorem unary). $\forall \Gamma, pc, \theta, e, \tau, \delta, n$.

$$\begin{aligned} & \Gamma \vdash_{pc} e : \tau \wedge \\ & (\theta, n, \delta) \in \lfloor \Gamma \rfloor_V \implies \\ & (\theta, n, e \delta) \in \lfloor \tau \rfloor_E^{pc} \end{aligned}$$

Proof. Proof by induction on FG typing derivation

1. FG-var:

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

To prove: $(\theta, n, x \delta) \in \lfloor \tau \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall j < n. (H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap H and $j < n$ s.t $(n, H) \triangleright \theta \wedge (H, x \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-V0}) \end{aligned}$$

In order to prove FU-V0 we instantiate θ' with θ . From reduction relation we know that $H' = H$, $v' = \delta(x)$ and $j = 1$

We need to prove the following:

- (a) $\theta \sqsubseteq \theta \wedge (n - 1, H) \triangleright \theta \wedge (\theta, n - 1, v') \in \lfloor \tau \rfloor_V$:
 - $\theta \sqsubseteq \theta$: From Definition 1.2
 - $(n - 1, H) \triangleright \theta$: From Lemma 1.19

- $(\theta, n-1, v') \in \lfloor \tau \rfloor_V$:
 Since we are given that $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$ and $v' = \delta(x)$
 Therefore $(\theta, n, v') \in \lfloor \Gamma(x) \rfloor_V$, where $\Gamma(x) = \tau$
 And finally from Lemma 1.15 we get $(\theta, n-1, v') \in \lfloor \tau \rfloor_V$
- (b) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$:
 Since $H' = H$, so we are done
- (c) $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$:
 Since $\theta' = \theta$, so we are done

2. FG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove: $(\theta, \lambda x. e_i \delta) \in \lfloor ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall j < n. (H, (\lambda x. e_i) \delta) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H') \triangleright \theta' \wedge (\theta', n-j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap H and $j < n$ s.t $(n, H) \triangleright \theta \wedge (H, (\lambda x. e_i) \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H') \triangleright \theta' \wedge (\theta', n-j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-L0}) \end{aligned}$$

IH1:

$$\forall \theta_i, v_x, n. (\theta_i, n, e_i \delta \cup \{x \mapsto v_x\}) \in \lfloor \tau_2 \rfloor_E^{\ell_e}, \text{ s.t } (\theta_i, n, v_x) \in \lfloor \tau_1 \rfloor_V$$

In order to prove FU-L0 we instantiate θ' with θ . From reduction relation we know that $H' = H$, $j = 0$ and $v' = \lambda x. e_i \delta$

$$(a) \theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in \lfloor ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \rfloor_V:$$

- $\theta \sqsubseteq \theta$: From Definition 1.2
- $(n, H) \triangleright \theta$: Given
- $(\theta, n, (\lambda x. e_i) \delta) \in \lfloor ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \rfloor_V$:

From Definition 1.6 it suffices to prove that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < n. \forall v. (\theta'', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta'', j, e_i[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

This means given some θ'' , j and v such that $\theta \sqsubseteq \theta''$, $j < n$ and $(\theta'', j, v) \in \lfloor \tau_1 \rfloor_V$

It suffices to prove that $(\theta'', j, e_i[v/x] \delta) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$

Since $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$ and $j < n$ therefore from Lemma 1.17 we have

$$(\theta, j, \delta) \in \lfloor \Gamma \rfloor_V$$

So we can apply IH1 instantiated with θ'' , v and j we get

$$(\theta'', j, e_i[v/x] \delta) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

(b) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$:

Since $H' = H$ so we are done

(c) $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$:

Since $\theta' = \theta$ so we are done

3. FG-app:

$$\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau_1 \quad \mathcal{L} \vdash \tau_2 \searrow \ell \quad \mathcal{L} \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove: $(\theta, n, (e_1 e_2) \delta) \in [\tau_2]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap H s.t. $(n, H) \triangleright \theta \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-P0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall n_1, H_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with n, H and since we know that $(n, H) \triangleright \theta \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-P1}) \end{aligned}$$

From evaluation rule we know that $v'_1 = \lambda x. e_i$. Since from FU-P1 we know that

$$(\theta'_1, n - i, \lambda x. e_i) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V$$

This means from Definition 1.6 we have

$$\forall \theta''. \theta'_1 \sqsubseteq \theta'' \wedge \forall j < (n - i). \forall v. (\theta'', j, v) \in [\tau_1]_V \implies (\theta'', j, e_i[v/x]) \in [\tau_2]_E^{\ell_e} \quad (2)$$

IH2:

$$\begin{aligned} & \forall n_2, \forall H_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall k < n_2. (H_2, (e_2) \delta) \Downarrow_k (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - k, v'_2) \in [(\tau_1)]_V \wedge \end{aligned}$$

$$\begin{aligned}
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc)
\end{aligned}$$

Instantiating IH2 with $n-i$, H'_1 and since we know that $(n-i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1 \ e_2) \ \delta) \Downarrow_{n'}$ (H', v') therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n-i-k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n-i-k, v'_2) \in [(\tau_1)]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \quad (\text{FU-P2})
\end{aligned}$$

Instantiating θ'' , j and v in Equation 2 with θ'_2 , $n-i-k$ and v'_2 from FU-P2 respectively, we get

$$(\theta'_2, n-i-k, e_i[v'_2/x]) \in [\tau_2]_E^{\ell_e}$$

This means from Definition 1.7 we have

$$\begin{aligned}
& \forall H_3. (n-i-k, H_3) \triangleright \theta'_2 \wedge \forall l < (n-i-k). (H_3, e_i[v'_2/x]) \Downarrow_l (H'_3, v'_3) \implies \\
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n-i-k-l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e)
\end{aligned}$$

Instantiating H_3 with H'_2 from FU-P2 and since we know that $((n-i-k), H'_2) \triangleright \theta'_2$ and since the reduction happens therefore we have

$$\begin{aligned}
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n-i-k-l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e) \quad (\text{FU-P3})
\end{aligned}$$

In order to prove FU-P0 we choose θ' as θ'_3 from FU-P3. Also we know that $H' = H'_3$, $v' = v'_3$ and $n' = i+k+l$. Now we are required to show

$$(a) \ \theta \sqsubseteq \theta'_3 \wedge ((n-i-k-l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V:$$

- $\theta \sqsubseteq \theta'_3$:
Since $\theta \sqsubseteq \theta'_1$ from FU-P1, $\theta'_1 \sqsubseteq \theta'_2$ from FU-P2 and $\theta'_2 \sqsubseteq \theta'_3$ from FU-P3 therefore from Definition 1.2 we get $\theta \sqsubseteq \theta'_3$
- $((n-i-k-l), H'_3) \triangleright \theta'_3$:
From FU-P3 we get $((n-i-k-l), H'_3) \triangleright \theta'_3$
- $(\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V$:
From FU-P3 we get $(\theta'_3, (n-i-k-l), v'_3) \in [\tau_2]_V$

$$(b) \ (\forall a \in \text{dom}(H). H(a) \neq H'_3(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

Since $pc \sqsubseteq \ell_e$ therefore we get the desired from FU-P1, FU-P2 and FU-P3

$$(c) \ (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta). \theta'_3(a) \searrow pc)$$

Since $pc \sqsubseteq \ell_e$ therefore we get the desired from FU-P1, FU-P2 and FU-P3

4. FG-prod:

$$\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \quad \Gamma \vdash_{pc} e_2 : \tau_2}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove: $(\theta, n, (e_1, e_2) \ \delta) \in [(\tau_1 \times \tau_2)^\perp]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall n' < n.(H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 \times \tau_2)^\perp]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap H s.t $H \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 \times \tau_2)^\perp]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-PA0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1.(n_1, H_1) \triangleright \theta \wedge \forall i < n_1.(H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau_1]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

We instantiate IH1 with H and n . And since we know that $(n, H) \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau_1]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-PA1}) \end{aligned}$$

IH2:

$$\begin{aligned} & \forall H_2, n_2.(n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2.(H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in [(\tau_2)]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \end{aligned}$$

We instantiate IH2 with H'_1 and $n - i$. And since we know that $(n - i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau_2)]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \quad (\text{FU-PA2}) \end{aligned}$$

In order to prove FU-PA0 we choose θ' as θ'_2 from FU-PA2. Also we know from the evaluation rule, that let $v' = (v'_1, v'_2)$, $H' = H'_2$ and $n' = i + j + 1$. Now we are required to show

$$(a) \theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v') \in [(\tau_1 \times \tau_2)^\perp]_V:$$

- $\theta \sqsubseteq \theta'_2$:
Since $\theta \sqsubseteq \theta'_1$ from FU-PA1 and $\theta'_1 \sqsubseteq \theta'_2$ from FU-PA2 therefore from Definition 1.2 we get $\theta \sqsubseteq \theta'_2$
- $(n - i - j - 1, H') \triangleright \theta'_2$:
From FU-PA2 we get $(n - i - j, H'_2) \triangleright \theta'_2$ therefore from Lemma 1.19 we get $(n - i - j - 1, H'_2) \triangleright \theta'_2$

- $(\theta'_2, n - i - j, v') \in [(\tau_1 \times \tau_2)^\perp]_V$:

From Definition 1.6 it suffices to show

- i. $(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1)]_V$:

Since from FU-PA1 we know that $(\theta'_1, n - i, v'_1) \in [(\tau_1)]_V$ and since $\theta'_1 \sqsubseteq \theta'_2$

(from FU-PA2) therefore from Lemma 1.15 we get

$$(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1)]_V$$

- ii. $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2)]_V$:

From FU-PA2 we know that $(\theta'_2, n - i - j, v'_2) \in [(\tau_2)]_V$ therefore from

Lemma 1.15 we get $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2)]_V$

$$(b) \ (\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

From FU-PA1 and FU-PA2

$$(c) \ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$$

From FU-PA1 and FU-PA2

5. FG-fst:

$$\frac{\Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$

To prove: $(\theta, n, \text{fst}(e_i) \delta) \in [\tau_1]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{fst}(e_i) \delta) \Downarrow_{n'} (H', v') \implies$$

$$\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$$

This means that given some heap H s.t $(n, H) \triangleright \theta \wedge (H, \text{fst}(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-F0})$$

IH1:

$$\forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies$$

$$\exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell]_V \wedge$$

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$$

Instantiating IH1 with H and n . Since we know that $H \triangleright \theta \wedge (H, \text{fst}(e_i) \delta) \Downarrow (H', v')$ therefore we have

$$\exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell]_V \wedge$$

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-F1})$$

From evaluation rule we know that $v'_1 = (v''_1, v''_2)$

In order to prove FU-F0 we choose θ' as θ'_1 from FU-P1. Also we know that $H' = H'_1$ and $v' = v''_1$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v'_1) \in \lfloor \tau_1 \rfloor_V$:
- $\theta \sqsubseteq \theta'_1$:
From FU-F1
 - $(n - i - 1, H'_1) \triangleright \theta'_1$:
From FU-F1 we know $(n - i, H'_1) \triangleright \theta'_1$ therefore from Lemma 1.19 we get $(n - i - 1, H'_1) \triangleright \theta'_1$
 - $(\theta'_1, n - i, v''_1) \in \lfloor \tau_1 \rfloor_V$:
Since from FU-F1 we know that $(\theta'_1, n - i, (v''_1, v''_2)) \in \lfloor (\tau_1 \times \tau_2) \rfloor_V$
Therefore from Definition 1.6 we know that $(\theta'_1, n - i, v''_1) \in \lfloor \tau_1 \rfloor_V$
From Lemma 1.15 we get $(\theta'_1, n - i - 1, v''_1) \in \lfloor \tau_1 \rfloor_V$
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$
From FU-F1
- (c) $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$
From FU-F1

6. FG-snd:

Symmetric case to FG-fst

7. FG-inl:

$$\frac{\Gamma \vdash_{pc} e_i : \tau_1}{\Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove: $(\theta, n, \text{inl}(e_i) \delta) \in \lfloor (\tau_1 + \tau_2)^\perp \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor (\tau_1 + \tau_2)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor (\tau_1 + \tau_2)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-LE0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in \lfloor \tau_1 \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with H and n . Since we know that $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in \lfloor \tau_1 \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-LE1}) \end{aligned}$$

In order to prove FU-LE0 we choose θ' as θ'_1 from FU-LE1. Also we know from the evaluation rule, that let $v' = \text{inl}(v'_1)$, $H' = H'_1$ and $n' = i + 1$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2)]_V$:
- $\theta \sqsubseteq \theta'_1$:
From FU-LE1
 - $(n - i - 1, H') \triangleright \theta'_1$:
From FU-LE1 we know that $(n - i, H') \triangleright \theta'_1$ therefore from Lemma 1.19 we get $(n - i - 1, H') \triangleright \theta'_1$
 - $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2)]_V$:
Since $v' = \text{inl}(v'_1)$ and from FU-LE1 we know that $(\theta'_1, n - i, v'_1) \in [\tau_1]_V$
Therefore from Definition 1.6 we get $(\theta'_1, n - i, v') \in [(\tau_1 + \tau_2)]_V$
From Lemma 1.15 we get $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2)]_V$
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$
From FU-LE1
- (c) $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$
From FU-LE1

8. FG-inr:

Symmetric case to FG-inl

9. FG-case:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau}$$

To prove: $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in [\tau]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-C0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_c) \delta) \Downarrow_i (H'_1, v'_c) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_c) \in [(\tau_1 + \tau_2)^\ell]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with H and n . Since we know that $H \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_c) \in \llbracket (\tau_1 + \tau_2)^\ell \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-C1}) \end{aligned}$$

2 cases arise:

(a) $v'_c = \text{inl}(v_{ci})$:

IH2:

$$\begin{aligned} & \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_1) \delta \cup \{x \mapsto v_{ci}\}) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in \llbracket (\tau) \rrbracket_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

Instantiating IH2 with H'_1 and $n - i$ since we know that $H'_1 \triangleright \theta'_1 \wedge (H', (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in \llbracket (\tau) \rrbracket_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell)) \quad (\text{FU-C2}) \end{aligned}$$

In order to prove FU-C0 we choose θ' as θ'_2 from FU-C2. Also we know that $H' = H'_2$, $v' = v'_2$ and $n' = i + j + 1$. Now we are required to show

i. $\theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v'_2) \in \llbracket \tau \rrbracket_V$:

• $\theta \sqsubseteq \theta'_2$:

Since $\theta \sqsubseteq \theta'_1$ from FU-C1 and $\theta'_1 \sqsubseteq \theta'_2$ from FU-C2 therefore from Definition 1.2 we get $\theta \sqsubseteq \theta'_2$

• $(n - i - j - 1, H'_2) \triangleright \theta'_2$:

From FU-C2 we know that $(n - i - j, H'_2) \triangleright \theta'_2$ therefore from Lemma 1.19 we get $(n - i - j - 1, H'_2) \triangleright \theta'_2$

• $(\theta'_2, n - i - j - 1, v'_2) \in \llbracket \tau \rrbracket_V$:

From FU-C2 we know that $(\theta'_2, n - i - j, v'_2) \in \llbracket \tau \rrbracket_V$ therefore from Lemma 1.15 we get $(\theta'_2, n - i - j - 1, v'_2) \in \llbracket \tau \rrbracket_V$

ii. $(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$:

Since from FU-C2 we know that

$$(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell')$$

therefore we also have

$$(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc) \sqsubseteq \ell')$$

and from FU-C1 we know that

$$(\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge (pc) \sqsubseteq \ell')$$

Combining the two we get

$$(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

iii. $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$:

Since from FU-C2 we know that

$$(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell))$$

therefore we also have

$$(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc))$$

and from FU-C1 we know that
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow (pc \sqcup \ell))$

Combining the two we get
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$

- (b) $v'_c = \text{inr}(v_{ci})$:
 Symmetric case as $v'_c = \text{inl}(v_{ci})$

10. FG-ref:

$$\frac{\Gamma \vdash_{pc} e_i : \tau \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove: $(\theta, n, \text{new } (e_i) \delta) \in [(\text{ref } \tau)^\perp]_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\text{ref } \tau)^\perp]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\text{ref } \tau)^\perp]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-R0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with H and n . Since we know that $(n, H) \triangleright \theta \wedge (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-R1}) \end{aligned}$$

From the evaluation rule we know that $H' = H'_1[a \mapsto v'_1]$ where $a \notin \text{dom}(H'_1)$, $v' = a$ and $n' = i + 1$. In order to prove FU-R0 we choose θ' as $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau\})$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_2 \wedge (n - i - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - 1, v') \in [(\text{ref } \tau)^\perp]_V$:

- $\theta \sqsubseteq \theta'_2$:

From FU-R1 we know that $\theta \sqsubseteq \theta'_1$ therefore from Definition 1.2 $\theta \sqsubseteq \theta'_2$

- $(n - i - 1, H') \triangleright \theta'_2$:
 From FU-R1 we know that $(n - i, H'_1) \triangleright \theta'_1$. Therefore from Lemma 1.19 we get $(n - i - 1, H'_1) \triangleright \theta'_1$.
 We also know that $(\theta'_1, n - i, v'_1) \in \lfloor \tau \rfloor_V$ (from FU-R1) therefore from Lemma 1.15 we get $(\theta'_1, n - i - 1, v'_1) \in \lfloor \tau \rfloor_V$
 Since $H' = H'_1[a \mapsto v'_1]$ and $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau\})$ therefore from Definition 1.8 we get $(n - i - 1, H') \triangleright \theta'_2$
 - $(\theta'_2, n - i - 1, a) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V$:
 Since $\theta'_2(a) = \tau$ therefore from Definition 1.6 we get $(\theta'_2, n - i - 1, a) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V$
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$
 From FU-R1
- (c) $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$:
 We get this from FU-R1 and $\tau \searrow pc$ (given)

11. FG-deref:

$$\frac{\Gamma \vdash_{pc} e_i : (\text{ref } \tau)^\ell \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e_i : \tau'}$$

To prove: $(\theta, n, (!e_i) \delta) \in \lfloor \tau' \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (!e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor \tau' \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, (!e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor \tau' \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-D0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in \lfloor ((\text{ref } \tau))^\ell \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with H and n . Since we know that $(n, H) \triangleright \theta \wedge (H, (!e_i) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in \lfloor ((\text{ref } \tau))^\ell \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-D1}) \end{aligned}$$

In order to prove FU-D0 we choose θ' as θ'_1 from FU-D1. Also we know from the evaluation rule, that $H' = H'_1$, $v' = H'_1(a)$, $v'_1 = a$ and $n' = i + 1$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in \lfloor \tau \rfloor_V$:
- $\theta \sqsubseteq \theta'_1$:
From FU-D1
 - $(n - i - 1, H') \triangleright \theta'_1$:
From FU-D1 we know that $(n - i, H') \triangleright \theta'_1$ therefore from Lemma 1.19 we get $(n - i - 1, H') \triangleright \theta'_1$
 - $(\theta'_1, n - i - 1, v') \in \lfloor \tau' \rfloor_V$:
Since from FU-D1 we know that $(n - i, H'_1) \triangleright \theta'_1$ therefore from the Definition 1.8 we get $(\theta'_1, n - i, H'_1(a)) \in \lfloor \tau \rfloor_V$
From Lemma 1.15 we get $(\theta'_1, n - i - 1, H'_1(a)) \in \lfloor \tau \rfloor_V$
Since $\tau <: \tau'$ therefore from Lemma 1.23 we get $(\theta'_1, n - i - 1, H'_1(a)) \in \lfloor \tau' \rfloor_V$
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$
From FU-D1
- (c) $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$
From FU-D1

12. FG-assign:

$$\frac{\Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau \quad \mathcal{L} \vdash \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_1 := e_2 : \text{unit}}$$

To prove: $(\theta, n, (e_1 := e_2) \delta) \in \lfloor \text{unit} \rfloor_E^{pc}$

This means that from Definition 1.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor \text{unit} \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor \text{unit} \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-A0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in \lfloor ((\text{ref } \tau))^\ell \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiating IH1 with H and n . Since we know that $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in \lfloor ((\text{ref } \tau))^\ell \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-A1}) \end{aligned}$$

IH2:

$$\begin{aligned}
& \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\
& \exists \theta'_2. \theta'_1 \sqsubseteq (n_2 - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in \llbracket (\tau) \rrbracket_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc)
\end{aligned}$$

Instantiating IH2 with H'_1 and since we know that $H'_1 \triangleright \theta'_1 \wedge (H', (e_1 := e_2) \delta) \Downarrow (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq (n - i - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in \llbracket (\tau) \rrbracket_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \quad \text{(FU-A2)}
\end{aligned}$$

In order to prove FU-A0 we choose θ' as θ'_2 from FU-A2. Also we know from the evaluation rule, assign, that let $v'_1 = a_1$, $H' = H'_2[a_1 \mapsto v'_2]$, $v' = ()$ and $n' = i + j + 1$. Now we are required to show

$$(a) \theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, ()) \in \llbracket \text{unit} \rrbracket_V:$$

- $\theta \sqsubseteq \theta'_2$:

Since $\theta \sqsubseteq \theta'_1$ from FU-A1 and $\theta'_1 \sqsubseteq \theta'_2$ from FU-A2 therefore from Definition 1.2 we get $\theta \sqsubseteq \theta'_2$

- $(n - i - j - 1, H') \triangleright \theta'_2$:

From Definition 1.8 it suffices to prove that

- i. $\text{dom}(\theta'_2) \subseteq \text{dom}(H')$: From FU-A2

- ii. $\forall a \in \text{dom}(\theta'_2). (\theta'_2, n - i - j - 1, H'(a)) \in \llbracket \theta'_2(a) \rrbracket_V$:
 $\forall a \in \text{dom}(\theta'_2).$

- $a = a_1$:

From FU-A2 (since we know that $(\theta'_2, n - i - j, v'_2) \in \llbracket (\tau) \rrbracket_V$)

Therefore from Lemma 1.15 we get $(\theta'_2, n - i - j - 1, v'_2) \in \llbracket (\tau) \rrbracket_V$

- $a \neq a_1$:

From FU-A2 (since we know that $(n - i - j, H'_2) \triangleright \theta'_2$ therefore from Lemma 1.19 we get $(n - i - j - 1, H'_2) \triangleright \theta'_2$)

- $(\theta'_2, n - i - j - 1, ()) \in \llbracket \text{unit} \rrbracket_V$:

From Definition 1.6

$$(b) (\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell')$$

$$\forall a \in \text{dom}(H).$$

- $a = a_1$:

Since we know that $H(a_1) \neq H'(a_1)$ and $\theta(a_1) = \tau = A^{\ell_i}$ (given)

It is given that $\tau \searrow pc$ therefore $pc \sqsubseteq \ell_i$

- $a \neq a_1$:

From FU-A2

$$(c) (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$$

From FU-A2

□

Lemma 1.22 (Expression subtyping). $\forall pc, pc', \tau.$

$$\mathcal{L} \models pc \sqsubseteq pc' \implies \llbracket \tau \rrbracket_E^{pc'} \subseteq \llbracket \tau \rrbracket_E^{pc}$$

Proof. Given: $\mathcal{L} \models pc \sqsubseteq pc'$

To prove: $\llbracket(\tau)\rrbracket_E^{pc'} \subseteq \llbracket(\tau)\rrbracket_E^{pc}$

This means we need to prove that

$$\forall(\theta, n, e) \in \llbracket(\tau)\rrbracket_E^{pc'} . (\theta, n, e) \in \llbracket(\tau)\rrbracket_E^{pc}$$

This means given $\forall(\theta, n, e) \in \llbracket(\tau)\rrbracket_E^{pc'}$

It suffices to prove that $(\theta, n, e) \in \llbracket(\tau)\rrbracket_E^{pc}$

From Definition 1.7 for the chosen θ, n, e we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \llbracket\tau\rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc') \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall k < n.(H_1, e) \Downarrow_k (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \llbracket\tau\rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

This means that we are given some H_1 and k such that $(n, H_1) \triangleright \theta$, $k < n$ and $(H_1, e) \Downarrow_k (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \llbracket\tau\rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned} \quad (\text{B})$$

Instantiate H with H_1 and j with k in (A) to get

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - k, H'_1) \triangleright \theta' \wedge (\theta', n - k, v') \in \llbracket\tau\rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc') \end{aligned} \quad (\text{C})$$

In order to prove (B) we choose θ'_1 as θ' and we need to prove

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - k, H'_1) \triangleright \theta' \wedge (\theta', n - k, v') \in \llbracket\tau\rrbracket_V$:

We get this directly from (C)

- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$:

Since $pc \sqsubseteq pc'$ and we are given

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell')$$

Therefore

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$:

We are given

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc')$$

and since $pc \sqsubseteq pc'$ Therefore

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$$

□

Lemma 1.23 (Subtyping unary). *The following holds:*

1. $\forall A, A', \mathcal{L}$.

$$(a) \mathcal{L} \vdash A <: A' \implies \llbracket (A) \rrbracket_V \subseteq \llbracket (A') \rrbracket_V$$

2. $\forall \tau, \tau', \mathcal{L}$.

$$(a) \mathcal{L} \vdash \tau <: \tau' \implies \llbracket (\tau) \rrbracket_V \subseteq \llbracket (\tau') \rrbracket_V$$

$$(b) \forall pc. \mathcal{L} \vdash \tau <: \tau' \implies \llbracket (\tau) \rrbracket_E^{pc} \subseteq \llbracket (\tau') \rrbracket_E^{pc}$$

Proof. Proof by simultaneous induction on $A <: A'$ and $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of $A <: A'$ in the last step:

1. FGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2 \quad \mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove: $\llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V \subseteq \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V$

IH1: $\llbracket (\tau'_1) \rrbracket_V \subseteq \llbracket (\tau_1) \rrbracket_V$ (Statement 2(a))

IH2: $\forall pc. \llbracket (\tau_2) \rrbracket_E^{pc} \subseteq \llbracket (\tau'_2) \rrbracket_E^{pc}$ (Statement 2(b))

It suffices to prove: $\forall (\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V. (\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V$

This means that given some θ, n and $\lambda x.e_i$ s.t $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V$

Therefore from Definition 1.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta_1, i, e_i[v/x]) \in \llbracket \tau_2 \rrbracket_E^{\ell_e} \quad (3)$$

And it suffices to prove: $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V$

Again from Definition 1.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V \implies (\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E^{\ell'_e}$$

This means that given some $\theta_2, j < n, v$ s.t $\theta \sqsubseteq \theta_2$ and $(\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V$

And we are required to prove: $(\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E^{\ell'_e}$

Since $(\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V$ therefore from IH1 we know that $(\theta_2, j, v) \in \llbracket \tau_1 \rrbracket_V$

As a result from Equation 3 we know that

$$(\theta_2, j, e_i[v/x]) \in \llbracket \tau_2 \rrbracket_E^{\ell_e}$$

From IH2, we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau_2']_E^{\ell_e}$$

Since $\mathcal{L} \models \ell'_e \sqsubseteq \ell_e$ therefore from Lemma 1.22 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau_2']_E^{\ell'_e}$$

2. FGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau_1' \quad \mathcal{L} \vdash \tau_2 <: \tau_2'}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau_1' \times \tau_2'} \text{FGsub-prod}$$

To prove: $[\!(\tau_1 \times \tau_2)\!]_V \subseteq [\!(\tau_1' \times \tau_2')\!]_V$

IH1: $[\!(\tau_1)\!]_V \subseteq [\!(\tau_1')\!]_V$ (Statement 2(a))

IH2: $[\!(\tau_2)\!]_V \subseteq [\!(\tau_2')\!]_V$ (Statement 2(a))

It suffices to prove: $\forall(\theta, n, (v_1, v_2)) \in [\!(\tau_1 \times \tau_2)\!]_V. (\theta, n, (v_1, v_2)) \in [\!(\tau_1' \times \tau_2')\!]_V$

This means that given some θ, n and (v_1, v_2) $(\theta, (v_1, v_2)) \in [\!(\tau_1 \times \tau_2)\!]_V$

Therefore from Definition 1.6 we are given:

$$(\theta, n, v_1) \in [\tau_1]_V \wedge (\theta, n, v_2) \in [\tau_2]_V \quad (4)$$

And it suffices to prove: $(\theta, (v_1, v_2)) \in [\!(\tau_1' \times \tau_2')\!]_V$

Again from Definition 1.6, it suffices to prove:

$$(\theta, n, v_1) \in [\tau_1']_V \wedge (\theta, n, v_2) \in [\tau_2']_V$$

Since from Equation 4 we know that $(\theta, n, v_1) \in [\tau_1]_V$ therefore from IH1 we have $(\theta, n, v_1) \in [\tau_1']_V$

Similarly since $(\theta, n, v_2) \in [\tau_2]_V$ from Equation 4 therefore from IH2 we have $(\theta, n, v_2) \in [\tau_2']_V$

3. FGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau_1' \quad \mathcal{L} \vdash \tau_2 <: \tau_2'}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau_1' + \tau_2'} \text{FGsub-sum}$$

To prove: $[\!(\tau_1 + \tau_2)\!]_V \subseteq [\!(\tau_1' + \tau_2')\!]_V$

IH1: $[\!(\tau_1)\!]_V \subseteq [\!(\tau_1')\!]_V$ (Statement 2(a))

IH2: $[\!(\tau_2)\!]_V \subseteq [\!(\tau_2')\!]_V$ (Statement 2(a))

It suffices to prove: $\forall(\theta, n, v_s) \in [\!(\tau_1 + \tau_2)\!]_V. (\theta, v_s) \in [\!(\tau_1' + \tau_2')\!]_V$

This means that given: $(\theta, n, v_s) \in [\!(\tau_1 + \tau_2)\!]_V$

And it suffices to prove: $(\theta, n, v_s) \in [\!(\tau_1' + \tau_2')\!]_V$

2 cases arise

(a) $v_s = \text{inl } v_i$:

From Definition 1.6 we are given:

$$(\theta, n, v_i) \in \llbracket \tau_1 \rrbracket_V \quad (5)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \llbracket \tau'_1 \rrbracket_V$$

From Equation 5 and IH1 we know that

$$(\theta, n, v_i) \in \llbracket \tau'_1 \rrbracket_V$$

(b) $v_s = \text{inr } v_i$:

From Definition 1.6 we are given:

$$(\theta, n, v_i) \in \llbracket \tau_2 \rrbracket_V \quad (6)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \llbracket \tau'_2 \rrbracket_V$$

From Equation 6 and IH2 we know that

$$(\theta, n, v_i) \in \llbracket \tau'_2 \rrbracket_V$$

4. FGsub-ref:

Given:

$$\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove: $\llbracket ((\text{ref } \tau)) \rrbracket_V \subseteq \llbracket ((\text{ref } \tau)) \rrbracket_V$

It suffices to prove: $\forall (\theta, n, a) \in \llbracket ((\text{ref } \tau)) \rrbracket_V. (\theta, n, a) \in \llbracket ((\text{ref } \tau)) \rrbracket_V$

Trivial

5. FGsub-base:

Given:

$$\frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove: $\llbracket ((\mathbf{b})) \rrbracket_V \subseteq \llbracket ((\mathbf{b})) \rrbracket_V$

Directly from Definition 1.6

6. FGsub-unit:

Given:

$$\frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove: $\llbracket ((\text{unit})) \rrbracket_V \subseteq \llbracket ((\text{unit})) \rrbracket_V$

Directly from Definition 1.6

Proof of statement 2(a)

Given:

$$\frac{\mathcal{L} \vdash \ell \sqsubseteq \ell' \quad \mathcal{L} \vdash A <: A'}{\mathcal{L} \vdash A^\ell <: A^{\ell'}} \text{FGsub-label}$$

To prove: $\llbracket (A^\ell) \rrbracket_V \subseteq \llbracket (A^{\ell'}) \rrbracket_V$

From Definition 1.6 it suffices to prove: $\llbracket (A) \rrbracket_V \subseteq \llbracket (A') \rrbracket_V$

This we get directly from IH (Statement 1(a))

Proof of statement 2(b)

Given: $\mathcal{L} \vdash \tau <: \tau'$

To prove: $\llbracket (\tau) \rrbracket_E^{pc} \subseteq \llbracket (\tau') \rrbracket_E^{pc}$

This means we need to prove that

$$\forall (\theta, n, e) \in \llbracket (\tau) \rrbracket_E^{pc}. (\theta, n, e) \in \llbracket (\tau') \rrbracket_E^{pc}$$

This means given $(\theta, n, e) \in \llbracket (\tau) \rrbracket_E^{pc}$

It suffices to prove that $(\theta, n, e) \in \llbracket (\tau') \rrbracket_E^{pc}$

From Definition 1.7 we know we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall i < n.(H, e) \Downarrow_i (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - i, H') \triangleright \theta' \wedge (\theta', n - i, v') \in \llbracket \tau \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall j < n.(H_1, e) \Downarrow_j (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - j, v') \in \llbracket \tau' \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

This means that we are given some H_1 and $j < n$ s.t $(n, H_1) \triangleright \theta \wedge (H_1, e) \Downarrow_j (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - j, v') \in \llbracket \tau' \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiate H in (A) with H_1 and i with j then we choose θ'_1 as θ'

Also we have IH1 as $\llbracket \tau \rrbracket_V \subseteq \llbracket \tau' \rrbracket_V$ (Statement 2(a))

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H'_1) \triangleright \theta' \wedge (\theta', n - j, v') \in \llbracket \tau' \rrbracket_V$:

We are given $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H'_1) \triangleright \theta' \wedge (\theta', n - j, v') \in \llbracket \tau \rrbracket_V$

From IH1 we know that $\llbracket \tau \rrbracket_V \subseteq \llbracket \tau' \rrbracket_V$

Therefore, $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H'_1) \triangleright \theta' \wedge (\theta', n - j, v') \in \llbracket \tau' \rrbracket_V$

- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell')$:

Given

- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$:

Given

□

Lemma 1.24 (Binary interpretation of Γ implies Unary interpretation of Γ). $\forall W, \gamma, \Gamma, n.$
 $(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

Proof. Given: $(W, n, \gamma) \in [\Gamma]_V^A$

To prove: $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 1.13 we know that we are given:

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And we are required to prove:

$$\forall i \in \{1, 2\}. \forall m.$$

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i) \wedge \forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$$

Case $i = 1$

Given some m we need to show:

- $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_1)$:

$$\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_1)$$

Therefore, $\text{dom}(\Gamma) \subseteq (\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_1))$ (Given)

- $\forall x \in \text{dom}(\Gamma). (W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$:

$$\text{We are given: } \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

Therefore from Lemma 1.14 we know that

$$\forall m'. (W.\theta_1, m', \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Instantiating m' with m we get

$$(W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Case $i = 2$

Symmetric case as $i = 1$

□

Theorem 1.25 (Fundamental theorem binary). $\forall \Gamma, pc, W, \mathcal{A}, e, \tau, \gamma, n.$

$$\Gamma \vdash_{pc} e : \tau \wedge (W, n, \gamma) \in [\Gamma]_V^A \implies$$

$$(W, n, e(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau]_E^A$$

Proof. Proof by induction on the typing derivation

1. FG-var:

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

$$\text{To prove: } (W, n, x(\gamma \downarrow_1), x(\gamma \downarrow_2)) \in [\tau]_E^A$$

$$\text{Say } e_1 = x(\gamma \downarrow_1) \text{ and } e_2 = x(\gamma \downarrow_2)$$

From Definition of $[\tau]_E^A$ it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall j < n. (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^A$$

This means given some H_1, H_2 and j s.t $(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$

We are required to prove: $\exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^A$

Here

- $H'_1 = H_1$ and $H'_2 = H_2$
- $e_1 = v'_1 = \gamma(x) \Downarrow_1$
- $e_2 = v'_2 = \gamma(x) \Downarrow_2$
- $j = 1$

We choose $W' = W$.

- $W \sqsubseteq W$: From Definition 1.3
- $(n - 1, H_1, H_2) \stackrel{A}{\triangleright} W$:
Since we know that $(n, H_1, H_2) \stackrel{A}{\triangleright} W$ therefore from Lemma 1.20 we get $(n - 1, H_1, H_2) \stackrel{A}{\triangleright} W$
- $(W, n - 1, \gamma(x) \Downarrow_1, \gamma(x) \Downarrow_2) \in [\tau]_{\mathcal{V}}^A$:
We are given that $(W, n, \gamma) \in [\Gamma]_{\mathcal{V}}^A$ therefore from Lemma 1.18 we get $(W, n - 1, \gamma) \in [\Gamma]_{\mathcal{V}}^A$
which means from Definition 1.13 we have $(W, n - 1, \gamma(x) \Downarrow_1, \gamma(x) \Downarrow_2) \in [\tau]_{\mathcal{V}}^A$

2. FG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove: $(W, n, \lambda x. e (\gamma \Downarrow_1), \lambda x. e (\gamma \Downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)]_E^A$

Say $e_1 = \lambda x. e (\gamma \Downarrow_1)$ and $e_2 = \lambda x. e (\gamma \Downarrow_2)$

From Definition of $[(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_E^A$ it suffices to prove that

$$\forall H_1, H_2, j < n. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_{\mathcal{V}}^A$$

This means that given H_1, H_2 and j s.t $(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$

It suffices to prove:

$$\exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_{\mathcal{V}}^A \quad (\text{FB-L0})$$

IH1:

$\forall W, n. (W, n, e (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2]_E^A$

s.t

$(W, n, (v_1, v_2)) \in [\tau_1]_V^A$

We know from the evaluation rules that $H'_1 = H_1$, $H'_2 = H_2$, $v'_1 = e_1 = \lambda x.e (\gamma \downarrow_1)$, $v'_2 = e_2 = \lambda x.e (\gamma \downarrow_2)$ and $j = 0$. In order to prove FB-L0 we choose $W' = W$ and we need to prove the following:

- $W \sqsubseteq W$: From Definition 1.3
- $(n, H_1, H_2) \triangleright^A W$: Given
- $(W, n, \lambda x.e (\gamma \downarrow_1), \lambda x.e (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_V^A$

From Definition 1.4 it suffices to prove that:

$\forall W'' \sqsupseteq W, k < n, v_1, v_2.$

$((W'', k, v_1, v_2) \in [\tau_1]_V^A \implies (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2]_E^A) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_1, k, v_c.$

$((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e (\gamma \downarrow_1)[v_c/x]) \in [\tau_2]_E^{\ell_e}) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_2, k, v_c.$

$((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e (\gamma \downarrow_2)[v_c/x]) \in [\tau_2]_E^{\ell_e})$

This means that we need to prove the following:

- $\forall W'' \sqsupseteq W, k < n, v_1, v_2. ((W'', k, v_1, v_2) \in [\tau_1]_V^A \implies (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2]_E^A):$

This means given $W'' \sqsupseteq W, k < n, v_1, v_2$ s.t $((W'', k, v_1, v_2) \in [\tau_1]_V^A$

We need to prove: $(W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2]_E^A$

We instantiate IH1 with W'' and k

And since $(W'', k, v_1, v_2) \in [\tau_1]_V^A$ therefore we get

$(W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2]_E^A$

- $\forall \theta_l \sqsupseteq W.\theta_1, k, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e (\gamma \downarrow_1)[v_c/x]) \in [\tau_2]_E^{\ell_e}):$

This means that we are given θ_l, k and v_c s.t

$\theta_l \sqsupseteq W.\theta_1$ and $(\theta_l, k, v_c) \in [\tau_1]_V$

And we are required to prove:

$(\theta_l, k, e (\gamma \downarrow_1)[v_c/x]) \in [\tau_2]_E^{\ell_e}$

It is given to us that

$\forall v_1, v_2. (W, n, \gamma \in [\Gamma]_V^A$

Therefore from Lemma 1.24 we know that

$\forall m. (W.\theta_1, m, (\gamma \downarrow_1) \in [\Gamma]_V$

Therefore, we can apply Theorem 1.21 to obtain

$\forall m. (W.\theta_1, m, \lambda x.e \gamma \downarrow_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_V$

From Definition 1.6 it means that we have

$$\forall m. \forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies \\ (\theta', j, e[v/x]\gamma \downarrow_1) \in [\tau_2]_E^{\ell_e}$$

We instantiate m with some $l > k$, θ' with θ_l , j with k and v with v_c to get $W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2]_E^{\ell_e}$

Since we thow that $W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1]_V$ therefore we get $(\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2]_E^{\ell_e}$

$$- \forall \theta_l \sqsupseteq W.\theta_2, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V \implies \\ (\theta_l, k, e(\gamma \downarrow_2)[v_c/x]) \in [\tau_2]_E^{\ell_e}): \\ \text{Symmetric case as above}$$

3. FG-app:

$$\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Gamma \vdash_{pc} e_2 : \tau_1 \quad \mathcal{L} \vdash \tau_2 \searrow \ell \quad \mathcal{L} \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove: $(W, n, (e_1 e_2) (\gamma \downarrow_1), (e_1 e_2) (\gamma \downarrow_2)) \in [(\tau_2)]_E^A$

This means from Definition 1.5 we need to prove:

$$\forall H_1, H_2, n' < n. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2)]_V^A$$

This further means that given $H_1, H_2, n' < n$ s.t

$$(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2)]_V^A \quad (\text{FB-A0})$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}, i < n. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge (H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V^A$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $(e_1 e_2)$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps. Therefore $\exists i < n' < n$ s.t $(H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1)$. $(H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$ is known because $(e_1 e_2)$ reduces to value with $\gamma \downarrow_2$. Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_V^A \quad (7)$$

$$\underline{\text{IH2}}: (W'_1, n - i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [(\tau_1)]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{j1}, H_{j2}, j < (n - i). (n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge (H_{j1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies \\ \exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1)]_V^A$$

Instantiating H_{j_1} with H'_1 and H_{j_2} with H'_2 in IH2. Since the $(e_1 e_2)$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps. Also, e_1 reduces to value $\gamma \downarrow_1$ in $i < n'$ steps. Therefore $\exists j < n' - i < n - i$ s.t $(H_{i_1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j_1}, v'_{j_1})$. $(H_{i_2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j_2}, v'_{j_2})$ is known because $(e_1 e_2)$ reduces to value with $\gamma \downarrow_2$. Hence we get

$$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in \lceil (\tau_1) \rceil_V^A \quad (8)$$

We case analyze on $(W'_1, n - i, v'_1, v'_2) \in \lceil (\tau_1 \xrightarrow{\ell_e} \tau_2) \rceil_V^A$ from Equation 7

- Case $\ell \sqsubseteq \mathcal{A}$:

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_1, v'_2) \in \lceil (\tau_1 \xrightarrow{\ell_e} \tau_2) \rceil_V^A$$

This means

$$(W'_1, n - i, v'_1, v'_2) \in \lceil (\tau_1 \xrightarrow{\ell_e} \tau_2) \rceil_V^A$$

Let $v'_1 = \lambda x.e_{h1}$ and $v'_2 = \lambda x.e_{h2}$

Again from Definition 1.4 it means that

$$\forall W'_{h1} \sqsupseteq W'_1, j_1 < (n - i), v_1, v_2.$$

$$((W'_{h1}, j_1, v_1, v_2) \in \lceil \tau_1 \rceil_V^A \implies (W'_{h1}, j_1, e_{h1}[v_1/x], e_{h2}[v_2/x]) \in \lceil \tau_2 \rceil_E^A) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_1, m_1, v_c.$$

$$\wedge ((\theta_{l1}, m_1, v_1) \in \lceil \tau_1 \rceil_V \implies (W'_{h1}.\theta_1, e_{h1}[v_1/x]) \in \lceil \tau_2 \rceil_E^{\ell_e}) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_2, m_1, v_c.$$

$$\wedge (\theta_{l1}, m_1, v_2) \in \lceil \tau_1 \rceil_V \implies (W'_{h1}.\theta_2, e_{h2}[v_2/x]) \in \lceil \tau_2 \rceil_E^{\ell_e}$$

We instantiate W'_{h1} with W'_2 obtained from Equation 8. Similarly we also instantiate v_1 and v_2 with v'_{j_1} and v'_{j_2} respectively from Equation 8, and j_1 with $n - i - j$. And we get

$$(W'_2, n - i - j, e_{h1}[v'_{j_1}/x], e_{h2}[v'_{j_2}/x]) \in \lceil \tau_2 \rceil_E^A$$

From Definition 1.5 we get

$$\forall H_1, H_2, k_e < (n - i - j).(n - i - j, H_1, H_2) \stackrel{A}{\triangleright} W'_2 \wedge$$

$$(H_1, e_{h1}[v'_{j_1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1}) \wedge (H_2, e_{h2}[v'_{j_2}/x]) \Downarrow (H'_{f2}, v_{f2}) \implies$$

$$\exists W' \sqsupseteq W'_2.(n - i - j - k_e, H'_{f1}, H'_{f2}) \stackrel{A}{\triangleright} W' \wedge (W', n - i - j - k_e, v_{f1}, v_{f2}) \in \lceil \tau_2 \rceil_V^A$$

Instantiating H_1 with H'_{j_1} and H_2 with H'_{j_2} obtained from Equation 8. And since we know that $e_1 e_2$ reduces with $\gamma \downarrow_1$ in $n' < n$ steps. And e_2 reduces to value $\gamma \downarrow_1$ in $j < n' - 1 < n - i$ steps. Therefore $\exists k_e = n' - i - j < n - i - j$ s.t $(H_1, e_{h1}[v'_{j_1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1})$. $(H_2, e_{h2}[v'_{j_2}/x]) \Downarrow (H'_{f2}, v_{f2})$ is known because $(e_1 e_2)$ reduces to value with $\gamma \downarrow_2$. Hence we get

$$\exists W' \sqsupseteq W'_2.((n - i - j - k_e), H'_{f1}, H'_{f2}) \stackrel{A}{\triangleright} W' \wedge (W', (n - i - j - k_e), v_{f1}, v_{f2}) \in \lceil \tau_2 \rceil_V^A \quad (9)$$

This concludes the proof in this case.

- Case $\ell \not\sqsubseteq \mathcal{A}$:

From FB-A0 we know that we need to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2)]_V^A$$

In this case since we know that $\ell \not\sqsubseteq \mathcal{A}$. Let $\tau_2 = A^{\ell_i}$ and since $\tau_2 \searrow \ell$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

Therefore from Definition 1.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau_2)]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau_2)]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2)]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau_2)]_V)$$

This means given m_1 and m_2 it suffices to prove:

$$(\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2)]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau_2)]_V \quad (10)$$

In this case from Definition 1.6 we know that

$$\forall m. (W'_1.\theta_1, m, \lambda x. e_{h1}) \in [(\tau_1 \xrightarrow{\ell_\xi} \tau_2)]_V \quad (11)$$

$$\forall m. (W'_1.\theta_2, m, \lambda x. e_{h2}) \in [(\tau_1 \xrightarrow{\ell_\xi} \tau_2)]_V \quad (12)$$

Applying Definition 1.6 on Equation 11 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall v. (\theta', j_1, v) \in [\tau_1]_V \implies (\theta', j_1, e_{h1}[v/x]) \in [\tau_2]_E^{\ell_e} \text{ where } \theta = W'_1.\theta_1$$

We instantiate m with $m_1 + 2 + t_1$ where t_1 is the number of steps in which e_{h1} reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 1 + t_1). \forall v. (\theta', j_1, v) \in [\tau_1]_V \implies (\theta', j_1, e_{h1}[v/x]) \in [\tau_2]_E^{\ell_e} \quad (\text{FB-AC1})$$

Since from Equation 8 we have

$$(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1)]_V^A$$

Therefore from Lemma 1.14 we get

$$\forall m. (W'_2.\theta_1, m, v'_{j1}) \in [\tau_1]_V$$

Instantiating m with $m_1 + 1 + t_1$ we get

$$(W'_2.\theta_1, m_1 + 1 + t_1, v'_{j1}) \in [\tau_1]_V$$

Instantiating θ' with $W'_2.\theta_1$, j_1 with $m_1 + t_1$ and v with v'_{j1} from Equation 8.

$$\text{Therefore we get } (W'_2.\theta_1, m_1 + 1 + t_1, e_{h1}[v'_{j1}/x]) \in [\tau_2]_E^{\ell_e}$$

From Definition 1.7, we get

$$\begin{aligned}
& \forall H.(m_1 + 1 + t_1, H) \triangleright W'_2.\theta_1 \wedge \forall k_c < (m_1 + 1 + t_1).(H, e_{h1}[v'_{j1}/x]) \Downarrow_{k_c} (H'_1, v'_1) \implies \\
& \exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1 + t_1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1 + t_1 - k_c), v'_1) \in [\tau_2]_V \wedge \\
& (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e))
\end{aligned}$$

Since from Equation 8 we have $(n - i - j, H'_{j1}, H'_{j1}) \triangleright W'_2$

Therefore from Lemma 1.26 we get $\forall m.(m, H'_{j1}) \triangleright W'_2.\theta_1$

Instantiating m with $m_1 + 1 + t_1$ we get $(m_1 + 1 + t_1, H'_{j1}) \triangleright W'_2.\theta_1$

Now instantiating H with H'_{j1} from Equation 8 and k_c with t_1 we get

$$\begin{aligned}
& \exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2]_V \wedge \\
& (\forall a. H'_{j1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e)) \quad (\text{R1})
\end{aligned}$$

Similarly we can apply Definition 1.6 on Equation 12 to get

$$\begin{aligned}
& \forall m. \forall \theta'_2.(m, W'_1.\theta_2) \sqsubseteq \theta'_2 \wedge \forall j_2 < m. \forall v. (\theta'_2, j_2, v) \in [\tau_1]_V \implies \\
& (\theta'_2, j_2, e_{h2}[v/x]) \in [\tau_2]_E^{\ell_e}
\end{aligned}$$

We instantiate m with $m_2 + 2 + t_2$ where t_2 is the number of steps in which e_{h2} reduces

$$\begin{aligned}
& \forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall v. (\theta', j_1, v) \in [\tau_1]_V \implies \\
& (\theta', j_1, e_{h2}[v/x]) \in [\tau_2]_E^{\ell_e} \quad (\text{FB-AC2})
\end{aligned}$$

Since from Equation 8 we have

$$(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1)]_V^A$$

Therefore from Lemma 1.14 we get

$$\forall m. (W'_2.\theta_2, m, v'_{j2}) \in [\tau_1]_V$$

Instantiating m with $m_2 + 1 + t_2$ we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, v'_{j2}) \in [\tau_1]_V$$

Instantiating θ' with $W'_2.\theta_2$, j_1 with $m_2 + 1 + t_2$ and v with v'_{j2} from Equation 8 in FB-AC2 we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, e_{h2}[v'_{j2}/x]) \in [\tau_2]_E^{\ell_e}$$

From Definition 1.7, we get

$$\begin{aligned}
& \forall H.(m_2 + 1 + t_2, H) \triangleright W'_2.\theta_2 \wedge \forall k_c < (m_2 + 1 + t_2).(H, e_{h2}[v'_{j1}/x]) \Downarrow_{k_c} (H'_2, v'_2) \implies \\
& \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1 + t_2 - k_c), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1 + t_2 - k_c), v'_2) \in [\tau_2]_V \wedge \\
& (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e))
\end{aligned}$$

Since from Equation 8 we have $(n - i - j, H'_{j1}, H'_{j1}) \triangleright W'_2$

Therefore from Lemma 1.26 we get $\forall m.(m, H'_{j2}) \triangleright W'_2.\theta_2$

Instantiating m with $m_2 + 1 + t_2$ we get $(m_2 + 1 + t_2, H'_{j2}) \triangleright W'_2.\theta_2$

Now Instantiating H with H'_{j2} from Equation 8 and k_c with t_2 .

$$\begin{aligned}
& \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2]_V \wedge \\
& (\forall a. H'_{j2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e)) \quad (\text{R2})
\end{aligned}$$

In order to prove FB-A0 we choose W' to be $(\theta'_1, \theta'_2, W'_2.\beta)$. Now we need to show two things:

(a) $(n - n', H'_1, H'_2) \triangleright W'$:

From Definition 1.9 it suffices to show that

– $dom(W'.\theta_1) \subseteq dom(H'_1) \wedge dom(W.\theta_2) \subseteq dom(H'_2)$:

From R1 we know that $(m_1 + 1, H'_1) \triangleright \theta'_1$, therefore from Definition 1.8 we get

$dom(W'.\theta_1) \subseteq dom(H'_1)$

Similarly, from R2 we know that $(m_2 + 1, H'_2) \triangleright \theta'_2$, therefore from Definition 1.8 we get $dom(W'.\theta_2) \subseteq dom(H'_2)$

– $(W'.\hat{\beta}) \subseteq (dom(W'.\theta_1) \times dom(W'.\theta_2))$:

Since from Equation 8 we know that $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$ therefore from

Definition 1.9 we know that $(W'_2.\hat{\beta}) \subseteq (dom(W'_2.\theta_1) \times dom(W'_2.\theta_2))$

From R1 and R2 we know that $W'_2.\theta_1 \sqsubseteq \theta'_1$ and $W'_2.\theta_2 \sqsubseteq \theta'_2$ therefore

$(W'_2.\hat{\beta}) \subseteq (dom(\theta'_1) \times dom(\theta'_2))$

– $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge$

$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

4 cases arise for each $(a_1, a_2) \in W'_2.\hat{\beta}$

i. $H'_{j_1}(a_1) = H'_1(a_1) \wedge H'_{j_2}(a_2) = H'_2(a_2)$:

* $W'.\theta_1(a_1) = W'.\theta_2(a_2)$:

We know from Equation 8 that $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$

Therefore from Definition 1.9 we have

$\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

Since $W'.\hat{\beta} = W'_2.\hat{\beta}$ by construction therefore

$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

From R1 and R2 we know that $W'_2.\theta_1 \sqsubseteq \theta'_1$ and $W'_2.\theta_2 \sqsubseteq \theta'_2$ respectively.

Therefore from Definition 1.2

$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$

* $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

From Equation 8 we know that $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$

This means from Definition 1.9 that

$\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$

Instantiating with a_1 and a_2 and since $W'_2 \sqsubseteq W'$ and $n - n' - 1 < n - i - j - 1$ (since $n' = i + j + t_1$ where t_1 is the number of steps taken by e_{h_1} , i is the number of steps taken by $e_1 \gamma \downarrow_1$ to reduce and j is the number of steps taken by $e_2 \gamma \downarrow_1$ to reduce) therefore from Lemma 1.16 we get

$(W', n - n' - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$

ii. $H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$:

* $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

Same reasoning as in the previous case

* $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

From R1 and R2 we know that

$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = A^{\ell'} \wedge (\ell_e \sqsubseteq \ell')$

$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sqsubseteq \ell')$

This means we have

$$\begin{aligned}\exists \ell'. W'_2.\theta_1(a_1) &= A^{\ell'} \wedge (\ell_e) \sqsubseteq \ell' \text{ and} \\ \exists \ell'. W'_2.\theta_2(a_2) &= A^{\ell'} \wedge (\ell_e) \sqsubseteq \ell'\end{aligned}$$

Since $pc \sqcup \ell \sqsubseteq \ell_e$ (given) and $\ell \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_e \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Also from R1 and R2, $(m_1 + 1, H'_1) \triangleright \theta'_1$ and $(m_2 + 1, H'_2) \triangleright \theta'_2$. Therefore from Definition 1.8 we have

$$\begin{aligned}(\theta'_1, m_1, H'_1(a_1)) &\in [\theta'_1(a_1)]_V \text{ and} \\ (\theta'_2, m_2, H'_2(a_1)) &\in [\theta'_2(a_2)]_V\end{aligned}$$

Since m_1 and m_2 are arbitrary indices therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$\text{iii. } H'_{j_1}(a_1) = H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2):$$

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2)$$

Same reasoning as in the previous case

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'.\theta_1(a_1)]_V^A:$$

From R2 we know that

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e) \sqsubseteq \ell')$$

This means that a_2 was protected at ℓ_e in the world before the modification. Since $pc \sqcup \ell \sqsubseteq \ell_e$ (given) and $\ell \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_e \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 8 we know that $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$ that means from Definition 1.9 that $(W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in [W'_2.\theta_1(a_1)]_V^A$. Since $(\ell_e) \sqsubseteq \ell'$ therefore from Definition 1.4 we know that $H'_{j_1}(a_1)$ must also be protected at some label $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_2.\theta_1, m, H'_{j_1}(a_1)) \in W'_2.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_2.\theta_2, m, H'_{j_2}(a_2)) \in W'_2.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with m_1 and using Lemma 1.15 we get

$$(\theta'_1, m_1, H'_{j_1}(a_1)) \in \theta'_1(a_1)$$

Since from R2 we know that $(m_2 + 1, H'_2) \triangleright \theta'_2$ therefore from Definition 1.8

$$\text{we know that } (\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$$

Therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$\text{iv. } H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) = H'_2(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V:$$

$$\underline{i = 1}$$

This means that given some m we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

Like before we instantiate Equation 11 and Equation 12 with $m + 2 + t_1$ and $m + 2 + t_2$ respectively. This will give us

$$\begin{aligned} & \exists \theta'_1. W'_2. \theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2]_V \wedge \\ & (\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2. \theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2. \theta_1). \theta'_1(a) \searrow (\ell_e)) \end{aligned}$$

and

$$\begin{aligned} & \exists \theta'_2. W'_2. \theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2]_V \wedge \\ & (\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2. \theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2. \theta_2). \theta'_2(a) \searrow (\ell_e)) \end{aligned}$$

Since we have $(m + 1, H'_1) \triangleright \theta'_1$ and $(m + 1, H'_2) \triangleright \theta'_2$ therefore we get the desired from Definition 1.8

$$\underline{i = 2}$$

Symmetric to $i = 1$

$$(b) (W', n - n' - 1, v'_1, v'_2) \in [\tau_2]_{\mathcal{V}}^{\mathcal{A}}:$$

Let $\tau_2 = \mathbf{A}^{\ell_i}$ Since $\tau_2 \searrow \ell$ and since $\ell \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

From R1 and R2 we and Definition 1.4 we get the desired.

4. FG-prod:

$$\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \quad \Gamma \vdash_{pc} e_2 : \tau_2}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove: $(W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\perp]_E^{\mathcal{A}}$

Say $e_1 = (e_1, e_2) (\gamma \downarrow_1)$ and $e_2 = (e_1, e_2) (\gamma \downarrow_2)$

From Definition of $[(\tau_1 \times \tau_2)^\perp]_E^{\mathcal{A}}$ it suffices to prove that

$$\begin{aligned} & \forall H_1, H_2. (n, H_1, H_2) \stackrel{\mathcal{A}}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ & \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp]_{\mathcal{V}}^{\mathcal{A}} \end{aligned}$$

This means that given some H_1, H_2 and $n' < n$ s.t

$$(n, H_1, H_2) \stackrel{\mathcal{A}}{\triangleright} W \wedge (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp]_{\mathcal{V}}^{\mathcal{A}} \quad (13)$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [\tau_1]_E^{\mathcal{A}}$$

This means from Definition 1.5 we get

$$\begin{aligned} & \forall H_{p11}, H_{p12}. (n, H_{p11}, H_{p12}) \stackrel{\mathcal{A}}{\triangleright} W \wedge \forall i < n. (H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11}) \wedge (H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12}) \implies \\ & \exists W'_1 \sqsupseteq W. (n - i, H'_{p11}, H'_{p12}) \stackrel{\mathcal{A}}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}} \end{aligned}$$

Instantiating H_{p11} with H_1 and H_{p22} with H_2 in IH1 and since the (e_1, e_2) reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore we know that $\exists i < n' < n$ s.t $(H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11})$. Similarly since we know that (e_1, e_2) reduces to value with $\gamma \downarrow_2$ therefore we know that $(H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12})$. Hence we get

$$\exists W'_1 \sqsubseteq W.(n-i, H'_{p11}, H'_{p12}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n-i, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^A \quad (14)$$

$$\underline{\text{IH2}} (W, n-i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_2]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{p21}, H_{p22}.(n-i, H_{p21}, H_{p22}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n-i. (H_{p21}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{p22}, v'_{p22}) \implies$$

$$\exists W'_2 \sqsubseteq W'_1.(n-i-j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n-i-j, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A$$

Instantiating H_{p21} with H'_{p11} and H_{p22} with H'_{p21} and in IH2. Since (e_1, e_2) reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps and e_1 has reduced with $i < n'$ steps. Therefore we know that $\exists j < n' - i < n - i$ s.t $(H_{p21}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{p21}, v'_{p21})$. Similarly since we know that (e_1, e_2) reduces to value with $\gamma \downarrow_2$ therefore we know that $(H_{p22}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{p22}, v'_{p22})$. Hence we get

since the (e_1, e_2) reduces to value with both $\gamma \downarrow_1$ and $\gamma \downarrow_2$ therefore we know that $(H_{p21}, e_2 (\gamma \downarrow_1)) \Downarrow (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_{p22}, v'_{p22})$. Hence we get

$$\exists W'_2 \sqsubseteq W'_1.(n-i-j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n-i-j, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A \quad (15)$$

In order to prove Equation 13 we instantiate W' in Equation 13 with W'_2 we are required to show the following:

- $W \sqsubseteq W'_2$:

Since $W \sqsubseteq W'_1$ from Equation 14 and $W'_1 \sqsubseteq W'_2$ from Equation 15

Therefore, $W \sqsubseteq W'_2$ from Definition 1.3

- $(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W'$:

Here $n' = i + j + 1$

From evaluation rule of products we know that $H'_1 = H'_{p21}$ and $H'_2 = H'_{p22}$

From Equation 15 we know that $(n-i-j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2$

Therefore from Lemma 1.20 we get $(n-i-j-1, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2$

- $(W', n-i-j-1, v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp]_{\mathcal{V}}^A$:

From evaluation rule of products we know that $v'_1 = (v'_{p11}, v'_{p21})$ and $v'_2 = (v'_{p12}, v'_{p22})$

We are required to show

$$- (W'_2, n-i-j-1, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^A \wedge (W'_2, n-i-j-1, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A:$$

From Equation 14 and Equation 15 we know that

$$(W'_2, n-i-j, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^A \wedge (W'_2, n-i-j, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A$$

Therefore from Lemma 1.16 we get

$$(W'_2, n-i-j-1, v'_{p11}, v'_{p12}) \in [\tau_1]_{\mathcal{V}}^A \wedge (W'_2, n-i-j-1, v'_{p21}, v'_{p22}) \in [\tau_2]_{\mathcal{V}}^A$$

5. FG-fst:

$$\frac{\Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$

To prove: $(W, n, (\text{fst}(e_i)) (\gamma \downarrow_1), (\text{fst}(e_i)) (\gamma \downarrow_2)) \in [\tau_1]_E^A$

Say $e_1 = (\text{fst}(e_i)) (\gamma \downarrow_1)$ and $e_2 = (\text{fst}(e_i)) (\gamma \downarrow_2)$

From Definition 1.5 it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1]_V^A \end{aligned}$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1]_V^A \quad (16)$$

IH1

$$(W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\ell]_E^A$$

This means from Definition 1.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow \\ (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell]_V^A \end{aligned}$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $\text{fst}(e_i)$ reduces to value reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore we know that $\exists i < n' < n$ s.t $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1})$. Similarly since we know that $\text{fst}(e_i)$ reduces to value with $\gamma \downarrow_2$ therefore we know that $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell]_V^A \quad (17)$$

We case analyze on $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell]_V^A$ from Equation 17

- Case $\ell \sqsubseteq \mathcal{A}$:

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)]_V^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)]_V^A$$

Let $v'_{i1} = (v_{i1}, v_{j1})$ and $v'_{i2} = (v_{i2}, v_{j2})$

Again from Definition 1.4 it means that

$$(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W'_1, n - i, v_{i2}, v_{j2}) \in [\tau_2]_V^A \quad (\text{F1})$$

In order to prove Equation 16 we choose W' as W'_1 and from the evaluation rule of fst we know that $H'_1 = H'_{i1}$ and $H'_2 = H'_{i2}$. Also, from reduction rules we know that $n' = i + 1$. And then we need to show:

- $W \sqsubseteq W'_1$:

Directly from Equation 17

– $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$:

Since from Equation 17 we know that $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 1.20 we get $(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

– $(W'_1, n - n', v'_1, v'_2) \in [\tau_1]_V^A$:

From the evaluation rule we know that $v'_1 = v_{i1}$ and $v'_2 = v_{j1}$

From F1 we know that $(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1]_V^A$

Therefore from Lemma 1.16 we get $(W'_1, n - i - 1, v_{i1}, v_{j1}) \in [\tau_1]_V^A$

• Case $\ell \not\sqsubseteq \mathcal{A}$:

In this case from Definition 1.6 we know that

(a) $\forall m. (W'_1.\theta_1, m, v'_{i1}) \in [(\tau_1 \times \tau_2)]_V$ and

(b) $\forall m. (W'_1.\theta_2, m, v'_{i2}) \in [(\tau_1 \times \tau_2)]_V$

where

$v'_{i1} = (v_{i1}, v_{i2})$ and $v'_{i2} = (v_{j1}, v_{j2})$

In order to prove Equation 16 we choose W' as W'_1 and from the evaluation rule of fst we know that $H'_1 = H'_{i1}$ and $H'_2 = H'_{i2}$. And then we need to show:

– $W \sqsubseteq W'_1$:

Directly from Equation 17

– $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$:

From Equation 17 we know that $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 1.20 we get

$(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

– $(W'_1, n - n', v'_1, v'_2) \in [\tau_1]_V^A$:

From the evaluation rule we know that $v'_1 = v_{i1}$ and $v'_2 = v_{j1}$

Let $\tau_1 = \mathbf{A}^{\ell_i}$ Since $\tau_1 \searrow \ell$ and since $\ell \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

Therefore from Definition 1.4 it suffices to prove that

$\forall m_1. (W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$

and

$\forall m_2. (W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1]_V$

This means given m_1 and it suffices to prove:

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1]_V \tag{18}$$

Similarly given m_2 , it suffices to prove:

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1]_V \tag{19}$$

Instantiating (a) with m_1

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1]_V \wedge (W'_1.\theta_1, m_1, v_{i2}) \in [\tau_2]_V \tag{20}$$

Instantiating (b) with m_2

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1]_V \wedge (W'_1.\theta_2, m_2, v_{j2}) \in [\tau_2]_V \tag{21}$$

From Equation 20 and Equation 21 we get

$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$ and $(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1]_V$

6. FG-snd:

Symmetric case as FG-fst

7. FG-inl:

$$\frac{\Gamma \vdash_{pc} e_i : \tau_1}{\Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove: $(W, n, (\text{inl}(e_i)) (\gamma \downarrow_1), (\text{inl}(e_i)) (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\perp]_E^A$

Say $e_1 = (\text{inl}(e_i)) (\gamma \downarrow_1)$ and $e_2 = (\text{inl}(e_i)) (\gamma \downarrow_2)$

From Definition of $[(\tau_1 + \tau_2)^\perp]_E^A$ it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp]_V^A \end{aligned}$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp]_V^A \quad (22)$$

$$\text{IH1 } (W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [\tau_1]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1]_V^A$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $\text{inl}(e_i)$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore we know that $\exists i < n' < n$ s.t $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$. Similarly since we know that $\text{inl}(e_i)$ reduces to value with $\gamma \downarrow_2$ therefore we know that $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1]_V^A \quad (23)$$

Instantiating W' in Equation 22 with W'_1 . Also from reduction relation we know that $n' = i + 1$ we are required to show the following:

- $W \sqsubseteq W'_1$:

Directly from Equation 23

- $(n - n', H'_1, H'_2) \triangleright^A W'_1$:

From Equation 23 we know that $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$

Therefore from Lemma 1.20 we get

$$(n - n', H'_1, H'_2) \triangleright^A W'_1$$

- $(W'_1, n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^+]^A_{\mathcal{V}}$:

From evaluation rule of inl we know that $v'_1 = \text{inl}(v'_{i1})$ and $v'_2 = \text{inl}(v'_{i2})$

We are required to show

- $(W'_1, n - n', v'_{i1}, v'_{i2}) \in [\tau_1]^A_{\mathcal{V}}$:

From Equation 23 we know that $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1]^A_{\mathcal{V}}$

Therefore from Lemma 1.16 we get

$$(W'_1, n - i - 1, v'_{i1}, v'_{i2}) \in [\tau_1]^A_{\mathcal{V}}$$

8. FG-inr:

Symmetric case to FG-inl.

9. FG-case:

$$\frac{\Gamma \vdash_{pc} e_i : (\tau_1 + \tau_2)^\ell \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{i1} : \tau \quad \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_{i2} : \tau \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \text{case}(e_i, x.e_{i1}, y.e_{i2}) : \tau}$$

To prove: $(W, (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1), (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)) \in [(\tau)]^A_E$

Say $e_1 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1)$ and $e_2 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)$

This means from Definition 1.5 we need to prove:

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau)]^A_{\mathcal{V}} \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau)]^A_{\mathcal{V}} \quad (24)$$

$$\underline{\text{IH1}} \quad (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\ell]^A_E$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell]^A_{\mathcal{V}}$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$ reduces to value with both $\gamma \downarrow_1$ and $\gamma \downarrow_2$ therefore we know that $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell]^A_{\mathcal{V}} \quad (25)$$

IH2:

$$(W'_1, n - i, (e_{i1}) (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\}), (e_{i1}) (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \in [(\tau)]^A_E$$

This means from Definition 1.5 we get

$$\forall H_{j_1}, H_{j_2}. (n - i, H_{j_1}, H_{j_2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_1, e_{i1} (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\})) \Downarrow_j (H'_{j_1}, v'_{j_1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \Downarrow (H'_{j_2}, v'_{j_2}) \implies$$

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau)]_V^A$$

Instantiating H_{j_1} with H'_1 and H_{j_2} with H'_2 in IH2. Also instantiating W with W'_1 . Since the $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$ reduces to value in both runs therefore we know that $(H_1, e_{i1} (\gamma \downarrow_1)) \Downarrow (H'_{j_1}, v'_{j_1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{j_2}, v'_{j_2})$. Hence we get

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau)]_V^A \quad (26)$$

IH3:

$$(W'_1, n - i, (e_{i2}) (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\}), (e_{i2}) (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \in [(\tau)]_E^A$$

This means from Definition 1.5 we get

$$\forall H_{k_1}, H_{k_2}. (n - i, H_{k_1}, H_{k_2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall k < n - i. (H_1, e_{i2} (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\})) \Downarrow_k (H'_{k_1}, v'_{k_1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \Downarrow (H'_{k_2}, v'_{k_2}) \implies$$

$$\exists W'_3 \sqsupseteq W'_1. (n - i - k, H'_{k_1}, H'_{k_2}) \stackrel{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k_1}, v'_{k_2}) \in [(\tau)]_V^A$$

Instantiating H_{k_1} with H'_1 and H_{k_2} with H'_2 in IH2. Also instantiating W with W'_1 . Since the $(\text{case}(e_i, x.e_{i2}, y.e_{i2}))$ reduces to value in both runs therefore we know that $(H_1, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{k_1}, v'_{k_1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{k_2}, v'_{k_2})$. Hence we get

$$\exists W'_3 \sqsupseteq W'_1. (n - i - k, H'_{k_1}, H'_{k_2}) \stackrel{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k_1}, v'_{k_2}) \in [(\tau)]_V^A \quad (27)$$

We case analyze $(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 + \tau_2)^\ell]_V^A$ from Equation 25

- Case $\ell \sqsubseteq \mathcal{A}$:

From Definition 1.4 2 further cases arise:

- $v'_1 = \text{inl}(v_{i1})$ and $v'_2 = \text{inl}(v_{i2})$:

In this case from Definition 1.4 we know that $(W, n - i, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

In order to prove Equation 24 we choose W' as W'_2 from Equation 26 and from the first evaluation rule of case we know that $H'_1 = H'_{j_1}$ and $H'_2 = H'_{j_2}$. Also we know from the evaluation rule that $n' = i + j + 1$. And then we need to show:

- * $W \sqsubseteq W'_2$:

Since $W \sqsubseteq W'_1$ from Equation 25 and $W'_1 \sqsubseteq W'_2$ from Equation 26

Therefore, $W \sqsubseteq W'_2$ from Definition 1.3

- * $(n - n', H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2$:

From Equation 26 we know that $(n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2$

Therefore from Lemma 1.20 we get

$$(n - i - j - 1, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2$$

* $(W'_2, n - n', v'_1, v'_2) \in [\tau]_V^A$:

From the evaluation rule we know that $v'_1 = v'_{j_1}$ and $v'_2 = v'_{j_2}$

From Equation 26 we know that $(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [\tau]_V^A$

Therefore from Lemma 1.16 we get

$(W'_2, n - i - j - 1, v'_{j_1}, v'_{j_2}) \in [\tau]_V^A$

- $v'_1 = \text{inr}(v_{i_1})$ and $v'_2 = \text{inr}(v_{i_2})$:

In this case from Definition 1.4 we know that $(W, v_{i_1}, v_{i_2}) \in [\tau_2]_V^A$

In order to prove Equation 24 we choose W' as W'_3 from Equation 27 and from the second evaluation rule of case we know that $H'_1 = H'_{k_1}$ and $H'_2 = H'_{k_2}$. Also we know from the evaluation rule that $n' = i + k + 1$. And then we need to show:

* $W \sqsubseteq W'_3$:

Since $W \sqsubseteq W'_1$ from Equation 25 and $W'_1 \sqsubseteq W'_3$ from Equation 27

Therefore, $W \sqsubseteq W'_3$ from Definition 1.3

* $(n - n', H'_1, H'_2) \triangleright^A W'_3$:

From Equation 27 we know that $(n - i - k, H'_{k_1}, H'_{k_2}) \triangleright^A W'_3$

Therefore from Lemma 1.20 we get

$(n - i - k - 1, H'_{k_1}, H'_{k_2}) \triangleright^A W'_3$

* $(W'_3, n - n', v'_1, v'_2) \in [\tau]_V^A$:

From the evaluation rule we know that $v'_1 = v'_{k_1}$ and $v'_2 = v'_{k_2}$

From Equation 27 we know that $(W'_3, n - i - k, v'_{k_1}, v'_{k_2}) \in [\tau]_V^A$

Therefore from Lemma 1.16 we get

$(W'_3, n - i - k - 1, v'_{k_1}, v'_{k_2}) \in [\tau]_V^A$

• Case $\ell \not\sqsubseteq \mathcal{A}$:

The following cases arise:

- Reduction of e_1 happens via Case1 and Reduction of e_2 happens via Case1 :
Exactly the same reasoning as in the $v'_1 = \text{inl}(v_{i_1})$ and $v'_2 = \text{inl}(v_{i_2})$ subcase of the $\ell \not\sqsubseteq \mathcal{A}$ case before.
- Reduction of e_1 happens via Case2 and Reduction of e_2 happens via Case2 :
Exactly the same reasoning as in the $v'_1 = \text{inr}(v_{i_1})$ and $v'_2 = \text{inr}(v_{i_2})$ subcase of the $\ell \not\sqsubseteq \mathcal{A}$ case before.
- Reduction of e_1 happens via Case1 and Reduction of e_2 happens via Case2 :

From Equation 24 we know that we need to prove

$\exists W' \sqsubseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau)]_V^A$

In this case since we know that $\ell \not\sqsubseteq \mathcal{A}$. Let $\tau = A^{\ell_i}$ and since $\tau \searrow \ell$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove $\exists W' \sqsubseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau)]_V^A$

From Definition 1.4 it will suffice to prove

$\exists W' \sqsubseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau)]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau)]_V)$

This means it suffices to prove

$(\forall m_1, m_2. \exists W' \sqsubseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau)]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau)]_V)$

This means given m_1 and m_2 it suffices to prove:

$$(\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau)]_V \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau)]_V) \quad (28)$$

Since we know that $(W, n, \gamma) \in [\Gamma]_V^A$ (given) therefore from Lemma 1.24 we know that $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

Therefore by instantiating it at $m_1 + 1 + j$ we know that

$$(W.\theta_1, m_1 + 1 + j, \gamma \downarrow_1) \in [\Gamma]_V \quad (29)$$

Next we apply Theorem 1.21 on $e_{i1} \gamma \downarrow_1$. Here j is the number of steps in which $e_{i1} \gamma \downarrow_1$ reduces. We use $\gamma \downarrow_1 \cup \{x \mapsto v'_{s1}\}$ as the unary substitution to get $(W.\theta_1, m_1 + 1 + j, e_{i1} \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \in [(\tau)]_E^{pc}$

This means from Definition 1.7 we get

$$\begin{aligned} & \forall H_{c2}. (m_1 + 1 + j, H_{c1}) \triangleright W_1.\theta_1 \wedge \forall l_c < (m_1 + 1 + j). (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \downarrow_{k_c} \\ & (H'_{c2}, v'_c) \implies \\ & \exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1 + j - l_c, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1 + j - l_c, v'_c) \in [(\tau)]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W_1.\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

Since from Equation 25 we know that $(n-i, H'_1, H'_2) \triangleright W'_1$ therefore from Lemma 1.26 we get $\forall m. (m, H'_1) \triangleright W'_1.\theta_1$

Instantiating m with $m_1 + 1 + j$ we get $(m_1 + 1 + j, H'_1) \triangleright W'_1.\theta_1$

Instantiating H_{c2} with H'_1 from Equation 25 and l_c with j we get

$$\begin{aligned} & \exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1, v'_c) \in [(\tau)]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell)) \quad (\text{CC1}) \end{aligned}$$

Similarly we apply Theorem 1.21 on $e_{i2} \gamma \downarrow_2$. Here j_2 is the number of steps in which $e_{i2} \gamma \downarrow_2$ reduces. We use $\gamma \downarrow_2 \cup \{y \mapsto v'_{s2}\}$ as the unary substitution to get $(W_1.\theta_2, m_2 + 1 + j_2, e_{i2} \gamma \downarrow_2 \cup \{y \mapsto v'_c\}) \in [(\tau)]_E^{pc}$

This means from Definition 1.7 we get

$$\begin{aligned} & \forall H_{c2}. (m_2 + 1 + j_2, H_{c1}) \triangleright W_1.\theta_2 \wedge \forall l_c < m_2 + 1 + j_2. (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \downarrow_{k_c} \\ & (H'_{c2}, v'_c) \implies \\ & \exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1 + j_2 - l_c, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1 + j_2 - l_c, v'_c) \in [(\tau)]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_1). \theta'_2(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

Since from Equation 25 we know that $(n-i, H'_1, H'_2) \triangleright W'_1$ therefore from Lemma 1.26 we get $\forall m. (m, H'_2) \triangleright W'_1.\theta_2$

Instantiating m with $m_2 + 1 + j_2$ we get $(m_2 + 1 + j_2, H'_2) \triangleright W'_1.\theta_2$

Instantiating H_{c2} with H'_2 (from Equation 25) and l_c with j_2 to get

$$\begin{aligned} & \exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1, v'_c) \in [(\tau)]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_1). \theta'_2(a) \searrow (pc \sqcup \ell)) \quad (\text{CC2}) \end{aligned}$$

We choose

$$W_n.\theta_1 = \theta'_1 \text{ (from CC1)}$$

$$W_n.\theta_2 = \theta'_2 \text{ (from CC2)}$$

$$W_n.\hat{\beta} = W'_1.\hat{\beta} \text{ (from Equation 25)}$$

In order to prove Equation 24 we choose W' as W_n

i. $(n - n', H'_1, H'_2) \triangleright W'$:

From Definition 1.9 it suffices to show that

$$- \text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H'_2):$$

From (CC1) we know that $(m_1 + 1, H'_1) \triangleright \theta'_1$, therefore from Definition 1.8 we get $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from (CC2) we know that $(m_2 + 1, H'_2) \triangleright \theta'_2$, therefore from Definition 1.8 we get $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

$$- (W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2)):$$

Since from Equation 25 we have $(n - i, H'_1, H'_2) \triangleright W'_1$ therefore from Definition 1.9 we get $(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$

From (CC1) and (CC2) we know that $W'_1.\theta_1 \sqsubseteq \theta'_1$ and $W'_1.\theta_2 \sqsubseteq \theta'_2$ therefore $(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

$$- \forall(a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge \\ (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

4 cases arise for each a_1 and a_2

A. $H'_{j1}(a_1) = H'_1(a_1) \wedge H'_{j2}(a_2) = H'_2(a_2):$

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2):}$$

We know from Equation 25 that $(n - i, H'_1, H'_2) \triangleright W'_1$

Therefore from Definition 1.9 we have

$$\forall(a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since $W'.\hat{\beta} = W'_1.\hat{\beta}$ by construction therefore

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From (CC1) and (CC2) we know that $W'_1.\theta_1 \sqsubseteq \theta'_1$ and $W'_1.\theta_2 \sqsubseteq \theta'_2$ respectively.

Therefore from Definition 1.2

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:}$$

From Equation 25 we know that $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

This means from Definition 1.9 that

$$\forall(a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) \wedge (W'_1, n - i - 1, H'_1(a_{i1}), H'_2(a_{i2})) \in \lceil W'_1.\theta_1(a_{i1}) \rceil_V^A$$

Instantiating with a_1 and a_2 and since $W'_1 \sqsubseteq W'$ and $n - n' - 1 < n - i - 1$ (since $n' = i + t_1 + 1$ where t_1 is the number of steps taken by e_{i1} , i is the number of steps taken by $e_1 \gamma \downarrow_1$ to reduce) therefore from Lemma 1.16 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

B. $H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2):$

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2):}$$

Same as before

$$\frac{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}{\text{From (CC1) and (CC2) we know that}}$$

$$(\forall a. H'_1(a) \neq H'_{c1}(a) \implies \exists \ell'. W'_1.\theta_1(a) = A^{\ell'} \wedge ((pc \sqcup \ell) \sqsubseteq \ell')$$

$$(\forall a. H'_2(a) \neq H'_{c2}(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge ((pc \sqcup \ell) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = A^{\ell'} \wedge ((pc \sqcup \ell) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = A^{\ell'} \wedge ((pc \sqcup \ell) \sqsubseteq \ell')$$

Since $\ell \not\sqsubseteq \mathcal{A}$. Therefore, $(pc \sqcup \ell) \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Also from (CC1) and (CC2), $(m_1 + 1, H'_{c1}) \triangleright \theta'_1$ and $(m_2 + 1, H'_{c2}) \triangleright \theta'_2$.

Therefore from Definition 1.8 we have

$$(\theta'_1, m_1, H'_{c1}(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V \text{ and}$$

$$(\theta'_2, m_2, H'_{c2}(a_1)) \in \lfloor \theta'_2(a_2) \rfloor_V$$

Since m_1 and m_2 are arbitrary indices therefore from Definition 1.4 we

get (here $H'_1 = H'_{c1}$ and $H'_2 = H'_{c2}$)

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

C. $H'_{j1}(a_1) = H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2)$:

$$\frac{W'.\theta_1(a_1) = W'.\theta_2(a_2):}{\text{Same as before}}$$

Same as before

$$\frac{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}{\text{From (CC2) we know that}}$$

From (CC2) we know that

$$(\forall a. H'_2(a) \neq H'_{c2}(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge ((pc \sqcup \ell) \sqsubseteq \ell')$$

This means that a_2 was protected at $(pc \sqcup \ell)$ in the world before the modification. Since $\ell \not\sqsubseteq \mathcal{A}$. Therefore, $(pc \sqcup \ell) \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 25 we know that $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$ that means from Definition 1.9 that $(W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$. Since $((pc \sqcup \ell) \sqsubseteq \ell'$ therefore from Definition 1.4 we know that $H'_1(a_1)$ must also be protected at some label $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_1.\theta_1, m, H'_1(a_1)) \in W'_1.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_1.\theta_2, m, H'_2(a_2)) \in W'_1.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with m_1 and using Lemma 1.15 we get

$$(\theta'_1, m_1, H'_1(a_1)) \in \theta'_1(a_1)$$

Since from (CC2) we know that $(m_2 + 1, H'_{c2}) \triangleright \theta'_2$ therefore from Definition 1.8 we know that $(\theta'_2, m_2, H'_{c2}(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 1.4 we get

$$(W', n - n' - 1, H'_{c1}(a_1), H'_{c2}(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

D. $H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) = H'_2(a_2)$:

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in \lfloor W'.\theta_i(a_i) \rfloor_V$$

$i = 1$

This means that given some m we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i).(W'.\theta_i, m, H'_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V$$

Like before we apply Theorem 1.21 on e_{i1} γ_1 and e_{i2} γ_2 but this time using $m + 1 + i$ and $m + 1 + j$ where i and j are the number of steps in which e_{i1} γ_1 and e_{i2} γ_2 reduces respectively. This will give us

$$\begin{aligned} \exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m + 1, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m + 1, v'_c) \in \llbracket (\tau) \rrbracket_V \wedge \\ (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

and

$$\begin{aligned} \exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m + 1, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m + 1, v'_c) \in \llbracket (\tau) \rrbracket_V \wedge \\ (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_2). \theta'_2(a) \searrow (pc \sqcup \ell)) \end{aligned}$$

Since we have $(m + 1, H'_{c1}) \triangleright \theta'_1$ and $(m + 1, H'_{c2}) \triangleright \theta'_2$ therefore we get the desired from Definition 1.8

$$i = 2$$

Symmetric to $i = 1$

$$\text{ii. } (W', n - n' - 1, v'_1, v'_2) \in \llbracket \tau_2 \rrbracket_V^A:$$

Let $\tau_2 = \mathbf{A}^{\ell_i}$ Since $\tau_2 \searrow \ell$ and since $\ell \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

From CC1 and CC2 we and Definition 1.4 we get the desired.

- (d) Reduction of e_1 happens via Case2 and Reduction of e_2 happens via Case1 :
Symmetric case as before

10. FG-ref:

$$\frac{\Gamma \vdash_{pc} e_i : \tau \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove: $(W, (\text{new } (e_i)) (\gamma \downarrow_1), (\text{new } (e_i)) (\gamma \downarrow_2)) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_E^A$

Say $e_1 = (\text{new } (e_i)) (\gamma \downarrow_1)$ and $e_2 = (\text{new } (e_i)) (\gamma \downarrow_2)$

From Definition of $\llbracket (\text{ref } \tau)^\perp \rrbracket_E^A$ it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_V^A \end{aligned}$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_V^A \quad (30)$$

$$\underline{\text{IH1}} (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in \llbracket \tau \rrbracket_E^A$$

This means from Definition 1.5 we get

$$\forall H_{i1}, H_{i2}.(n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau]_V^A$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $\text{ref}(e_i)$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore $\exists i < n' < n$. s.t $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$. Similarly since $\text{ref}(e_i)$ reduces with $\gamma \downarrow_2$ therefore we know that $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau]_V^A \quad (31)$$

From the evaluation rule of ref we know that $H'_1 = H'_{i1} \cup \{a_{n1} \mapsto v_{i1}\}$ and $H'_2 = H'_{i2} \cup \{a_{n2} \mapsto v_{i2}\}$

In order to prove Equation 30 we instantiate W' with W_n where W_n is

$$W_n.\theta_1 = W'_1.\theta_1 \cup \{a_{n1} \mapsto \tau\}$$

$$W_n.\theta_2 = W'_1.\theta_2 \cup \{a_{n2} \mapsto \tau\}$$

$$W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$$

Also we know that $n' = i + 1$

We are now required to prove

- $W \sqsubseteq W_n$:

From Equation 31 we know that $W \sqsubseteq W'_1$ and $W'_1 \sqsubseteq W_n$ by construction. Therefore from Definition 1.3, $W \sqsubseteq W_n$

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W_n$:

From Definition 1.9 it suffices to show that

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2):$$

From Equation 31 and by construction of W_n

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)):$$

From Equation 31 and by construction of W_n

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, n - n', H'_1(a_1), H'_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A:$$

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

From Equation 31 and by construction of W_n

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A:$$

From Equation 31 since we know that $(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$ that means

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in [W'_1.\theta_1(a_1)]_V^A$$

Therefore from Lemma 1.16 we get $(n - i - 2 = n - n' - 1, \text{ since } n' = i + 1)$

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 2, H'_1(a_1), H'_2(a_2)) \in [W'_1.\theta_1(a_1)]_V^A$$

Since $W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$ and from Equation 31 we know that $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau]_V^A$

Therefore combining the two we get

$$\forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

– $\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V$:

From Equation 31 we have $(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$ that means from Definition 1.9 we have

$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W'_1.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V$

Also from Equation 31 we know that $(W'_1, n - i, v'_{i1}, v'_{i2}) \in \llbracket \tau \rrbracket_V^A$

Therefore from Lemma 1.14 and Lemma 1.15 we get

$\forall m. (W'_1.\theta_1, m, v'_{i1}) \in \llbracket \tau \rrbracket_V$

and

$\forall m. (W'_1.\theta_2, m, v'_{i2}) \in \llbracket \tau \rrbracket_V$

Combining the two we get

$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V$

• $(W_n, n - n', v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_V^A$:

Here $v'_1 = a_{n1}$ and $v'_2 = a_{n2}$

Since $(a_{n1}, a_{n2}) \in W_n$ and also $W_n.\theta_1(a_{n1}) = W_n.\theta_1(a_{n1}) = \tau$

Therefore from Definition 1.4 $(W_n, v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\perp \rrbracket_V^A$

11. FG-deref:

$$\frac{\Gamma \vdash_{pc} e_i : (\text{ref } \tau)^\ell \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e_i : \tau'}$$

To prove: $(W, n, !(e_i) (\gamma \downarrow_1), !(e_i) (\gamma \downarrow_2)) \in \llbracket (\tau') \rrbracket_E^A$

Say $e_1 = !(e_i) (\gamma \downarrow_1)$ and $e_2 = !(e_i) (\gamma \downarrow_2)$

This means from Definition 1.5 we need to prove:

$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !(e_i)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !(e_i)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$

$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in \llbracket (\tau') \rrbracket_V^A$

This further means that given

$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !(e_i)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !(e_i)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in \llbracket (\tau') \rrbracket_V^A \quad (32)$$

IH1 $(W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in \llbracket (\text{ref } \tau)^\ell \rrbracket_E^A$

This means from Definition 1.5 we get

$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$

$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in \llbracket (\text{ref } \tau)^\ell \rrbracket_V^A$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $!(e_i)$ reduces to value with both $\gamma \downarrow_1$ in $n' < n$ steps therefore $\exists i < n' < n$ s.t $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1)$. Similarly since $!e_i$ reduces to value with $\gamma \downarrow_2$ therefore $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$. Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell]_V^A \quad (33)$$

We case analyze on $(W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell]_V^A$ from Equation 33

- Case $\ell \sqsubseteq \mathcal{A}$:

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)]_V^A$$

This means

$$(W'_1, n - i, v'_1, v'_2) \in [(\text{ref } (\tau))]_V^A$$

Let $v'_{i1} = a_{i1}$ and $v'_{i2} = a_{i2}$

Again from Definition 1.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1.\hat{\beta} \wedge W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau \quad (\text{D1})$$

In order to prove Equation 32 we instantiate W' with W'_1 . Also we know that $n' = i + 1$

- $W'_1 \sqsupseteq W$:

From Equation 33

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$:

From Equation 33 we know that

$$(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$$

Therefore from Lemma 1.20 we get

$$(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$$

- $(W'_1, n - n', v'_1, v'_2) \in [(\tau')]_V^A$:

From the evaluation rule of deref we know that $v'_1 = H'_1(a_{i1})$ and $v'_2 = H'_2(a_{i2})$

Since from Equation 33 we know that $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$, therefore from Definition 1.9 we know that

$$(W'_1, n - i - 1, H'_1(a_{i1}), H'_2(a_{i2})) \in [W'_1.\theta_1(a_{i1})]_V^A$$

And from D1 we know that $W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau$

Therefore $(W'_1, v'_1, v'_2) \in [(\tau)]_V^A$

Since $\tau <: \tau'$ Therefore from Lemma 1.27, we get

$$(W'_1, n - i - 1, v'_1, v'_2) \in [(\tau')]_V^A$$

- Case $\ell \not\sqsubseteq \mathcal{A}$:

From the evaluation rule of deref we know that $v'_{i1} = a_1$ and $v'_{i2} = a_2$

In this case from Definition 1.4 we know that

$$\forall m_1.(W'_1.\theta_1, m_1, a_1) \in [(\text{ref } \tau)]_V \quad (34)$$

and

$$\forall m_2.(W'_1.\theta_2, m_2, a_2) \in [(\text{ref } \tau)]_V \quad (35)$$

In order to prove Equation 32 we choose W' as W'_1 . And then we need to show:

- $W \sqsubseteq W'_1$:
Directly from Equation 33
- $(n - n', H'_1, H'_2) \triangleright^A W'_1$:
From Equation 33 we know that $(n - i, H'_1, H'_2) \triangleright^A W'_1$
Therefore from Lemma 1.20 we get
 $(n - i - 1, H'_1, H'_2) \triangleright^A W'_1$
- $(W'_1, n - n', v'_1, v'_2) \in \lceil \tau' \rceil_V^A$:
Let $\tau' = A^{\ell_i}$ Since $\tau' \searrow \ell$ and since $\ell \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

Therefore from Definition 1.4 it suffices to prove that

$$\forall m_1. (W'_1.\theta_1, m_1, v'_1) \in \lfloor \tau' \rfloor_V$$

and

$$\forall m_2. (W'_1.\theta_2, m_2, v'_2) \in \lfloor \tau' \rfloor_V$$

This means given m_1 and it suffices to prove:

$$(W'_1.\theta_1, m_1, v'_1) \in \lfloor \tau' \rfloor_V \quad (36)$$

Similarly given m_2 , it suffices to prove:

$$(W'_1.\theta_2, m_2, v'_2) \in \lfloor \tau' \rfloor_V \quad (37)$$

Since from Equation 33 we know that $(n - i, H'_1, H'_2) \triangleright W'_1$ therefore from Lemma 1.26 we get

$$\forall m_{h1}. (m_{h1}, H'_1) \triangleright W'_1.\theta_1 \quad (38)$$

$$\forall m_{h2}. (m_{h2}, H'_2) \triangleright W'_1.\theta_2 \quad (39)$$

Instantiating m_{h1} in Equation 38 with $m_1 + 1$ we get $(m_1, H'_1) \triangleright W'_1.\theta_1$

Therefore from Definition 1.8, we get

$$\forall a \in \text{dom}(W'_1.\theta_1). (W'_1.\theta_1, m_1, H'_1(a)) \in \lfloor W'_1.\theta_1(a) \rfloor_V$$

Instantiating a with a_1 we get $(W'_1.\theta_1, m_1, H'_1(a_1)) \in \lfloor W'_1.\theta_1(a) \rfloor_V$

Since $W'_1.\theta_1(a_{i1}) = \tau$ therefore we get

$$(W'_1.\theta_1, m_1, v'_1) \in \lfloor \tau \rfloor_V$$

and since $\tau <: \tau'$ therefore from Lemma 1.23 we get

$$(W'_1.\theta_1, m_1, v'_1) \in \lfloor \tau' \rfloor_V$$

Similarly we also get

$$(W'_1.\theta_2, m_2, v'_2) \in \lfloor \tau' \rfloor_V$$

Finally from Definition 1.4 we get

$$(W'_1, v'_1, v'_2) \in \lceil (\tau') \rceil_V^A$$

12. FG-assign:

$$\frac{\Gamma \vdash_{pc} e_{i1} : (\text{ref } \tau)^\ell \quad \Gamma \vdash_{pc} e_{i2} : \tau \quad \mathcal{L} \vdash \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_{i1} := e_{i2} : \text{unit}}$$

To prove: $(W, n, (e_{i1} := e_{i2}) (\gamma \downarrow_1), (e_{i1} := e_{i2}) (\gamma \downarrow_2)) \in [(\mathbf{unit})]_E^A$

Say $e_1 = (e_{i1} := e_{i2}) (\gamma \downarrow_1)$ and $e_2 = (e_{i1} := e_{i2}) (\gamma \downarrow_2)$

This means from Definition 1.5 we need to prove:

$$\begin{aligned} & \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies \\ & \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\mathbf{unit})]_V^A \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\mathbf{unit})]_V^A \quad (40)$$

IH1 $(W, n, (e_{i1}) (\gamma \downarrow_1), (e_{i1}) (\gamma \downarrow_2)) \in [(\mathbf{ref} \ \tau)^\ell]_E^A$

This means from Definition 1.5 we get

$$\begin{aligned} & \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies \\ & \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\mathbf{ref} \ \tau)^\ell]_V^A \end{aligned}$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $(e_{i1} := e_{i2})$ reduces to value with both $\gamma \downarrow_1$ in $n' < n$ steps therefore $\exists i < n' < n$ s.t $(H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow (H'_{i1}, v'_{i1})$. Similarly since $(e_{i1} := e_{i2})$ reduces to value with $\gamma \downarrow_2$ therefore we also have $(H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\mathbf{ref} \ \tau)^\ell]_V^A \quad (41)$$

IH2 $(W, n - i, (e_{i2}) (\gamma \downarrow_1), (e_{i2}) (\gamma \downarrow_2)) \in [(\tau)]_E^A$

This means from Definition 1.5 we get

$$\begin{aligned} & \forall H_{j1}, H_{j2}. (n - i, H_{j1}, H_{j2}) \triangleright^A W'_1 \wedge \forall j < n - i. (H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies \\ & \exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \triangleright^A W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau)]_V^A \end{aligned}$$

Instantiating H_{j1} with H'_{i1} and H_{j2} with H'_{i2} in IH2 and since the $(e_{i1} := e_{i2})$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps and e_1 reduces $\gamma \downarrow_1$ with $i < n'$ steps therefore $\exists j < (n' - i) < (n - i)$ s.t $(H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{j1}, v'_{j1})$. Similarly we also have $(H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$. Hence we get

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \triangleright^A W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau)]_V^A \quad (42)$$

We case analyze on $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\mathbf{ref} \ \tau)^\ell]_V^A$ from Equation 41

- Case $\ell \sqsubseteq \mathcal{A}$:

From Definition 1.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau)]_V^{\mathcal{A}}$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } (\tau))]_V^{\mathcal{A}}$$

Let $v'_{i1} = a_{i1}$ and $v'_{i2} = a_{i2}$

Again from Definition 1.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1.\hat{\beta} \wedge W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau \quad (\text{A1})$$

In order to prove Equation 40 we instantiate W' with W'_2

- $W'_2 \sqsupseteq W$:

Since $W'_1 \sqsupseteq W$ from Equation 41 and $W'_2 \sqsupseteq W'_1$ from Equation 42

Therefore from Definition 1.3 we get $W'_2 \sqsupseteq W$

- $(n - n', H'_1, H'_2) \triangleright^{\mathcal{A}} W'_2$:

From the evaluation rule assign we know that

$$H'_1 = H'_{j1}[a_{i1} \mapsto v'_{j1}] \text{ and } H'_2 = H'_{j2}[a_{i2} \mapsto v'_{j2}]$$

In order to prove $(n - n', H'_1, H'_2) \triangleright^{\mathcal{A}} W'_2$ we need to show:

$$* \text{ dom}(W'_2.\theta_1) \subseteq \text{ dom}(H'_1) \wedge \text{ dom}(W'_2.\theta_2) \subseteq \text{ dom}(H'_2):$$

Directly from Equation 42

$$* W'_2.\hat{\beta} \subseteq (\text{ dom}(W'_2.\theta_1) \times \text{ dom}(W'_2.\theta_2)):$$

Directly from Equation 42

$$* \forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge \\ (W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'_2.\theta_1(a_1)]_V^{\mathcal{A}}:$$

$$(a) \forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2):$$

$$\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$$

- i. When $a_1 = a_{i1}$ and $a_2 = a_{i2}$:

From A1 we know that $W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$

and since $W'_1 \sqsubseteq W'_2$ therefore from Lemma 1.15 we get $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$

- ii. When $a_1 = a_{i1}$ and $a_2 \neq a_{i2}$: This case cannot arise

- iii. When $a_1 \neq a_{i1}$ and $a_2 = a_{i2}$: This case cannot arise

- iv. When $a_1 \neq a_{i1}$ and $a_2 \neq a_{i2}$: From Equation 42 and Lemma 1.16

$$(b) \forall (a_1, a_2) \in (W'_2.\hat{\beta}). (W'_2, n - n', H'_1(a_1), H'_2(a_2)) \in [W'_2.\theta_1(a_1)]_V^{\mathcal{A}}:$$

$$\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$$

- i. When $a_1 = a_{i1}$ and $a_2 = a_{i2}$:

Since $H'_1(a_{i1}) = v'_{j1}$ and $H'_1(a_{i2}) = v'_{j2}$

From A1 we know that $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$

And since from Equation 42 we know that $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau)]_V^{\mathcal{A}}$

Therefore from Lemma 1.16 we get

$$(W'_2, n - j - i - 1, H'_1(a_1), H'_2(a_2)) \in [W'_2.\theta_1(a_1)]_V^{\mathcal{A}}$$

- ii. When $a_1 = a_{i1}$ and $a_2 \neq a_{i2}$: This case cannot arise

- iii. When $a_1 \neq a_{i1}$ and $a_2 = a_{i2}$: This case cannot arise

- iv. When $a_1 \neq a_{i1}$ and $a_2 \neq a_{i2}$: From Equation 42 and from Lemma 1.16

* $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'_2.\theta_i). (W'_2.\theta_i, m, H'_i(a_i)) \in \llbracket W'_2.\theta_i(a_i) \rrbracket_V$:

When $i = 1$

Given some m

$\forall a_1 \in \text{dom}(W'_2.\theta_1)$.

· when $a_1 = a_{i1}$:

From Equation 42 we know that $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in \llbracket (\tau) \rrbracket_V^A$ thus from Lemma 1.14 we know that

$\forall m_1. (W'_2.\theta_1, m_1, H'_1(a_1)) \in \llbracket W'_2.\theta_1(a_1) \rrbracket_V$

Instantiating with m we get

$(W'_2.\theta_1, m, H'_1(a_1)) \in \llbracket W'_2.\theta_1(a_1) \rrbracket_V$

· Otherwise:

From Equation 42 and Lemma 1.26

When $i = 2$

Similar reasoning as with $i = 1$

– $(W'_1, n - n', \text{val}'_1, v'_2) \in \llbracket (\text{unit}) \rrbracket_V^A$:

From evaluation rule assign we know that $v'_1 = v'_2 = ()$

Directly from Definition 1.4

• Case $\ell \not\sqsubseteq \mathcal{A}$:

From Definition 1.4 we know that this would mean that

$$\forall m_1. (W'_1.\theta_1, m_1, a_{i1}) \in \llbracket (\text{ref } \tau) \rrbracket_V \quad (43)$$

$$\forall m_2. (W'_1.\theta_2, m_2, a_{i2}) \in \llbracket (\text{ref } \tau) \rrbracket_V \quad (44)$$

In order to prove Equation 40 we instantiate W' with W'_2 and then we need to show that:

– $W'_2 \sqsupseteq W$:

Since $W'_1 \sqsupseteq W$ from Equation 41 and $W'_2 \sqsupseteq W'_1$ from Equation 42

Therefore from Definition 1.3 we get $W'_2 \sqsupseteq W$

– $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$:

From the evaluation rule assign we know that

$H'_1 = H'_{j1}[a_{i1} \mapsto v'_{j1}]$ and $H'_2 = H'_{j2}[a_{i2} \mapsto v'_{j2}]$

In order to prove $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$ we need to show:

* $\text{dom}(W'_2.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'_2.\theta_2) \subseteq \text{dom}(H'_2)$:

Directly from Equation 42

* $W'_2.\hat{\beta} \subseteq (\text{dom}(W'_2.\theta_1) \times \text{dom}(W'_2.\theta_2))$:

Directly from Equation 42

* $\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \llbracket W'_2.\theta_1(a_1) \rrbracket_V^A$:

(a) When $(a_{i1}, a_{i2}) \in W'_2.\hat{\beta}$:

$\forall (a_1, a_2) \in (W'_2.\hat{\beta})$.

i. When $a_1 = a_{i1}$ and $a_2 = a_{i2}$:

Instantiating Equation 43 and Equation 44 with $n - n' - 1$ we get

$W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$

and since $W'_1 \sqsubseteq W'_2$ therefore from Definition 1.3 we get $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$

From Equation 42 we know that $(W'_2, v'_{j_1}, v'_{j_2}) \in [(\tau)]_V^A$

Therefore $(W'_2, H'_1(a_{i_1}), H'_2(a_{i_2})) \in [(\tau)]_V^A$

- ii. When $a_1 = a_{i_1}$ and $a_2 \neq a_{i_2}$: This case cannot arise
 - iii. When $a_1 \neq a_{i_1}$ and $a_2 = a_{i_2}$: This case cannot arise
 - iv. When $a_1 \neq a_{i_1}$ and $a_2 \neq a_{i_2}$: From Equation 42
- (b) When $(a_{i_1}, a_{i_2}) \notin W'_2.\hat{\beta}$:

$\forall (a_1, a_2) \in (W'_2.\hat{\beta})$.

i. When $a_1 = a_{i_1}$ and $a_2 = a_{i_2}$: This case cannot arise

ii. When $a_1 = a_{i_1}$ and $a_2 \neq a_{i_2}$:

From Equation 42 we know that $(n - i - j, H'_{j_1}, H'_{j_2}) \stackrel{A}{\triangleright} W'_2$ and since $(a_{i_1}, a_2) \in W'_2.\hat{\beta}$ therefore from Definition 1.9 we know that

$$(W'_2.\theta_1(a_{i_1}) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j_1}(a_{i_1}), H'_{j_2}(a_2)) \in [W'_2.\theta_1(a_{i_1})]_V^A) \quad (45)$$

Instantiating Equation 43 and Equation 44 with $n - i - j - 1$ we get $W'_1.\theta_1(a_{i_1}) = \tau$ therefore from monotonicity we also have $W'_2.\theta_1(a_{i_1}) = \tau$. As a result from Equation 45 we get $W'_2.\theta_2(a_2) = \tau$

Also since from Equation 45 $(W'_2, n - i - j - 1, H'_{j_1}(a_{i_1}), H'_{j_2}(a_2)) \in [(\tau)]_V^A$ and $\tau \searrow \ell, \ell \not\sqsubseteq \mathcal{A}$ therefore from Lemma 1.14 we know that

$$\forall m.(W'_2.\theta_1, m, H'_{j_1}(a_{i_1})) \in [(\tau)]_V \quad (46)$$

$$\forall m.(W'_2.\theta_2, m, H'_{j_2}(a_2)) \in [(\tau)]_V \quad (47)$$

Instantiating m with $n - i - j - 1$ in Equation 46 and Equation 47 to get

$$(W'_2.\theta_1, n - i - j - 1, H'_{j_1}(a_{i_1})) \in [(\tau)]_V$$

and

$$(W'_2.\theta_2, n - i - j - 1, H'_{j_2}(a_2)) \in [(\tau)]_V$$

Since $H'_1(a_{i_1}) = v'_{j_1}$ and $H'_2(a_2) = H'_{j_2}(a_2)$

Again from Equation 42 we know that $(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau)]_V^A$. This means from Lemma 1.14 and instantiating it with $n - i - j - 1$ we get

$$(W'_2.\theta_1, n - i - j - 1, v'_{j_1}) \in [(\tau)]_V \quad (48)$$

Therefore from Equation 47 and Equation 48 we have $(W'_2, n - i - j - 1, H'_1(a_{i_1}), H'_2(a_2)) \in [(\tau)]_V^A$

iii. When $a_1 \neq a_{i_1}$ and $a_2 = a_{i_2}$:

Symmetric case as (ii)

iv. When $a_1 \neq a_{i_1}$ and $a_2 \neq a_{i_2}$:

From Equation 42 and Definition 1.9

* $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'_2.\theta_i). (W'_2.\theta_i, m, H'_i(a_i)) \in [W'_2.\theta_i(a_i)]_V$:
When $i = 1$
 Given some m
 $\forall a_1 \in \text{dom}(W'_2.\theta_1)$.
 · when $a_1 = a_{i1}$:
 From Equation 42 we know that $(W'_2, v'_{j1}, v'_{j2}) \in [(\tau)]_V^A$ thus from Lemma 1.14 we know that
 $(W'_2.\theta_1, H'_1(a_1)) \in [W'_2.\theta_1(a_1)]_V$
 · Otherwise:
 From Equation 42 and Lemma 1.26
When $i = 2$
 Similar reasoning as with $i = 1$
 – $(W'_1, n - n', v'_1, v'_2) \in [(\text{unit})]_V^A$:
 From evaluation rule assign we know that $v'_1 = v'_2 = ()$
 Directly from Definition 1.4

□

Lemma 1.26 (Binary heap well formedness implies unary heap well formedness). $\forall H_1, H_2, W. (n, H_1, H_2) \triangleright W \implies \forall i \in \{1, 2\}. \forall m. (m, H_i) \triangleright W.\theta_i$

Proof. Directly from Definition 1.9

□

Lemma 1.27 (Subtyping binary). *The following holds:*

1. $\forall A, A'$.

$$(a) \mathcal{L} \vdash A <: A' \implies [(A)]_V^A \subseteq [(A')]_V^A$$

2. $\forall \tau, \tau'$.

$$(a) \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_V^A \subseteq [(\tau')]_V^A$$

$$(b) \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_E^A \subseteq [(\tau')]_E^A$$

Proof. Proof by simultaneous induction on $A <: A'$ and $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of A in the last step:

1. FGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2 \quad \mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove: $[((\tau_1 \xrightarrow{\ell_e} \tau_2))]_V^A \subseteq [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2))]_V^A$

IH1: $[(\tau'_1)]_V^A \subseteq [(\tau_1)]_V^A$

IH2: $[(\tau_2)]_E^A \subseteq [(\tau'_2)]_E^A$

It suffices to prove:

$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2))]_{\mathcal{V}}^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2))]_{\mathcal{V}}^A$

This means that given: $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2))]_{\mathcal{V}}^A$

And it suffices to prove: $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2))]_{\mathcal{V}}^A$

From Definition 1.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_{\mathcal{V}}^A \implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_{\mathcal{E}}^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} \implies (\theta_l, j, e_1[v_1/x]) \in [\tau_2]_{\mathcal{E}}^{\ell_e}) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_{\mathcal{E}}^{\ell_e}) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 1.4 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ [\tau'_2]_{\mathcal{E}}^A) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}) \end{aligned}$$

This means given some $W'' \sqsupseteq W, k < n$ and v'_1, v'_2 we need to prove:

$$(a) \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_{\mathcal{E}}^A) :$$

Given: $W'' \sqsupseteq W, k < n$ and v'_1, v'_2 . We are also given $(W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A$

To prove: $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_{\mathcal{E}}^A$

Instantiating the first conjunct of Sub-A1 with W'', k, v'_1 and v'_2 we get

$$((W'', k, v'_1, v'_2) \in [\tau_1]_{\mathcal{V}}^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_{\mathcal{E}}^A) \quad (49)$$

Since $(W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A$ therefore from IH1 we know that $(W'', k, v'_1, v'_2) \in [\tau_1]_{\mathcal{V}}^A$

Thus from Equation 49 we get $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_{\mathcal{E}}^A$

Finally using IH2 we get $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_{\mathcal{E}}^A$

$$(b) \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}) :$$

Given: $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$. We are also given $(\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}}$

To prove: $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}$

Since we are given $(\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}}$ and since $\tau'_1 <: \tau_1$ therefore from Lemma 1.23 we get

$$(\theta'_l, k, v'_c) \in [\tau_1]_{\mathcal{V}} \quad (50)$$

Instantiating the second conjunct of Sub-A1 with θ'_l, k, v'_1 and v'_2 we get

$$((\theta'_l, k, v'_c) \in [\tau_1]_{\mathcal{V}} \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2]_{\mathcal{E}}^{\ell_e}) \quad (51)$$

Therefore from Equation 50 and 51 we get $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_{\mathcal{E}}^{\ell_e}$

Since $\tau_2 <: \tau'_2$ and $\ell'_e \sqsubseteq \ell_e$ therefore from Lemma 1.23 and 1.22 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_{\mathcal{E}}^{\ell'_e}$$

(c) $\forall \theta'_i \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_i, k, v'_c) \in \lceil \tau'_1 \rceil_V \implies (\theta'_i, k, e_2[v'_c/x]) \in \lceil \tau'_2 \rceil_E^{\ell'_e})$:

Similar reasoning as in the previous case

2. FGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove: $\lceil ((\tau_1 \times \tau_2)) \rceil_V^A \subseteq \lceil ((\tau'_1 \times \tau'_2)) \rceil_V^A$

IH1: $\lceil (\tau_1) \rceil_V^A \subseteq \lceil (\tau'_1) \rceil_V^A$

IH2: $\lceil (\tau_2) \rceil_V^A \subseteq \lceil (\tau'_2) \rceil_V^A$

It suffices to prove: $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau_1 \times \tau_2)) \rceil_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau'_1 \times \tau'_2)) \rceil_V^A$

This means that given: $(W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau_1 \times \tau_2)) \rceil_V^A$

Therefore from Definition 1.4 we are given:

$$(W, n, v_1, v'_1) \in \lceil \tau_1 \rceil_V^A \wedge (W, n, v_2, v'_2) \in \lceil \tau_2 \rceil_V^A \quad (52)$$

And it suffices to prove: $(W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau'_1 \times \tau'_2)) \rceil_V^A$

Again from Definition 1.4, it suffices to prove:

$$(W, n, v_1, v'_1) \in \lceil \tau'_1 \rceil_V^A \wedge (W, n, v_2, v'_2) \in \lceil \tau'_2 \rceil_V^A$$

Since from Equation 52 we know that $(W, n, v_1, v'_1) \in \lceil \tau_1 \rceil_V^A$ therefore from IH1 we have $(W, n, v_1, v'_1) \in \lceil \tau'_1 \rceil_V^A$

Similarly since $(W, n, v_2, v'_2) \in \lceil \tau_2 \rceil_V^A$ from Equation 52 therefore from IH2 we have $(W, n, v_2, v'_2) \in \lceil \tau'_2 \rceil_V^A$

3. FGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove: $\lceil ((\tau_1 + \tau_2)) \rceil_V^A \subseteq \lceil ((\tau'_1 + \tau'_2)) \rceil_V^A$

IH1: $\lceil (\tau_1) \rceil_V^A \subseteq \lceil (\tau'_1) \rceil_V^A$

IH2: $\lceil (\tau_2) \rceil_V^A \subseteq \lceil (\tau'_2) \rceil_V^A$

It suffices to prove: $\forall (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2)) \rceil_V^A. (W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2)) \rceil_V^A$

This means that given: $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau_1 + \tau_2)) \rceil_V^A$

And it suffices to prove: $(W, n, v_{s1}, v_{s2}) \in \lceil ((\tau'_1 + \tau'_2)) \rceil_V^A$

2 cases arise

- (a) $v_{s1} = \text{inl } v_{i1}$ and $v_{s1} = \text{inl } v_{i2}$:
From Definition 1.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_1 \rceil_V^A \quad (53)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_1' \rceil_V^A$$

From Equation 53 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_1' \rceil_V^A$$

- (b) $v_s = \text{inr } v_{i1}$ and $v_{s2} = \text{inr } v_{i2}$:
From Definition 1.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_2 \rceil_V^A \quad (54)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_2' \rceil_V^A$$

From Equation 54 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil \tau_2' \rceil_V^A$$

4. FGsub-ref:

Given:

$$\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove: $\lceil ((\text{ref } \tau)) \rceil_V^A \subseteq \lceil ((\text{ref } \tau)) \rceil_V^A$

Directly from Definition 1.4

5. FGsub-base:

Given:

$$\frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove: $\lceil ((\mathbf{b})) \rceil_V^A \subseteq \lceil ((\mathbf{b})) \rceil_V^A$

Directly from Definition 1.4

6. FGsub-unit:

Given:

$$\frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove: $\lceil ((\text{unit})) \rceil_V^A \subseteq \lceil ((\text{unit})) \rceil_V^A$

Directly from Definition 1.4

Proof of statement 2(a)

Given:

$$\frac{\mathcal{L} \vdash \ell \subseteq \ell' \quad \mathcal{L} \vdash A <: A'}{\mathcal{L} \vdash A^\ell <: A^{\ell'}} \text{FGsub-label}$$

To prove: $\lceil ((A^\ell)) \rceil_V^A \subseteq \lceil ((A^{\ell'})) \rceil_V^A$

2 cases arise

1. $\ell \sqsubseteq \ell'$:

From Definition 1.4 it suffices to prove: $\lceil ((A)) \rceil_V^A \subseteq \lceil ((A')) \rceil_V^A$

This we get directly from IH (Statement (1))

2. $\ell \not\sqsubseteq \ell'$:

We need to prove that

$$\forall (W, n, v_1, v_2) \in \lceil A \rceil_V^A. (W, n, v_1, v_2) \in \lceil A' \rceil_V^A$$

From Definition 1.4 it suffices to prove:

$$\forall i \in \{1, 2\}. \forall m. (W(n). \theta_i, m, v_i) \in \lceil A \rceil_V. (W(n). \theta_i, m, v_i) \in \lceil A \rceil_V \in \lceil A' \rceil_V$$

Since $A <: A'$ therefore from Lemma 1.23 we get the desired

Proof of statement 2(b)

Given: $\mathcal{L} \vdash \tau <: \tau'$

To prove: $\lceil (\tau) \rceil_E^A \subseteq \lceil (\tau') \rceil_E^A$

This means we need to prove that

$$\forall (W, n, e_1, e_2) \in \lceil (\tau) \rceil_E^A. (W, n, e_1, e_2) \in \lceil (\tau') \rceil_E^A$$

This means given $\forall (W, n, e_1, e_2) \in \lceil (\tau) \rceil_E^A$

It suffices to prove that $(W, n, e_1, e_2) \in \lceil (\tau') \rceil_E^A$

From Definition 1.5 we know we are given:

$$\begin{aligned} & \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in \lceil \tau \rceil_V^A \quad (\text{Sub-exp1}) \end{aligned}$$

And we need prove that

$$\begin{aligned} & \forall H_{21}, H_{22}, k < n. (n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22}) \implies \\ \exists W'' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in \lceil \tau \rceil_V^A \end{aligned}$$

This means that we are given some H_{21}, H_{22} and $k < n$ such that $(n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22})$

It suffices to prove:

$$\exists W'' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in \lceil \tau \rceil_V^A \quad (55)$$

Instantiating (Sub-exp1) with H_{21}, H_{22} and k we get

$$\exists W' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W' \wedge (W', n - k, v'_{21}, v'_{22}) \in \lceil \tau \rceil_V^A \quad (56)$$

We choose W'' in Equation 55 as W' from Equation 56 and we are done

□

Theorem 1.28 (NI for FG). *Say* $\text{bool} = (\text{unit} + \text{unit})$

$\forall v_1, v_2, e, \tau, n_1.$

$\emptyset \vdash_{\perp} v_1 : \text{bool}^{\top} \wedge \emptyset \vdash_{\perp} v_2 : \text{bool}^{\top}$

$x : \text{bool}^{\top} \vdash_{\perp} e : \text{bool}^{\perp} \wedge$

$(\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_{-} (-, v'_2) \implies$

$v'_1 = v'_2$

Proof. Given some

$$\begin{aligned} & \emptyset \vdash_{\perp} v_1 : \mathbf{bool}^{\top} \wedge \emptyset \vdash_{\perp} v_2 : \mathbf{bool}^{\top} \\ & x : \mathbf{bool}^{\top} \vdash_{\perp} e : \mathbf{bool}^{\perp} \wedge \\ & (\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow (-, v'_2) \end{aligned}$$

We need to prove

$$\overline{v'_1 = v'_2}$$

From Theorem 1.25 we have

$$\forall n. (\emptyset, n, v_1, v_2) \in \lceil \mathbf{bool}^{\top} \rceil_E^{\perp}$$

Therefore from Theorem 1.25 and from Definition 1.13 we have

$$\forall n. (\emptyset, n, e[v_1/x], e[v_1/x]) \in \lceil \mathbf{bool}^{\perp} \rceil_E^{\perp}$$

Therefore from Definition 1.5 we know that

$$\begin{aligned} & \forall n. (\forall H_1, H_2, j < n. (n, H_1, H_2) \overset{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \exists W' \sqsupseteq \\ & W.(n - j, H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in \lceil (\mathbf{unit} + \mathbf{unit})^{\perp} \rceil_V^A) \end{aligned}$$

Instantiating with $n_1 + 1$ and then with $\emptyset, \emptyset, n_1$ we get

$$\exists W' \sqsupseteq W.(1, H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', 1, v'_1, v'_2) \in \lceil (\mathbf{unit} + \mathbf{unit})^{\perp} \rceil_V^A$$

Since we have $(W', 1, v'_1, v'_2) \in \lceil (\mathbf{unit} + \mathbf{unit})^{\perp} \rceil_V^A$ therefore from Definition 1.4 we get $v'_1 = v'_2$ \square

1.2 Coarse-grained IFC enforcement (CG)

1.2.1 CG type system

Term, type, constraint syntax:

Expressions	$e ::= x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, y.e) \mid \text{new } e \mid !e \mid e := e \mid () \mid \text{Lb}(e) \mid \text{unlabel}(e) \mid \text{toLabeled}(e) \mid \text{ret}(e) \mid \text{bind}(e, x.e)$
Labels	$\ell ::= \perp \mid \top \mid l \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
Types	$\tau ::= \mathbf{b} \mid \text{unit} \mid \tau \rightarrow \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \ell \tau \mid \text{Labeled } \ell \tau \mid \mathbb{C} \ell_1 \ell_2 \tau$

Type system: $\boxed{\Gamma \vdash e : \tau}$

(All rules of the simply typed lambda-calculus pertaining to the types $\mathbf{b}, \tau \rightarrow \tau, \tau \times \tau, \tau + \tau, \text{unit}$ are included.)

$$\begin{array}{c}
\frac{\Gamma \vdash e : \tau}{\Gamma \vdash \text{Lb}(e) : \text{Labeled } \ell \tau} \text{CG-label} \qquad \frac{\Gamma \vdash e : \text{Labeled } \ell \tau}{\Gamma \vdash \text{unlabel}(e) : \mathbb{C} \top \ell \tau} \text{CG-unlabel} \\
\\
\frac{\Gamma \vdash e : \mathbb{C} \ell \ell' \tau}{\Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell \perp (\text{Labeled } \ell' \tau)} \text{CG-toLabeled} \qquad \frac{\Gamma \vdash e : \tau}{\Gamma \vdash \text{ret}(e) : \mathbb{C} \ell \ell' \tau} \text{CG-ret} \\
\\
\frac{\Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \quad \Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau'} \text{CG-bind} \\
\\
\frac{\Gamma \vdash e : \tau' \quad \mathcal{L} \vdash \tau' <: \tau}{\Gamma \vdash e : \tau} \text{CG-sub} \qquad \frac{\Gamma \vdash e : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } e : \mathbb{C} \ell \perp (\text{ref } \ell' \tau)} \text{CG-ref} \\
\\
\frac{\Gamma \vdash e : \text{ref } \ell' \tau}{\Gamma \vdash !e : \mathbb{C} \top \perp (\text{Labeled } \ell' \tau)} \text{CG-deref} \\
\\
\frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit}} \text{CG-assign}
\end{array}$$

Figure 4: Type system of CG.

1.2.2 CG semantics

Judgement: $e \Downarrow_i v$ and $(H, e) \Downarrow_i^f (H', v)$

1.2.3 Logical relation for CG

$W : ((\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \mapsto \text{Type})) \times (\text{Loc} \leftrightarrow \text{Loc})$

Definition 1.29 (θ_2 extends θ_1). $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

$$\begin{array}{c}
\frac{}{\mathcal{L} \vdash \tau <: \tau} \text{CGsub-refl} \qquad \frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2} \text{CGsub-arrow} \\
\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{CGsub-prod} \qquad \frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{CGsub-sum} \\
\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'} \text{CGsub-labeled} \\
\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i \quad \mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o}{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'} \text{CGsub-monad}
\end{array}$$

Figure 5: CG subtyping

$$\begin{array}{c}
\frac{e_1 \Downarrow_i \lambda x. e_i \quad e_2 \Downarrow_j v_2 \quad e_i[v_2/x] \Downarrow_k v_3}{e_1 e_2 \Downarrow_{i+j+k+1} v_3} \text{cg-app} \qquad \frac{e_1 \Downarrow_i v_1 \quad e_2 \Downarrow_j v_2}{(e_1, e_2) \Downarrow_{i+j+1} (v_1, v_2)} \text{cg-prod} \\
\frac{e \Downarrow_i (v_1, v_2)}{\text{fst}(e) \Downarrow_{i+1} v_1} \text{cg-fst} \qquad \frac{e \Downarrow_i (v_1, v_2)}{\text{snd}(e) \Downarrow_{i+1} v_2} \text{cg-snd} \qquad \frac{e \Downarrow_i v}{\text{inl}(e) \Downarrow_{i+1} \text{inl}(v)} \text{cg-inl} \\
\frac{e \Downarrow_i v}{\text{inr}(e) \Downarrow_{i+1} \text{inr}(v)} \text{cg-inr} \qquad \frac{e \Downarrow_i \text{inl } v \quad e_1[v/x] \Downarrow_j v_1}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_1} \text{cg-case1} \\
\frac{e \Downarrow_i \text{inr } v \quad e_2[v/x] \Downarrow_j v_2}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_2} \text{cg-case2} \qquad \frac{e \Downarrow_i v}{\text{Lb}(e) \Downarrow_{i+1} \text{Lb}(v)} \text{cg-Lb} \qquad \frac{e \Downarrow_i v}{(H, \text{ret}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-ret} \\
\frac{e_1 \Downarrow_i v_1 \quad (H, v_1) \Downarrow_j^f (H', v'_1) \quad e_2[v'_1/x] \Downarrow_k v_2 \quad (H', v_2) \Downarrow_l^f (H'', v'_2)}{(H, \text{bind}(e_1, x.e_2)) \Downarrow_{i+j+k+l+1}^f (H'', v'_2)} \text{cg-bind} \\
\frac{e \Downarrow_i \text{Lb}(v)}{(H, \text{unlabel}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-unlabel} \qquad \frac{e \Downarrow_i v \quad (H, v) \Downarrow_j^f (H', v')}{(H, \text{toLabeled}(e)) \Downarrow_{i+j+1}^f (H', \text{Lb}(v'))} \text{cg-toLabeled} \\
\frac{e \Downarrow_i \text{Lb}v \quad a \notin \text{dom}(H)}{(H, \text{new } (e)) \Downarrow_{i+1}^f (H[a \mapsto \text{Lb}v], a)} \text{cg-ref} \qquad \frac{e \Downarrow_i a}{(H, !e) \Downarrow_{i+1}^f (H, H(a))} \text{cg-deref} \\
\frac{e_1 \Downarrow_i a \quad e_2 \Downarrow_j \text{Lb}v}{(H, e_1 := e_2) \Downarrow_{i+j+1}^f (H[a \mapsto \text{Lb}v], ())} \text{cg-assign} \\
\frac{e \in \{x, \lambda y. -, \text{ret } -, \text{bind}(-, -, -), \text{unlabel}(-), \text{toLabeled}(-), \text{new } (-), !-, - := -\}}{e \Downarrow_0 e} \text{cg-val}
\end{array}$$

Figure 6: CG semantics

Definition 1.30 (W_2 extends W_1). $W_1 \sqsubseteq W_2 \triangleq$

1. $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2. $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

Definition 1.31 (Value Equivalence).

$$\text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau) \triangleq \begin{cases} (W, n, v_1, v_2) \in [\tau]_V^{\mathcal{A}} & \ell \sqsubseteq \mathcal{A} \\ \forall j. (W.\theta_1, j, v_1) \in [\tau]_V \wedge (W.\theta_2, j, v_2) \in [\tau]_V & \ell \not\sqsubseteq \mathcal{A} \end{cases}$$

Definition 1.32 (Binary value relation).

$$\begin{aligned} [\mathbf{b}]_V^{\mathcal{A}} &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\ [\mathbf{unit}]_V^{\mathcal{A}} &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{unit} \rrbracket\} \\ [\tau_1 \times \tau_2]_V^{\mathcal{A}} &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^{\mathcal{A}} \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^{\mathcal{A}}\} \\ [\tau_1 + \tau_2]_V^{\mathcal{A}} &\triangleq \{(W, n, \text{inl } v, \text{inl } v') \mid (W, n, v, v') \in [\tau_1]_V^{\mathcal{A}}\} \cup \\ &\quad \{(W, n, \text{inr } v, \text{inr } v') \mid (W, n, v, v') \in [\tau_2]_V^{\mathcal{A}}\} \\ [\tau_1 \rightarrow \tau_2]_V^{\mathcal{A}} &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\ &\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\ &\quad ((W', j, v_1, v_2) \in [\tau_1]_V^{\mathcal{A}} \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^{\mathcal{A}}) \wedge \\ &\quad \forall \theta_l \sqsupseteq W.\theta_1, v_c, j. \\ &\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E) \wedge \\ &\quad \forall \theta_l \sqsupseteq W.\theta_2, v_c, j. \\ &\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)\} \\ [\text{ref } \ell \tau]_V^{\mathcal{A}} &\triangleq \{(W, n, a_1, a_2) \mid \\ &\quad (a_1, a_2) \in W.\hat{\beta} \wedge W.\theta_1(a_1) = W.\theta_2(a_2) = \text{Labeled } \ell \tau\} \\ [\text{Labeled } \ell \tau]_V^{\mathcal{A}} &\triangleq \{(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \mid \text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau)\} \\ [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\mathcal{A}} &\triangleq \{(W, n, v_1, v_2) \mid \\ &\quad (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \\ &\quad \forall v'_1, v'_2, j. (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ &\quad \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\ &\quad \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\quad \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ &\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ &\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1))\} \end{aligned}$$

Definition 1.33 (Binary expression relation).

$$[\tau]_E^{\mathcal{A}} \triangleq \{(W, n, e_1, e_2) \mid \forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow_i v_2 \implies (W, n - i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}\}$$

Definition 1.34 (Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
\llbracket \mathbf{unit} \rrbracket_V &\triangleq \{(\theta, m, v \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V \wedge (\theta, m, v_2) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \mathbf{inl} \ v) \mid (\theta, m, v) \in \llbracket \tau_1 \rrbracket_V\} \cup \{(\theta, m, \mathbf{inr} \ v) \mid (\theta, m, v) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta' \sqsupseteq \theta, v, j < m. (\theta', j, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta', j, e[v/x]) \in \llbracket \tau_2 \rrbracket_E\} \\
\llbracket \mathbf{ref} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \mathbf{Labeled} \ \ell \ \tau\} \\
\llbracket \mathbf{Labeled} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, \mathbf{Lb}(v) \mid (\theta, m, v) \in \llbracket \tau \rrbracket_V\} \\
\llbracket \mathbb{C} \ \ell_1 \ \ell_2 \ \tau \rrbracket_V &\triangleq \{(\theta, m, e) \mid \\
&\quad \forall k \leq m, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k \implies \\
&\quad \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \mathbf{Labeled} \ \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)\}
\end{aligned}$$

Definition 1.35 (Unary expression relation).

$$\llbracket \tau \rrbracket_E \triangleq \{(\theta, n, e) \mid \forall i < n. e \Downarrow_i v \implies (\theta, n - i, v) \in \llbracket \tau \rrbracket_V\}$$

Definition 1.36 (Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in \llbracket \theta(a) \rrbracket_V$$

Definition 1.37 (Binary heap well formedness).

$$\begin{aligned}
(n, H_1, H_2) \overset{A}{\triangleright} W &\triangleq \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\
&\quad (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\
&\quad \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2) \wedge \\
&\quad (W, n - 1, H_1(a_1), H_2(a_2)) \in \llbracket W.\theta_1(a_1) \rrbracket_V^A) \wedge \\
&\quad \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V
\end{aligned}$$

Definition 1.38 (Binary substitution). $\gamma : \text{Var} \mapsto (\text{Val}, \text{Val})$

Definition 1.39 (Unary substitution). $\delta : \text{Var} \mapsto \text{Val}$

Definition 1.40 (Unary interpretation of Γ).

$$\llbracket \Gamma \rrbracket_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V\}$$

Definition 1.41 (Binary interpretation of Γ).

$$\llbracket \Gamma \rrbracket_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A\}$$

1.2.4 Soundness proof for CG

Lemma 1.42 (Binary value relation subsumes unary value relation). $\forall W, v_1, v_2, \mathcal{A}, n, \tau.$

$$(W, n, v_1, v_2) \in \llbracket \tau \rrbracket_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \llbracket \tau \rrbracket_V$$

Proof. Proof by induction on τ

1. Case \mathbf{b} , \mathbf{unit} :

From Definition 1.34

2. Case $\tau_1 \times \tau_2$:

Given: $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$ (P01)

and

$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$ (P02)

From Definition 1.32 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$ (P1)

IH1a: $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$ and

IH1b: $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a: $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$ and

IH2b: $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some m we need to prove

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly from (P02) we know that given some m we need to prove

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

We instantiate IH1a and IH2a with the given m from (P01) to get

$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$ and $(W.\theta_1, m, v_{i2}) \in [\tau_2]_V$

Then from Definition 1.34, we get

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly we instantiate IH1b and IH2b with the given m from (P02) to get

$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$ and $(W.\theta_2, m, v_{j2}) \in [\tau_2]_V$

Then from Definition 1.34, we get

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

3. Case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v_1 = \text{inl}(v_{i1})$ and $v_2 = \text{inl}(v_{j1})$

Given: $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$ (S01)

and

$\forall m. (W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$ (S02)

From Definition 1.32 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A$ (S0)

IH1: $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$ and

IH2: $\forall m_2. (W.\theta_2, m_2, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

From (S01) we know that given some m and we are required to prove:

$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

Also from (S02) we know that given some m and we are required to prove:

$(W.\theta_2, m, \text{inl}(v_{i2})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

We instantiate IH1 with m from (S01) to get

$(W.\theta_1, m, v_{i1}) \in \lfloor \tau_1 \rfloor_V$

Therefore from Definition 1.34, we get

$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

We instantiate IH2 with m from (S02) to get

$(W.\theta_2, m, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

Therefore from Definition 1.34, we get

$(W.\theta_2, m, \text{inl}(v_{j1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

(b) $v_1 = \text{inr}(v_{i2})$ and $v_2 = \text{inr}(v_{j2})$

Symmetric reasoning as in the (a) case above

4. Case $\tau_1 \rightarrow \tau_2$:

Given: $(W, n, \lambda x.e_1, \lambda x.e_2) \in \lceil \tau_1 \rightarrow \tau_2 \rceil_V^A$

This means from Definition 1.32 we know that

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in \lceil \tau_1 \rceil_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in \lceil \tau_2 \rceil_E^A)$
 $\wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in \lfloor \tau_1 \rfloor_V \implies (\theta_l, i, e_1[v_c/x]) \in \lfloor \tau_2 \rfloor_E)$
 $\wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in \lfloor \tau_1 \rfloor_V \implies (\theta_l, k, e_2[v_c/x]) \in \lfloor \tau_2 \rfloor_E)$ (L0)

To prove:

(a) $\forall m. (W.\theta_1, m, \lambda x.e_1) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$:

This means from Definition 1.34 we need to prove:

$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E$

This further means that we have some θ', j and v s.t

$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in \lfloor \tau_1 \rfloor_V$

And we need to prove: $(\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E$

Instantiating θ_l, i and v_c in the second conjunct of L0 with θ', j and v respectively and since we know that $W.\theta_1 \sqsubseteq \theta'$ and $(\theta', j, v) \in \lfloor \tau_1 \rfloor_V$

Therefore we get $(\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E$

(b) $\forall m. (W.\theta_2, m, \lambda x.e_2) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$:

Similar reasoning with e_2

5. Case ref $\ell \tau$:

From Definition 1.32 and 1.34

6. Case Labeled $\ell \tau$:

Given $(W, n, \text{Lb}v_1, \text{Lb}v_2) \in [\text{Labeled } \ell \tau]_V^A$

2 cases arise:

(a) $\ell \sqsubseteq \mathcal{A}$:

From Definition 1.31 we know that

$$(W, n, v_1, v_2) \in [\tau]_V^A$$

Therefore from IH we get $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$ and $\forall m. (W.\theta_2, m, v_2) \in [\tau]_V$

(b) $\ell \not\sqsubseteq \mathcal{A}$:

Directly from Definition 1.31

7. Case $\mathbb{C} \ell_1 \ell_2 \tau$:

Given: $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$

This means from Definition 1.32 we know that

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \quad (\text{CG0}) \end{aligned}$$

To prove: $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V$

This means from Definition 1.34 we need to prove

$$\begin{aligned} & \forall l \in \{1, 2\}. \forall m. (\forall k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j \text{ s.t. } (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \end{aligned}$$

Case $l = 1$

And given some m and $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t. $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove that

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating (CG0) with $l = 1$ and the given $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$ we get the desired.

Case $l = 2$

Symmetric reasoning as in the previous case above

□

Lemma 1.43 (Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m', \tau.$

$(\theta, m, v) \in [\tau]_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in [\tau]_V$

Proof. Proof by induction on τ

1. case **b**, unit:

Directly from Definition 1.34

2. case $\tau_1 \times \tau_2$:

Given: $(\theta, m, (v_1, v_2)) \in [\tau_1 \times \tau_2]_V$

To prove: $(\theta', m', (v_1, v_2)) \in [\tau_1 \times \tau_2]_V$

This means from Definition 1.34 we know that

$(\theta, m, v_1) \in [\tau_1]_V \wedge (\theta, m, v_2) \in [\tau_2]_V$

IH1 : $(\theta', m', v_1) \in [\tau_1]_V$

IH2 : $(\theta', m', v_2) \in [\tau_2]_V$

We get the desired from IH1, IH2 and Definition 1.34

3. case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v = \text{inl}(v_1)$:

Given: $(\theta, m, (\text{inl } v_1)) \in [\tau_1 + \tau_2]_V$

To prove: $(\theta', m', \text{inl } v_1) \in [\tau_1 + \tau_2]_V$

This means from Definition 1.34 we know that

$(\theta, m, v_1) \in [\tau_1]_V$

IH : $(\theta', m', v_1) \in [\tau_1]_V$

Therefore from IH and Definition 1.34 we get the desired

(b) $v = \text{inr}(v_2)$

Symmetric case

4. case $\tau_1 \rightarrow \tau_2$:

Given: $(\theta, m, (\lambda x. e_1)) \in [\tau_1 \rightarrow \tau_2]_V$

To prove: $(\theta', m', (\lambda x. e_1)) \in [\tau_1 \rightarrow \tau_2]_V$

This means from Definition 1.34 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in [\tau_1]_V \implies (\theta'', j, e_1[v/x]) \in [\tau_2]_E \quad (57)$$

Similarly from Definition 1.34 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall v_1. (\theta''', k, v_1) \in [\tau_1]_V \implies (\theta''', k, e_1[v_1/x]) \in [\tau_2]_E$$

This means that given some θ''', k and v_1 such that $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in [\tau_1]_V$

And we are required to prove $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E$

Instantiating Equation 57 with θ''' , k and v_1 and since we know that $\theta' \sqsubseteq \theta'''$ and $\theta \sqsubseteq \theta'$ therefore we have $\theta \sqsubseteq \theta'''$. Also, we know that $k < m' < m$ and $(\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V$

Therefore we get $(\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$

5. case ref $\ell \tau$:

From Definition 1.34 and Definition 1.29

6. case Labeled $\ell \tau$:

Given: $(\theta, m, (\text{Lb } v)) \in \lfloor \text{Labeled } \ell \tau \rfloor_V$

To prove: $(\theta', m', (\text{Lb } v)) \in \lfloor \text{Labeled } \ell \tau \rfloor_V$

This means from Definition 1.34 we know that $(\theta, m, v) \in \lfloor \tau \rfloor_V$

IH: $(\theta', m', v) \in \lfloor \tau \rfloor_V$

Therefore from IH and Definition 1.34 we get the desired

7. case $\mathbb{C} \ell_1 \ell_2 \tau$:

Given: $(\theta, m, e) \in \lfloor \mathbb{C} \ell_1 \ell_2 \tau \rfloor_V$

To prove: $(\theta', m', e) \in \lfloor \mathbb{C} \ell_1 \ell_2 \tau \rfloor_V$

This means from Definition 1.34 we know that

$$\begin{aligned} \forall k \leq m, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \rfloor_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) & \quad (\text{LB0}) \end{aligned}$$

Similarly from Definition 1.34 we are required to prove

$$\begin{aligned} \forall k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1. (k_1, H_1) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1 &\implies \\ \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H'_1) \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v'_1) \in \lfloor \tau \rfloor_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

This means we are given

$$k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1 \text{ s.t. } (k_1, H_1) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1$$

And we are required to prove:

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H'_1) \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v'_1) \in \lfloor \tau \rfloor_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

Instantiating (LB0), k with k_1 , θ_e with θ_{e1} , H with H_1 and j with j_1 . We know that $k_1 < m' < m$, $\theta \sqsubseteq \theta' \sqsubseteq \theta_{e1}$, $(k_1, H_1) \triangleright \theta_{e1}$, $(H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1)$ and $j_1 < k_1$. Therefore we get

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_{e1}. (k_1 - j_1, H'_1) \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v'_1) \in \lfloor \tau \rfloor_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

□

Lemma 1.44 (Monotonicity binary). *The following holds:*

$$\forall W, W', v_1, v_2, \mathcal{A}, n, n', \tau.$$

$$(W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}} \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$$

Proof. Proof by induction on τ

1. Case **b**, unit:

From Definition 1.32

2. Case $\tau_1 \times \tau_2$:

Given: $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

To prove: $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 1.32 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}} \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_{\mathcal{V}}^{\mathcal{A}}$$

$$\text{IH1} : (W', n', v_{i1}, v_{j1}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}}$$

$$\text{IH2} : (W', n', v_{i2}, v_{j2}) \in [\tau_2]_{\mathcal{V}}^{\mathcal{A}}$$

From IH1, IH2 and Definition 1.32 we get the desired.

3. Case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v_1 = \text{inl } v_{i1}$ and $v_2 = \text{inl } v_{i2}$:

Given: $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

To prove: $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 1.32 we know that we are given

$$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}}$$

$$\text{IH} : (W', n', v_{i1}, v_{i2}) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}}$$

Therefore from Definition 1.32 we get

$$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_{\mathcal{V}}^{\mathcal{A}}$$

(b) $v_1 = \text{inr}(v_{i1})$ and $v_2 = \text{inr}(v_{i2})$:

Symmetric case

4. Case $\tau_1 \rightarrow \tau_2$:

Given: $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

To prove: $(\theta', n', (\lambda x.e_1), (\lambda x.e_1)) \in [\tau_1 \rightarrow \tau_2]_{\mathcal{V}}^{\mathcal{A}}$

This means from Definition 1.32 we know that the following holds

$$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_{\mathcal{V}}^{\mathcal{A}} \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_{\mathcal{E}}^{\mathcal{A}}) \quad (\text{BM-A0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_{\mathcal{E}}) \quad (\text{BM-A1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_{\mathcal{E}}) \quad (\text{BM-A2})$$

Similarly from Definition 1.32 we know that we are required to prove

- (a) $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$:

This means that we are given some $W'' \sqsupseteq W', k < n'$ and v'_1, v'_2 s.t

$$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$$

And we are required to prove: $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with W'', k and v'_1, v'_2 we get

$$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$$

- (b) $\forall \theta'_i \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_i, k, v'_c) \in [\tau_1]_V \implies (\theta'_i, k, e_1[v'_c/x]) \in [\tau_2]_E)$:

This means that we are given some $\theta'_i \sqsupseteq W'.\theta_1, k$ and v'_c s.t

$$(\theta'_i, k, v'_c) \in [\tau_1]_V$$

And we are required to prove: $(\theta'_i, k, e_1[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with θ'_i, k and v'_c we get

$$(\theta'_i, k, e_1[v'_c/x]) \in [\tau_2]_E$$

- (c) $\forall \theta'_i \sqsupseteq W'.\theta_2, k, v'_c. ((\theta'_i, k, v'_c) \in [\tau_1]_V \implies (\theta'_i, k, e_2[v'_c/x]) \in [\tau_2]_E)$:

This means that we are given some $\theta'_i \sqsupseteq W'.\theta_2, k$ and v'_c s.t

$$(\theta'_i, k, v'_c) \in [\tau_1]_V$$

And we are required to prove: $(\theta'_i, k, e_2[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with θ'_i, k and v'_c we get

$$(\theta'_i, k, e_2[v'_c/x]) \in [\tau_2]_E$$

5. Case ref $\ell \tau$:

From Definition 1.32 and Definition 1.30

6. Case Labeled $\ell \tau$:

Given: $(W, n, (\text{Lb } v_1), (\text{Lb } v_2)) \in [\text{Labeled } \ell \tau]_V^A$

To prove: $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in [\text{Labeled } \ell \tau]_V^A$

From Definition 1.32 2 cases arise:

- (a) $\ell \sqsubseteq \mathcal{A}$:

In this case we know that $(W, n, v_1, v_2) \in [\tau]_V^A$

Therefore from IH we know that $(W', n', v_1, v_2) \in [\tau]_V^A$

Hence from Definition 1.32 we get $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in [\text{Labeled } \ell \tau]_V^A$

- (b) $\ell \not\sqsubseteq \mathcal{A}$:

In this case we know that $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$ and $(W.\theta_2, m, v_2) \in [\tau]_V$

Since $W.\theta_1 \sqsubseteq W'.\theta_1$ (from Definition 1.30). Therefore from Lemma 1.43 we know that $\forall m' < m. (W'.\theta_1, m', v_1) \in [\tau]_V$

Similarly since $W.\theta_2 \sqsubseteq W'.\theta_2$ (from Definition 1.30). Therefore from Lemma 1.43 we know that

$$\forall m' < m. (W'.\theta_2, m', v_2) \in [\tau]_V$$

Finally from Definition 1.32 we get $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in [\text{Labeled } \ell \tau]_V^A$

7. Case $\mathbb{C} \ell_1 \ell_2 \tau$:

Given: $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$

To prove: $(W', n', v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$

From Definition 1.32 we are given that

$$\begin{aligned}
& \left(\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\
& \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \right) \wedge \\
& \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \quad (\text{BM-M0})
\end{aligned}$$

Similarly from Definition 1.32 it suffices to prove that

$$\begin{aligned}
& \text{(a) } \left(\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\
& \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \right): \\
& \text{This means that given some } k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j \text{ s.t} \\
& (k, H_1, H_2) \triangleright W_e \wedge (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k
\end{aligned}$$

It suffices to prove that

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)$$

Instantiating the first conjunct of (BM-M0) with the given $k, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$ and since we know that $n' \leq n$ and $W \sqsubseteq W'$ we get the desired

$$\begin{aligned}
& \text{(b) } \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right):
\end{aligned}$$

Similar reasoning as in the previous case but using Lemma 1.43

□

Lemma 1.45 (Unary monotonicity for Γ). $\forall \theta, \theta', \delta, \Gamma, n, n'.$
 $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in [\Gamma]_V$

Proof. Given: $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$

To prove: $(\theta', n', \delta) \in [\Gamma]_V$

From Definition 1.40 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V$$

And again from Definition 1.40 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$$

- $dom(\Gamma) \subseteq dom(\delta)$:

Given

- $\forall x \in dom(\Gamma).(\theta', n', \delta(x)) \in [\Gamma(x)]_V$:

Since we know that $\forall x \in dom(\Gamma).(\theta, n, \delta(x)) \in [\Gamma(x)]_V$ (given)

Therefore from Lemma 1.43 we get

$$\forall x \in dom(\Gamma).(\theta', n', \delta(x)) \in [\Gamma(x)]_V$$

□

Lemma 1.46 (Binary monotonicity for Γ). $\forall W, W', \delta, \Gamma, n, n'$.

$$(W, n, \gamma) \in [\Gamma]_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in [\Gamma]_V$$

Proof. Given: $(W, n, \gamma) \in [\Gamma]_V \wedge n' < n \wedge W \sqsubseteq W'$

To prove: $(W', n', \gamma) \in [\Gamma]_V$

From Definition 1.41 it is given that

$$dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And again from Definition 1.40 we are required to prove that

$$dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

- $dom(\Gamma) \subseteq dom(\gamma)$:

Given

- $\forall x \in dom(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$:

Since we know that $\forall x \in dom(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$ (given)

Therefore from Lemma 1.44 we get

$$\forall x \in dom(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

□

Lemma 1.47 (Unary monotonicity for H). $\forall \theta, H, n, n'$.

$$(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$$

Proof. Given: $(n, H) \triangleright \theta \wedge n' < n$

To prove: $(n', H) \triangleright \theta$

From Definition 1.36 it is given that

$$dom(\theta) \subseteq dom(H) \wedge \forall a \in dom(\theta).(\theta, n - 1, H(a)) \in [\theta(a)]_V$$

And again from Definition 1.40 we are required to prove that

$$dom(\theta) \subseteq dom(H) \wedge \forall a \in dom(\theta).(\theta, n' - 1, H(a)) \in [\theta'(a)]_V$$

- $dom(\theta) \subseteq dom(H)$:

Given

- $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$:

Since we know that $\forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$ (given)

Therefore from Lemma 1.43 we get

$$\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

□

Lemma 1.48 (Binary monotonicity for heaps). $\forall W, H_1, H_2, n, n'$.

$$(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$$

Proof. Given: $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$

To prove: $(n', H_1, H_2) \triangleright W$

From Definition 1.37 it is given that

$$\begin{aligned} \text{dom}(W.\theta_1) &\subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ (W.\hat{\beta}) &\subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ \forall (a_1, a_2) \in (W.\hat{\beta}). & (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ (W, n - 1, H_1(a_1), H_2(a_2)) &\in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). & (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

And again from Definition 1.37 we are required to prove:

- $\text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2)$:
Given
- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$:
Given
- $\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2))$ and $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$:
 $\forall (a_1, a_2) \in (W.\hat{\beta})$.
 - $(W.\theta_1(a_1) = W.\theta_2(a_2))$: Given
 - $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$:
Given and from Lemma 1.44
- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$:
Given

□

Theorem 1.49 (Fundamental theorem unary). $\forall \Gamma, \theta, e, \tau, \delta, n$.

$$\begin{aligned} \Gamma \vdash e : \tau \wedge \\ (\theta, n, \delta) \in \lfloor \Gamma \rfloor_V &\implies \\ (\theta, n, e \delta) \in \lfloor \tau \rfloor_E \end{aligned}$$

Proof. Proof by induction on CG typing derivation

1. CG-var:

$$\frac{}{\Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, x \delta) \in [\tau]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. x \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau]_V$$

This means that given some $i < n$ s.t $x \delta \Downarrow_i v$

(from cg-val we know that $v = x \delta$ and $i = 0$)

It suffices to prove $(\theta, n, x \delta) \in [\tau]_V$ (FU-V0)

Since $(\theta, n, \delta) \in [\Gamma']_V$ where $\Gamma' = \Gamma \cup \{x : \tau\}$. Therefore from Definition 1.40 we know that $(\theta, n, \delta(x)) \in [\Gamma'(x)]_V$

So we are done.

2. CG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash e' : \tau_2}{\Gamma \vdash \lambda x. e' : (\tau_1 \rightarrow \tau_2)}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, \lambda x. e_i \delta) \in [(\tau_1 \rightarrow \tau_2)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \lambda x. e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 \rightarrow \tau_2)]_V$$

This means that given some $i < n$ s.t $\lambda x. e' \delta \Downarrow_i v$

(from cg-val we know that $v = \lambda x. e' \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \lambda x. e' \delta) \in [(\tau_1 \rightarrow \tau_2)]_V \quad (\text{FU-L0})$$

From Definition 1.34 it further suffices to prove

$$\forall \theta'' \sqsupseteq \theta, v', j < n. (\theta'', j, v') \in [\tau_1]_V \implies (\theta'', j, (e' \delta)[v'/x]) \in [\tau_2]_E$$

This means given some θ'', v', j s.t $\theta'' \sqsupseteq \theta, j < n$ and $(\theta'', j, v') \in [\tau_1]_V$ (FU-L1)

We are required to prove

$$(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2]_E$$

Since $(\theta, n, \delta) \in [\Gamma]_V$ therefore from Lemma 1.45 we know that $(\theta, j, \delta) \in [\Gamma]_V$ where $j < n$ (from FU-L1)

IH:

$$\forall \theta_h, v_x. (\theta_h, j, e' \delta \cup \{x \mapsto v_x\}) \in [\tau_2]_E, \text{ s.t } (\theta_h, j, v_x) \in [\tau_1]_V$$

Instantiating IH with θ'' and v' from (FU-L1) we get $(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2]_E$

3. CG-app:

$$\frac{\Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, (e_1 e_2) \delta) \in [\tau_2]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (e_1 e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_2]_V$$

This means that given some $i < n$ s.t $(e_1 e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau_2]_V \quad (\text{FU-P0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2)]_V$$

Since we know that $(e_1 e_2) \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t $e_1 \delta \Downarrow_j v_1$. This means we have

$$(\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2)]_V$$

From cg-app we know that $v_1 = \lambda x. e'$. Therefore we have

$$(\theta, n - j, \lambda x. e') \in [(\tau_1 \rightarrow \tau_2)]_V \quad (\text{FU-P1})$$

This means from Definition 1.34 we have

$$\forall \theta'' \sqsupseteq \theta \wedge I < (n - j), v. (\theta'', I, v) \in [\tau_1]_V \implies (\theta'', I, e'[v/x]) \in [\tau_2]_E \quad (58)$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in [\tau_1]_V$$

Since we know that $(e_1 e_2) \delta \Downarrow_i v$ therefore $\exists k < i - j$ (since $i < n$ therefore $i - j < n - j$) s.t $e_2 \delta \Downarrow_k v_2$. This means we have

$$(\theta, n - j - k, v_2) \in [\tau_1]_V \quad (\text{FU-P2})$$

Instantiating Equation 58 with $\theta, (n - j - k), v_2$ and since we know that $(\theta, n - j - k, v_2) \in [\tau_1]_V$ therefore we get

$$(\theta, n - j - k, e'[v_2/x]) \in [\tau_2]_E$$

This means from Definition 1.35 we have

$$\forall J < n - j - k. e'[v_2/x] \Downarrow_J v_f \implies (\theta, n - j - k - J, v_f) \in [\tau_2]_E$$

Since we know that $(e_1 e_2) \delta \Downarrow_i v$ therefore we know that $\exists J < i < n$ s.t $i = j + k + J$ (since $j + k + J < n$ therefore $J < n - j - k$) and $e'[v_2/x] \Downarrow_J v_f$

Therefore we have $(\theta, n - j - k - J, v_f) \in [\tau_2]_E$

Since we know that $i = j + k + J$ and $v = v_f$ therefore we get $(\theta, n - i, v_f) \in [\tau_2]_E$ (so FU-P0 is proved)

4. CG-prod:

$$\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, (e_1, e_2) \delta) \in [(\tau_1 \times \tau_2)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (e_1, e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 \times \tau_2)]_V$$

This means that given some $i < n$ s.t. $(e_1, e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [(\tau_1 \times \tau_2)]_V \quad (\text{FU-PA0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [\tau_1]_V$$

Since we know that $(e_1, e_2) \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t. $e_1 \delta \Downarrow_j v_1$. This means we have

$$(\theta, n - j, v_1) \in [\tau_1]_V \quad (\text{FU-PA1})$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in [\tau_2]_V$$

Since we know that $(e_1, e_2) \delta \Downarrow_i v$ therefore $\exists k < i - j$ (since $i < n$ therefore $i - j < n - j$) s.t. $e_2 \delta \Downarrow_k v_2$. This means we have

$$(\theta, n - j - k, v_2) \in [\tau_2]_V \quad (\text{FU-PA2})$$

In order to prove (FU-PA0) from cg-prod we know that $i = j + k + 1$ and $v = (v_1, v_2)$ therefore from Definition 1.34 it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in [\tau_1]_V \text{ and } (\theta, n - j - k - 1, v_2) \in [\tau_2]_V$$

We get this from (FU-PA1) and Lemma 1.43 and from (FU-PA2) and Lemma 1.43

5. CG-fst:

$$\frac{\Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Gamma \vdash \text{fst}(e') : \tau_1}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, \text{fst}(e') \delta) \in [\tau_1]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{fst}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_1]_V$$

This means that given some $i < n$ s.t $\text{fst}(e') \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor \tau_1 \rfloor_V \quad (\text{FU-F0})$$

IH1:

$$\forall j < n. e' \delta \Downarrow_j (v_1, v_2) \implies (\theta, n - j, (v_1, v_2)) \in \lfloor (\tau_1 \times \tau_2) \rfloor_V$$

Since we know that $\text{fst}(e') \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t $e' \delta \Downarrow_j (v_1, v_2)$. This means we have

$$(\theta, n - j, (v_1, v_2)) \in \lfloor (\tau_1 \times \tau_2) \rfloor_V$$

From Definition 1.34 we know the following holds

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \rfloor_V \text{ and } (\theta, n - j, v_2) \in \lfloor \tau_2 \rfloor_V \quad (\text{FU-F1})$$

From cg-fst we know that $v = v_1$ and $i = j + 1$. Therefore from (FU-F0), we are required to prove

$$(\theta, n - j - 1, v_1) \in \lfloor \tau_1 \rfloor_V$$

We get this from (FU-F1) and Lemma 1.43

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Gamma \vdash e' : \tau_1}{\Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

Also given is $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$

To prove: $(\theta, n, \text{inl}(e') \delta) \in \lfloor (\tau_1 + \tau_2) \rfloor_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{inl}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\tau_1 + \tau_2) \rfloor_V$$

This means that given some $i < n$ s.t $\text{inl}(e') \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor (\tau_1 + \tau_2) \rfloor_V \quad (\text{FU-LE0})$$

IH1:

$$\forall j < n. e' \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in \lfloor \tau_1 \rfloor_V$$

Since we know that $\text{inl}(e') \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t $e' \delta \Downarrow_j v_1$. This means we have

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \rfloor_V \quad (\text{FU-LE1})$$

From cg-inl we know that $v = v_1$ and $i = j + 1$. Therefore from (FU-LE0) we are required to prove

$$(\theta, n - j - 1, v_1) \in [(\tau_1 + \tau_2)]_V$$

From Definition 1.34 it suffices to prove

$$(\theta, n - j - 1, v_1) \in [\tau_1]_V$$

We get this from (FU-LE1) and Lemma 1.43

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Gamma \vdash \text{case}(e, x.e_1, y.e_2) : \tau}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in [\tau]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau]_V$$

This means that given some $i < n$ s.t $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau]_V \quad (\text{FU-C0})$$

IH1:

$$\forall j < n. e_c \delta \Downarrow_j v_c \implies (\theta, n - j, v_1) \in [(\tau_1 + \tau_2)]_V$$

Since we know that $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t $e_c \delta \Downarrow_j v_c$. This means we have

$$(\theta, n - j, v_c) \in [(\tau_1 + \tau_2)]_V \quad (\text{FU-C1})$$

2 cases arise:

(a) $v_c = \text{inl}(v_l)$:

IH2:

$$\forall k < (n - j). e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1 \implies (\theta, n - j - k, v_1) \in [\tau]_V$$

Since we know that $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$ therefore $\exists k < i - j$ (since $i < n$ therefore $i - j < n - j$) s.t $e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1$. This means we have

$$(\theta, n - j - k, v_1) \in [\tau]_V \quad (\text{FU-C2})$$

From cg-case1 we know that $i = j + k + 1$ and $v = v_1$. Therefore from (FU-C0) it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in [\tau]_V$$

We get this from (FU-C2) and Lemma 1.43

(b) $v_c = \text{inr}(v_r)$:

Symmetric reasoning as in the previous case

10. CG-ref:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } (e') : \mathbb{C} \ell \perp (\text{ref } \ell' \tau)}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, \text{new } (e') \delta) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{new } (e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V$$

This means that given some $i < n$ s.t $\text{new } (e') \delta \Downarrow_i v$

(from cg-val we know that $v = \text{new } (e') \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \text{new } (e') \delta) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{ref } \ell' \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k$.

Also from cg-ref we know that $v' = a$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, a) \in [(\text{ref } \ell' \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IIH:

$$(\theta_e, k, e' \delta) \in [(\text{Labeled } \ell' \tau)]_E$$

From Definition 1.35 this means we have

$$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, n - l, v_h) \in [(\text{Labeled } \ell' \tau)]_V$$

Since we know that $(H, \text{new } (e')) \Downarrow_j^f (H', a)$ therefore from cg-ref we know that

$$\exists l < j < k \text{ s.t } e' \delta \Downarrow_l v_h$$

Therefore we have

$$(\theta_e, n - l, v_h) \in [(\text{Labeled } \ell' \tau)]_V \quad (\text{FU-R2})$$

In order to prove (FU-R0) we choose θ' as $\theta_n = \theta_e \cup \{a \mapsto \text{Labeled } \ell' \tau\}$

Now we need to prove:

(a) $(k - j, H') \triangleright \theta_n$:

From Definition 1.36 it suffices to prove that

$$\text{dom}(\theta_n) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in [\theta_n(a)]_V$$

- $dom(\theta_n) \subseteq dom(H')$:
 We know that $dom(H') = dom(H) \cup \{a\}$
 We know that $dom(\theta_n) = dom(\theta_e) \cup \{a\}$
 And $(k, H) \triangleright \theta_e$ therefore from Definition 1.36 we know that $dom(\theta_e) \subseteq dom(H)$
 So we are done
 - $\forall a \in dom(\theta_n).(\theta_n, (k-j)-1, H'(a)) \in [\theta_n(a)]_V$:
 Since from (FU-R2) we know that $(\theta_h, n-l, v_h) \in [(\text{Labeled } \ell' \tau)]_V$
 Since $\theta_h \sqsubseteq \theta_n$ and $k-j-1 < n-l$ (since $k < n$ and $l < j$) therefore from Lemma 1.43 we know that $(\theta_n, k-j-1, v_h) \in [(\text{Labeled } \ell' \tau)]_V$
- (b) $(\theta_n, k-j-1, a) \in [(\text{ref } \ell' \tau)]_V$:
 From Definition 1.34 it suffices to prove that $\theta_n(a) = \text{Labeled } \ell' \tau$
 We get this by construction of θ_n
- (c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$:
 Holds vacuously
- (d) $(\forall a \in dom(\theta_n) \setminus dom(\theta_e). \theta_n(a) \searrow \ell)$:
 From CG-ref we know that $\ell \sqsubseteq \ell'$

11. CG-deref:

$$\frac{\Gamma \vdash e' : \text{ref } \ell \tau}{\Gamma \vdash !e' : \mathbb{C} \top \perp (\text{Labeled } \ell \tau)}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, (!e') \delta) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \tau)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (!e') \delta \Downarrow_i v \implies (\theta, n-i, v) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \tau)]_V$$

(From cg-val we know that $v = !e' \delta$ and $i = 0$)

This means that given some $i < n$ s.t $!e' \delta \Downarrow_i !e' \delta$

It suffices to prove

$$(\theta, n, !e' \delta) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \tau)]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} & \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k-j, H') \triangleright \theta' \wedge (\theta', k-j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \top \sqsubseteq \ell'') \wedge \\ & (\forall a \in dom(\theta') \setminus dom(\theta_e). \theta'(a) \searrow \top) \end{aligned}$$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k$.

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k-j, H') \triangleright \theta' \wedge (\theta', k-j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \top \sqsubseteq \ell'') \wedge \\ & (\forall a \in dom(\theta') \setminus dom(\theta_e). \theta'(a) \searrow \top) \quad (\text{FU-D0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [(\text{ref } \ell \tau)]_E$$

From Definition 1.35 this means we have

$$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, k - l, v_h) \in [(\text{ref } \ell \tau)]_V$$

Since we know that $(H, !(e')) \Downarrow_j^f (H', a)$ therefore from cg-deref we know that

$$\exists l < j < k \text{ s.t. } e' \delta \Downarrow_l v_h, v_h = a$$

Therefore we have

$$(\theta_e, k - l, a) \in [(\text{ref } \ell \tau)]_V \quad (\text{FU-D1})$$

In order to prove (FU-D0) we choose θ' as θ_e

Now we need to prove:

(a) $(k - j, H') \triangleright \theta_e$:

From Definition 1.36 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

- $\text{dom}(\theta_e) \subseteq \text{dom}(H')$:

And $(k, H) \triangleright \theta_e$ therefore from Definition 1.36 we know that $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

And since $H' = H$ (from cg-deref) so we are done

- $\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$:

Since we know that $(k, H) \triangleright \theta_e$ therefore from Definition 1.36 we know that

$$\forall a \in \text{dom}(\theta_e). (\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since $H' = H$ and from Lemma 1.43 we get

$$\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

(b) $(\theta_e, k - j, v') \in [(\text{Labeled } \ell \tau)]_V$:

From cg-deref we know that $H = H'$ and $v' = H(a)$

From (FU-D1) and Definition 1.34 we know that $\theta_e(a) = \text{Labeled } \ell \tau$

Since we know that $(k, H) \triangleright \theta_e$ therefore from Definition 1.36 we know that

$$\forall a \in \text{dom}(\theta_e). (\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since from cg-deref we know that $j \geq 1$. Therefore from Lemma 1.43 we get $(\theta_e, k - j, H(a)) \in [(\text{Labeled } \ell \tau)]_V$

(c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell')$:

Holds vacuously

(d) $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \top)$:

Holds vacuously

12. CG-assign:

$$\frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit}}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, (e_1 := e_2) \delta) \in [(\mathbb{C} \ell \perp \text{unit})]_E^{pc}$

This means that from Definition 1.35 we need to prove

$$\forall i < n. (e_1 := e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell \perp \text{unit})]_V$$

This means that given some $i < n$ s.t $(e_1 := e_2) \delta \Downarrow_i v$.

It suffices to prove

$$(\theta, n - i, ()) \in [(\mathbb{C} \ell \perp \text{unit})]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} & \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{ref } \ell' \tau)]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$. Also from cg-assgn we know that $v' = ()$

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, ()) \in [\text{unit}]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-A0}) \end{aligned}$$

IH1:

$$\forall l < k. e_1 \delta \Downarrow_l v_1 \implies (\theta, k - l, a) \in [(\text{ref } \ell' \tau)]_V$$

Since we know that $(e_1 := e_2) \delta \Downarrow_j^f v$ therefore $\exists l < j < k$ s.t $e_1 \delta \Downarrow_l a$. This means we have

$$(\theta, k - l, a) \in [(\text{ref } \ell' \tau)]_V \quad (\text{FU-A1})$$

IH2:

$$\forall m < (k - l). e_2 \delta \Downarrow_m v_2 \implies (\theta, k - l - m, v_2) \in [\text{Labeled } \ell' \tau]_V$$

Since we know that $(e_1 := e_2) \delta \Downarrow_j^f v$ therefore $\exists m < j - l$ (since $j < k$ therefore $j - l < k - l$) s.t $e_2 \delta \Downarrow_m v_2$. This means we have

$$(\theta, k - l - m, v_2) \in [(\text{Labeled } \ell' \tau)]_V \quad (\text{FU-A2})$$

In order to prove (FU-A0) we choose θ' as θ_e

Now we need to prove:

(a) $(k - j, H') \triangleright \theta_e$:

From Definition 1.36 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

• $\text{dom}(\theta_e) \subseteq \text{dom}(H')$:

We know that $\text{dom}(H') = \text{dom}(H)$

And $(k, H) \triangleright \theta_e$ therefore from Definition 1.36 we know that $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

- $\forall a \in \text{dom}(\theta_e). (\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$:
 $\forall a \in \text{dom}(\theta_e)$.

i. $H(a) = H'(a)$:

Since $(k, H) \triangleright \theta_e$ therefore from Definition 1.36 we know that

$$(\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Therefore from Lemma 1.43 we get

$$(\theta_e, k - 1 - j, H(a)) \in [\theta_e(a)]_V$$

ii. $H(a) \neq H'(a)$:

From cg-assign we know that $H'(a) = v_2$

From (FU-A1) we know that $\theta_e(a) = \text{Labeled } \ell' \tau$

Also we know that $j = l + m + 1$

Since from (FU-A2) we know that

$$(\theta, k - l - m, v_2) \in [(\text{Labeled } \ell' \tau)]_V$$

Therefore we get

$$(\theta, k - j + 1, v_2) \in [(\text{Labeled } \ell' \tau)]_V$$

Therefore from Lemma 1.43 we get

$$(\theta, k - j - 1, v_2) \in [(\text{Labeled } \ell' \tau)]_V$$

(b) $(\theta_e, k - j - 1, ()) \in [\text{unit}]_V$:

From Definition 1.34

(c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell')$:

From CG-assign we know that $\ell \sqsubseteq \ell'$

(d) $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \ell)$:

Holds vacuously

13. CG-label:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, \text{Lb}(e') \delta) \in [\text{Labeled } \ell \tau]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{Lb}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\text{Labeled } \ell \tau]_V$$

This means we are given some $i < n$ s.t $\text{Lb}(e') \delta \Downarrow_i v$ and we are required to prove

$$(\theta, n - i, v) \in [\text{Labeled } \ell \tau]_V$$

Let $v = \text{Lb}(v_i)$. This means from Definition 1.34 we are required to prove

$$(\theta, n - i, v_i) \in [\tau]_V$$

IH: $(\theta, n, e' \delta) \in [\tau]_E$

This means from Definition 1.35 we have

$$\forall j < n. e' \delta \Downarrow_j v_i \implies (\theta, n - j, v_i) \in [\tau]_V$$

Since we know that $\text{Lb}(e') \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t $e' \delta \Downarrow_j v_i$

Therefore we have $(\theta, n - j, v_i) \in \lfloor \tau \rfloor_V$

From cg-label we know that $i = j + 1$ therefore from Lemma 1.43 we have

$$(\theta, n - i, v_i) \in \lfloor \tau \rfloor_V$$

14. CG-unlabel:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell \tau}{\Gamma \vdash \text{unlabel}(e') : \mathbb{C} \top \ell \tau}$$

Also given is $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$

To prove: $(\theta, n, \text{unlabel}(e') \delta) \in \lfloor (\mathbb{C} \top \ell \tau) \rfloor_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{unlabel}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\mathbb{C} \top \ell \tau) \rfloor_V$$

This means that given some $i < n$ s.t $\text{unlabel}(e') \delta \Downarrow_i v$

(from cg-val we know that $v = \text{unlabel}(e') \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \text{unlabel}(e') \delta) \in \lfloor (\mathbb{C} \top \ell \tau) \rfloor_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \end{aligned}$$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$. Also from cg-unlabel we know that $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \quad (\text{FU-U0}) \end{aligned}$$

IIH:

$$(\theta_e, k, e' \delta) \in \lfloor (\text{Labeled } \ell \tau) \rfloor_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in \lfloor (\text{Labeled } \ell \tau) \rfloor_V$$

Since we know that $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$ therefore from cg-unlabel we know that

$$\exists h_1 < j < k \text{ s.t } e' \delta \Downarrow_{h_1} \text{Lb } v'$$

This means we have

$$(\theta_e, k - h_1, \text{Lb } v') \in \lfloor (\text{Labeled } \ell \tau) \rfloor_V$$

This means from Definition 1.34 we have

$$(\theta_e, k - h_1, v') \in [\tau]_V \quad (\text{FU-U1})$$

In order to prove (FU-U0) we choose θ' as θ_e . And we are required to prove:

(a) $(k - j, H) \triangleright \theta_e$:

Since we have $(k, H) \triangleright \theta_e$ therefore from Lemma 1.47 we get $(k - j, H) \triangleright \theta_e$

(b) $(\theta', k - j, v') \in [\tau]_V$:

Since from (FU-U1) we know that $(\theta_e, k - h_1, v') \in [\tau]_V$

And since $j = h_1 + 1$, therefore from Lemma 1.43 we get $(\theta_e, k - j, v') \in [\tau]_V$

(c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell')$:

Holds vacuously

(d) $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top)$:

Holds vacuously

15. CG-ret:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{ret}(e') : \mathbb{C} \ell \ell' \tau}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell \ell' \tau]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{ret}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell \ell' \tau]_V$$

This means we are given some $i < n$ s.t. $\text{ret}(e') \delta \Downarrow_i v$ and we are required to prove

$$(\theta, n - i, v) \in [\mathbb{C} \ell \ell' \tau]_V$$

(from cg-val we know that $v = \text{ret}(e') \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell \ell' \tau]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t. $(k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$.

Also from cg-ret we know that $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in \lfloor \tau \rfloor_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in \lfloor \tau \rfloor_V$$

Since we know that $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$ therefore from cg-ret we know that $\exists h_1 < j < k$ s.t $e' \delta \Downarrow_{h_1} v'$

This means we have

$$(\theta_e, k - h_1, v') \in \lfloor \tau \rfloor_V \quad (\text{FU-R1})$$

In order to prove (FU-U0) we choose θ' as θ_e . And we are required to prove:

(a) $(k - j, H) \triangleright \theta_e$:

Since have $(k, H) \triangleright \theta_e$ therefore from Lemma 1.47 we get $(k - j, H) \triangleright \theta_e$

(b) $(\theta', k - j, v') \in \lfloor \tau \rfloor_V$:

Since from (FU-R1) we know that $(\theta_e, k - h_1, v') \in \lfloor \tau \rfloor_V$

And since $j = h_1 + 1$, therefore from Lemma 1.43 we get $(\theta_e, k - j, v') \in \lfloor \tau \rfloor_V$

(c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell \sqsubseteq \ell'')$:

Holds vacuously

(d) $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$:

Holds vacuously

16. CG-bind:

$$\frac{\Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \quad \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau'}$$

Also given is $(\theta, n, \delta) \in \lfloor \Gamma \rfloor_V$

To prove: $(\theta, n, \text{bind}(e_1, x.e_2) \delta) \in \lfloor \mathbb{C} \ell \ell' \tau' \rfloor_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{bind}(e_1, x.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \mathbb{C} \ell \ell' \tau' \rfloor_V$$

This means we are given some $i < n$ s.t $\text{bind}(e_1, x.e_2) \delta \Downarrow_i v$ and we are required to prove

$$(\theta, n - i, v) \in \lfloor \mathbb{C} \ell \ell' \tau' \rfloor_V$$

(from cg-val we know that $v = \text{bind}(e_1, x.e_2) \delta$ and $i = 0$)

Therefore we need to prove

$$(\theta, n, v) \in \lfloor \mathbb{C} \ell \ell' \tau' \rfloor_V$$

From Definition 1.34 it suffices to prove

$$\forall k \leq n. \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies$$

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \rfloor_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$$

This means we are given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$.

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-B0}) \end{aligned}$$

IH1:

$$(\theta_e, k, e_1 \delta) \in \lfloor (\mathbb{C} \ell_1 \ell_2 \tau) \rfloor_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_1 < k. e_1 \delta \Downarrow_{h_1} v_1 \implies (\theta_e, k - h_1, v_1) \in \lfloor (\mathbb{C} \ell_1 \ell_2 \tau) \rfloor_V$$

Since we know that $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$ therefore from cg-bind we know that

$$\exists h_1 < j < k \text{ s.t } e_1 \delta \Downarrow_{h_1} v_1$$

This means we have

$$(\theta_e, k - h_1, v_1) \in \lfloor (\mathbb{C} \ell_1 \ell_2 \tau) \rfloor_V$$

From Definition 1.34 we know that

$$\begin{aligned} & \forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H, J.(k_{h_1}, H) \triangleright \theta'_e \wedge (H, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies \\ & \exists \theta'' \sqsupseteq \theta'_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_1 \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell_1) \end{aligned}$$

Instantiating k_{h_1} with $k - h_1, \theta'_e$ with θ_e . Since we know that $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$ therefore $\exists J < j - h_1 < k - h_1$ s.t $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$. And since we already know that $(k, H) \triangleright \theta_e$ therefore from Lemma 1.47 we get $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned} & \exists \theta'' \sqsupseteq \theta_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_1 \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta_e). \theta''(a) \searrow \ell_1) \quad (\text{FU-B1}) \end{aligned}$$

IH2:

$$(\theta'', k - h_1 - J, e_2 \delta \cup \{x \mapsto v'\}) \in \lfloor (\mathbb{C} \ell_3 \ell_4 \tau') \rfloor_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_2 < k - h_1 - J. e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v'' \implies (\theta'', k - h_1 - J - h_2, v'') \in \lfloor (\mathbb{C} \ell_3 \ell_4 \tau') \rfloor_V$$

Since we know that $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H, v_1)$ therefore from cg-bind we know that

$$\exists h_2 < j - h_1 - J < k - h_1 - J \text{ s.t } e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v''$$

This means we have

$$(\theta'', k - h_1 - J - h_2, v'') \in \lfloor (\mathbb{C} \ell_3 \ell_4 \tau') \rfloor_V$$

From Definition 1.34 we know that

$$\begin{aligned} \forall k_{h_2} \leq (k - h_1 - J - h_2), \theta'_e \sqsupseteq \theta'', H, J'.(k_{h_2}, H) \triangleright \theta'_e \wedge (H, v'') \Downarrow_{J'}^f (H'', v'_{h_2}) \wedge J' < k_{h_2} \implies \\ \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v') \in [\tau']_V \wedge \\ (\forall a. H(a) \neq H''(a) \implies \exists \ell''. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_3 \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell_3) \end{aligned}$$

Since we know that $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$ therefore $\exists v_{h_2}, i$ s.t $(v'' \Downarrow_i v_{h_2})$. From cg-val we know that $v_{h_2} = v''$ and $i = 0$. Instantiating k_{h_2} with $k - h_1 - J - h_2$, θ'_e with θ'' , H with H' (from FU-B1) and $\exists J' < j - h_1 - J - h_2 < k - h_1 - J - h_2$ s.t $(H', v_{h_2}) \Downarrow_{J'}^f (H'', v'_{h_2})$. And since we already know that $(k - h_1, H') \triangleright \theta''$ therefore from Lemma 1.47 we get $(k - h_1 - J - h_2, H') \triangleright \theta''$

This means we have

$$\begin{aligned} \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v') \in [\tau]_V \wedge \\ (\forall a. H(a) \neq H''(a) \implies \exists \ell''. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell_3 \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell_3) \quad (\text{FU-B2}) \end{aligned}$$

We get (FU-B0) by choosing θ' as θ''' (from FU-B2)

17. CG-toLabeled:

$$\frac{\Gamma \vdash e' : \mathbb{C} \ell_1 \ell_2 \tau}{\Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau)}$$

Also given is $(\theta, n, \delta) \in [\Gamma]_V$

To prove: $(\theta, n, \text{toLabeled}(e') \delta) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_E$

This means that from Definition 1.35 we need to prove

$$\forall i < n. \text{toLabeled}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$$

This means that given some $i < n$ s.t $\text{toLabeled}(e') \delta \Downarrow_i v$

(from cg-val we know that $v = \text{toLabeled}(e') \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \text{toLabeled}(e') \delta) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$$

From Definition 1.34 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

And given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$. Also from cg-tolabeled we know that $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \quad (\text{FU-TL0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_E$$

This means that from Definition 1.35 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_1 \implies (\theta, k - h_1, v_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V$$

Since $H, \text{toLabeled}(e') \Downarrow_j^f H', v'$ therefore from cg-tolabeled we know that $\exists h_1 < j < k$ s.t $e' \delta \Downarrow_{h_1} v_1$

Therefore we get $(\theta, k - h_1, v_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V$

From Definition 1.34 we know that

$$\begin{aligned} \forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H_h, J.(k_{h_1}, H_h) \triangleright \theta'_e \wedge (H_h, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies \\ \exists \theta'' \sqsupseteq \theta'_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v_1) \in [\tau]_V \wedge \\ (\forall a. H_h(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell_1) \end{aligned}$$

Instantiating k_{h_1} with $k - h_1$, H_h with H , θ'_e with θ_e . Since we know that $(H, \text{toLabeled}(e')) \Downarrow_j^f (H', v_1)$ therefore $\exists J < j - h_1 < k - h_1$ s.t $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$. And since we already know that $(k, H) \triangleright \theta_e$ therefore from Lemma 1.47 we get $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned} \exists \theta'' \sqsupseteq \theta'_e.(k - h_1 - J, H') \triangleright \theta'' \wedge (\theta'', k - h_1 - J, v_1) \in [\tau]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell_1) \quad (\text{FU-TL1}) \end{aligned}$$

In order to prove (FU-TL0) we choose θ' as θ'' . Now we need to prove the following

(a) $(k - j, H') \triangleright \theta''$:

Since $(k - h_1 - J, H') \triangleright \theta''$ and $j = h_1 + J + 1$ therefore from Lemma 1.47 we get $(k - j, H') \triangleright \theta''$

(b) $(\theta'', k - j - 1, v') \in [(\text{Labeled } \ell_o \tau)]_V$:

From cg-tolabeled we know that $v' = \text{toLabeled}(v_1)$

From Definition 1.32 it suffices to prove that $(\theta'', k - j - 1, v_1) \in [\tau]_V$

We get this from (FU-TL1) and Lemma 1.43

(c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$:

Directly from (FU-TL1)

(d) $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow \ell)$:

Directly from (FU-TL1)

□

Lemma 1.50 (Subtyping unary). *The following holds:*

$$\forall \mathcal{L}, \tau, \tau'.$$

$$1. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_V \subseteq [(\tau')]_V$$

$$2. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_E \subseteq [(\tau')]_E$$

Proof. Proof of Statement (1)

Proof by induction on $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove: $\llbracket ((\tau_1 \rightarrow \tau_2)) \rrbracket_V \subseteq \llbracket ((\tau'_1 \rightarrow \tau'_2)) \rrbracket_V$

IH1: $\llbracket (\tau'_1) \rrbracket_V \subseteq \llbracket (\tau_1) \rrbracket_V$ (Statement (1))

$\llbracket (\tau_2) \rrbracket_E \subseteq \llbracket (\tau'_2) \rrbracket_E$ (Sub-A0, From Statement (2))

It suffices to prove: $\forall (\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \rightarrow \tau_2)) \rrbracket_V. (\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \rightarrow \tau'_2)) \rrbracket_V$

This means that given some θ, n and $\lambda x.e_i$ s.t $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau_1 \rightarrow \tau_2)) \rrbracket_V$

Therefore from Definition 1.34 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta_1, i, e_i[v/x]) \in \llbracket \tau_2 \rrbracket_E \quad (59)$$

And it suffices to prove: $(\theta, n, \lambda x.e_i) \in \llbracket ((\tau'_1 \rightarrow \tau'_2)) \rrbracket_V$

Again from Definition 1.34, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V \implies (\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E$$

This means that given some $\theta_2, j < n, v$ s.t $\theta \sqsubseteq \theta_2$ and $(\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V$

And we are required to prove: $(\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E$

Since $(\theta_2, j, v) \in \llbracket \tau'_1 \rrbracket_V$ therefore from IH1 we know that $(\theta_2, j, v) \in \llbracket \tau_1 \rrbracket_V$

As a result from Equation 59 we know that

$$(\theta_2, j, e_i[v/x]) \in \llbracket \tau_2 \rrbracket_E$$

From (Sub-A0), we know that

$$(\theta_2, j, e_i[v/x]) \in \llbracket \tau'_2 \rrbracket_E$$

2. CGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove: $\llbracket ((\tau_1 \times \tau_2)) \rrbracket_V \subseteq \llbracket ((\tau'_1 \times \tau'_2)) \rrbracket_V$

IH1: $\llbracket (\tau_1) \rrbracket_V \subseteq \llbracket (\tau'_1) \rrbracket_V$ (Statement (1))

IH2: $\llbracket (\tau_2) \rrbracket_V \subseteq \llbracket (\tau'_2) \rrbracket_V$ (Statement (1))

It suffices to prove: $\forall (\theta, n, (v_1, v_2)) \in \llbracket ((\tau_1 \times \tau_2)) \rrbracket_V. (\theta, n, (v_1, v_2)) \in \llbracket ((\tau'_1 \times \tau'_2)) \rrbracket_V$

This means that given some θ, n and (v_1, v_2) $(\theta, (v_1, v_2)) \in [((\tau_1 \times \tau_2))]_V$

Therefore from Definition 1.34 we are given:

$$(\theta, n, v_1) \in [\tau_1]_V \wedge (\theta, n, v_2) \in [\tau_2]_V \quad (60)$$

And it suffices to prove: $(\theta, (v_1, v_2)) \in [((\tau'_1 \times \tau'_2))]_V$

Again from Definition 1.34, it suffices to prove:

$$(\theta, n, v_1) \in [\tau'_1]_V \wedge (\theta, n, v_2) \in [\tau'_2]_V$$

Since from Equation 60 we know that $(\theta, n, v_1) \in [\tau_1]_V$ therefore from IH1 we have $(\theta, n, v_1) \in [\tau'_1]_V$

Similarly since $(\theta, n, v_2) \in [\tau_2]_V$ from Equation 60 therefore from IH2 we have $(\theta, n, v_2) \in [\tau'_2]_V$

3. CGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove: $[((\tau_1 + \tau_2))]_V \subseteq [((\tau'_1 + \tau'_2))]_V$

IH1: $[(\tau_1)]_V \subseteq [(\tau'_1)]_V$ (Statement (1))

IH2: $[(\tau_2)]_V \subseteq [(\tau'_2)]_V$ (Statement (1))

It suffices to prove: $\forall(\theta, n, v_s) \in [((\tau_1 + \tau_2))]_V. (\theta, v_s) \in [((\tau'_1 + \tau'_2))]_V$

This means that given: $(\theta, n, v_s) \in [((\tau_1 + \tau_2))]_V$

And it suffices to prove: $(\theta, n, v_s) \in [((\tau'_1 + \tau'_2))]_V$

2 cases arise

(a) $v_s = \text{inl } v_i$:

From Definition 1.34 we are given:

$$(\theta, n, v_i) \in [\tau_1]_V \quad (61)$$

And we are required to prove that:

$$(\theta, n, v_i) \in [\tau'_1]_V$$

From Equation 61 and IH1 we know that

$$(\theta, n, v_i) \in [\tau'_1]_V$$

(b) $v_s = \text{inr } v_i$:

From Definition 1.34 we are given:

$$(\theta, n, v_i) \in [\tau_2]_V \quad (62)$$

And we are required to prove that:

$$(\theta, n, v_i) \in [\tau'_2]_V$$

From Equation 62 and IH2 we know that

$$(\theta, n, v_i) \in [\tau'_2]_V$$

4. CGsub-label:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove: $\llbracket ((\text{Labeled } \ell \tau)) \rrbracket_V \subseteq \llbracket ((\text{Labeled } \ell' \tau')) \rrbracket_V$

IH: $\llbracket (\tau) \rrbracket_V \subseteq \llbracket (\tau') \rrbracket_V$ (Statement (1))

It suffices to prove:

$\forall (\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell \tau)) \rrbracket_V. (\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell' \tau')) \rrbracket_V$

This means that given some θ, n and $\text{Lb}(e_i)$ s.t $(\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell \tau)) \rrbracket_V$

Therefore from Definition 1.34 we are given:

$$(\theta, n, v_i) \in \llbracket (\tau) \rrbracket_V \quad (\text{SL})$$

And we are required to prove that

$$(\theta, n, \text{Lb}(v_i)) \in \llbracket ((\text{Labeled } \ell' \tau')) \rrbracket_V$$

From Definition 1.34 it suffices to prove

$$(\theta, n, v_i) \in \llbracket (\tau') \rrbracket_V$$

We get this directly from (SL) and IH

5. CGsub-CG:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i \quad \mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o}{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove: $\llbracket ((\mathbb{C} \ell_i \ell_o \tau)) \rrbracket_V \subseteq \llbracket ((\mathbb{C} \ell'_i \ell'_o \tau')) \rrbracket_V$

IH: $\llbracket (\tau) \rrbracket_V \subseteq \llbracket (\tau') \rrbracket_V$ (Statement (1))

It suffices to prove:

$\forall (\theta, n, e) \in \llbracket ((\mathbb{C} \ell_i \ell_o \tau)) \rrbracket_V. (\theta, n, e) \in \llbracket ((\mathbb{C} \ell'_i \ell'_o \tau')) \rrbracket_V$

This means that given some θ, n and e s.t $(\theta, n, e) \in \llbracket ((\mathbb{C} \ell_i \ell_o \tau)) \rrbracket_V$

Therefore from Definition 1.34 we are given:

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i) & \quad (\text{SC0}) \end{aligned}$$

And we are required to prove

$$(\theta, n, e) \in \llbracket ((\mathbb{C} \ell'_i \ell'_o \tau')) \rrbracket_V$$

So again from Definition 1.34 we need to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau' \rrbracket_V \wedge & \end{aligned}$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$$

This means we are given some $k \leq n, \theta_e \sqsupseteq \theta, H, j < k$ s.t. $(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v')$
(SC1)

And we need to prove

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \rfloor_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$$

We instantiate (SC0) with k, θ_e, H, j from (SC1) and we get

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \rfloor_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)$$

Since $\tau <: \tau'$ therefore from IH we get

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \rfloor_V$$

And since $\ell'_i \sqsubseteq \ell_i$ therefore we also have

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$$

6. CGsub-base:

Trivial

Proof of Statement(2)

It suffice to prove that

$$\forall (\theta, n, e) \in \lfloor (\tau) \rfloor_E. (\theta, n, e) \in \lfloor (\tau') \rfloor_E$$

This means that we are given $(\theta, n, e) \in \lfloor (\tau) \rfloor_E$

From Definition 1.35 it means we have

$$\forall i < n.e \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau \rfloor_V \quad (\text{Sub-E0})$$

And we need to prove

$$(\theta, n, e) \in \lfloor (\tau') \rfloor_E$$

From Definition 1.35 we need to prove

$$\forall i < n.e \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau' \rfloor_V$$

This further means that given some $i < n$ s.t. $e \Downarrow_i v$, it suffices to prove that

$$(\theta, n - i, v) \in \lfloor \tau' \rfloor_V$$

Instantiating (Sub-E0) with the given i we get $(\theta, n - i, v) \in \lfloor \tau \rfloor_V$

Finally from Statement(1) we get $(\theta, n - i, v) \in \lfloor \tau' \rfloor_V$

□

Lemma 1.51 (Binary interpretation of Γ implies Unary interpretation of Γ). $\forall W, \gamma, \Gamma, n.$

$$(W, n, \gamma) \in \lfloor \Gamma \rfloor_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \Downarrow_i) \in \lfloor \Gamma \rfloor_V$$

Proof. Given: $(W, n, \gamma) \in [\Gamma]_V^A$

To prove: $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 1.41 we know that we are given:

$dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$

And we are required to prove:

$\forall i \in \{1, 2\}. \forall m.$

$dom(\Gamma) \subseteq dom(\gamma \downarrow_i) \wedge \forall x \in dom(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Case $i = 1$

Given some m we need to show:

- $dom(\Gamma) \subseteq dom(\gamma \downarrow_1)$:

$$dom(\gamma) = dom(\gamma \downarrow_1)$$

Therefore, $dom(\Gamma) \subseteq (dom(\gamma) = dom(\gamma \downarrow_1))$ (Given)

- $\forall x \in dom(\Gamma). (W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$:

We are given: $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$

Therefore from Lemma 1.42 we know that

$$\forall m'. (W.\theta_1, m', \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Instantiating m' with m we get

$$(W.\theta_1, m, \gamma \downarrow_1(x)) \in [\Gamma(x)]_V$$

Case $i = 2$

Symmetric reasoning as in the $i = 1$ case above

□

Theorem 1.52 (Fundamental theorem binary). $\forall \Gamma, pc, W, \mathcal{A}, e, \tau, \gamma, n.$

$\Gamma \vdash e : \tau \wedge$

$(W, n, \gamma) \in [\Gamma]_V^A \implies$

$(W, n, e(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau]_E^A$

Proof. Proof by induction on the typing derivation

1. CG-var:

$$\frac{}{\Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

To prove: $(W, n, x(\gamma \downarrow_1), x(\gamma \downarrow_2)) \in [\tau]_E^A$

Say $e_1 = x(\gamma \downarrow_1)$ and $e_2 = x(\gamma \downarrow_2)$

From Definition 1.33 it suffices to prove that

$$\forall i < n. e_1 \downarrow_i v'_1 \wedge e_2 \downarrow_i v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

This means given some $i < n$ s.t $e_1 \downarrow_i v'_1 \wedge e_2 \downarrow_i v'_2$

We are required to prove: $(W, n - i, v'_1, v'_2) \in [\tau]_V^A$

From cg-val we know that $x (\gamma \downarrow_1) \Downarrow x (\gamma \downarrow_1)$ and $x (\gamma \downarrow_2) \Downarrow x (\gamma \downarrow_2)$

This means $v'_1 = x (\gamma \downarrow_1)$ and $v'_2 = x (\gamma \downarrow_2)$

Since $(W, n, \gamma) \in [\tau]_V^A$. Therefore from Definition 1.41 we know that

$(W, n, v'_1, v'_2) \in [\tau]_V^A$

From Lemma 1.44 we get

$(W, n - i, v'_1, v'_2) \in [\tau]_V^A$

2. CG-lam:

$$\frac{\Gamma, x : \tau_1 \vdash e_i : \tau_2}{\Gamma \vdash \lambda x. e_i : (\tau_1 \rightarrow \tau_2)}$$

To prove: $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2)]_E^A$

Say $e_1 = \lambda x. e (\gamma \downarrow_1)$ and $e_2 = \lambda x. e (\gamma \downarrow_2)$

From Definition of $[(\tau_1 \rightarrow \tau_2)]_E^A$ it suffices to prove that

$\forall i < n. e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$

This means given some $i < n$ s.t $e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2$

From cg-val we know that $v'_1 = (\lambda x. e_i) \gamma \downarrow_1$ and $v'_2 = (\lambda x. e_i) \gamma \downarrow_2$

We are required to prove:

$(W, n - i, (\lambda x. e_i) \gamma \downarrow_1, (\lambda x. e_i) \gamma \downarrow_2) \in [\tau]_V^A$

From Definition 1.32 it suffices to prove

$\forall W' \sqsupseteq W, j < n, v_1, v_2.$

$((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_1) \in [\tau_2]_E^A) \wedge$

$\forall \theta_l \sqsupseteq W. \theta_1, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2]_E) \wedge$

$\forall \theta_l \sqsupseteq W. \theta_2, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2]_E) \quad \text{(FB-L0)}$

IH:

$\forall W, n. (W, n, e_i (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e_i (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2]_E^A$

s.t

$(W, n, (\gamma \cup \{x \mapsto (v_1, v_2)\})) \in [\Gamma]_V^A$

In order to prove (FB-L0) we need to prove the following:

- (a) $\forall W' \sqsupseteq W, j < n, v_1, v_2.$
 $((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2]_E^A):$

This means given some $W' \sqsupseteq W, j < n, v_1, v_2$ s.t. $(W', j, v_1, v_2) \in [\tau_1]_V^A$
We need to prove $(W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2]_E^A$

We get this by instantiating IH with W' and j

- (b) $\forall \theta_l \sqsupseteq W.\theta_1, v_c, j.$
 $((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2]_E):$
This means given some $\theta_l \sqsupseteq W.\theta_1, v_c, j$ s.t. $(\theta_l, j, v_c) \in [\tau_1]_V$
We need to prove: $(\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2]_E$

It is given to us that

$$(W, n, \gamma) \in [\Gamma]_V^A$$

Therefore from Lemma 1.51 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$$

Instantiating m with j we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$$

From Lemma 1.46 we know that

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Since we know that $(\theta_l, j, v_c) \in [\tau_1]_V$

Therefore we also have

$$(\theta_l, j, \gamma \downarrow_1 \cup \{x \mapsto v_c\}) \in [\Gamma \cup \{x \mapsto \tau_1\}]_V$$

Therefore, we can apply Theorem 1.49 to obtain

$$(\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2]_V$$

- (c) $\forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$
 $((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2]_E):$
Similar reasoning as in the previous case

3. CG-app:

$$\frac{\Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2}$$

To prove: $(W, n, (e_1 e_2) (\gamma \downarrow_1), (e_1 e_2) (\gamma \downarrow_2)) \in [(\tau_2)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. (e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau_2]_V^A$$

This further means that given some $i < n$ s.t. $(e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2}$

It suffices to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau_2]_V^A$$

IH1: $(W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2)]_E^A$

This means from Definition 1.33 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \downarrow_j v_{h_1} \wedge e_1 \gamma \downarrow_2 \downarrow v_{h_2} \implies (W, n - j, v_{h_1}, v_{h_2}) \in [(\tau_1 \rightarrow \tau_2)]_V^A$$

Since we know that $(e_1 \ e_2) \gamma \downarrow_1 \downarrow_i v_{f_1}$. Therefore $\exists j < i < n$ s.t $e_1 \gamma \downarrow_1 \downarrow_j v_{h_1}$. Similarly since $(e_1 \ e_2) \gamma \downarrow_2 \downarrow v_{f_2}$ therefore $e_1 \gamma \downarrow_2 \downarrow v_{h_2}$

This means we have $(W, n - j, v_{h_1}, v_{h_2}) \in [(\tau_1 \rightarrow \tau_2)]_V^A$

From cg-app we know that $val_{h_1} = \lambda x. e_{h_1}$ and $val_{h_2} = \lambda x. e_{h_2}$

From Definition 1.32 this further means

$$\begin{aligned} \forall W' \sqsupseteq W, J < (n - j), v_1, v_2. \\ ((W', J, v_1, v_2) \in [\tau_1]_V^A \implies (W', J, e_{h_1}[v_1/x], e_{h_2}[v_2/x]) \in [\tau_2]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, v_c, j. \\ ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, v_c, j. \\ ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E) \end{aligned} \quad (\text{FB-A1})$$

$$\underline{\text{IH2}}: (W, n - j, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_1]_E^A$$

This means from Definition 1.33 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \downarrow_j v_{h_1'} \wedge e_2 \gamma \downarrow_2 \downarrow v_{h_2'} \implies (W, n - j - k, v_{h_1'}, v_{h_2'}) \in [\tau_1]_V^A$$

Since we know that $(e_1 \ e_2) \gamma \downarrow_1 \downarrow_i v_{f_1}$. Therefore $\exists k < i - j < n - j$ s.t $e_2 \gamma \downarrow_1 \downarrow_k v_{h_1'}$. Similarly since $(e_1 \ e_2) \gamma \downarrow_2 \downarrow v_{f_2}$ therefore $e_2 \gamma \downarrow_2 \downarrow v_{h_2'}$

This means we have $(W, n - j - k, v_{h_1'}, v_{h_2'}) \in [\tau_1]_V^A$ (FB-A2)

Instantiating the first conjunct of (FB-A1) as follows W' with W , J with $n - j - k$, v_1 and v_2 with v'_{h_1} and v'_{h_2} respectively, we obtain

$$(W, n - j - k, e_{h_1}[v'_{h_1}/x], e_{h_2}[v'_{h_2}/x]) \in [\tau_2]_E^A$$

From Definition 1.33

$$\forall l < n - j - k. (e_{h_1}[v'_{h_1}/x]) \gamma \downarrow_l v_{f_1} \wedge e_{h_2}[v'_{h_2}/x] \downarrow v_{f_2} \implies (W, n - j - k - l, v_{f_1}, v_{f_2}) \in [\tau_2]_V^A$$

Since we know that $(e_1 \ e_2) \gamma \downarrow_1 \downarrow_i v_{f_1}$. Therefore $\exists l < i - j - k < n - j - k$ s.t $e_{h_1}[v'_{h_1}/x] \downarrow_l v_{f_1}$. Similarly since $(e_1 \ e_2) \gamma \downarrow_2 \downarrow v_{f_2}$ therefore $e_{h_2}[v'_{h_2}/x] \downarrow v_{f_2}$

Therefore we have $(W, n - j - k - l, v_{f_1}, v_{f_2}) \in [\tau_2]_V^A$

Since $i = j + k + l$ threfore we are done

4. CG-prod:

$$\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

To prove: $(W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} \forall i < n. (e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2}) \wedge (e_1, e_2) \gamma \downarrow_2 \downarrow (v'_{f_1}, v'_{f_2}) \implies \\ (W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A \end{aligned}$$

This means that given some $i < n$ s.t $(e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2}) \wedge (e_1, e_2) \gamma \downarrow_2 \downarrow (v'_{f_1}, v'_{f_2})$

We are required to prove

$$(W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A \quad (\text{FB-P0})$$

$$\underline{\text{IH1}}: (W, n, e_1 (\gamma \downarrow_1), e_1 (\gamma \downarrow_2)) \in [\tau_1]_E^A$$

This means from Definition 1.33 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \downarrow_i v_{f_1} \wedge e_1 \gamma \downarrow_2 \downarrow v'_{f_1} \implies (W, n - j, (v_{f_1}, v'_{f_1})) \in [\tau_1]_V^A$$

Since we know that $(e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2})$. Therefore $\exists j < i < n$ s.t $e_1 \gamma \downarrow_1 \downarrow_j v_{f_1}$. Similarly since $(e_1, e_2) \gamma \downarrow_2 \downarrow v_{f_2}$ therefore $e_1 \gamma \downarrow_2 \downarrow v'_{f_1}$

This means we have

$$(W, n - j, (v_{f_1}, v'_{f_1})) \in [\tau_1]_V^A \quad (\text{FB-P1})$$

$$\underline{\text{IH2}}: (W, n - j, e_2 (\gamma \downarrow_1), e_2 (\gamma \downarrow_2)) \in [\tau_2]_E^A$$

This means from Definition 1.33 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \downarrow_i v_{f_2} \wedge e_2 \gamma \downarrow_2 \downarrow v'_{f_2} \implies (W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2]_V^A$$

Since we know that $(e_1, e_2) \gamma \downarrow_1 \downarrow_i (v_{f_1}, v_{f_2})$. Therefore $\exists k < i - j < n - j$ s.t $e_2 \gamma \downarrow_1 \downarrow_j v_{f_2}$. Similarly since $(e_1, e_2) \gamma \downarrow_2 \downarrow v_{f_2}$ therefore $e_2 \gamma \downarrow_2 \downarrow v'_{f_2}$

This means we have

$$(W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2]_V^A \quad (\text{FB-P2})$$

In order to prove (FB-P0) from Definition 1.32 it suffices to prove that

$$(W, n - i, (v_{f_1}, v'_{f_1})) \in [\tau_1]_V^A \text{ and } (W, n - i, (v_{f_2}, v'_{f_2})) \in [\tau_2]_V^A$$

Since $i = j + k + 1$ therefore from (FB-P1) and (FB-P2) and from Lemma 1.44 we get

$$(W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A$$

5. CG-fst:

$$\frac{\Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Gamma \vdash \text{fst}(e') : \tau_1}$$

To prove: $(W, n, \text{fst}(e') (\gamma \downarrow_1), \text{fst}(e') (\gamma \downarrow_2)) \in [(\tau_1)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{fst}(e') \gamma \downarrow_1 \downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \downarrow v'_{f_1} \implies (W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1]_V^A$$

This means that given some $i < n$ s.t $\text{fst}(e') \gamma \downarrow_1 \downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \downarrow v'_{f_1}$

We are required to prove

$$(W, n - i, v_{f_1}, v_{f_1}) \in [\tau_1]_V^A \quad (\text{FB-F0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)]_E^A$$

This means from Definition 1.33 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2}) \wedge e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, v'_{f_2}) \implies \\ (W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A$$

Since we know that $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1}$. Therefore $\exists j < i < n$ s.t $e' \gamma \downarrow_1 \Downarrow_j (v_{f_1}, -)$. Similarly since $\text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$ therefore $e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, -)$

This means we have

$$(W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2)]_V^A$$

From Definition 1.32 we know that

$$(W, n - j, v_{f_1}, v'_{f_1}) \in [\tau_1]_V^A$$

Since from cg-fst $i = j + 1$ therefore from Lemma 1.44 we get

$$(W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1]_V^A$$

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Gamma \vdash e' : \tau_1}{\Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

To prove: $(W, n, \text{inl}(e') (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{inl}(e') \gamma \downarrow_1 \Downarrow_i \text{inl}(v_{f_1}) \wedge \text{inl}(e') \gamma \downarrow_2 \Downarrow \text{inl}(v'_{f_1}) \implies \\ (W, n - i, \text{inl}(v_{f_1}), \text{inl}(v'_{f_1})) \in [(\tau_1 + \tau_2)]_V^A$$

This means that given some $i < n$ s.t $\text{inl}(e') \gamma \downarrow_1 \Downarrow_i \text{inl}(v_{f_1}) \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow \text{inl}(v'_{f_1})$

We are required to prove

$$(W, n - i, \text{inl}(v_{f_1}), \text{inl}(v'_{f_1})) \in [(\tau_1 + \tau_2)]_V^A \quad (\text{FB-IL0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)]_E^A$$

This means from Definition 1.33 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{f_1} \implies \\ (W, n - j, v_{f_1}, v'_{f_1}) \in [\tau_1]_V^A$$

Since we know that $\text{inl}(e') \gamma \downarrow_1 \Downarrow_i \text{inl}(v_{f_1})$. Therefore $\exists j < i < n$ s.t $e' \gamma \downarrow_1 \Downarrow_j v_{f_1}$. Similarly since $\text{fst}(e') \gamma \downarrow_2 \Downarrow \text{inl}(v'_{f_1})$ therefore $e' \gamma \downarrow_2 \Downarrow v'_{f_1}$

This means we have

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau_1]_V^A \quad (\text{FB-IL1})$$

In order to prove (FB-IL0) from Definition 1.32 it suffices to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau_1]_V^A$$

From cg-inl since $i = j + 1$ therefore from (FB-IL1) and Lemma 1.44 we get (FB-IL0)

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Gamma \vdash \text{case}(e_c, x.e_1, y.e_2) : \tau}$$

To prove: $(W, n, \text{case}(e_c, x.e_1, y.e_2) (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau]_V^A$$

This means that given some $i < n$ s.t $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v_{f2}$

We are required to prove

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau]_V^A \quad (\text{FB-C0})$$

IH1:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)]_E^A$$

This means from Definition 1.33 we have:

$$\forall j < n. e_c \gamma \downarrow_1 \Downarrow_j v_{h1} \wedge e_c \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2)]_V^A$$

Since we know that $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$. Therefore $\exists j < i < n$ s.t $e_c \gamma \downarrow_1 \Downarrow_j v_{h1}$. Similarly since $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v'_{h1}$ therefore $e_c \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2)]_V^A \quad (\text{FB-C1})$$

2 cases arise

(a) $v_{h1} = \text{inl}(v_1)$ and $v'_{h1} = \text{inl}(v'_1)$:

IH2:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)]_E^A$$

This means from Definition 1.33 we have:

$$\forall k < n - j. e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_i v_{h2} \wedge e_1 \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2} \implies (W, n - j - k, v_{h2}, v'_{h2}) \in [\tau]_V^A$$

Since we know that $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$. Therefore $\exists k < i - j < n - j$ s.t $e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_j v_{h2}$. Similarly since $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2}$ therefore $e_1 \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$$(W, n - j - k, v_{h2}, v'_{h2}) \in [\tau]_V^A$$

From cg-case1 we know that $i = j + k + 1$ therefore from Lemma 1.44 we get (FB-C0)

(b) $v_{h1} = \text{inr}(v_1)$ and $v'_{h1} = \text{inr}(v'_1)$:

Symmetric case

10. CG-label:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

To prove: $(W, n, \text{Lb}(e') (\gamma \downarrow_1), \text{Lb}(e') (\gamma \downarrow_2)) \in [\text{Labeled } \ell \tau]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{Lb}(e') \gamma \downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{f1}) \implies (W, n - i, \text{Lb}(v_{f1}), \text{Lb}(v'_{f1})) \in [\text{Labeled } \ell \tau]_V^A$$

This means that given some $i < n$ s.t $\text{Lb}(e') \gamma \downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{f1})$

We are required to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\text{Labeled } \ell \tau]_V^A \quad (\text{FB-LB0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau]_E^A$$

This means from Definition 1.33 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{f1} \implies (W, n - j, v_{f1}, v'_{f1}) \in [\tau]_V^A$$

Since we know that $\text{Lb}(e') \gamma \downarrow_1 \Downarrow_i \text{Lb}(v_{f1})$. Therefore $\exists j < i < n$ s.t $e' \gamma \downarrow_1 \Downarrow_j v_{f1}$. Similarly since $\text{Lb}(e') \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{f1})$ therefore $e' \gamma \downarrow_2 \Downarrow v'_{f1}$

This means we have

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau]_V^A \quad (\text{FB-LB1})$$

In order to prove (FB-LB0) from Definition 1.32 it suffices to prove that

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau]_V^A$$

From cg-label we know that $i = j + 1$. Therefore we get the desired from (FB-LB1) and Lemma 1.44

11. CG-unlabel:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell \tau}{\Gamma \vdash \text{unlabel}(e') : \mathbb{C} \top \ell \tau}$$

To prove: $(W, n, \text{unlabel}(e') (\gamma \downarrow_1), \text{unlabel}(e') (\gamma \downarrow_2)) \in [(\mathbb{C} \top \ell \tau)]_E^A$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f1} \implies (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \top \ell \tau)]_V^A$$

This means that given some $i < n$ s.t $\text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that $v_{f1} = \text{unlabel}(e') \gamma \downarrow_1$ and $v'_{f1} = \text{unlabel}(e') \gamma \downarrow_2$. Also $i = 0$

We are required to prove

$$(W, n, \text{unlabel}(e') \gamma \downarrow_1, \text{unlabel}(e') \gamma \downarrow_2) \in [(\mathbb{C} \top \ell \tau)]_V^A$$

This means from Definition 1.32 we need to prove

Let $e_1 = \text{unlabel}(e') \gamma \downarrow_1$ and $e_2 = \text{unlabel}(e') \gamma \downarrow_2$

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau' \rfloor_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \right) \end{aligned}$$

We need to show

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \tau): \end{aligned}$$

Also given is some $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$ s.t $(k, H_1, H_2) \triangleright W_e$ and $(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k$

And we are required to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \tau) \quad (\text{FB-U0})$$

$$\underline{\text{IH}}: (W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{Labeled } \ell \tau)]_E^A$$

This means from Definition 1.33 we are given

$$\forall I < k. e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1}) \implies (W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in [(\text{Labeled } \ell \tau)]_V^A$$

Since we know that

$(H_1, \text{unlabel}(e') \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{unlabel}(e') \gamma \downarrow_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k$ therefore
 $\exists I < j < k$ s.t $e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1})$

Therefore we have

$$(W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in [(\text{Labeled } \ell \tau)]_V^A$$

This means from Definition 1.32 we have

$$\text{ValEq}(\mathcal{A}, W_e, k - I, \ell, v_{h1}, v'_{h1}, \tau) \quad (\text{FB-U1})$$

In order to prove (FB-U0) we choose W' as W_e and from cg-unlabel we know that $H'_1 = H_1$ and $H'_2 = H_2$. And we already know that $(k, H_1, H_2) \triangleright W_e$. Therefore from Lemma 1.48 we get $(k - j, H_1, H_2) \triangleright W_e$

From cg-unlabel we know that v'_1, v'_2 in (FB-U0) is v_{h1}, v'_{h1} respectively. And since from (FB-U1) we know that $\text{ValEq}(\mathcal{A}, W_e, k - I, \ell, v_{h1}, v'_{h1}, \tau)$. Therefore from Lemma 1.53 we get

$$\text{ValEq}(\mathcal{A}, W_e, k - j, \ell, v_{h1}, v'_{h1}, \tau)$$

$$(b) \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \right):$$

Case $l = 1$

$$\text{Given some } k, \theta_e \sqsupseteq W.\theta_l, H, j \text{ s.t } (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k$$

We need to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \top \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top) \end{aligned}$$

Since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 1.51 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

$$\text{Instantiating } m \text{ with } k \text{ we get } (W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$$

Now we can apply Theorem 1.49 to get

$$(W.\theta_1, k, (\text{unlabel } e') \gamma \downarrow_1) \in [(\mathbb{C} \top \ell \tau)]_E$$

This means from Definition 1.35 we get

$$\forall c < k. (\text{unlabel } e') \gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \top \ell \tau)]_V$$

This further means that given some $c < k$ s.t $(\text{unlabel } e') \gamma \downarrow_1 \Downarrow_c v$. From cg-val we know that $c = 0$ and $v = (\text{unlabel } e') \gamma \downarrow_1$

$$\text{And we have } (W.\theta_1, k, (\text{unlabel } e') \gamma \downarrow_1) \in [(\mathbb{C} \top \ell \tau)]_V$$

From Definition 1.34 we have

$$\begin{aligned} & \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{unlabel } e') \gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies \\ & \exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\tau]_V \wedge \\ & (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \top \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \top) \end{aligned}$$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

12. CG-tolabeled:

$$\frac{\Gamma \vdash e' : \mathbb{C} \ell_1 \ell_2 \tau}{\Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_1 \perp (\text{Labeled} \ell_2 \tau)}$$

To prove: $(W, n, \text{toLabeled}(e') (\gamma \downarrow_1), \text{toLabeled}(e') (\gamma \downarrow_2)) \in [\mathbb{C} \ell_1 \perp (\text{Labeled} \ell_2 \tau)]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} \forall i < n. \text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f_1} &\implies \\ (W, n - i, v_{f_1}, v'_{f_1}) \in [\mathbb{C} \ell_1 \perp (\text{Labeled} \ell_2 \tau)]_V^A & \end{aligned}$$

This means that given some $i < n$ s.t $\text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$

From cg-val we know that $v_{f_1} = \text{toLabeled}(e') \gamma \downarrow_1$, $v_{f_2} = \text{toLabeled}(e') \gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, \text{toLabeled}(e') \gamma \downarrow_1, \text{toLabeled}(e') \gamma \downarrow_2) \in [\mathbb{C} \ell_1 \perp (\text{Labeled} \ell_2 \tau)]_V^A$$

Let $v_1 = \text{toLabeled}(e') \gamma \downarrow_1$ and $v_2 = \text{toLabeled}(e') \gamma \downarrow_2$

This means from Definition 1.32 we are required to prove

$$\begin{aligned} &(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled} \ell_2 \tau))) \wedge \\ &\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled} \ell_o \tau)]_V \wedge \\ &(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled} \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ &(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \end{aligned}$$

We need to prove:

$$\begin{aligned} \text{(a)} \quad &\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled} \ell_2 \tau)): \end{aligned}$$

This means that we are given some $k \leq n$, $W_e \sqsupseteq W$, $H_1, H_2, v'_1, v'_2, j < k$ s.t

$$(k, H_1, H_2) \triangleright W_e \text{ and } (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$$

And we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled} \ell_2 \tau))$$

From Definition 1.31 it suffices to prove that

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge (W', k - j, v'_1, v'_2) \in [(\text{Labeled} \ell_2 \tau)]_V^A$$

Further from Definition 1.32 it suffices to prove

$$\begin{aligned} & \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v''_1, v''_2, \tau) \quad (\text{FB-TL0}) \\ & \text{where } v''_1 = \text{Lb } v''_1 \text{ and } v''_2 = \text{Lb } v''_2 \end{aligned}$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\mathbb{C} \ell_1 \ell_2 \tau]_E^A$$

This means from Definition 1.33 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \downarrow v'_{h1} \implies (W_e, n - J, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$$

Since we know that $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \downarrow_j (H'_1, v'_1)$ and $(H_2, \text{toLabeled}(e')\gamma \downarrow_1) \downarrow_j (H'_2, v'_2)$. Therefore from cg-val we know that $\exists J < j < k \leq n$ s.t $e' \gamma \downarrow_1 \downarrow_J v_{h1}$ and similarly we also know that $e' \gamma \downarrow_2 \downarrow v'_{h1}$

This means we have

$$(W_e, k - J, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$$

From Definition 1.32 we know that

$$\begin{aligned} & (\forall k_1 \leq (k - J), W_e'' \sqsupseteq W_e. \forall H''_1, H''_2. (k_1, H''_1, H''_2) \triangleright W_e'' \wedge \forall v''_1, v''_2, m. \\ & (H''_1, v_{h1}) \downarrow_m^f (H''_1, v''_1) \wedge (H''_2, v'_{h1}) \downarrow^f (H''_2, v''_2) \wedge m < k_1 \implies \\ & \exists W' \sqsupseteq W_e''. (k_1 - m, H''_1, H''_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k_1 - m, \ell_2, v''_1, v''_2, \tau)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \quad (\text{FB-TL1}) \end{aligned}$$

We instantiate W_e'' with W_e , H''_1 with H_1 , H''_2 with H_2 and k_1 with k in (FB-TL1).

Since we know that $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{toLabeled}(e')\gamma \downarrow_2) \downarrow_j^f (H'_2, v'_2)$, therefore $\exists m < j < k \leq n$ s.t $(H_1, v_{h1}) \downarrow_m^f (H'_1, v'_1) \wedge (H_2, v'_{h1}) \downarrow^f (H'_2, v'_2)$

This means we have

$$\exists W' \sqsupseteq W_e.(k - m, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - m, \ell_2, v''_1, v''_2, \tau) \quad (\text{FB-TL2})$$

In order to prove (FB-TL0) we choose W' as W' from (FB-TL2). Since from cg-tolabeled we know that $v''_1 = \text{Lb}(v''_1)$, $v''_2 = \text{Lb}(v''_2)$ and $j = m + 1$ (therefore from Lemma 1.48 we get $(k - j, H'_1, H'_2) \triangleright W'$) and from (FB-TL2) and Lemma 1.53 we get $\text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v''_1, v''_2, \tau)$

$$\begin{aligned} \text{(b)} \quad & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1))): \end{aligned}$$

Case $l = 1$

Given some $k, \theta_e \sqsupseteq W. \theta_l, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, v_l) \downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_2 \tau)]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 1.51 we know that
 $\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$ and $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$$(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_E$$

This means from Definition 1.35 we get

$$\forall c < k. (\text{toLabeled } e')\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$$

Instantiating c with 0 and from cg-val we know $v = (\text{toLabeled } e')\gamma \downarrow_1$

And we have $(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \perp \text{Labeled } \ell_2 \tau)]_V$

From Definition 1.34 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{toLabeled } e')\gamma \downarrow_1) \Downarrow_j^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\text{Labeled } \ell_2 \tau]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

13. CG-ret:

$$\frac{\Gamma \vdash e' : \tau}{\Gamma \vdash \text{ret}(e') : \mathbb{C} \ell_1 \ell_2 \tau}$$

To prove: $(W, n, \text{ret}(e') (\gamma \downarrow_1), \text{ret}(e') (\gamma \downarrow_2)) \in [\mathbb{C} \ell_1 \ell_2 \tau]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} \forall i < n. \text{ret}(e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A \end{aligned}$$

This means that given some $i < n$ s.t $\text{ret}(e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \downarrow v'_{f1}$

From cg-val we know that $v_{f1} = \text{ret}(e')\gamma \downarrow_1$, $v_{f2} = \text{ret}(e')\gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, \text{ret}(e')\gamma \downarrow_1, \text{ret}(e')\gamma \downarrow_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$$

Let $v_1 = \text{ret}(e')\gamma \downarrow_1$ and $v_2 = \text{ret}(e')\gamma \downarrow_2$

From Definition 1.32 it suffices to prove

$$\begin{aligned} (\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \end{aligned}$$

$$\begin{aligned}
& \forall l \in \{1, 2\}. \left(\forall v, i. (e_l \Downarrow_i v_l) \implies \right. \\
& \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\
& \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right)
\end{aligned}$$

It suffices to prove:

$$\begin{aligned}
\text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\
& (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau):
\end{aligned}$$

We are given is some $k \leq n$, $W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$ s.t $(k, H_1, H_2) \triangleright W_e$ and $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

From cg-ret we know that $H'_1 = H_1$ and $H'_2 = H_2$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H_1, H_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \quad (\text{FB-R0})$$

$$\underline{\text{IH}}: (W_e, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau]_E^A$$

This means from Definition 1.33 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \Downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, k - J, v_{h1}, v'_{h1}) \in [\tau]_V^A$$

Since we know that $(H_1, \text{ret}(e')\gamma \downarrow_1) \Downarrow_j^f (H_1, v'_1) \wedge (H_2, \text{ret}(e')\gamma \downarrow_2) \Downarrow_j^f (H_2, v'_2)$, therefore $\exists J < j < k$ s.t $e' \gamma \downarrow_1 \Downarrow_J v_{h1}$ and similarly $e' \gamma \downarrow_2 \Downarrow v'_{h1}$.

$$\text{Therefore we have } (W_e, k - J, v_{h1}, v'_{h1}) \in [\tau]_V^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose W' as W_e and from cg-ret we know that $v'_1 = v_{h1}$ and $v'_2 = v'_{h1}$. We need to prove the following:

i. $(k - j, H_1, H_2) \triangleright W_e$:

Since we have $(k, H_1, H_2) \triangleright W_e$ therefore from Lemma 1.48 we get

$$(k - j, H_1, H_2) \triangleright W_e$$

ii. $\text{ValEq}(\mathcal{A}, W_e, k - j, \ell_2, v'_1, v'_2, \tau)$:

2 cases arise:

A. $\ell_2 \sqsubseteq \mathcal{A}$:

In this case from Definition 1.31 it suffices to prove

$$(W_e, k - j, v'_1, v'_2) \in [\tau]_V^A$$

Since $j = J + 1$ therefore we get this from (FB-R1) and Lemma 1.44

B. $\ell_2 \not\sqsubseteq \mathcal{A}$:

In this case from Definition 1.31 it suffices to prove that

$$\forall m. (W_e, m, v'_1) \in [\tau]_V \text{ and } \forall m. (W_e, m, v'_2) \in [\tau]_V$$

We get this From (FB-R1) and Lemma 1.42

(b) $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1):$

Case $l = 1$

Given some $k, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t. $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_o \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_o)$

Since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 1.51 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$ and $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_E$

This means from Definition 1.35 we get

$\forall c < k. (\text{ret } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V$

Instantiating c with 0 and from cg-val we know that $v = (\text{ret } e')\gamma \downarrow_1$

And we have $(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V$

From Definition 1.34 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies$
 $\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\tau]_V \wedge$
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_1)$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

14. CG-bind:

$$\frac{\Gamma \vdash e_l : \mathbb{C} \ell_1 \ell_2 \tau \quad \Gamma, x : \tau \vdash e_b : \mathbb{C} \ell_3 \ell_4 \tau' \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_l, x.e_b) : \mathbb{C} \ell \ell' \tau'}$$

To prove: $(W, n, \text{bind}(e_l, x.e_b) (\gamma \downarrow_1), \text{bind}(e_l, x.e_b) (\gamma \downarrow_2)) \in [\mathbb{C} \ell \ell' \tau']_E^A$

This means from Definition 1.33 we need to prove:

$\forall i < n. \text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1} \implies$
 $(W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell \ell' \tau']_V^A$

This means that given some $i < n$ s.t. $\text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that $v_{f_1} = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$, $v_{f_2} = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, \text{bind}(e_l, x.e_b)\gamma \downarrow_1, \text{bind}(e_l, x.e_b)\gamma \downarrow_2) \in [\mathbb{C} \ell \ell' \tau']_V^A$$

Let $v_1 = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$ and $v_2 = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$

This means from Definition 1.32 we need to prove

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell', v'_1, v'_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \rfloor_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell', v'_1, v'_2, \tau): \end{aligned}$$

This means we are given some $k \leq n$, $W_e \sqsupseteq W, H_1, H_2$ s.t $(k, H_1, H_2) \triangleright W_e$

Also given some $v'_1, v'_2, j < k$ s.t $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell', v'_1, v'_2, \tau) \quad (\text{FB-B0})$$

IH1:

$$(W_e, k, e_l (\gamma \downarrow_1), e_l (\gamma \downarrow_2)) \in [\mathbb{C} \ell_1 \ell_2 \tau]_E^A$$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} & \forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1} \implies \\ & (W_e, k - f, v_{h_1}, v'_{h_1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A \end{aligned}$$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$ therefore $\exists f < j < k$ s.t $e_l \gamma \downarrow_f \Downarrow_j v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1}$

This means we have

$$(W_e, k - f, v_{h_1}, v'_{h_1}) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$$

This means from Definition 1.32 we have

$$\begin{aligned} & \left(\forall K \leq (k - f), W_e \sqsupseteq W_e. \forall H''_1, H''_2. (K, H''_1, H''_2) \triangleright W_e \wedge \forall v''_1, v''_2, J. \right. \\ & (H''_1, v_{h_1}) \Downarrow_J^f (H''_1, v''_1) \wedge (H''_2, v'_{h_1}) \Downarrow_J^f (H''_2, v''_2) \wedge J < K \implies \\ & \left. \exists W'' \sqsupseteq W_e. (K - J, H''_1, H''_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', K - J, \ell_2, v''_1, v''_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \rfloor_V \wedge \right. \end{aligned}$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)$$

Instantiating K with $(k - f)$, W'_e with W_e , H_1'' with H_1 and H_2'' with H_2 in the first conjunct of the above equation. Since we know that $(k, H_1, H_2) \triangleright W_e$ therefore from Lemma 1.48 we also have $(k - f, H_1, H_2) \triangleright W_e$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$ therefore $\exists J < j - f < k - f$ s.t. $(H_1, v_{h1}) \Downarrow_J^f (H'_1, v''_1) \wedge (H_2, v'_{h1}) \Downarrow_J^f (H'_2, v''_2)$

This means we have

$$\exists W'' \sqsupseteq W'_e. (k - f - J, H'_1, H'_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', k - f - J, \ell_2, v''_1, v''_2, \tau) \quad (\text{FB-B1})$$

From Definition 1.31 two cases arise:

i. $\ell_2 \sqsubseteq \mathcal{A}$:

In this case we know that $(W'', k - f - J, v''_1, v''_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

IH2:

$$(W'', k - f - J, e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}), e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\})) \in [\mathbb{C} \ell_3 \ell_4 \tau']_{\mathcal{E}}^{\mathcal{A}}$$

This means from Definition 1.33 we need to prove:

$$\forall s < k - f - J. e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \Downarrow_s v_{h2} \wedge e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \Downarrow v'_{h2} \implies \\ (W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\mathbb{C} \ell_3 \ell_4 \tau']_{\mathcal{V}}^{\mathcal{A}}$$

Since we know that $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow_j^f (H'_2, v'_2)$ therefore $\exists s < j - f - J < k - f - J$ s.t. $e_b (\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \Downarrow_s v_{h2} \wedge e_b (\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \Downarrow v'_{h2}$

This means we have

$$(W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\mathbb{C} \ell_3 \ell_4 \tau']_{\mathcal{V}}^{\mathcal{A}}$$

This means from Definition 1.32 we know that

$$\left(\forall K_s \leq (k - f - J - s), W_s \sqsupseteq W'' . \forall H_1, H_2. (K_s, H_1, H_2) \triangleright W_s \wedge \forall v'_{s1}, v'_{s2}, J_s. \right.$$

$$(H_1, v_{h2}) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H_2, v'_{h2}) \Downarrow_{J_s}^f (H'_{s2}, v'_{s2}) \wedge J_s < K_s \implies$$

$$\left. \exists W'_s \sqsupseteq W_s. (K_s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, K_s - J_s, \ell_4, v'_1, v'_2, \tau') \right) \wedge$$

$$\forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$$

$$\left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_{\mathcal{V}} \wedge \right.$$

$$\left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_3 \sqsubseteq \ell') \wedge \right.$$

$$\left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_3) \right)$$

Instantiating K_s with $(k - f - J - s)$, W_s with W'' , H_1 with H'_1 and H'_2 with H_2 . Since we know that $(k - f - J, H'_1, H'_2) \triangleright W''$ therefore from Lemma 1.48 we also have $(k - f - J - s, H'_1, H'_2) \triangleright W''$

Since we know that $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow_j^f (H'_2, v'_2)$ therefore $\exists J_s < j - f - J - s < k - f - J - s$ s.t. $(H'_1, v'_1) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H'_2, v'_2) \Downarrow_{J_s}^f (H'_{s2}, v'_{s2})$

This means we have

$$\exists W'_s \sqsupseteq W_s.(k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_4, v'_{s1}, v'_{s2}, \tau') \quad (\text{FB-B2})$$

In order to prove (FB-B0) we choose W' as W'_s . From cg-bind we know that $H'_1 = H'_{s1}$, $H'_2 = H'_{s2}$, $v'_1 = v'_{s1}$, $v'_2 = v'_{s2}$ and $j = f + J + s + J_s + 1$. And we need to prove:

A. $(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$:

Since from (FB-B2) we know that $(k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s$ therefore from Lemma 1.48 we get

$$(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$$

B. $\text{ValEq}(\mathcal{A}, W'_s, k - j, \ell', v'_{s1}, v'_{s2}, \tau')$:

Since from (FB-B2) we know that $\text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_4, v'_{s1}, v'_{s2}, \tau')$ therefore from Lemma 1.53 we get

$$\text{ValEq}(\mathcal{A}, W'_s, k - j, \ell', v'_{s1}, v'_{s2}, \tau')$$

ii. $\ell_2 \not\sqsubseteq \mathcal{A}$:

From (FB-B0) we know that we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau')$$

Since $\ell_2 \sqsubseteq \ell_4 \sqsubseteq \ell'$ and $\ell \not\sqsubseteq \mathcal{A}$ therefore we have $\ell_4 \not\sqsubseteq \mathcal{A}$ and $\ell' \not\sqsubseteq \mathcal{A}$

This means that from Definition 1.31 it suffices to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \forall m_{u1}. (W'.\theta_1, m_{u1}, v'_1) \in [\tau']_V \wedge \forall m_{u2}. (W'.\theta_2, m_{u2}, v'_2) \in [\tau']_V$$

This means given some m_{u1}, m_{u2} and we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge (W'.\theta_1, m_{u1}, v'_1) \in [\tau']_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in [\tau']_V \quad (\text{FB-B01})$$

In this case from (FB-B1) and Definition 1.31 we know that

$$\forall m. (W''.\theta_1, m, v''_1) \in [\tau]_V \text{ and } \forall m. (W''.\theta_2, m, v''_2) \in [\tau]_V \quad (\text{FB-B3})$$

Since $\text{bind}(e_l, x.e_b)\gamma \downarrow_1 \downarrow_j v'_1$ therefore $\exists J_1 < j - f - J < k - f - J$ s.t $(e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \downarrow_{J_1} v'_1$. Similarly, $\exists J'_1 < j - f - J - J_1 < k - f - J - J_1$ s.t $(H'_1, v'_1) \downarrow_{J'_1}^f -$

Instantiating m with $m_{u1} + 1 + J_1 + J'_1$ in the first conjunct of (FB-B3)

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, v''_1) \in [\tau]_V$$

Since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 1.51 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$$

Instantiating m with $m_{u1} + 1 + J_1 + J'_1$ we get $(W.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in [\Gamma]_V$

From Lemma 1.45 we know that

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in [\Gamma]_V \quad (\text{FB-B4})$$

Now we can apply Theorem 1.49 to get

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_E$$

This means from Definition 1.35 we get

$$\forall c_1 < m_{u1} + 1 + J_1 + J'_1. (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \downarrow_{c_1} v_{o1} \implies (W''.\theta_1, m_{u1} + 1 + J_1 + J'_1 - c_1, v_{o1}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V \quad (\text{FB-B5})$$

Instantiating c_1 with J_1 in (FB-B5)

$$\text{Therefore we have } (W''.\theta_1, m_{u1} + 1 + J'_1, v_{o1}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V$$

From Definition 1.34 we have

$$\forall K \leq (m_{u1} + 1 + J'_1), \theta'_e \sqsupseteq W''.\theta_1, H_1, J_2. (K, H_1) \triangleright \theta'_e \wedge (H_1, v_{o1}) \Downarrow_{J_2}^f (H''_1, v'_1) \wedge J_2 < K \implies$$

$$\begin{aligned} \exists \theta'_1 \sqsupseteq \theta'_e. (K - J_2, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, K - J_2, v'_1) \in [\tau']_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_3) \end{aligned}$$

Instantiating K with $m_{u1} + 1 + J'_1$, θ'_e with $W''.\theta_1$, H_1 with H'_1 (from FB-B1) and J_2 with J'_1 we get

$$\begin{aligned} \exists \theta'_1 \sqsupseteq W''.\theta_1. (m_{u1} + 1, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in [\tau']_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_3) \quad (\text{FB-B6}) \end{aligned}$$

Since we know that $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \Downarrow v'_2$. Say this reduction happens in t steps. Therefore $\exists t_1 < t < k \leq n$ s.t. $(e_l)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{t_1} v_{l2}$ and similarly $\exists t_2 < t - t_1 < k - t_1$ s.t. $(H, v_{l2})\gamma \downarrow_2 \Downarrow_{t_2}^f (H''_2, v''_2)$

Again since $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \Downarrow v'_2$ therefore $\exists J_2 < t - t_1 - t_2 < k - t_1 - t_2$ s.t. $(e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{J_2} v'_2$. Similarly $\exists J'_2 < t - t_1 - t_2 - J_2 < k - t_1 - t_2 - J_2$ s.t. $(H'_2, v'_2) \downarrow_{J'_2}^f -$

Instantiating the second conjunct of (FB-B3) with $m_{u2} + 1 + J_2 + J'_2$ we get $(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, v''_2) \in [\tau]_V$

Again since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 1.51 we know that $\forall m. (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating m with $m_{u2} + 1 + J_2 + J'_2$ we get $(W.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V$

From Lemma 1.45 we know that

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V \quad (\text{FB-B7})$$

Now we can apply Theorem 1.49 to get

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, (e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_E$$

This means from Definition 1.35 we get

$$\forall c_2 < (m_{u2} + 1 + J_2 + J'_2). (e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{c_2} v_{o2} \implies (W''.\theta_2, m_{u2} + 1 + J_2 - c_2, v_{o2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V \quad (\text{FB-B8})$$

Instantiating c_2 with J_2 in (FB-B8) we get

$$(W''.\theta_2, m_{u2} + 1 + J'_2, v_{o2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V$$

From Definition 1.34 we have

$$\forall K \leq (m_{u2} + 1 + J'_2), \theta'_e \sqsupseteq W''.\theta_2, H_2, J_3. (K, H_2) \triangleright \theta'_e \wedge (H_2, v_{o2}) \Downarrow_{J_3}^f (H''_2, v'_2) \wedge J_3 < K \implies$$

$$\begin{aligned} \exists \theta'_2 \sqsupseteq \theta'_e. (K - J_3, H''_2) \triangleright \theta'_2 \wedge (\theta'_2, K - J_3, v'_2) \in [\tau']_V \wedge \\ (\forall a. H_2(a) \neq H''_2(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_e). \theta'_2(a) \searrow \ell_3) \end{aligned}$$

Instantiating K with $m_{u2} + 1 + J'_2$, θ'_e with $W''.\theta_2$, H_2 with H'_2 (from FB-B1) and J_3 with J'_2 , we get

$$\begin{aligned} & \exists \theta'_2 \sqsupseteq W''.\theta_2.(m_{u2} + 1, H''_2) \triangleright \theta'_2 \wedge (\theta'_2, m_{u2} + 1, v'_2) \in [\tau']_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_e). \theta'_2(a) \searrow \ell_3) \quad (\text{FB-B9}) \end{aligned}$$

In order to prove (FB-B01) we chose W' as W_n where W_n is defined as follows:

$$\begin{aligned} W_n.\theta_1 &= \theta'_1 \quad (\text{From (FB-B6)}) \\ W_n.\theta_2 &= \theta'_2 \quad (\text{From (FB-B9)}) \\ W_n.\hat{\beta} &= W''.\hat{\beta} \quad (\text{From (FB-B1)}) \end{aligned}$$

It suffices to prove

- $(k - j, H'_1, H'_2) \triangleright W_n$:

From Definition 1.37 we need to prove the following

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2):$$

From (FB-B6) we know that $(m_{u1} + 1, H''_1) \triangleright \theta'_1$ therefore from Definition 1.36 we know that $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly from (FB-B9) we know that $(m_{u2} + 1, H''_2) \triangleright \theta'_2$ therefore from Definition 1.36 we know that $\text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2)$

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)):$$

Since from (FB-B1) we know that $(k - f - J, H'_1, H'_2) \triangleright W''$ therefore from Definition 1.37 we know that $(W''.\hat{\beta}) \subseteq (\text{dom}(W''.\theta_1) \times \text{dom}(W''.\theta_2))$

Since from (FB-B6) and (FB-B9) we know that $W''.\theta_1 \sqsubseteq W_n.\theta_1$ and $W''.\theta_2 \sqsubseteq W_n.\theta_2$

Therefore we get

$$(W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$$

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, k - j - 1, H'_1(a_1), H'_2(a_2)) \in [\tau]_V^A):$$

4 cases arise for each $(a_1, a_2) \in W_n.\hat{\beta}$

$$A. H'_1(a_1) = H''_1(a_1) \wedge H'_2(a_2) = H''_2(a_2):$$

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

We know from that $(k - f - J, H'_1, H'_2) \triangleright W''$

Therefore from Definition 1.37 we have

$$\forall (a'_1, a'_2) \in (W''.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$$

Since $W_n.\hat{\beta} = W''.\hat{\beta}$ by construction therefore

$$\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$$

From (FB-B6) and (FB-B9) we know that $W''.\theta_1 \sqsubseteq \theta'_1$ and $W''.\theta_2 \sqsubseteq \theta'_2$ respectively.

Therefore from Definition 1.29

$$\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

To prove:

$$\overline{(W_n, k - j - 1, H_1''(a_1), H_2''(a_2))} \in \lceil W_n.\theta_1(a_1) \rceil_V^A:$$

From (FB-B1) we know that $(k - f - J, H_1', H_2') \triangleright^A W''$

This means from Definition 1.37 we know that

$$\forall (a_{i1}, a_{i2}) \in (W''.\hat{\beta}). W''.\theta_1(a_{i1}) = W''.\theta_2(a_{i2}) \wedge \\ (W'', k - f - J - 1, H_1'(a_{i1}), H_2'(a_{i2})) \in \lceil W''.\theta_1(a_{i1}) \rceil_V^A$$

Instantiating with a_1 and a_2 and since $W'' \sqsubseteq W_n$ and $k - j - 1 < k - f - J - 1$ (since $j = f + J + J_1 + 1$ therefore from Lemma 1.44 we get

$$\overline{(W_n, k - j - 1, H_1'(a_1), H_2'(a_2))} \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

B. $H_1'(a_1) \neq H_1''(a_1) \wedge H_2'(a_2) \neq H_2''(a_2)$:

To prove:

$$\overline{W_n.\theta_1(a_1)} = \overline{W_n.\theta_2(a_2)}$$

Same reasoning as in the previous case

To prove:

$$\overline{(W_n, k - j - 1, H_1''(a_1), H_2''(a_2))} \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

From (FB-B6) and (FB-B9) we know that

$$(\forall a. H_1'(a) \neq H_1''(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell') \\ (\forall a. H_2'(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W''.\theta_1(a_1) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell' \text{ and} \\ \exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell'$$

Since $\ell_2 \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_3 \not\sqsubseteq \mathcal{A}$.

Also from (FB-B6) and (FB-B9), $(m_{u1} + 1, H_1'') \triangleright \theta_1'$ and $(m_{u2} + 1, H_2'') \triangleright \theta_2'$.

Therefore from Definition 1.36 we have

$$(\theta_1', m_{u1}, H_1''(a_1)) \in \lceil \theta_1'(a_1) \rceil_V \text{ and} \\ (\theta_2', m_{u2}, H_2''(a_1)) \in \lceil \theta_2'(a_2) \rceil_V$$

Since m_{u1} and m_{u2} are arbitrary indices therefore from Definition 1.32 we get

$$\overline{(W_n, k - j - 1, H_1''(a_1), H_2''(a_2))} \in \lceil \theta_1'(a_1) \rceil_V^A$$

C. $H_1'(a_1) = H_1''(a_1) \wedge H_2'(a_2) \neq H_2''(a_2)$:

To prove:

$$\overline{W_n.\theta_1(a_1)} = \overline{W_n.\theta_2(a_2)}$$

Same reasoning as in the previous case

To prove:

$$\overline{(W_n, k - j - 1, H_1''(a_1), H_2''(a_2))} \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

From (FB-B9) we know that

$$(\forall a. H_2'(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell')$$

This means we have

$\exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell_3) \sqsubseteq \ell'$
 Since $\ell_2 \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_3 \not\sqsubseteq \mathcal{A}$.

Since from (FB-B1) we know that $(k - f - J, H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W''$ that means from Definition 1.37 that $(W'', k - f - J - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W''.\theta_1(a_1) \rceil_V^{\mathcal{A}}$. Since $W''.\theta_1(a_1) = W''.\theta_2(a_2) = \text{Labeled } \ell' \tau''$ and since $\ell' \not\sqsubseteq \mathcal{A}$ therefore from Definition 1.32 and Definition 1.31 we know that

Therefore
 $\forall m. (W''.\theta_1, m, H'_1(a_1)) \in W''.\theta_1(a_1) \quad (\text{F})$

Instantiating the (F) with m_{u1} and using Lemma 1.43 we get
 $(\theta'_1, m_{u1}, H'_1(a_1)) \in \theta'_1(a_1)$

Since from (FB-B9) we know that $(m_{u2} + 1, H''_2) \triangleright \theta'_2$ therefore from Definition 1.36 we know that $(\theta'_2, m_{u2}, H''_2(a_2)) \in \theta'_2(a_2)$
 Therefore from Definition 1.32 we get
 $(W', k - j - 1, H''_1(a_1), H''_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^{\mathcal{A}}$

D. $H'_1(a_1) \neq H''_1(a_1) \wedge H'_2(a_2) = H''_2(a_2)$:
 Symmetric reasoning as in the previous case

– $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V$:

Case $i = 1$

Given some m we need to prove

$\forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V$

This further means that given some $a_1 \in \text{dom}(W_n.\theta_1)$ we need to show
 $(W_n.\theta_1, m, H''_1(a_1)) \in \lfloor W_n.\theta_1(a_1) \rfloor_V$

Since $W_n.\theta_1 = \theta'_1$, it suffices to prove
 $(\theta'_1, m, H''_1(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V$

Like before we apply Theorem 1.49 on $e_b \gamma \downarrow_1 \cup \{x \mapsto v''_1\}$ but this time at $m + 1 + J_1 + J'_1$ to get

$\exists \theta'_1 \sqsupseteq W''.\theta_1.(m + 1, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in \lceil \tau' \rceil_V \wedge$
 $(\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_3 \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_3)$

Since we have $\ell \sqsubseteq \ell_3$ and $(m + 1, H''_1) \triangleright \theta'_1$ therefore from Definition 1.36 we get the desired.

Case $i = 2$

Similar reasoning as in the $i = 1$ case

• $(W'.\theta_1, m_{u1}, v'_1) \in \lceil \tau' \rceil_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in \lceil \tau' \rceil_V$:

We get this from (FB-B6), (FB-B9) and Lemma 1.43 we get the desired

15. CG-ref:

$$\frac{\Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new}(e') : \mathbb{C} \ell \perp (\text{ref } \ell' \tau)}$$

To prove: $(W, n, \text{new}(e')(\gamma \downarrow_1), \text{new}(e')(\gamma \downarrow_2)) \in \lceil (\mathbb{C} \ell \perp (\text{ref } \ell' \tau)) \rceil_E^{\mathcal{A}}$

This means from Definition 1.33 we need to prove:

$$\forall i < n. \text{new } (e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{new } (e') \gamma \downarrow_2 \Downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_{\mathcal{V}}^A$$

This means that given some $i < n$ s.t $\text{new } (e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{new } (e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that $v_{f1} = \text{new } (e') \gamma \downarrow_1$, $v_{f2} = \text{new } (e') \gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, \text{new } (e') \gamma \downarrow_1, \text{new } (e') \gamma \downarrow_2) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_{\mathcal{V}}^A$$

Let $v_1 = \text{new } (e') \gamma \downarrow_1$ and $v_2 = \text{new } (e') \gamma \downarrow_2$

From Definition 1.32 we are required to prove

$$\left(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ \left. (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \right. \\ \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{ref } \ell' \tau)) \right) \wedge \\ \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau)]_{\mathcal{V}} \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right)$$

This means we need to prove the following:

$$(a) \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{ref } \ell' \tau)):$$

This means we are given some $k \leq n$, $W_e \sqsupseteq W, H_1, H_2$ s.t $(k, H_1, H_2) \triangleright W_e$

Also we are given some $v'_1, v'_2, j < k$ s.t $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{ref } \ell' \tau))$$

Further from Definition 1.31 it suffices to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge (W', k - j, v'_1, v'_2) \in [(\text{ref } \ell' \tau)]_{\mathcal{V}}^A \quad (\text{FB-R0})$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau]_E^A$$

This means from Definition 1.33 we need to prove:

$$\forall f < k. e' \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, k - f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau]_{\mathcal{V}}^A$$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$ therefore $\exists f < j < k$ s.t $e' \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau]_{\mathcal{V}}^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose W' as W_n where

$$W_n.\theta_1 = W_e.\theta_1 \cup \{a_1 \mapsto (\text{Labeled } \ell' \tau)\}$$

$$W_n.\theta_2 = W_e.\theta_2 \cup \{a_2 \mapsto (\text{Labeled } \ell' \tau)\}$$

$$W_n.\hat{\beta} = W_e.\hat{\beta} \cup \{a_1, a_2\}$$

Now we need to prove:

i. $(k - j, H'_1, H'_2) \triangleright W_n$:

From Definition 1.37 it suffices to prove:

$$\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2) \wedge$$

$$(W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)) \wedge$$

$$\forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a_1) = W_n.\theta_2(a_2)) \wedge$$

$$(W_n, (k - j) - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A \wedge$$

$$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V$$

This means we need to prove

- $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H'_2) \wedge (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$:

We know that $\text{dom}(W_n.\theta_1) = \text{dom}(W_e.\theta_1) \cup \{a_1\}$ and $\text{dom}(W_n.\theta_2) = \text{dom}(W_e.\theta_2) \cup \{a_2\}$

Also $\text{dom}(H'_1) = \text{dom}(H_1) \cup \{a_1\}$ and $\text{dom}(H'_2) = \text{dom}(H_2) \cup \{a_2\}$

Therefore from $(k, H_1, H_2) \triangleright W_e$ and from construction of W_n we get the desired.

- $\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a'_1) = W_n.\theta_2(a'_2)) \wedge (W_n, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \lceil W_n.\theta_1(a'_1) \rceil_V^A$:

$\forall (a'_1, a'_2) \in (W_n.\hat{\beta})$.

A. When $a'_1 = a_1$ and $a'_2 = a_2$:

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$$

Since from (FB-R1) we know that $(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$

And since from cg-ref we know that $H'_1(a_1) = v_{h1}$, $H'_2(a_2) = v'_{h1}$ and $j = f + 1$ thfore from Lemma 1.44 we get

$$(W_n, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

B. When $a'_1 = a_1$ and $a'_2 \neq a_2$: This case cannot arise

C. When $a'_1 \neq a_1$ and $a'_2 = a_2$: This case cannot arise

D. When $a'_1 \neq a_1$ and $a'_2 \neq a_2$:

Since $(k, H_1, H_2) \triangleright W_e$ therefore the desired is obtained directly from Definition 1.37

- $\forall i \in \{1, 2\}. \forall m. \forall a'_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H_i(a'_i)) \in \lfloor W_n.\theta_i(a'_i) \rfloor_V$:

When $i = 1$

Given some m

$\forall a'_1 \in \text{dom}(W_n.\theta_1)$.

– when $a'_1 = a_1$:

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$$

And from (FB-R1) we know that $(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$

Therefore from Lemma 1.42 get the desired

– Otherwise:

Since $(k, H_1, H_2) \triangleright W_e$ therefore the desired is obtained directly from Definition 1.37

When $i = 2$

Similar reasoning as with $i = 1$

ii. $(W', k - j, v'_1, v'_2) \in [(\text{ref } \ell' \tau)]_V^A$:

From cg-ref we know that $v'_1 = a_1$ and $v'_2 = a_2$

From Definition 1.32 it suffices to prove

$(a_1, a_2) \in W_n.\hat{\beta} \wedge W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau)$

This holds from construction of W_n

(b) $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau)]_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell):$

Case $l = 1$

Given some $k, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau)]_V \wedge$

$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sqsubseteq \ell') \wedge$

$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$

Since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 1.51 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$ and $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$(W.\theta_1, k, (\text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_E$

This means from Definition 1.35 we get

$\forall c < k. \text{ref } (e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_V$

This further means that given some $c < k$ s.t $\text{ref } (e')\gamma \downarrow_1 \Downarrow_c v$. From cg-val we know that $c = 0$ and $v = \text{ref } (e')\gamma \downarrow_1$

And we have $(W.\theta_1, k, \text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \perp (\text{ref } \ell' \tau))]_V$

From Definition 1.34 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, \text{ref } (e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies$

$\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{ref } \ell' \tau)]_V \wedge$

$(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sqsubseteq \ell') \wedge$

$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell)$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

16. CG-deref:

$$\frac{\Gamma \vdash e' : \text{ref } \ell \ \tau}{\Gamma \vdash !e' : \mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)}$$

To prove: $(W, n, !e' (\gamma \downarrow_1), !e' (\gamma \downarrow_2)) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} \forall i < n. !e' \ \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge !e' \ \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) &\in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_V^A \end{aligned}$$

This means that given some $i < n$ s.t $!e' \ \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge !e' \ \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that $v_{f1} = !e' \ \gamma \downarrow_1$, $v_{f2} = !e' \ \gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, !e' \ \gamma \downarrow_1, !e' \ \gamma \downarrow_2) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_V^A$$

Let $v_1 = !e' \ \gamma \downarrow_1$ and $v_2 = !e' \ \gamma \downarrow_2$

From Definition 1.32 it suffices to prove

$$\begin{aligned} &(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell \ \tau))) \wedge \\ &\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \ \tau)]_V \wedge \\ &(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \ \tau' \wedge \top \sqsubseteq \ell'') \wedge \\ &(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top)) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad &\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell \ \tau)): \end{aligned}$$

This means we are given is some $k \leq n$, $W_e \sqsupseteq W, H_1, H_2$ s.t $(k, H_1, H_2) \triangleright W_e$

Also given some $v'_1, v'_2, j < k$ s.t $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, (\text{Labeled } \ell \ \tau))$$

This means from Definition 1.31 it suffices to prove $\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge (W', k - j, v'_1, v'_2) \in [(\text{Labeled } \ell \ \tau)]_V^A$ (FB-D0)

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{ref } \ell \ \tau)]_E^A$$

This means from Definition 1.33 we need to prove:

$$\forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1} \implies \\ (W_e, k - f, v_{h_1}, v'_{h_1}) \in [(\text{ref } \ell \tau)]_V^A$$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$ therefore $\exists f < j < k$ s.t $e_l \gamma \downarrow_f \Downarrow_j v_{h_1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h_1}$

This means we have

$$(W_e, k - f, v_{h_1}, v'_{h_1}) \in [(\text{ref } \ell \tau)]_V^A \quad (\text{FB-D1})$$

In order to prove (FB-D0) we choose W' as W_e . Also from cg-deref we know that $H'_1 = H_1$ and $H'_2 = H_2$. Also we know that $v_{h_1} = a_1$ and $v'_{h_1} = a_2$.

- $(k - j, H_1, H_2) \triangleright W_e$:

Since we know that $(k, H_1, H_2) \triangleright W_e$ therefore from Lemma 1.48 we get

$$(k - j, H_1, H_2) \triangleright W_e$$

- $(W', k - j, v'_1, v'_2) \in [(\text{Labeled } \ell \tau)]_V^A$:

Since from (FB-D1) we know that $(W_e, k - f, a_1, a_2) \in [\text{ref } \ell \tau]_V^A$

Therefore from Definition 1.32 we know that $(a_1, a_2) \in W_e.\hat{\beta} \wedge W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = \text{Labeled } \ell \tau$

And since we know that $(k, H_1, H_2) \triangleright W_e$ therefore from Definition we know that

$$(W_e, k, H_1(a_1), H_2(a_2)) \in [\text{Labeled } \ell \tau]_V^A$$

Also from cg-ref we know that $v'_1 = H_1(a_1)$ and $v'_2 = H_2(a_2)$

From Lemma 1.44 we get $(W', k - j, H_1(a_1), H_2(a_2)) \in [(\text{Labeled } \ell \tau)]_V^A$

- (b) $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \top \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \top):$

Case $l = 1$

Given some $k, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau)]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell')$$

Since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 1.51 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$$(W.\theta_1, k, (!e' \gamma \downarrow_1)) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau))]_E$$

This means from Definition 1.35 we get

$$\forall c < k. !e' \gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau))]_V$$

Instantiating c with 0 and from cg-val we know that $v = !e' \gamma \downarrow_1$

And we have $(W.\theta_1, k, !e' \gamma \downarrow_1) \in [(\mathbb{C} \top \perp (\text{Labeled } \ell \tau))]_V$

From Definition 1.34 we have

$$\begin{aligned} & \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ & \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau)]_V \wedge \\ & (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \top \sqsubseteq \ell'') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \top) \end{aligned}$$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

17. CG-assign:

$$\frac{\Gamma \vdash e_l : \text{ref } \ell' \tau \quad \Gamma \vdash e_r : \text{Labeled } \ell' \tau \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_l := e_r : \mathbb{C} \ell \perp \text{unit}}$$

To prove: $(W, n, (e_l := e_r) (\gamma \downarrow_1), (e_l := e_r) (\gamma \downarrow_2)) \in [\mathbb{C} \ell \perp \text{unit}]_E^A$

This means from Definition 1.33 we need to prove:

$$\begin{aligned} & \forall i < n. (e_l := e_r) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \Downarrow v'_{f1} \implies \\ & (W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell \perp \text{unit}]_V^A \end{aligned}$$

This means that given some $i < n$ s.t $(e_l := e_r) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \Downarrow v'_{f1}$

From cg-val we know that $v_{f1} = (e_l := e_r) \gamma \downarrow_1$, $v_{f2} = (e_l := e_r) \gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, (e_l := e_r) \gamma \downarrow_1, (e_l := e_r) \gamma \downarrow_2) \in [\mathbb{C} \ell \ell \text{unit}]_V^A$$

Let $e_1 = (e_l := e_r) \gamma \downarrow_1$ and $e_2 = (e_l := e_r) \gamma \downarrow_2$

From Definition 1.32 it suffices to prove

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, \text{unit}) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\text{unit}]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} & \text{(a) } \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, \text{unit}): \end{aligned}$$

This means we are given some $k \leq n$, $W_e \sqsupseteq W$, H_1, H_2 s.t $(k, H_1, H_2) \triangleright W_e$

And finally given some $v'_1, v'_2, j < k$ s.t $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

And we are required to prove:

$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \perp, v'_1, v'_2, \text{unit})$
(FB-A0)

IH1:

$(W_e, k, e_l(\gamma \downarrow_1), e_l(\gamma \downarrow_2)) \in [\text{ref } \ell' \tau]_E^A$

This means from Definition 1.33 we need to prove:

$\forall f < k. e_l \gamma \downarrow_1 \downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \downarrow v'_{h1} \implies$
 $(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau]_V^A$

Since we know that $(H_1, v_1) \downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \downarrow^f (H'_2, v'_2)$ therefore $\exists f < j < k$ s.t
 $e_l \gamma \downarrow_f \downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \downarrow v'_{h1}$

This means we have

$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau]_V^A$ (FB-A1)

IH2:

$(W_e, k - f, e_r(\gamma \downarrow_1), e_r(\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau]_E^A$

This means from Definition 1.33 we need to prove:

$\forall s < k - f. e_r \gamma \downarrow_1 \downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \downarrow v'_{h2} \implies$
 $(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau]_V^A$

Since we know that $(H_1, v_1) \downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \downarrow^f (H'_2, v'_2)$ therefore $\exists s < j - f < k - f$ s.t $e_r \gamma \downarrow_1 \downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \downarrow v'_{h2}$

This means we have

$(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau]_V^A$ (FB-A2)

In order to prove (FB-A0) we choose W' as W_e . Also from cg-assign we know that $H'_1 = H_1[v_{h1} \mapsto v_{h2}]$ and $H'_2 = H_2[v'_{h1} \mapsto v'_{h2}]$, and $j = f + s + 1$

We need to prove the following:

i. $(k - j, H'_1, H'_2) \triangleright W_e$:

Say $v_{h1} = a_1$ and $v'_{h1} = a_2$

From Definition 1.37 it suffices to prove:

$\text{dom}(W_e.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge$

$(W_e.\hat{\beta}) \subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2)) \wedge$

$\forall (a_1, a_2) \in (W_e.\hat{\beta}). (W_e.\theta_1(a_1) = W_e.\theta_2(a_2)) \wedge$

$(W_e, (k - j) - 1, H'_1(a_1), H'_2(a_2)) \in [W_e.\theta_1(a_1)]_V^A \wedge$

$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_e.\theta_i). (W_e.\theta_i, m, H_i(a_i)) \in [W_e.\theta_i(a_i)]_V$

This means we need to prove

- $\text{dom}(W_e.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge (W_e.\hat{\beta}) \subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2))$:

Since $\text{dom}(H_1) = \text{dom}(H'_1)$ and $\text{dom}(H_2) = \text{dom}(H'_2)$, and also we know that $(k, H_1, H_2) \triangleright W_e$. Therefore we obtain the desired directly from Definition 1.37

- $\forall (a'_1, a'_2) \in (W_e.\hat{\beta}).(W_e.\theta_1(a'_1) = W_e.\theta_2(a'_2) \wedge (W_e, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \lceil W_e.\theta_1(a'_1) \rceil_V^A)$
 $\forall (a'_1, a'_2) \in (W_e.\hat{\beta}).$
 - A. When $a'_1 = a_1$ and $a'_2 = a_2$:
From (FB-A1) and from Definition 1.32 we get
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$
Since from (FB-A2) we know that $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$
And since from cg-assign we know that $H'_1(a_1) = v_{h2}$, $H'_2(a_2) = v'_{h2}$ and $j = f + s + 1$ threfore from Lemma 1.44 we get
 $(W_e, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_e.\theta_1(a_1) \rceil_V^A$
 - B. When $a'_1 = a_1$ and $a'_2 \neq a_2$: This case cannot arise
 - C. When $a'_1 \neq a_1$ and $a'_2 = a_2$: This case cannot arise
 - D. When $a'_1 \neq a_1$ and $a'_2 \neq a_2$:
Since $(k, H_1, H_2) \triangleright W_e$ therefore the desired is obtained directly from Definition 1.37
- $\forall i \in \{1, 2\}.\forall m.\forall a'_i \in \text{dom}(W_e.\theta_i).(W_e.\theta_i, m, H_i(a'_i)) \in \lfloor W_e.\theta_i(a'_i) \rfloor_V$:
When $i = 1$
Given some m
 $\forall a'_1 \in \text{dom}(W_e.\theta_1).$
 - when $a'_1 = a_1$:
From (FB-A1) and from Definition 1.32 we get
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau))$
Since from (FB-A2) we know that $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \lceil \text{Labeled } \ell' \tau \rceil_V^A$
Therefore from Lemma 1.42 get the desired
 - Otherwise:
Since $(k, H_1, H_2) \triangleright W_e$ therefore the desired is obtained directly from Definition 1.37
- When $i = 2$
Similar reasoning as with $i = 1$

ii. $\text{ValEq}(\mathcal{A}, W_e, k - j, \perp, (), (), \text{unit})$:

Holds directly from Definition 1.31 and Definition 1.32

- (b) $\forall l \in \{1, 2\}.\left(\forall k, \theta_e \sqsupseteq \theta, H, j, (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$
 $\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \text{unit} \rfloor_V \wedge$
 $(\forall a.H(a) \neq H'(a) \implies \exists \ell'.\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e).\theta'(a) \searrow \ell)$:

Case $l = 1$

Given some $k, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor (\text{unit}) \rfloor_V \wedge$
 $(\forall a.H(a) \neq H'(a) \implies \exists \ell'.\theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell \sqsubseteq \ell'') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e).\theta'(a) \searrow \ell)$

Since $(W, n, \gamma) \in \lceil \Gamma \rceil_V^A$ therefore from Lemma 1.51 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in \lfloor \Gamma \rfloor_V$ and $(W.\theta_2, m, \gamma \downarrow_2) \in \lfloor \Gamma \rfloor_V$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 1.49 to get

$$(W.\theta_1, k, ((e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \perp (\text{unit}))]_E$$

This means from Definition 1.35 we get

$$\forall c < k.(e_l := e_r)\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \perp (\text{unit}))]_V$$

Instantiating c with 0 and from cg-val we know that $v = (e_l := e_r)\gamma \downarrow_1$

And we have $(W.\theta_1, k, (e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{unit}))]_V$

From Definition 1.34 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau)]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell') \end{aligned}$$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

□

Lemma 1.53. $\forall \mathcal{A}, W, W', \ell, \ell', v_1, v_2, \tau, i, j.$

$$\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau) \wedge j < i \wedge \ell \sqsubseteq \ell' \wedge W \sqsubseteq W' \implies$$

$$\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$$

Proof. Given that $\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau)$. From Definition 1.31 two cases arise

1. $\ell \sqsubseteq \mathcal{A}$:

In this case we know that $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$

2 cases arise

(a) $\ell' \sqsubseteq \mathcal{A}$:

Since $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$ therefore from Lemma 1.44 we know that $(W', j, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$

And thus from Definition 1.31 we know that $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

(b) $\ell' \not\sqsubseteq \mathcal{A}$:

Since $(W, i, v_1, v_2) \in [\tau]_V^{\mathcal{A}}$ therefore from Lemma 1.42 we know that $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 1.43 we know that $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 1.31 we know that $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

2. $\ell \not\sqsubseteq \mathcal{A}$:

Given is $\ell \sqsubseteq \ell' \not\sqsubseteq \mathcal{A}$

In this case we know that $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 1.43 we know that $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 1.31 we know that $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

□

Lemma 1.54 (Subtyping binary). *The following holds:*

\forall, τ, τ' .

$$1. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_{\mathcal{V}}^A \subseteq [(\tau')]_{\mathcal{V}}^A$$

$$2. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_E^A \subseteq [(\tau')]_E^A$$

Proof. Proof of statement (1)

Proof by induction on the $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove: $[((\tau_1 \rightarrow \tau_2))]_{\mathcal{V}}^A \subseteq [((\tau'_1 \rightarrow \tau'_2))]_{\mathcal{V}}^A$

IH1: $[(\tau'_1)]_{\mathcal{V}}^A \subseteq [(\tau_1)]_{\mathcal{V}}^A$ (Statement 1)

$[(\tau_2)]_E^A \subseteq [(\tau'_2)]_E^A$ (Sub-A0 From Statement 2)

It suffices to prove:

$$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2))]_{\mathcal{V}}^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2))]_{\mathcal{V}}^A$$

This means that given: $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2))]_{\mathcal{V}}^A$

And it suffices to prove: $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2))]_{\mathcal{V}}^A$

From Definition 1.32 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_{\mathcal{V}}^A &\implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A &\wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} &\implies (\theta_l, j, e_1[v_1/x]) \in [\tau_2]_E) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_{\mathcal{V}} &\implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 1.32 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A &\implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ [\tau'_2]_E^A) &\wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} &\implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_E) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_{\mathcal{V}} &\implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2]_E) \end{aligned}$$

This means need to prove:

$$(a) \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_E^A) :$$

Given: $W'' \sqsupseteq W, k < n$ and v'_1, v'_2 . We are also given $(W'', k, v'_1, v'_2) \in [\tau'_1]_{\mathcal{V}}^A$

To prove: $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_E^A$

Instantiating the first conjunct of Sub-A1 with W'', k, v'_1 and v'_2 we get

$$((W'', k, v'_1, v'_2) \in [\tau_1]_{\mathcal{V}}^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A) \quad (63)$$

Since $(W'', k, v'_1, v'_2) \in [\tau'_1]_V^A$ therefore from IH1 we know that $(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

Thus from Equation 63 we get $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Finally using (Sub-A0) we get $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2]_E^A$

(b) $\forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_E)$:

Given: $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$. We are also given $(\theta'_l, k, v'_c) \in [\tau'_1]_V$

To prove: $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_E$

Since we are given $(\theta'_l, k, v'_c) \in [\tau'_1]_V$ and since $\tau'_1 <: \tau_1$ therefore from Lemma 1.50 we get

$$(\theta'_l, k, v'_c) \in [\tau_1]_V \tag{64}$$

Instantiating the second conjunct of Sub-A1 with θ'_l, k, v'_1 and v'_2 we get

$$((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2]_E) \tag{65}$$

Therefore from Equation 64 and 65 we get $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$

Since $\tau_2 <: \tau'_2$ therefore from Lemma 1.50 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2]_E$$

(c) $\forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2]_E)$:

Similar reasoning as in the previous case

2. CGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove: $[\tau_1 \times \tau_2]_V^A \subseteq [(\tau'_1 \times \tau'_2)]_V^A$

IH1: $[\tau_1]_V^A \subseteq [\tau'_1]_V^A$ (Statement (1))

IH2: $[\tau_2]_V^A \subseteq [\tau'_2]_V^A$ (Statement (1))

It suffices to prove: $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in [(\tau_1 \times \tau_2)]_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in [(\tau'_1 \times \tau'_2)]_V^A$

This means that given: $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [(\tau_1 \times \tau_2)]_V^A$

Therefore from Definition 1.32 we are given:

$$(W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A \tag{66}$$

And it suffices to prove: $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [(\tau'_1 \times \tau'_2)]_V^A$

Again from Definition 1.32, it suffices to prove:

$$(W, n, v_1, v'_1) \in [\tau'_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau'_2]_V^A$$

Since from Equation 66 we know that $(W, n, v_1, v'_1) \in [\tau_1]_V^A$ therefore from IH1 we have $(W, n, v_1, v'_1) \in [\tau'_1]_V^A$

Similarly since $(W, n, v_2, v'_2) \in [\tau_2]_V^A$ from Equation 66 therefore from IH2 we have $(W, n, v_2, v'_2) \in [\tau'_2]_V^A$

3. CGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove: $\lceil((\tau_1 + \tau_2))\rceil_{\mathcal{V}}^A \subseteq \lceil((\tau'_1 + \tau'_2))\rceil_{\mathcal{V}}^A$

IH1: $\lceil(\tau_1)\rceil_{\mathcal{V}}^A \subseteq \lceil(\tau'_1)\rceil_{\mathcal{V}}^A$ (Statement (1))

IH2: $\lceil(\tau_2)\rceil_{\mathcal{V}}^A \subseteq \lceil(\tau'_2)\rceil_{\mathcal{V}}^A$ (Statement (1))

It suffices to prove: $\forall (W, n, v_{s1}, v_{s2}) \in \lceil((\tau_1 + \tau_2))\rceil_{\mathcal{V}}^A. (W, n, v_{s1}, v_{s2}) \in \lceil((\tau'_1 + \tau'_2))\rceil_{\mathcal{V}}^A$

This means that given: $(W, n, v_{s1}, v_{s2}) \in \lceil((\tau_1 + \tau_2))\rceil_{\mathcal{V}}^A$

And it suffices to prove: $(W, n, v_{s1}, v_{s2}) \in \lceil((\tau'_1 + \tau'_2))\rceil_{\mathcal{V}}^A$

2 cases arise

(a) $v_{s1} = \text{inl } v_{i1}$ and $v_{s1} = \text{inl } v_{i2}$:

From Definition 1.32 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau_1\rceil_{\mathcal{V}}^A \tag{67}$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau'_1\rceil_{\mathcal{V}}^A$$

From Equation 67 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau'_1\rceil_{\mathcal{V}}^A$$

(b) $v_s = \text{inr } v_{i1}$ and $v_{s2} = \text{inr } v_{i2}$:

From Definition 1.32 we are given:

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau_2\rceil_{\mathcal{V}}^A \tag{68}$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau'_2\rceil_{\mathcal{V}}^A$$

From Equation 68 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in \lceil\tau'_2\rceil_{\mathcal{V}}^A$$

4. CGsub-label:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove: $\lceil((\text{Labeled } \ell \tau))\rceil_{\mathcal{V}}^A \subseteq \lceil((\text{Labeled } \ell' \tau'))\rceil_{\mathcal{V}}^A$

IH: $\lceil(\tau)\rceil_{\mathcal{V}}^A \subseteq \lceil(\tau')\rceil_{\mathcal{V}}^A$

It suffices to prove: $\forall (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil((\text{Labeled } \ell \tau))\rceil_{\mathcal{V}}^A. (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil((\text{Labeled } \ell' \tau'))\rceil_{\mathcal{V}}^A$

This means we are given $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in \lceil((\text{Labeled } \ell \tau))\rceil_{\mathcal{V}}^A$

From Definition 1.32 it means we have $\text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau)$ (Sub-L0)

and it suffices to prove $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell' \tau'))]_{\mathcal{V}}^{\mathcal{A}}$

Again from Definition 1.32 it means we need to prove that

$$\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau')$$

Since we have (Sub-L0) and $\ell \sqsubseteq \ell'$ therefore from Lemma 1.53 we have

$$\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau)$$

2 cases arise:

(a) $\ell' \sqsubseteq \mathcal{A}$:

In this case from Definition 1.31 we know that $(W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

From IH we also know that $(W, n, v_1, v_2) \in [\tau']_{\mathcal{V}}^{\mathcal{A}}$

And from Definition 1.32 we get $\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau')$

(b) $\ell' \not\sqsubseteq \mathcal{A}$:

In this case from Definition 1.31 we know that $\forall j. (W.\theta_1, j, v_1) \in [\tau]_{\mathcal{V}}$ and $(W.\theta_2, j, v_2) \in [\tau]_{\mathcal{V}}$

Since $\tau <: \tau'$ therefore from Lemma 1.50 we get $(W.\theta_1, j, v_1) \in [\tau']_{\mathcal{V}}$ and $(W.\theta_2, j, v_2) \in [\tau']_{\mathcal{V}}$

And from Definition 1.32 we get $\text{ValEq}(\mathcal{A}, W, \ell', n, \text{Lb}(v_1), \text{Lb}_\ell(v_2), \tau')$

5. CGsub-CG:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i \quad \mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o}{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove: $[((\mathbb{C} \ell_i \ell_o \tau))]_{\mathcal{V}}^{\mathcal{A}} \subseteq [((\mathbb{C} \ell'_i \ell'_o \tau'))]_{\mathcal{V}}^{\mathcal{A}}$

IH: $[(\tau)]_{\mathcal{V}}^{\mathcal{A}} \subseteq [(\tau')]_{\mathcal{V}}^{\mathcal{A}}$

It suffices to prove: $\forall (W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau))]_{\mathcal{V}}^{\mathcal{A}}. (W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau'))]_{\mathcal{V}}^{\mathcal{A}}$

This means we are given $(W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau))]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 1.32 it means we have

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_{\mathcal{V}} \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i) \right) \quad (\text{Sub-CG0}) \end{aligned}$$

And we need to prove

$$(W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau'))]_{\mathcal{V}}^{\mathcal{A}}$$

Again from Definition 1.32 it means we need to prove

$$\begin{aligned}
& \left(\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \right. \\
& (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau') \right) \wedge \\
& \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau' \rfloor_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i) \right)
\end{aligned}$$

It means we need to prove:

$$\begin{aligned}
& \text{(a) } \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\
& (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
& \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau') :
\end{aligned}$$

This means we are given $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$ s.t
 $(k, H_1, H_2) \triangleright W_e, (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2)$

And we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau')$$

Instantiating the first conjunct of (Sub-CG0) to get

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau) \quad (\text{Sub-CG1})$$

Since from (Sub-CG1) $\text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau)$

Therefore from Lemma 1.53 we get $\text{ValEq}(\mathcal{A}, W', k - j, \ell'_o, v'_1, v'_2, \tau)$

$$\begin{aligned}
& \text{(b) } \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau' \rfloor_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i) \right) :
\end{aligned}$$

Case $l = 1$

Here we are given $k, \theta_e \sqsupseteq \theta, H, j < k$ s.t $(k, H) \triangleright \theta_e \wedge (H, e_1) \Downarrow_j^f (H', v'_1)$

And we need to prove

$$\text{i. } \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in \lfloor \tau' \rfloor_V :$$

Instantiating the second conjunct of (Sub-CG0) with the given k, θ_e, H, j to get
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in \lfloor \tau \rfloor_V$

Since $\tau <: \tau'$ therefore from Lemma 1.50 we get $(\theta', k - j, v'_1) \in \lfloor \tau' \rfloor_V$

$$\text{ii. } (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell') :$$

Instantiating the second conjunct of (Sub-CG0) with the given v, i, k, θ_e, H, j to get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell')$$

Since $\ell'_i \sqsubseteq \ell_i$ therefore we also get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sqsubseteq \ell')$$

$$\text{iii. } (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i) :$$

Instantiating the second conjunct of (Sub-CG0) with the given v, i, k, θ_e, H, j to get

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$$

Since $\ell'_i \sqsubseteq \ell_i$ therefore we also get
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i)$

Case $l = 2$

Symmetric reasoning as in the previous $l = 1$ case

6. CGsub-base:

Trivial

Proof of Statement (2)

It suffice to prove that

$$\forall (W, n, e_1, e_2) \in [(\tau)]_E^A. (W, n, e_1, e_2) \in [(\tau')]_E^A$$

This means given $(W, n, e_1, e_2) \in [(\tau)]_E^A$

From Definition 1.33 it means we have

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau]_V^A \quad (\text{Sub-E0})$$

And it suffices to prove $(W, n, e_1, e_2) \in [(\tau')]_E^A$

Again from Definition 1.33 it means we need to prove

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [(\tau')]_V^A$$

This means that given $i < n$ s.t $e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2$ we need to prove $(W, n - i, v_1, v_2) \in [(\tau')]_V^A$

Instantiating (Sub-E0) with the given i we get $(W, n - i, v_1, v_2) \in [\tau]_V^A$

From Statement (1) we get $(W, n - i, v_1, v_2) \in [(\tau')]_V^A$ □

Theorem 1.55 (NI for CG). *Say $\text{bool} = (\text{unit} + \text{unit})$*

$\forall v_1, v_2, e, n'$.

$\emptyset \vdash v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash v_2 : \text{Labeled } \top \text{ bool} \wedge$

$x : \text{Labeled } \top \text{ bool} \vdash e : \mathbb{C} \perp \perp \text{ bool} \wedge$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_-^f (-, v'_2) \implies v'_1 = v'_2$$

Proof. Given some

$\emptyset \vdash v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash v_2 : \text{Labeled } \top \text{ bool} \wedge$

$x : \text{Labeled } \top \text{ bool} \vdash e : \mathbb{C} \perp \perp \text{ bool} \wedge$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow_-^f (-, v'_2)$$

And we need to prove

$$v'_1 = v'_2$$

From Theorem 1.52 we know that

$$\forall n. (\emptyset, n, v_1, v_2) \in [\text{Labeled } \top \text{ bool}]_E^\perp$$

Similarly from Theorem 1.52 and Definition 1.41 we also get

$$\forall n. (\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \text{ bool}]_E^\perp$$

From Definition 1.33 we get

$$\forall n. \forall i < n. e_1[v_1/x] \Downarrow_i v_{11} \wedge e_2 \Downarrow v_{22} \implies (\emptyset, n - i, v_{11}, v_{22}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\perp$$

Instantiating it with $n' + 1$ and then with 0 , from CG-val we have $v_{11} = e[v_1/x]$ and $v_{22} = e[v_2/x]$

Therefore we have

$$(\emptyset, n' + 1, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathbb{V}}^{\perp}$$

From Definition 1.34 we have

$$\begin{aligned} & \left(\forall k \leq n' + 1, W_e \sqsupseteq \emptyset, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v_1'', v_2'', j. (H_1, e[v_1/x]) \Downarrow_j^f (H_1', v_1'') \wedge (H_2, e[v_2/x]) \Downarrow_j^f (H_2', v_2'') \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', k - j, \perp, v_1', v_2', \mathbf{b}) \Big) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \right. \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v_l') \in [\mathbf{b}]_{\mathbb{V}} \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \perp \sqsubseteq \ell') \wedge \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \perp) \right) \end{aligned}$$

Instantiating the first conjunct with $n' + 1, \emptyset, \emptyset, \emptyset$. And then with v_1', v_2', n' we get $\exists W' \sqsupseteq \emptyset. (1, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', 1, \perp, v_1', v_2', \text{bool})$

From Definition 1.31 and Definition 1.34 we get $v_1' = v_2'$

□

1.3 CG to FG translation

1.3.1 Type directed translation from CG to FG

CG types are translated into FG types by the following definition of $\llbracket \cdot \rrbracket$

$$\begin{aligned} \llbracket \mathbf{b} \rrbracket &= \mathbf{b}^\perp \\ \llbracket \tau_1 \rightarrow \tau_2 \rrbracket &= (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^\perp & \llbracket \text{ref } \ell \tau \rrbracket &= (\text{ref } (\llbracket \tau \rrbracket) + \text{unit})^\ell^\perp \\ \llbracket \tau_1 \times \tau_2 \rrbracket &= (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp & \llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket &= (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket) + \text{unit})^{\ell_2}^\perp \\ \llbracket \tau_1 + \tau_2 \rrbracket &= (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp \\ \llbracket \text{Labeled } \ell \tau \rrbracket &= (\llbracket \tau \rrbracket + \text{unit})^\ell \end{aligned}$$

The translation judgment for expressions is of the form $\boxed{\Gamma \vdash_{pc} e_C : \tau_C \rightsquigarrow e_F}$.

The translation for the pure calculus is omitted as it is straightforward.

$$\begin{array}{c} \frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{Lb}_\ell(e) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_F)} \text{label} \qquad \frac{\Gamma \vdash e : \text{Labeled } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash \text{unlabel}(e) : \mathbb{C} \top \ell \tau \rightsquigarrow \lambda_{\cdot}.e_F} \text{unlabel} \\ \\ \frac{\Gamma \vdash e : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_F}{\Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau) \rightsquigarrow \lambda_{\cdot}.\text{inl}(e_F ())} \text{toLabeled} \\ \\ \frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{ret}(e) : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow \lambda_{\cdot}.\text{inl}(e_F)} \text{ret} \\ \\ \frac{\Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_{F1} \quad \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \rightsquigarrow e_{F2} \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau' \rightsquigarrow \lambda_{\cdot}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}())} \text{bind} \\ \\ \frac{\Gamma \vdash e : \text{Labeled } \ell' \tau \rightsquigarrow e_F \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } e : \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \rightsquigarrow \lambda_{\cdot}.\text{inl}(\text{new } (e_F))} \text{ref} \\ \\ \frac{\Gamma \vdash e : \text{ref } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash !e : \mathbb{C} \top \perp (\text{Labeled } \ell \tau) \rightsquigarrow \lambda_{\cdot}.\text{inl}(e_F)} \text{deref} \\ \\ \frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \rightsquigarrow e_{F1} \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \rightsquigarrow e_{F2} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit} \rightsquigarrow \lambda_{\cdot}.\text{inl}(e_{F1} := e_{F2})} \text{assign} \end{array}$$

Figure 7: Expression translation from CG to FG

1.3.2 Type preservation for CG to FG translation

Theorem 1.56 (Type preservation, CG \rightsquigarrow FG). $\forall \Gamma, e_C, \tau.$

$\Gamma \vdash e_C : \tau$ is a valid typing derivation in CG \implies

$\exists e_F.$

$\Gamma \vdash e_C : \tau \rightsquigarrow e_F \wedge$

$\llbracket \Gamma \rrbracket \vdash_{\top} e_F : \llbracket \tau \rrbracket$ is a valid typing derivation in FG

Proof. Proof by induction on the translation judgment. We show selected cases below.

1. label:

$$\frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{Lb}_\ell(e) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_F)} \text{ label}$$

$$\frac{\frac{\overline{\Gamma \vdash_\top e_F : \llbracket \tau \rrbracket}}{\Gamma \vdash_\top \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^\perp} \text{IH}}{\Gamma \vdash_\top \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^\ell} \text{FG-inl}}{\Gamma \vdash_\top \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^\ell} \text{FG-sub}$$

2. unlabel:

$$\frac{\Gamma \vdash e : \text{Labeled } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash \text{unlabel}(e) : \mathbb{C} \top \ell \tau \rightsquigarrow \lambda_{\cdot} e_F} \text{ unlabel}$$

Main derivation:

$$\frac{\overline{\Gamma, - : \text{unit} \vdash_\top e_F : (\llbracket \tau \rrbracket + \text{unit})^\ell} \text{IH}}{\Gamma, - : \text{unit} \vdash_\top \lambda_{\cdot} e_F : (\text{unit} \xrightarrow{\top} (\llbracket \tau \rrbracket + \text{unit})^\ell)^\perp} \text{FG-lam}$$

3. toLabeled:

$$\frac{\Gamma \vdash e : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_F}{\Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau) \rightsquigarrow \lambda_{\cdot} \text{inl}(e_F ())} \text{ toLabeled}$$

P2:

$$\frac{\overline{\Gamma, - : \text{unit} \vdash_\top e_F : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^\perp} \text{IH, Weakening} \quad \mathcal{L} \vdash \ell_1 \sqsubseteq \top}{\Gamma, - : \text{unit} \vdash_{\ell_1} e_F : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^\perp} \text{FG-sub}$$

P1:

$$\frac{P2 \quad \overline{\Gamma, - : \text{unit} \vdash_{\ell_1} () : \text{unit}} \quad \mathcal{L} \vdash \ell_1 \sqcup \perp \sqsubseteq \ell_1 \quad \mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell_2} \searrow \perp}{\Gamma, - : \text{unit} \vdash_{\ell_1} e_F() : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{FG-app}$$

Main derivation:

$$\frac{P1 \quad \overline{\Gamma, - : \text{unit} \vdash_{\ell_1} \text{inl}(e_F()) : ((\llbracket \tau \rrbracket + \text{unit})^{\ell_2} + \text{unit})^\perp} \text{FG-inl}}{\Gamma \vdash_\top \lambda_{\cdot} \text{inl}(e_F()) : (\text{unit} \xrightarrow{\ell_1} ((\llbracket \tau \rrbracket + \text{unit})^{\ell_2} + \text{unit})^\perp)^\perp} \text{FG-lam}$$

4. ret:

$$\frac{\Gamma \vdash e : \tau \rightsquigarrow e_F}{\Gamma \vdash \text{ret}(e) : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow \lambda_{-}.\text{inl}(e_F)} \text{ ret}$$

$$\frac{\frac{\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_F : \llbracket \tau \rrbracket}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_F : \llbracket \tau \rrbracket} \text{ IH, Weakening} \quad \mathcal{L} \vdash \ell_1 \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{ FG-sub} \quad \mathcal{L} \vdash \perp \sqsubseteq \ell_2}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} \text{inl}(e_F) : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{ FG-sub, FG-inl}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(e_F) : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^{\perp}} \text{ FG-lam}$$

5. bind:

$$\frac{\Gamma \vdash e_1 : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_{F1} \quad \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell_3 \ell_4 \tau' \rightsquigarrow e_{F2} \quad \ell \sqsubseteq \ell_1 \quad \ell \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell'}{\Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell \ell' \tau' \rightsquigarrow \lambda_{-}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}())} \text{ bind}$$

P1.1:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F1} : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^{\perp}} \text{ IH1, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F1} : (\text{unit} \xrightarrow{\ell_1} (\llbracket \tau \rrbracket + \text{unit})^{\ell_2})^{\perp}} \text{ FG-sub}$$

P1:

$$P1.1 \quad \frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} () : \text{unit}} \text{ FG-var} \quad \mathcal{L} \vdash (\ell \sqcup \perp) \sqsubseteq \ell_1 \quad \frac{\mathcal{L} \vdash \perp \sqsubseteq \ell_2}{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell_2} \searrow \perp}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell_1} e_{F1}() : (\llbracket \tau \rrbracket + \text{unit})^{\ell_2}} \text{ FG-app}$$

P2.1:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\top} e_{F2} : (\text{unit} \xrightarrow{\ell_3} (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4})^{\perp}} \text{ IH2, Weakening} \quad \mathcal{L} \vdash \ell \sqcup \ell_2 \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\ell \sqcup \ell_2} e_{F2} : (\text{unit} \xrightarrow{\ell_3} (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4})^{\perp}} \text{ FG-sub}$$

P2:

$$P2.1 \quad \frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\ell \sqcup \ell_2} () : \text{unit}} \text{ FG-var} \quad \mathcal{L} \vdash (\ell \sqcup \ell_2 \sqcup \perp) \sqsubseteq \ell_3 \quad \frac{\mathcal{L} \vdash \perp \sqsubseteq \ell_4}{\mathcal{L} \vdash (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4} \searrow \perp}}{\llbracket \Gamma \rrbracket, - : \text{unit}, x : \llbracket \tau \rrbracket \vdash_{\ell \sqcup \ell_2} e_{F2}() : (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4}} \text{ FG-app}$$

P3:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit}, y : \text{unit} \vdash_{\ell \sqcup \ell_2} () : \text{unit}} \text{ FG-var} \quad \mathcal{L} \vdash \perp \sqsubseteq \ell_4}{\llbracket \Gamma \rrbracket, - : \text{unit}, y : \text{unit} \vdash_{\ell \sqcup \ell_2} \text{inr}() : (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4}} \text{ FG-sub, FG-inr}$$

Main derivation:

$$\begin{array}{c}
P1 \quad P2 \quad P3 \quad \frac{\overline{\mathcal{L} \vdash \ell_2 \sqsubseteq \ell_4} \text{ Given}}{\mathcal{L} \vdash (\llbracket \tau' \rrbracket + \text{unit})^{\ell_4} \searrow \ell_2} \quad \frac{\overline{\ell_4 \sqsubseteq \ell'}}{\text{Given}} \\
\hline
\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : (\llbracket \tau' \rrbracket + \text{unit})^{\ell'}}{\text{FG-case, FG-sub}} \\
\hline
\overline{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : (\text{unit} \xrightarrow{\ell} (\llbracket \tau' \rrbracket + \text{unit})^{\ell'})^{\perp}} \text{FG-lam}
\end{array}$$

6. ref:

$$\frac{\Gamma \vdash e : \text{Labeled } \ell' \tau \rightsquigarrow e_F \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } e : \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \rightsquigarrow \lambda_{-}.\text{inl}(\text{new } (e_F))} \text{ref}$$

P1:

$$\begin{array}{c}
\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_F : (\llbracket \tau \rrbracket + \text{unit})^{\ell'}} \text{IH, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_F : (\llbracket \tau \rrbracket + \text{unit})^{\ell'}}} \text{FG-sub} \\
\frac{\overline{\mathcal{L} \vdash \ell \sqsubseteq \ell'}}{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell'} \searrow \ell} \\
\hline
\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{new } e_F : (\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp}} \text{FG-ref}
\end{array}$$

Main derivation:

$$\begin{array}{c}
P1 \\
\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{inl}(\text{new } e_F) : ((\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\perp}} \text{FG-inl}}{\overline{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(\text{new } e_F) : (\text{unit} \xrightarrow{\ell} ((\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\perp})^{\perp}} \text{FG-lam}}
\end{array}$$

7. deref:

$$\frac{\Gamma \vdash e : \text{ref } \ell \tau \rightsquigarrow e_F}{\Gamma \vdash !e : \mathbb{C} \top \perp (\text{Labeled } \ell \tau) \rightsquigarrow \lambda_{-}.\text{inl}(e_F)} \text{deref}$$

P2:

$$\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_F : (\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell})^{\perp}} \text{IH}$$

P1:

$$\frac{P2 \quad \frac{\overline{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell} < : (\llbracket \tau \rrbracket + \text{unit})^{\ell}} \text{Lemma 1.1} \quad \overline{\mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell} \searrow \perp}}{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} !e_F : (\llbracket \tau \rrbracket + \text{unit})^{\ell}} \text{FG-deref}}$$

Main derivation:

$$\begin{array}{c}
P1 \\
\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} \text{inl}(!e_F) : ((\llbracket \tau \rrbracket + \text{unit})^{\ell} + \text{unit})^{\perp}} \text{FG-inl}}{\overline{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{-}.\text{inl}(!e_F) : (\text{unit} \xrightarrow{\top} ((\llbracket \tau \rrbracket + \text{unit})^{\ell} + \text{unit})^{\perp})^{\perp}} \text{FG-lam}}
\end{array}$$

8. assign:

$$\frac{\Gamma \vdash e_1 : \text{ref } \ell' \tau \rightsquigarrow e_{F1} \quad \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \rightsquigarrow e_{F2} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \perp \text{unit} \rightsquigarrow \lambda_{\cdot} \text{inl}(e_{F1} := e_{F2})} \text{ assign}$$

P3:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F2} : (\llbracket \tau \rrbracket + \text{unit})^{\ell'}} \text{ IH2, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F2} : (\llbracket \tau \rrbracket + \text{unit})^{\ell'}} \text{ FG-sub}$$

P2:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\top} e_{F1} : (\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp}} \text{ IH1, Weakening} \quad \mathcal{L} \vdash \ell \sqsubseteq \top}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F1} : (\text{ref}(\llbracket \tau \rrbracket + \text{unit})^{\ell'})^{\perp}} \text{ FG-sub}$$

P1:

$$\frac{\begin{array}{c} P2 \quad P3 \quad \overline{\mathcal{L} \vdash \ell \sqsubseteq \ell'} \text{ Given} \\ \mathcal{L} \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell'} \searrow (\ell \sqcup \perp) \end{array}}{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} e_{F1} := e_{F2} : \text{unit}} \text{ FG-assign}$$

Main derivation:

$$\frac{\overline{\llbracket \Gamma \rrbracket, - : \text{unit} \vdash_{\ell} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} + \text{unit})^{\perp}} \text{ P1} \quad \text{FG-inl}}{\llbracket \Gamma \rrbracket \vdash_{\top} \lambda_{\cdot} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} \xrightarrow{\ell} (\text{unit} + \text{unit})^{\perp})^{\perp}} \text{ FG-lam}}$$

9. sub:

$$\frac{\overline{\llbracket \Gamma \rrbracket \vdash_{\top} e_F : \llbracket \tau' \rrbracket} \text{ IH} \quad \mathcal{L} \vdash \top \sqsubseteq \top \quad \frac{\mathcal{L} \vdash \tau' <: \tau}{\mathcal{L} \vdash \llbracket \tau' \rrbracket <: \llbracket \tau \rrbracket} \text{ Lemma 1.57}}{\llbracket \Gamma \rrbracket \vdash_{\top} e_F : \llbracket \tau \rrbracket} \text{ FG-sub}$$

□

Lemma 1.57 (Subtyping type preservation: CG to FG). *For any CG types τ and τ' , Σ , and Ψ , if $\mathcal{L} \vdash \tau <: \tau'$, then $\mathcal{L} \vdash \llbracket \tau \rrbracket <: \llbracket \tau' \rrbracket$.*

Proof. Proof by induction on CG's subtyping relation

1. CGsub-base:

$$\overline{\mathcal{L} \vdash \llbracket \tau \rrbracket <: \llbracket \tau \rrbracket} \text{ Lemma 1.1}$$

2. CGsub-arrow:

$$\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1' \rrbracket <: \llbracket \tau_1 \rrbracket} \text{ IH1} \quad \overline{\mathcal{L} \vdash \llbracket \tau_2' \rrbracket <: \llbracket \tau_2 \rrbracket} \text{ IH2} \quad \mathcal{L} \vdash \top \sqsubseteq \top}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^{\perp} <: (\llbracket \tau_1' \rrbracket \xrightarrow{\top} \llbracket \tau_2' \rrbracket)^{\perp}} \text{ FGsub-arrow}$$

$$\frac{}{\mathcal{L} \vdash \llbracket (\tau_1 \xrightarrow{\ell_e} \tau_2) \rrbracket <: \llbracket (\tau_1' \xrightarrow{\ell'_e} \tau_2') \rrbracket} \text{ Definition of } \llbracket \cdot \rrbracket$$

3. CGsub-prod:

$$\frac{\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\mathcal{L} \vdash \llbracket \tau_2 \rrbracket <: \llbracket \tau'_2 \rrbracket}}{\text{IH2}}}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp <: (\llbracket \tau'_1 \rrbracket \times \llbracket \tau'_2 \rrbracket)^\perp} \text{FGsub-arrow} \\ \hline \mathcal{L} \vdash \llbracket (\tau_1 \times \tau_2) \rrbracket <: \llbracket (\tau'_1 \times \tau'_2) \rrbracket \quad \text{Definition of } \llbracket \cdot \rrbracket$$

4. CGsub-sum:

$$\frac{\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\mathcal{L} \vdash \llbracket \tau_2 \rrbracket <: \llbracket \tau'_2 \rrbracket}}{\text{IH2}}}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp <: (\llbracket \tau'_1 \rrbracket + \llbracket \tau'_2 \rrbracket)^\perp} \text{FGsub-arrow} \\ \hline \mathcal{L} \vdash \llbracket (\tau_1 + \tau_2) \rrbracket <: \llbracket (\tau'_1 + \tau'_2) \rrbracket \quad \text{Definition of } \llbracket \cdot \rrbracket$$

5. CGsub-labeled:

$$\frac{\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\mathcal{L} \vdash \text{unit} <: \text{unit}}}{\text{FGsub-unit}}}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})} \text{FGsub-sum} \\ \frac{\overline{\text{Labeled } \ell_1 \tau_1 <: \text{Labeled } \ell'_1 \tau'_1}}{\text{Given}}}{\ell_1 \sqsubseteq \ell'_1} \text{By inversion} \\ \hline \mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_1} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_1} \quad \text{FGsub-arrow} \\ \hline \mathcal{L} \vdash \llbracket \text{Labeled } \ell_1 \tau_1 \rrbracket <: \llbracket \text{Labeled } \ell'_1 \tau'_1 \rrbracket \quad \text{Definition of } \llbracket \cdot \rrbracket$$

6. CGsub-monad:

P3:

$$\frac{\frac{\overline{\mathcal{L} \vdash \llbracket \tau_1 \rrbracket <: \llbracket \tau'_1 \rrbracket}}{\text{IH}} \quad \frac{\overline{\mathcal{L} \vdash \text{unit} <: \text{unit}}}{\text{FGsub-unit}}}{\mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit}) <: (\llbracket \tau'_1 \rrbracket + \text{unit})} \text{FGsub-sum}$$

P2:

$$\frac{\frac{\overline{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau_1 <: \mathbb{C} \ell'_i \ell'_o \tau'_1}}{\text{Given}}}{\mathcal{L} \vdash \ell_o \sqsubseteq \ell'_o} \text{By inversion} \\ \hline \mathcal{L} \vdash (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o} <: (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o} \quad \text{FGsub-label}$$

P1:

$$\frac{\overline{\mathcal{L} \vdash \text{unit} <: \text{unit}} \quad \frac{\overline{\mathcal{L} \vdash \mathbb{C} \ell_i \ell_o \tau_1 <: \mathbb{C} \ell'_i \ell'_o \tau'_1}}{\text{Given}}}{\mathcal{L} \vdash \ell'_i \sqsubseteq \ell_i} \text{FGsub-arrow} \\ \hline \mathcal{L} \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o}) <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})$$

Main derivation:

$$\frac{\frac{\overline{\mathcal{L} \vdash \perp \sqsubseteq \perp}}{P1}}{\mathcal{L} \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})^{\ell_o})^\perp <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})^{\ell'_o})^\perp} \text{FGsub-label} \\ \hline \mathcal{L} \vdash \llbracket \mathbb{C} \ell_i \ell_o \tau_1 \rrbracket <: \llbracket \mathbb{C} \ell'_i \ell'_o \tau'_1 \rrbracket \quad \text{Definition of } \llbracket \cdot \rrbracket$$

□

1.3.3 Logical relation for CG to FG translation

Definition 1.58 (${}^s\theta_2$ extends ${}^s\theta_1$). ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$

$$\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$$

Definition 1.59 ($\hat{\beta}_2$ extends $\hat{\beta}_1$). $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$

$$\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$$

Definition 1.60 (Unary value relation).

$$\begin{aligned} \llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket \wedge {}^s v = {}^t v\} \\ \llbracket \text{unit} \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \text{unit} \rrbracket \wedge {}^t v \in \llbracket \text{unit} \rrbracket\} \\ \llbracket \tau_1 \times \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \mid \\ &\quad ({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\ \llbracket \tau_1 + \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \text{inl } {}^s v, \text{inl } {}^t v) \mid ({}^s\theta, m, {}^s v, {}^t v) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}}\} \cup \\ &\quad \{({}^s\theta, m, \text{inr } {}^s v, \text{inr } {}^t v) \mid ({}^s\theta, m, {}^s v, {}^t v) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\ \llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \lambda x. e_s, \lambda x. e_t) \mid \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v, {}^t v, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^s v, {}^t v) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'} \\ &\quad \implies ({}^s\theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}'}\} \\ \llbracket \text{ref } \ell \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s a, {}^t a) \mid {}^s\theta({}^s a) = \text{Labeled } \ell \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}\} \\ \llbracket \text{Labeled } \ell \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid \\ &\quad \exists {}^s v', {}^t v'. {}^s v = \text{Lb}({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, m, {}^s v', {}^t v') \in \llbracket \tau \rrbracket_V^{\hat{\beta}}\} \\ \llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^s v, {}^t v) \mid \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ &\quad (k, H_s, H_t) \hat{\triangleright}^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H_s, {}^s v') \wedge i < k \implies \\ &\quad \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge \\ &\quad \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in \llbracket \tau \rrbracket_V^{\hat{\beta}''}\} \end{aligned}$$

Definition 1.61 (Unary expression relation).

$$\begin{aligned} \llbracket \tau \rrbracket_E^{\hat{\beta}} &\triangleq \{({}^s\theta, n, e_s, e_t) \mid \\ &\quad \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. e_s \Downarrow_i {}^s v \implies \\ &\quad \exists H'_t, {}^t v. (H_t, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in \llbracket \tau \rrbracket_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta\} \end{aligned}$$

Definition 1.62 (Unary heap well formedness).

$$\begin{aligned} (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta &\triangleq \text{dom}({}^s\theta) \subseteq \text{dom}(H_s) \wedge \\ &\quad \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \\ &\quad \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n - 1, H_s(a_1), H_t(a_2)) \in \llbracket {}^s\theta(a) \rrbracket_V^{\hat{\beta}} \end{aligned}$$

Definition 1.63 (Value substitution). $\delta^s : \text{Var} \mapsto \text{Val}$, $\delta^t : \text{Var} \mapsto \text{Val}$

Definition 1.64 (Unary interpretation of Γ).

$$\begin{aligned} \llbracket \Gamma \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\ &\quad \forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in \llbracket \Gamma(x) \rrbracket_V^{\hat{\beta}}\} \end{aligned}$$

1.3.4 Soundness proof for CG to FG translation

Lemma 1.65 (Monotonicity). $\forall^s \theta, {}^s \theta', n, {}^s v, {}^t v, n', \beta, \beta'$.

$$({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s \theta', n', {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$$

Proof. Proof by induction on τ

1. Case **b**:

Given:

$$({}^s \theta, n, {}^s v, {}^t v) \in [\mathbf{b}]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s \theta', n', {}^s v, {}^t v) \in [\mathbf{b}]_V^{\hat{\beta}'}$$

Since $({}^s \theta, n, {}^s v, {}^t v) \in [\mathbf{b}]_V^{\hat{\beta}}$ therefore from Definition 1.60 we know that ${}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket$

Therefore from Definition 1.60 ${}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket$ we get the desired

2. Case **unit**:

Given:

$$({}^s \theta, n, {}^s v, {}^t v) \in [\mathbf{unit}]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s \theta', n', {}^s v, {}^t v) \in [\mathbf{unit}]_V^{\hat{\beta}'}$$

Since $({}^s \theta, n, {}^s v, {}^t v) \in [\mathbf{unit}]_V^{\hat{\beta}}$ therefore from Definition 1.60 we know that ${}^s v \in \llbracket \mathbf{unit} \rrbracket \wedge {}^t v \in \llbracket \mathbf{unit} \rrbracket$

Therefore from Definition 1.60 ${}^s v \in \llbracket \mathbf{unit} \rrbracket \wedge {}^t v \in \llbracket \mathbf{unit} \rrbracket$ we get the desired

3. Case $\tau_1 \times \tau_2$:

Given:

$$({}^s \theta, n, {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s \theta', n', {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.60 we know that ${}^s v = ({}^s v_1, {}^s v_2)$ and ${}^t v = ({}^t v_1, {}^t v_2)$.

We also know that $({}^s \theta, n, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$ and $({}^s \theta, n, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}$

$$\underline{\text{IH1:}} ({}^s \theta', n', {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'}$$

$$\underline{\text{IH2:}} ({}^s \theta', n', {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}'}$$

Therefore from Definition 1.60, IH1 and IH2 we get

$$({}^s \theta', n', {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case $\tau_1 + \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.60 two cases arise

(a) ${}^sv = \text{inl}({}^sv')$ and ${}^tv = \text{inl}({}^tv')$:

$$\text{IH: } ({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$$

Therefore from Definition 1.60 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b) ${}^sv = \text{inr}({}^sv')$ and ${}^tv = \text{inr}({}^tv')$:

Symmetric reasoning as in the previous case

5. Case $\tau_1 \rightarrow \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.60 we know that

$$\forall {}^s\theta'' \sqsubseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta'', j, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}} \implies ({}^s\theta'', j, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau_2]_E^{\hat{\beta}'} \quad (\text{A0})$$

Similarly from Definition 1.60 we are required to prove

$$\forall {}^s\theta'_1 \sqsubseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}'} \implies ({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

This means we are given some ${}^s\theta'_1 \sqsubseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''$ s.t $({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}'}$ and we are required to prove

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}'}$$

Instantiating (A0) with ${}^s\theta'_1, {}^sv_2, {}^tv_2, j, \hat{\beta}''$ since ${}^s\theta'_1 \sqsubseteq {}^s\theta' \sqsubseteq {}^s\theta, j < n' < n$ and $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$ therefore we get

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

6. Case $\text{ref } \ell \tau$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\text{ref } \ell \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \ell \ \tau]_V^{\hat{\beta}'}$$

From Definition 1.60 we know that ${}^s v = {}^s a$ and ${}^t v = {}^t a$. We also know that

$${}^s\theta({}^s a) = \text{Labeled } \ell \ \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}$$

From Definition 1.60, Definition 1.58 and Definition 1.59 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \ell \ \tau]_V^{\hat{\beta}'}$$

7. Case Labeled $\ell \ \tau$:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}'}$$

From Definition 1.60 it means

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, n, {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}}$$

$$\text{IH: } ({}^s\theta', n', {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}'}$$

Similarly from Definition 1.60 we need to prove that

$$\exists {}^s v'', {}^t v''. {}^s v = \text{Lb}_\ell({}^s v'') \wedge {}^t v = \text{inl } {}^t v'' \wedge ({}^s\theta', n', {}^s v'', {}^t v'') \in [\tau]_V^{\hat{\beta}'}$$

We choose ${}^s v''$ as ${}^s v'$ and ${}^t v''$ as ${}^t v'$ and since from IH we know that $({}^s\theta', n', {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}'}$

Therefore from Definition 1.60 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}'}$$

8. Case $\mathbb{C} \ \ell_1 \ \ell_2 \ \tau$:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\mathbb{C} \ \ell_1 \ \ell_2 \ \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\mathbb{C} \ \ell_1 \ \ell_2 \ \tau]_V^{\hat{\beta}'}$$

This means from Definition 1.60 we know that

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}_1.$$

$$(k, H_s, H_t) \stackrel{\hat{\beta}_1}{\triangleright} ({}^s\theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies$$

$$\exists {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}_1 \sqsubseteq \hat{\beta}_2. (k - i, H'_s, H'_t) \stackrel{\hat{\beta}_2}{\triangleright} {}^s\theta' \wedge$$

$$\exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', {}^t\theta', k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}_2} \wedge$$

$$(\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s\theta_e(a) = \text{Labeled } \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e). {}^s\theta'(a) \searrow \ell_1) \quad (\text{CG0})$$

Similarly from Definition 1.60 we need to prove

$$\begin{aligned}
& \forall^s \theta'_e \sqsupseteq {}^s \theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1. \\
& (k', H'_s, H'_t) \triangleright^{\hat{\beta}'_1} ({}^s \theta'_e) \wedge (H'_s, {}^s v) \Downarrow_i^f (H''_s, {}^s v'') \wedge (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge i' < k' \implies \\
& \exists {}^t v''. (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} {}^s \theta'' \wedge \\
& \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k' - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\
& (\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}({}^s \theta') / \text{dom}({}^s \theta_e). {}^s \theta'(a) \searrow \ell_1)
\end{aligned}$$

This means we are given some ${}^s \theta'_e \sqsupseteq {}^s \theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$ s.t. $(k', H'_s, H'_t) \triangleright ({}^s \theta'_e) \wedge (H'_s, {}^s v) \Downarrow_i^f (H''_s, {}^s v'') \wedge i' < k'$

And we need to prove

$$\begin{aligned}
& \exists {}^t v''. (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} {}^s \theta'' \wedge \\
& \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta'', k' - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\
& (\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}({}^s \theta') / \text{dom}({}^s \theta_e). {}^s \theta'(a) \searrow \ell_1)
\end{aligned}$$

Instantiating (CG0) with ${}^s \theta'_e \sqsupseteq {}^s \theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$ we get the desired

□

Lemma 1.66 (Unary monotonicity for Γ). $\forall \theta, \theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$.

$$(\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies (\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

Proof. Given: $(\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove: $(\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$

From Definition 1.64 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s \theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$$

And again from Definition 1.64 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s \theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t)$:

Given

- $\forall x \in \text{dom}(\Gamma). ({}^s \theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$:

Since we know that $\forall x \in \text{dom}(\Gamma). ({}^s \theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$ (given)

Therefore from Lemma 1.65 we get

$$\forall x \in \text{dom}(\Gamma). ({}^s \theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

□

Lemma 1.67 (Unary monotonicity for H). $\forall {}^s \theta, H_s, H_t, n, n', \hat{\beta}, \hat{\beta}'$.

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s \theta$$

Proof. Given: $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta \wedge n' < n$

To prove: $(n', H_s, H_t) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta$

From Definition 1.62 it is given that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

And again from Definition 1.62 we are required to prove that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

- $\text{dom}({}^s\theta) \subseteq \text{dom}(H_S)$:

Given

- $\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t))$:

Given

- $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$:

Since we know that $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$ (given)

Therefore from Lemma 1.65 we get

$$\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

□

Theorem 1.68 (Fundamental theorem). $\forall \Gamma, \tau, e, \delta^s, \delta^t, {}^s\theta, n$.

$$\Gamma \vdash e_s : \tau \rightsquigarrow e_t \wedge$$

$$({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$$

\implies

$$({}^s\theta, n, e_s, \delta^s, e_t, \delta^t) \in [\tau]_E^{\hat{\beta}}$$

Proof. Proof by induction on the \rightsquigarrow relation

1. CF-var:

$$\frac{}{\Gamma, x : \tau \vdash x : \tau \rightsquigarrow x} \text{CF-var}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau\}]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, x, \delta^s, x, \delta^t) \in [\tau]_E^{\hat{\beta}}$

From Definition 1.61 it suffices to prove that

$$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. x \delta^s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta$$

This means given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $x \delta^s \Downarrow_i {}^s v$

From cg-val we know that $i = 0$, ${}^s v = x \delta^s$.

And we are required to prove

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-V0})$$

From fg-val we know that ${}^t v = x \delta^t$ and $H'_t = H_t$. So we are left with proving

$$({}^s \theta, n, x \delta^s, x \delta^t) \in [\tau]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we are given $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau\}]_{\hat{V}}^{\hat{\beta}}$, therefore from Definition 1.64 we get

$$({}^s \theta, n, x \delta^s, x \delta^t) \in [\tau]_{\hat{V}}^{\hat{\beta}}. \text{ And we have } (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \text{ in the context. So we are done.}$$

2. CF-lam:

$$\frac{\Gamma, x : \tau_1 \vdash e_s : \tau_2 \rightsquigarrow e_t}{\Gamma \vdash \lambda x. e_s : \tau_1 \rightarrow \tau_2 \rightsquigarrow \lambda x. e_t} \text{ lam}$$

Also given is: $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{V}}^{\hat{\beta}}$

To prove: $({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [\tau]_{\hat{E}}^{\hat{\beta}}$

From Definition 1.61 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (\lambda x. e_s) \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, (\lambda x. e_t) \delta^t) \Downarrow (H'_t, {}^t v) ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \rightarrow \tau_2)]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ and given some $i < n, {}^s v$ s.t $(\lambda x. e_s) \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that ${}^s v = (\lambda x. e_s) \delta^s$, ${}^t v = (\lambda x. e_t) \delta^t$, $H'_t = H_t$ and $i = 0$

It suffices to prove that

$$({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2)]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

We know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ from the context. So, we are only left to prove

$$({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2)]_{\hat{V}}^{\hat{\beta}}$$

From Definition 1.60 it suffices to prove

$$\begin{aligned} \forall {}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1]_{\hat{V}}^{\hat{\beta}'} \\ \implies ({}^s \theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2]_{\hat{E}}^{\hat{\beta}'} \end{aligned}$$

This means that we are given ${}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t $({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1]_{\hat{V}}^{\hat{\beta}'}$

And we need to prove

$$({}^s \theta', j, e_s[{}^s v/x] \delta^s, e_t[{}^t v/x] \delta^t) \in [\tau_2]_{\hat{E}}^{\hat{\beta}'} \quad (\text{F-L0})$$

Since $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{V}}^{\hat{\beta}}$ therefore from Lemma 1.66 we also have

$$({}^s\theta', j, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

IH:

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_1\}, e_t \cup \{x \mapsto {}^t v_1\}) \in [\tau_2]_E^{\hat{\beta}'}$$
 s.t

$$({}^s\theta', j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'}$$

We get (F-L0) directly from IH

3. CF-app:

$$\frac{\Gamma \vdash e_{s1} : (\tau_1 \rightarrow \tau_2) \rightsquigarrow e_{t1} \quad \Gamma \vdash e_{s2} : \tau_1 \rightsquigarrow e_{t2}}{\Gamma \vdash e_{s1} e_{s2} : \tau_2 \rightsquigarrow e_{t1} e_{t2}} \text{ app}$$

$$\text{Also given is: } ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$$

$$\text{To prove: } ({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, (e_{t1} e_{t2}) \delta^t) \in [\tau_2]_E^{\hat{\beta}}$$

This means from Definition 1.61 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This further means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n, {}^s v$ s.t $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A0})$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \rightarrow \tau_2)]_E^{\hat{\beta}}$$

This means from Definition 1.61 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n$ s.t $e_{s1} \delta^s \Downarrow_j {}^s v_1$.

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A1})$$

IH2:

$$({}^s\theta, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1]_E^{\hat{\beta}}$$

This means from Definition 1.61 it suffices to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall k < n - j, {}^s v_2. e_{s2} \Downarrow_i {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j - k, H_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta' \end{aligned}$$

Instantiating with H_s, H'_{t1} and since we know that $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$ therefore $\exists k < i - j < n - j$ s.t $e_{s2} \delta^s \Downarrow_k {}^s v_2$.

And we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta \quad (\text{F-A2})$$

Since from (F-A1) we know that $({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2)]_V^{\hat{\beta}}$ where ${}^s v_1 = \lambda x. e'_s$ and ${}^t v_1 = \lambda x. e'_t$

From Definition 1.60 we have

$$\begin{aligned} & \forall {}^s\theta'_3 \supseteq {}^s\theta, {}^s v, {}^t v, l < n - j, \hat{\beta}_3 \supseteq \hat{\beta}. ({}^s\theta'_3, l, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}_3} \\ & \implies ({}^s\theta'_3, l, e'_s[{}^s v/x], e'_t[{}^t v/x]) \in [\tau_2]_E^{\hat{\beta}_3} \end{aligned}$$

Instantiating with ${}^s\theta, {}^s v_2, {}^t v_2, n - j - k, \hat{\beta}$ we get

$$({}^s\theta, n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s4}, H_{t4}. (n - j - k, H_{s4}, H_{t4}) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall k' < n - j - k, {}^s v_4. e'_s[{}^s v_2/x] \Downarrow_{k'} {}^s v_4 \implies \\ & \exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s\theta, n - j - k - k', {}^s v_4, {}^t v_4) \in [\tau_2]_V^{\hat{\beta}} \wedge \\ & (n - j - k - k', H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with H_s, H'_{t2} , from (F-A2) we know that $(n - j - k, H_s, H'_{t2}) \hat{\triangleright}^{\hat{\beta}} s\theta$. Instantiating ${}^s v_4$ with ${}^s v$ and since we know that $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$ therefore $\exists k' < i - j - k < n - j - k$ s.t $e'_s[{}^s v_2/x] \delta^s \Downarrow_{k'} {}^s v$. therefore we have

$$\exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s\theta, n - j - k - k', {}^s v, {}^t v_4) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - j - k - k', H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}} s\theta \quad (\text{F-A3})$$

Since from cg-app we know that $i = j + k + k'$ and $H'_t = H'_{t4}$, ${}^t v = {}^t v_4$ therefore we get (F-A0) from (F-A3) and Lemma 1.65 and Lemma 1.67

4. CF-prod:

$$\frac{\Gamma \vdash e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Gamma \vdash e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Gamma \vdash (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2) \rightsquigarrow (e_{t1}, e_{t2})} \text{ prod}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, (e_{s1}, e_{s2}) \delta^s, (e_{t1}, e_{t2}) \delta^t) \in [(\tau_1 \times \tau_2)]_E^{\hat{\beta}}$

From Definition 1.61 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t, \hat{\beta}. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

This means that we are given some $H_s, H_t, \hat{\beta}$ s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ and given some $i < n$ s.t $(e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \quad (\text{F-P0})$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$ therefore $\exists j < i < n$ s.t $e_{s1} \delta^s \Downarrow_j {}^s v_1$.

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} s\theta \quad (\text{F-P1})$$

IH2:

$$({}^s\theta, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_2]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} s\theta \wedge \forall k < n - j. e_{s2} \delta^s \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

Instantiating with $H_s, H'_{t1}, \hat{\beta}'_1$ and since we know that $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$ therefore $\exists k < i - j < n - j$ s.t $e_{s2} \delta^s \Downarrow_k {}^s v_2$.

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} s\theta \quad (\text{F-P2})$$

From cg-prod we know that $i = j + k + 1$, $H'_t = H'_{t2}$ and ${}^t v = ({}^t v_1, {}^t v_2)$ therefore from Definition 1.60 and Lemma 1.65 we get $({}^s\theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}}$

And since we have $(n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} s\theta$ therefore from Lemma 1.67 we also get

$$(n - i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta$$

5. CF-fst:

$$\frac{\Gamma \vdash e_s : \tau_1 \times \tau_2 \rightsquigarrow e_t}{\Gamma \vdash \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{fst}(e_t)} \text{fst}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{fst}(e_t) \delta^t) \in [\tau_1]_E^{\hat{\beta}}$ (F-F0)

This means from Definition 1.61 we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{fst}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^s) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some $H_s, H_t, \hat{\beta}$ s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n, {}^s v$ s.t $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^s) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-F0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2)]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j ({}^s v_1, -) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v_1, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge \\ (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H_t and ${}^s v_1$ with ${}^s v$ since we know that $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n$ s.t $e_s \delta^s \Downarrow_j ({}^s v, -)$.

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}} \wedge \\ (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-F1}) \end{aligned}$$

From cg-fst we know that $i = j + 1$, $H'_t = H'_{t1}$ and ${}^t v = {}^t v_1$. Since we know $({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2)]_V^{\hat{\beta}}$ therefore from Definition 1.60 and Lemma 1.65 we get

$$({}^s\theta, n - i, {}^s v, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$$

And since from (F-F1) we have $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$ therefore from Lemma 1.67 we get

$$(n - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

6. CF-snd:

Symmetric reasoning as in the CF-fst case

7. CF-inl:

$$\frac{\Gamma \vdash e_s : \tau_1 \rightsquigarrow e_t}{\Gamma \vdash \text{inl}(e_s) : (\tau_1 + \tau_2) \rightsquigarrow \text{inl}(e_t)} \text{CF-inl}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{inl}(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\tau_1 + \tau_2)]_E^{\hat{\beta}}$

From Definition 1.61 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some $H_s, H_t, \hat{\beta}$ s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n, {}^s v$ s.t $\text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-IL0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge \\ ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $\text{inl}(e_s) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n$ s.t $e_s \delta^s \Downarrow_j {}^s v$.

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-IL1})$$

From cg-inl we know that $i = j + 1$ and $H'_t = H'_{t1}, {}^t v = {}^t v_1$. Since we know $({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$ therefore from Definition 1.60 and Lemma 1.65 we get

$$({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v_1)) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}}$$

And since from (F-IL1) we have $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$ therefore from Lemma 1.67 we get

$$(n - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

8. CF-inr:

Symmetric reasoning as in the CF-inl case

9. CF-case:

$$\frac{\Gamma \vdash e_s : \tau_1 + \tau_2 \rightsquigarrow e_t \quad \Gamma, x : \tau_1 \vdash e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Gamma, y : \tau_2 \vdash e_{s2} : \tau \rightsquigarrow e_{t2}}{\Gamma \vdash \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{case}(e_t, x.e_{t1}, y.e_{t2})} \text{CF-case}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \in [\tau]_E^{\hat{\beta}}$

This means from Definition 1.61 we need to prove

$$\begin{aligned} &\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ &\exists H'_t, {}^t v. (H_t, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some H_s, H_t s.t. $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n$ s.t. $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \text{ (F-C0)}$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2)]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} &\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ &\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n$ s.t. $e_s \delta^s \Downarrow_j {}^s v_1$.

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \text{ (F-C1)}$$

Two cases arise:

(a) ${}^s v_1 = \text{inl}({}^s v'_1)$ and ${}^t v_1 = \text{inl}({}^t v'_1)$:

IH2:

$$({}^s\theta, n - j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \in [\tau]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} &\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k < n - j, {}^s v_2. e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v_2 \implies \\ &\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}} \wedge (n - \\ &j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H'_{t1} and since we know that $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i^s v$ therefore $\exists k < i - j < n - j$ s.t $e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k^s v$.

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v, {}^t v_2) \in [\tau]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$$

From cg-case1 we know that $i = j + k + 1$ and $H'_t = H'_{t2}, {}^t v = {}^t v_2$. Since we know $({}^s \theta, n - j - k, {}^s v, {}^t v_2) \in [\tau]_V^{\hat{\beta}}$ therefore from Definition 1.60 and Lemma 1.65 we get $({}^s \theta, n - i, {}^s v, {}^t v_2) \in [\tau]_V^{\hat{\beta}}$

And since from (F-C2) we have $(n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$ therefore from Lemma 1.67 we get $(n - i, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$

(b) ${}^s v_1 = \text{inr}({}^s v'_1)$ and ${}^t v_1 = \text{inr}({}^t v'_1)$:

Symmetric reasoning as in the previous case

10. CF-ret:

$$\frac{\Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Gamma \vdash \text{ret}(e_s) : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow \lambda_.\text{inl}(e_t)} \text{ret}$$

Also given is: $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda_.\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_1 \ell_2 \tau]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{ret}(e_s) \Downarrow_i^s v \implies \exists H'_t, {}^t v. (H_t, \lambda_.\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

This means that given some $H_s, H_t, \hat{\beta}$ s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ and given some $i < n$ s.t $\text{ret}(e_s) \delta^s \Downarrow_i^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_.\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

From CG-ret and FG-lam we know that $i = 0, {}^s v = \text{ret}(e_s) \delta^s, {}^t v = \lambda_.\text{inl}(e_t) \delta^t$ and $H'_t = H_t$.

So we need to prove

$$({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda_.\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ from the context so we are left with proving

$$({}^s \theta, n, \text{ret}(e_s) \delta^s, \lambda_.\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}' \\ & (k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \exists H'_t, {}^t v'. (H_t, (\lambda_{-} \text{inl}(e_t) ()) \delta^t) \Downarrow \\ & (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau]_{V'}^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k. \text{ Also from cg-ret we know that } H'_s = H_s$$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v'. (H_t, (\lambda_{-} \text{inl}(e_t) ()) \delta^t) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau]_{V'}^{\hat{\beta}''} \quad (\text{F-R0}) \end{aligned}$$

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [\tau]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k. e_s \delta^s \Downarrow_f {}^s v \implies \\ & \exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s\theta_e, k - f, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating H_{s1} with H_s and H_{t1} with H_t . And since we know that $(H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists f < i < k \leq n$ s.t $e_s \delta^s \Downarrow_f {}^s v_h$. Therefore we have

$$\exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s\theta_e, k - f, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}'} \wedge (k - f, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-R1})$$

In order to prove (F-R0) we choose H'_t as H'_{t1} , ${}^t v'$ as $\text{inl}({}^t v)$, ${}^s\theta'$ as ${}^s\theta_e$, $\hat{\beta}''$ as $\hat{\beta}'$. Since from cg-ret we know that $i = f + 1$ therefore from (F-R1) and Lemma 1.67 we know that

$$(k - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$$

Next we choose ${}^t v''$ as ${}^t v$ (from F-R1) and from Lemma 1.65 we get $({}^s\theta_e, k - i, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}'}$ (we know from cg-ret that ${}^s v' = {}^s v$)

11. CF-bind:

$$\frac{\Gamma \vdash e_{s1} : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_{t1} \quad \Gamma, x : \tau \vdash e_{s2} : \mathbb{C} \ell_3 \ell_4 \tau' \rightsquigarrow e_{t2} \quad \ell_i \sqsubseteq \ell_1 \quad \ell_i \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_3 \quad \ell_2 \sqsubseteq \ell_4 \quad \ell_4 \sqsubseteq \ell_o}{\Gamma \vdash \text{bind}(e_{s1}, x.e_{s2}) : \mathbb{C} \ell_i \ell_o \tau' \rightsquigarrow \lambda_{-} \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}())} \text{ bind}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_{-} \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau')]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau')]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright}^s \theta$ and given some $i < n, {}^s v$ s.t $\text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \lambda_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau')]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

From cg-val and fg-val we know that $i = 0, {}^s v = \text{bind}(e_{s1}, x.e_{s2}) \delta^s,$
 ${}^t v = \lambda_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau')]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^s \theta$$

Since we already know $(n, H_s, H_t) \hat{\triangleright}^s \theta$ from the context so we are left with proving

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau')]_{V}^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \\ & \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}'''. (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau']_{V}^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_ \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}'''. (k - \\ & i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau']_{V}^{\hat{\beta}''} \quad (\text{F-B0}) \end{aligned}$$

IH1:

$$({}^s \theta, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_E^{\hat{\beta}}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_j {}^s v_{h1} \implies \\ & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V^{\hat{\beta}} \wedge (k - j, H_{s2}, H'_{t2}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists j < i < k \leq n$ s.t $e_{s1} \delta^s \Downarrow_j {}^s v_{h1}$.

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k-j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_1 \ell_2 \tau)]_V^{\hat{\beta}} \wedge (k-j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-B1.1})$$

From Definition 1.60 we know have

$$\begin{aligned} \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k-j, \hat{\beta} \sqsubseteq \hat{\beta}' \\ (m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies \\ \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m-b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m-b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating ${}^s \theta_e$ with ${}^s \theta$, H_{s3} with H_{s1} , H_{t3} with H'_{t2} , m with $k-j$ and $\hat{\beta}'$ with $\hat{\beta}$. Since we know that $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists b < i-j < k-j$ s.t $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$.

Therefore we have

$$\begin{aligned} \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k-j-b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k-j-b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau]_V^{\hat{\beta}''} \quad (\text{F-B1}) \end{aligned}$$

IH2:

$$({}^s \theta'', k-j-b, e_{s2} \delta^s \cup \{x \mapsto {}^s v'_{h1}\}, e_{t2} \delta^t \cup \{x \mapsto {}^t v''_{h1}\}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_E^{\hat{\beta}''}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} \forall H_{s4}, H_{t4}. (k, H_{s4}, H_{t4}) \triangleright^{\hat{\beta}''} {}^s \theta \wedge \forall c < (k-j-b), {}^s v_{h2}. e_{s2} \delta^s \Downarrow_j {}^s v_{h2} \implies \\ \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k-j-b-c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V^{\hat{\beta}''} \wedge (k-j-b-c, H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}''} {}^s \theta'' \end{aligned}$$

Instantiating H_{s4} with H'_{s3} and H_{t4} with H'_{t3} . And since we know that $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists c < i-j-b < k-j-b$ s.t $e_{s2} \delta^s \Downarrow_c {}^s v_{h2}$.

Therefore we have

$$\begin{aligned} \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k-j-b-c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell_3 \ell_4 \tau')]_V^{\hat{\beta}''} \wedge (k-j-b-c, H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}''} {}^s \theta'' \quad (\text{F-B2.1}) \end{aligned}$$

From Definition 1.60 we know have

$$\begin{aligned} \forall {}^s \theta_e \sqsupseteq {}^s \theta'', H_{s5}, H_{t5}, d, {}^s v'_{h2}, {}^t v'_{h2}, m \leq k-j-b-c, \hat{\beta}'' \sqsubseteq \hat{\beta}''_1 \\ (m, H_{s5}, H_{t5}) \triangleright^{\hat{\beta}''_1} ({}^s \theta_e) \wedge (H_{s5}, {}^s v_{h2}) \Downarrow_d^f (H'_{s5}, {}^s v'_{h2}) \wedge d < m \implies \\ \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s \theta''' \sqsupseteq {}^s \theta_e, \hat{\beta}''_1 \sqsubseteq \hat{\beta}''_2 . (m-d, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}''_2} {}^s \theta''' \wedge \\ \exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', m-d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau']_V^{\hat{\beta}''_2} \end{aligned}$$

Instantiating ${}^s\theta_e$ with ${}^s\theta''$, H_{s5} with H'_{s3} , H_{t5} with H'_{t3} , m with $k - j - b - c$ and $\hat{\beta}'_1$ with $\hat{\beta}''$. Since we know that $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists d < i - j - b - c < k - j - b - c$ s.t. $(H'_{s3}, {}^s v_{h2}) \delta^s \Downarrow_d (H'_{s5}, {}^s v'_{h2})$.

Therefore we have

$$\begin{aligned} & \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s\theta''' \sqsupseteq {}^s\theta_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_2. (k - j - b - c - d, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}''_2} {}^s\theta''' \wedge \\ & \exists {}^t v'' . {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s\theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau']_{V}^{\hat{\beta}''_2} \quad (\text{F-B2}) \end{aligned}$$

In order to prove (F-B0) we choose H'_{t1} as H'_{t5} and ${}^t v'$ as ${}^t v'_{h2}$. Next we choose ${}^s\theta'$ as ${}^s\theta'''$ and $\hat{\beta}''$ as $\hat{\beta}''_2$ (both chosen from (F-B2)). Also from cg-bind we know that in (F-B0) H'_{s1} will be H'_{s5} .

Since $(k - j - b - c - d, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}''_2} {}^s\theta'''$ therefore Lemma 1.65 we get $(k - i, H'_{s5}, H'_{t5}) \triangleright^{\hat{\beta}''_2} {}^s\theta'''$

Also since from (F-B2) we have $\exists {}^t v'' . {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s\theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau']_{V}^{\hat{\beta}''_2}$

Since $i = j + b + c + d + 1$ therefore from Lemma 1.65 we get

$$\exists {}^t v'' . {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s\theta''', k - i, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau']_{V}^{\hat{\beta}''_2}$$

12. CF-label:

$$\frac{\Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Gamma \vdash \text{Lb}_\ell(e_s) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_t)} \text{ label}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{Lb}_\ell(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\text{Labeled } \ell \tau)]_E^{\hat{\beta}}$

From Definition 1.61 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some $H_s, H_t, \hat{\beta}$ s.t. $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n$ s.t. $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$.

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-LB0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau]_E^{\hat{\beta}}$$

From Definition 1.61 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$ therefore $\exists j < i < n$ s.t $e_s \delta^s \Downarrow_j {}^s v$.

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [(\tau)]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta \quad (\text{F-LB1})$$

Since from (F-LB0) we are required to prove $({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}}$. Since from cg-label we know that $i = j + 1$, ${}^s v = {}^s v_1$ and ${}^t v = {}^t v_1$. Therefore we get this from Definition 1.60, (F-LB1) and Lemma 1.65.

From Lemma 1.65 we get $(n - i, H_s, H'_{t1}) \hat{\triangleright}^s \theta$

13. CF-toLabeled:

$$\frac{\Gamma \vdash e_s : \mathbb{C} \ell_1 \ell_2 \tau \rightsquigarrow e_t}{\Gamma \vdash \text{toLabeled}(e_s) : \mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau) \rightsquigarrow \lambda_. \text{inl}(e_t ())} \text{toLabeled}$$

Also given is: $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda_. \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright}^s \theta$ and given some $i < n$ s.t $\text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\lambda_. \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

From cg-val and fg-val we know that $i = 0$, ${}^s v = \text{toLabeled}(e_s) \delta^s$, ${}^t v = (\lambda_. \text{inl } e_t()) \delta^t$, $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^s \theta$$

Since we already know $(n, H_s, H_t) \hat{\triangleright}^s \theta$ from the context so we are left with proving

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_1 \perp (\text{Labeled } \ell_2 \tau))]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\forall^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} & (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda _ . \text{inl } e_t()) () \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket (\text{Labeled } \ell_2 \tau) \rrbracket_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda _ . \text{inl } e_t()) () \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket (\text{Labeled } \ell_2 \tau) \rrbracket_V^{\hat{\beta}''} \quad (\text{F-TL0}) \end{aligned}$$

IH:

$$({}^s \theta, k, e_s \delta^s, e_t \delta^t) \in \llbracket (\mathbb{C} \ell_1 \ell_2 \tau) \rrbracket_E^{\hat{\beta}}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_{h1}. e_s \delta^s \Downarrow_j {}^s v_{h1} \implies \\ & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in \llbracket (\mathbb{C} \ell_1 \ell_2 \tau) \rrbracket_V^{\hat{\beta}} \wedge (k - j, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \end{aligned}$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists j < i < k \leq n$ s.t $e_s \delta^s \Downarrow_j {}^s v_{h1}$.

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in \llbracket (\mathbb{C} \ell_1 \ell_2 \tau) \rrbracket_V^{\hat{\beta}} \wedge (k - j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \quad (\text{F-TL1.1}) \end{aligned}$$

From Definition 1.60 we know have

$$\forall^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k - j, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} & (m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies \\ & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m - b, {}^s v'_{h1}, {}^t v''_{h1}) \in \llbracket \tau \rrbracket_V^{\hat{\beta}''} \end{aligned}$$

Instantiating ${}^s \theta_e$ with ${}^s \theta$, H_{s3} with H_{s1} , H_{t3} with H'_{t2} , m with $k - j$ and $\hat{\beta}'$ with $\hat{\beta}$. Since we know that $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists b < i - j < k - j$ s.t $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$.

Therefore we have

$$\begin{aligned} & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - j - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k - j - b, {}^s v'_{h1}, {}^t v''_{h1}) \in \llbracket \tau \rrbracket_V^{\hat{\beta}''} \quad (\text{F-TL1}) \end{aligned}$$

In order to prove (F-TL0) we choose ${}^s\theta'$ as ${}^s\theta''$ and $\hat{\beta}'$ as $\hat{\beta}''$ (both chosen from (F-TL2))
 Also from cg-toLabeled and fg-inl, fg-app we know that $H'_s = H'_{s3}$ and $H'_t = H'_{t3}$, and
 ${}^s v' = {}^s v'_{h1}$, ${}^t v' = {}^t v'_{h1}$

Therefore we get the desired from (F-TL1) and Lemma 1.65

14. CF-unlabel:

$$\frac{\Gamma \vdash e_s : \text{Labeled } \ell \tau \rightsquigarrow e_t}{\Gamma \vdash \text{unlabel}(e_s) : \mathbb{C} \top \ell \tau \rightsquigarrow \lambda_{-} e_t} \text{ unlabel}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{-} e_t \delta^t) \in [\mathbb{C} \top (\ell) \tau]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ \exists H'_t, {}^t v. (H_t, \lambda_{-} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \top (\ell) \tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n, {}^s v$ s.t
 $\text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v. (H_t, \lambda_{-} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \top (\ell) \tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \\ \text{From cg-val and fg-val we know that } i = 0, {}^s v = \text{unlabel}(e_s) \delta^s, {}^t v = \lambda_{-} e_t \delta^t, H'_t = H_t \end{aligned}$$

And we need to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [\mathbb{C} \top (\ell) \tau]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we already know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ from the context so we are left with proving

$$({}^s\theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{-} e_t \delta^t) \in [\mathbb{C} \top (\ell) \tau]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} e_t)() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} e_t)() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}''} \quad (\text{F-U0}) \end{aligned}$$

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{Labeled } \ell \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_h \cdot e_s \delta^s \Downarrow_f {}^s v_h \implies \\ & \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} \\ & {}^s\theta_e \end{aligned}$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists f < i < k \leq n$ s.t. $e_s \delta^s \Downarrow_f {}^s v_h$.

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} \\ & {}^s\theta_e \quad (\text{F-U1}) \end{aligned}$$

In order to prove (F-U0) we choose H'_{t1} as H'_{t2} , ${}^t v'$ as ${}^t v_h$, ${}^s\theta'$ as ${}^s\theta_e$ and $\hat{\beta}''$ as $\hat{\beta}'$

From cg-unlabel and fg-app we also know that $H'_{s1} = H_{s1}$ and $H'_{t1} = H'_{t2}$

We need to prove

$$(a) (k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e:$$

Since from (F-U1) we know that $(k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$

Therefore from Lemma 1.67 we also get $(k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$

$$(b) \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'}:$$

Since from (F-U1) we have

$$({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'}$$

This means from Definition 1.60 we know that

$$\exists {}^s v_i, {}^t v_i. {}^s v_h = \text{Lb}_\ell({}^s v_i) \wedge {}^t v_h = \text{inl } {}^t v_i \wedge ({}^s\theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau]_V^{\hat{\beta}'} \quad (\text{F-U2})$$

Since we know that ${}^t v' = {}^t v_h$ and since from (F-U2) we have ${}^t v_h = \text{inl } {}^t v_i$. Therefore from we choose ${}^t v''$ as ${}^t v_i$ to get the first conjunct

From cg-unlabel we know that ${}^s v = {}^s v_i$ and since we know that $({}^s\theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau]_V^{\hat{\beta}'}$

Therefore from Lemma 1.65 we also get $({}^s\theta_e, k - i, {}^s v_i, {}^t v_i) \in [\tau]_V^{\hat{\beta}'}$

15. CF-ref:

$$\frac{\Gamma \vdash e_s : \text{Labeled } \ell' \tau \rightsquigarrow e_t \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash \text{new } e_s : \mathbb{C} \ell \perp (\text{ref } \ell' \tau) \rightsquigarrow \lambda_. \text{inl}(\text{new } (e_t))} \text{ref}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{new } e_s \delta^s, \lambda_. \text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. \text{new } e_s \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(\text{new } (e_t)) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ and given some $i < n, {}^s v$ s.t $\text{new } e_s \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that $i = 0, {}^s v = \text{new } e_s \delta^s, {}^t v = \lambda_{-}.\text{inl}(\text{new } (e_t)) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, \text{new } e_s \delta^s, \lambda_{-}.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$$

Since we already know $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ from the context so we are left with proving

$$({}^s\theta, n, \text{new } e_s \delta^s, \lambda_{-}.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \perp (\text{ref } \ell' \tau)]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, \text{new } e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(\text{new } e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(\text{new } e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}''} \quad (\text{F-N0}) \end{aligned}$$

From cg-ref we know that ${}^s v' = a_s$ and from fg-ref, fg-inl we know that ${}^t v' = \text{inl } a_t$.

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{Labeled } \ell' \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_h. e_s \delta^s \Downarrow_f {}^s v_h \implies \\ & \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} \\ & {}^s\theta_e \end{aligned}$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists f < i < k \leq n$ s.t $e_s \delta^s \Downarrow_f {}^s v_h$.

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k-f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_{V}^{\hat{\beta}'} \wedge (k-f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-N1})$$

In order to prove (F-N0) we choose H'_{t1} as $H'_{t2} \cup \{a_t \mapsto {}^t v_h\}$, ${}^t v$ as a_t , ${}^s \theta'$ as ${}^s \theta_n$ where ${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau)\}$

And we choose $\hat{\beta}''$ as $\hat{\beta}_n$ where $\hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$

From cg-ref and fg-ref we also know that $H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$

We need to prove

$$(a) (k-i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}_n} {}^s \theta_n:$$

From Definition 1.62 it suffices to prove that

- $\text{dom}({}^s \theta_n) \subseteq \text{dom}(H'_{s1})$:

Since $\text{dom}({}^s \theta_e) \subseteq \text{dom}(H_{s1})$ (given that we have $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$)

And since we know that

$${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau)\} \text{ and } H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^s v_h\}$$

Therefore we get $\text{dom}({}^s \theta_n) \subseteq \text{dom}(H'_{s1})$

- $\hat{\beta}_n \subseteq (\text{dom}({}^s \theta_n) \times \text{dom}(H'_{t1}))$:

Since $\hat{\beta}' \subseteq (\text{dom}({}^s \theta_e) \times \text{dom}(H_{t1}))$ (given that we have $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$)

And since we know that

$${}^s \theta_n = {}^s \theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau)\}, H'_{t1} = H_{t1} \cup \{a_t \mapsto {}^t v_h\} \text{ and } \hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$$

Therefore we get $\hat{\beta}_n \subseteq (\text{dom}({}^s \theta_n) \times \text{dom}(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}_n. ({}^s \theta_n, k-i-1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s \theta_n(a)]_{V}^{\hat{\beta}_n}$:

$$\forall (a_1, a_2) \in \hat{\beta}_n$$

- $(a_1, a_2) = (a_s, a_t)$:

Since from (F-N1) we know that $({}^s \theta_e, k-f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_{V}^{\hat{\beta}'}$

From Lemma 1.65 we get $({}^s \theta_n, k-i-1, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau)]_{V}^{\hat{\beta}_n}$

- $(a_1, a_2) \neq (a_s, a_t)$:

Since we have $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$ therefore

from Definition 1.62 we get

$$({}^s \theta_e, k-1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s \theta_e(a_1)]_{V}^{\hat{\beta}'}$$

From Lemma 1.65 we get

$$({}^s \theta_n, k-i-1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s \theta_n(a_1)]_{V}^{\hat{\beta}'}$$

$$(b) \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_n, k-i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau)]_{V}^{\hat{\beta}_n}:$$

We choose ${}^t v''$ as ${}^t v_h$ from (F-N1), fg-inl and fg-ref we know that ${}^t v' = \text{inl } {}^t v_h$

In order to prove $({}^s \theta_n, k-i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau)]_{V}^{\hat{\beta}_n}$, from Definition 1.60 it suffices to prove that

$${}^s \theta_n(a_s) = (\text{Labeled } \ell' \tau) \wedge (a_s, a_t) \in \hat{\beta}_n$$

We get this by construction of ${}^s \theta_n$ and $\hat{\beta}_n$

16. CF-deref:

$$\frac{\Gamma \vdash e_s : \text{ref } \ell \ \tau \rightsquigarrow e_t}{\Gamma \vdash !e_s : \mathbb{C} \top \perp (\text{Labeled } \ell \ \tau) \rightsquigarrow \lambda_{-}.\text{inl}(e_t)} \text{deref}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}}$

To prove: $({}^s\theta, n, !e_s \ \delta^s, \lambda_{-}.\text{inl}(e_t) \ \delta^t) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{E}}^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v. !e_s \ \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_t) \ \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright} {}^s\theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$ and given some $i < n$ s.t $!e_s \ \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_t) \ \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s\theta$$

From cg-val and fg-val we know that $i = 0, {}^s v = !e_s \ \delta^s, {}^t v = \lambda_{-}.\text{inl}(e_t) \ \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright} {}^s\theta$$

Since we already know $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$ from the context so we are left with proving

$$({}^s\theta, n, !e_s \ \delta^s, \lambda_{-}.\text{inl}(e_t) \ \delta^t) \in [\mathbb{C} \top \perp (\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \hat{\triangleright} ({}^s\theta_e) \wedge (H_{s1}, !e_s \ \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}''.(k - i, H'_{s1}, H'_{t1}) \hat{\triangleright} {}^s\theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright} {}^s\theta_e \wedge (H_{s1}, !e_s \ \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}''.(k - i, H'_{s1}, H'_{t1}) \hat{\triangleright} {}^s\theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell \ \tau)]_{\hat{V}}^{\hat{\beta}''} \quad (\text{F-D0}) \end{aligned}$$

IH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{ref } \ell \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_h \cdot e_s \delta^s \Downarrow_f {}^s v_h \implies \\ \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{ref } \ell \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists f < i < k \leq n$ s.t $e_s \delta^s \Downarrow_f {}^s v_h$.

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{ref } \ell \tau)]_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-D1})$$

In order to prove (F-D0) we choose H'_{t1} as H'_{t2} , ${}^t v'_1$ as $H'_{t2}(a)$ (where ${}^t v_h = a_t$ from fg-deref), ${}^s\theta'$ as ${}^s\theta_e$ and we choose $\hat{\beta}''$ as $\hat{\beta}'$.

From cg-deref we also know that $H'_{s1} = H_{s1}$

We need to prove

$$(a) (k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e:$$

Since from (F-D1) we have $(k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$ and since $f < i$ thfore from Lemma 1.67 we get $(k - i, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$

$$(b) \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'}$$

Since from cg-deref and fg-deref we know that ${}^s v_h = a_s$ and ${}^t v_h = a_t$.

Therefore from (F-D1) and from Definition 1.60 we know that

$${}^s\theta_e(a_s) = (\text{Labeled } \ell \tau) \wedge (a_s, a_t) \in \hat{\beta}'$$

Since from (F-D1) we know that $(k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$ which means from Definition 1.62 we know that

$$({}^s\theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'} \quad (\text{F-D2})$$

This means from Definition 1.60 we know that

$$\exists {}^s v_i, {}^t v_i. H_{s1}(a_s) = \text{Lb}_\ell({}^s v_i) \wedge H'_{t2}(a_t) = \text{inl } {}^t v_i \wedge ({}^s\theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau]_V^{\hat{\beta}'}$$

We choose ${}^t v''$ as ${}^t v_i$ and we know that ${}^t v' = H'_{t2}(a_t) = \text{inl } {}^t v_i$. This proves the first conjunct.

Since from (F-D2) we have $({}^s\theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'}$ therefore from Lemma 1.65 we get

$$({}^s\theta, k - i - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau)]_V^{\hat{\beta}'}$$

This proves the second conjunct.

17. CF-assign:

$$\frac{\Gamma \vdash e_{s1} : \text{ref } \ell' \tau \rightsquigarrow e_{t1} \quad \Gamma \vdash e_{s2} : \text{Labeled } \ell' \tau \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\Gamma \vdash e_{s1} := e_{s2} : \mathbb{C} \ell \perp \text{unit} \rightsquigarrow \lambda_{-}.\text{inl}(e_{t1} := e_{t2})} \text{ assign}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \perp \text{unit}]_E^{\hat{\beta}}$

It means from Definition 1.61 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s\theta \wedge \forall i < n, {}^s v. (e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright} {}^s\theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$ and given some $i < n, {}^s v$ s.t $(e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s\theta$$

From cg-val and fg-val we know that $i = 0, {}^s v = (e_{s1} := e_{s2}) \delta^s, {}^t v = \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright} {}^s\theta$$

Since we already know $(n, H_s, H_t) \hat{\triangleright} {}^s\theta$ from the context so we are left with proving

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \perp \text{unit}]_V^{\hat{\beta}}$$

From Definition 1.60 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'. \\ & (k, H_{s1}, H_{t1}) \hat{\triangleright} ({}^s\theta_e) \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright} \\ & {}^s\theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \hat{\triangleright} {}^s\theta_e \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t)) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . \\ & (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright} {}^s\theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}''} \quad (\text{F-S0}) \end{aligned}$$

IH1:

$$({}^s\theta_e, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{ref } \ell' \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_f {}^s v_{h1} \implies$$

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta_e, k-f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists f < i < k \leq n$ s.t $e_s \delta^s \Downarrow_f {}^s v_{h1}$.

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s\theta_e, k-f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-S1})$$

IH2:

$$({}^s\theta_e, k-f, e_{s2} \delta^s, e_{t2} \delta^t) \in [(\text{Labeled } \ell' \tau)]_E^{\hat{\beta}'}$$

It means from Definition 1.61 that we need to prove

$$\forall H_{s3}, H_{t3}. (k, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall l < k-f, {}^s v_{h2}. e_{s2} \delta^s \Downarrow_l {}^s v_{h2} \implies$$

$$\exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k-f-l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f-l, H_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e$$

Instantiating H_{s3} with H_{s1} and H_{t3} with H'_{t2} . And since we know that $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists l < i-f < k-f$ s.t $e_{s2} \delta^s \Downarrow_l {}^s v_{h2}$.

Therefore we have

$$\exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k-f-l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'} \wedge (k-f-l, H_{s1}, H'_{t3}) \triangleright^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-S2})$$

In order to prove (F-S0) we choose H'_{t1} as $H'_{t3}[a_t \mapsto {}^t v_{h3}]$, ${}^t v'$ as $(\)$, ${}^s\theta'$ as ${}^s\theta_e$ and $\hat{\beta}''$ as $\hat{\beta}'$

From cg-assign and fg-assign we also know that ${}^s v_{h2} = a_s$, ${}^t v_{h2} = a_t$, $H'_{s1} = H_{s1}[a_s \mapsto {}^s v_{h3}]$ and $H'_{t1} = H'_{t3}[a_t \mapsto {}^t v_{h3}]$

We need to prove

$$(a) (k-i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e:$$

From Definition 1.62 it suffices to prove that

- $\text{dom}({}^s\theta_e) \subseteq \text{dom}(H'_{s1})$:

Since $\text{dom}({}^s\theta_e) \subseteq \text{dom}(H_{s1})$ (given that we have $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$)

And since $\text{dom}(H_{s1}) = \text{dom}(H'_{s1})$ therefore we also get

$$\text{dom}({}^s\theta_e) \subseteq \text{dom}(H'_{s1})$$

- $\hat{\beta}' \subseteq (\text{dom}({}^s\theta_e) \times \text{dom}(H'_{t1}))$:

Since $\hat{\beta}' \subseteq (\text{dom}({}^s\theta_e) \times \text{dom}(H_{t1}))$ (given that we have $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$)

And since $\text{dom}(H_{t1}) \subseteq \text{dom}(H'_{t1})$ therefore we also have $\hat{\beta}' \subseteq (\text{dom}({}^s\theta_e) \times \text{dom}(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}' . ({}^s\theta_e, k - i - 1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$
 $\forall (a_1, a_2) \in \hat{\beta}_n$

$$- (a_1, a_2) = (a_s, a_t):$$

Since from (F-S2) we know that $({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'}$

From Lemma 1.65 we get $({}^s\theta_e, k - i - 1, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau)]_V^{\hat{\beta}'}$

$$- (a_1, a_2) \neq (a_s, a_t):$$

Since we have $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$ therefore
from Definition 1.62 we get

$$({}^s\theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$$

From Lemma 1.65 we get

$$({}^s\theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$$

- (b) $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in [\text{unit}]_V^{\hat{\beta}_n}$:

We choose ${}^t v''$ as () from (F-S1), fg-inl and fg-assign we know that ${}^t v' = \text{inl } ()$

To prove: $({}^s\theta_n, k - i, (), ()) \in [\text{unit}]_V^{\hat{\beta}_n}$,

We get this directly from Definition 1.60

□

Lemma 1.69 (Subtyping). *The following holds:*

\forall, τ, τ' .

$$1. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_V^{\hat{\beta}} \subseteq [(\tau')]_V^{\hat{\beta}}$$

$$2. \mathcal{L} \vdash \tau <: \tau' \implies [(\tau)]_E^{\hat{\beta}} \subseteq [(\tau')]_E^{\hat{\beta}}$$

Proof. Proof of Statement (1)

Proof by induction on $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove: $[((\tau_1 \rightarrow \tau_2))]_V^{\hat{\beta}} \subseteq [((\tau'_1 \rightarrow \tau'_2))]_V^{\hat{\beta}}$

It suffices to prove: $\forall ({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2))]_V^{\hat{\beta}} . ({}^s\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2))]_V^{\hat{\beta}}$

This means that given some ${}^s\theta, n$ and $\lambda x.e_i$ s.t $({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2))]_V^{\hat{\beta}}$

Therefore from Definition 1.60 we are given:

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}' .$$

$$({}^s\theta', j, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}'} \implies ({}^s\theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2]_E^{\hat{\beta}'} \quad (\text{S-A0})$$

And it suffices to prove: $({}^s\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2))]_V^{\hat{\beta}}$

Again from Definition 1.60 it suffices to prove:

$$\forall^s \theta'_1 \sqsupseteq {}^s \theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1.$$

$$({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}'_1} \implies ({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau'_2]_E^{\hat{\beta}'_1}$$

This means that given some ${}^s \theta'_1 \sqsubseteq {}^s \theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$ s.t. $({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}'_1}$

And we are required to prove: $({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau'_2]_E^{\hat{\beta}'_1}$

$$\text{IH: } [(\tau'_1)]_V^{\hat{\beta}'_1} \subseteq [(\tau_1)]_V^{\hat{\beta}'_1} \text{ (Statement (1))}$$

$$[(\tau_2)]_E^{\hat{\beta}'_1} \subseteq [(\tau'_2)]_E^{\hat{\beta}'_1} \text{ (Sub-A0, From Statement (2))}$$

Instantiating (S-A0) with ${}^s \theta'_1, {}^s v_1, {}^t v_1, k, \hat{\beta}'_1$

Since $({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}'_1}$ therefore from IH1 we know that $({}^s \theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'_1}$

As a result we get

$$({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau_2]_E^{\hat{\beta}'_1}$$

From (Sub-A0), we know that

$$({}^s \theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau'_2]_E^{\hat{\beta}'_1}$$

2. CGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove: $[((\tau_1 \times \tau_2))]_V^{\hat{\beta}} \subseteq [((\tau'_1 \times \tau'_2))]_V^{\hat{\beta}}$

$$\text{IH1: } [(\tau_1)]_V^{\hat{\beta}} \subseteq [(\tau'_1)]_V^{\hat{\beta}} \text{ (Statement (1))}$$

$$\text{IH2: } [(\tau_2)]_V^{\hat{\beta}} \subseteq [(\tau'_2)]_V^{\hat{\beta}} \text{ (Statement (1))}$$

It suffices to prove:

$$\forall ({}^s \theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2))]_V^{\hat{\beta}}. ({}^s \theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2))]_V^{\hat{\beta}}$$

This means that given $({}^s \theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2))]_V^{\hat{\beta}}$

Therefore from Definition 1.60 we are given:

$$({}^s \theta, n, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge ({}^s \theta, n, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}} \quad (\text{S-P0})$$

And it suffices to prove: $({}^s \theta, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2))]_V^{\hat{\beta}}$

Again from Definition 1.60, it suffices to prove:

$$({}^s \theta, n, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge ({}^s \theta, n, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}$$

Since from (S-P0) we know that $({}^s \theta, n, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$ therefore from IH1 we have

$$({}^s \theta, n, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}}$$

Similarly since from (S-P0) we have $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}$ therefore from IH2 we get $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau'_2]_V^{\hat{\beta}}$

3. CGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove: $[\tau_1 + \tau_2]_V^{\hat{\beta}} \subseteq [(\tau'_1 + \tau'_2)]_V^{\hat{\beta}}$

IH1: $[\tau_1]_V^{\hat{\beta}} \subseteq [\tau'_1]_V^{\hat{\beta}}$ (Statement (1))

IH2: $[\tau_2]_V^{\hat{\beta}} \subseteq [\tau'_2]_V^{\hat{\beta}}$ (Statement (1))

It suffices to prove: $\forall ({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in [(\tau'_1 + \tau'_2)]_V^{\hat{\beta}}$

This means that given: $({}^s\theta, n, {}^sv, {}^tv) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}}$

And it suffices to prove: $({}^s\theta, n, {}^sv, {}^tv) \in [(\tau'_1 + \tau'_2)]_V^{\hat{\beta}}$

2 cases arise

(a) ${}^sv = \text{inl } {}^sv_i$ and ${}^tv = \text{inl } {}^tv_i$:

From Definition 1.60 we are given:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau_1]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau'_1]_V^{\hat{\beta}}$$

From (S-S0) and IH1 we get

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau'_1]_V^{\hat{\beta}}$$

(b) ${}^sv = \text{inr } {}^sv_i$ and ${}^tv = \text{inr } {}^tv_i$:

Symmetric reasoning

4. CGsub-label:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell \sqsubseteq \ell'}{\mathcal{L} \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove: $[\text{Labeled } \ell \tau]_V^{\hat{\beta}} \subseteq [\text{Labeled } \ell' \tau']_V^{\hat{\beta}}$

IH: $[\tau]_V^{\hat{\beta}} \subseteq [\tau']_V^{\hat{\beta}}$ (Statement (1))

It suffices to prove:

$$\forall ({}^s\theta, n, {}^sv, {}^tv) \in [\text{Labeled } \ell \tau]_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in [\text{Labeled } \ell' \tau']_V^{\hat{\beta}}$$

This means that given some $({}^s\theta, n, {}^sv, {}^tv) \in [\text{Labeled } \ell \tau]_V^{\hat{\beta}}$

Therefore from Definition 1.60 we are given:

$$\exists s v', t v'. s v = \text{Lb}_\ell(s v') \wedge t v = \text{inl } t v' \wedge (s \theta, m, s v', t v') \in [\tau]_V^{\hat{\beta}} \quad (\text{S-L0})$$

And we are required to prove that

$$(s \theta, n, s v, t v) \in [((\text{Labeled } \ell' \tau'))]_V^{\hat{\beta}}$$

From Definition 1.60 it suffices to prove

$$\exists s v', t v'. s v = \text{Lb}_\ell(s v') \wedge t v = \text{inl } t v' \wedge (s \theta, m, s v', t v') \in [\tau']_V^{\hat{\beta}}$$

We get this directly from (S-L0) and IH

5. CGsub-CG:

$$\frac{\mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \ell'_1 \sqsubseteq \ell_1 \quad \mathcal{L} \vdash \ell_2 \sqsubseteq \ell'_2}{\mathcal{L} \vdash \mathbb{C} \ell_1 \ell_2 \tau <: \mathbb{C} \ell'_1 \ell'_2 \tau'}$$

$$\text{To prove: } [((\mathbb{C} \ell_1 \ell_2 \tau))]_V^{\hat{\beta}} \subseteq [((\mathbb{C} \ell'_1 \ell'_2 \tau'))]_V^{\hat{\beta}}$$

It suffices to prove:

$$\forall (s \theta, n, s v, t v) \in [((\mathbb{C} \ell_1 \ell_2 \tau))]_V^{\hat{\beta}}. (s \theta, n, s v, t v) \in [((\mathbb{C} \ell'_1 \ell'_2 \tau'))]_V^{\hat{\beta}}$$

$$\text{This means that given } (s \theta, n, s v, t v) \in [((\mathbb{C} \ell_1 \ell_2 \tau))]_V^{\hat{\beta}}$$

Therefore from Definition 1.60 we are given:

$$\forall s \theta_e \sqsupseteq s \theta, H_s, H_t, i, s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(k, H_s, H_t) \triangleright^{\hat{\beta}'} (s \theta_e) \wedge (H_s, s v) \Downarrow_i^f (H'_s, s v') \wedge i < k \implies$$

$$\exists H'_t, t v'. (H_t, t v()) \Downarrow (H'_t, t v') \wedge \exists s \theta' \sqsupseteq s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} s \theta' \wedge$$

$$\exists t v'' . t v' = \text{inl } t v'' \wedge (s \theta', k - i, s v', t v'') \in [\tau']_V^{\hat{\beta}''} \quad (\text{S-M0})$$

And we are required to prove

$$(s \theta, n, s v, t v) \in [((\mathbb{C} \ell'_1 \ell'_2 \tau'))]_V^{\hat{\beta}}$$

So again from Definition 1.60 we need to prove

$$\forall s \theta_{e1} \sqsupseteq s \theta, H_{s1}, H_{t1}, i_1, s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1.$$

$$(k_1, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'_1} (s \theta_{e1}) \wedge (H_{s1}, s v) \Downarrow_{i_1}^f (H'_{s1}, s v'_1) \wedge i_1 < k_1 \implies$$

$$\exists H'_{t1}, t v'_1. (H_{t1}, t v()) \Downarrow (H'_{t1}, t v'_1) \wedge \exists s \theta' \sqsupseteq s \theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''_1} s \theta' \wedge$$

$$\exists t v''_1 . t v'_1 = \text{inl } t v''_1 \wedge (s \theta', k_1 - i_1, s v'_1, t v''_1) \in [\tau']_V^{\hat{\beta}''_1}$$

This means we are given some $s \theta_{e1} \sqsupseteq s \theta, H_{s1}, H_{t1}, i_1, s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$ s.t. $(k_1, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'_1} (s \theta_{e1}) \wedge (H_{s1}, s v) \Downarrow_{i_1}^f (H'_{s1}, s v'_1) \wedge i_1 < k_1$

And we need to prove

$$\exists H'_{t1}, t v'_1. (H_{t1}, t v()) \Downarrow (H'_{t1}, t v'_1) \wedge \exists s \theta' \sqsupseteq s \theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''_1} s \theta' \wedge$$

$$\exists t v''_1 . t v'_1 = \text{inl } t v''_1 \wedge (s \theta', k_1 - i_1, s v'_1, t v''_1) \in [\tau']_V^{\hat{\beta}''_1}$$

We instantiate (S-M0) with ${}^s\theta_{e1}, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1, \hat{\beta}'_1$ we get

$$\begin{aligned} & \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_{e1}, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau]_{V'}^{\hat{\beta}''} \end{aligned}$$

$$\text{IH: } [(\tau)]_{V'}^{\hat{\beta}''} \subseteq [(\tau')]_{V'}^{\hat{\beta}} \hat{\beta}'' \text{ (Statement (1))}$$

Since we have $({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau]_{V'}^{\hat{\beta}''}$ therefore from IH we get $({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\tau')]_{V'}^{\hat{\beta}''}$

6. CGsub-base:

Trivial

Proof of Statement(2)

It suffice to prove that

$$\forall ({}^s \theta, n, e_s, e_t) \in [(\tau)]_E^{\hat{\beta}} . ({}^s \theta, n, e_s, e_t) \in [(\tau')]_E^{\hat{\beta}}$$

This means that we are given $({}^s \theta, n, e_s, e_t) \in [(\tau)]_E^{\hat{\beta}}$

From Definition 1.61 it means we have

$$\forall H_s, H_t . (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v . e_s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v . (H_t, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{Sub-E0})$$

And we need to prove

$$({}^s \theta, n, e_s, e_t) \in [(\tau')]_E^{\hat{\beta}}$$

From Definition 1.61 we need to prove

$$\forall H_{s1}, H_{t1} . (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1 . e_s \Downarrow_j {}^s v_1 \implies$$

$$\exists H'_{t1}, {}^t v_1 . (H_{t1}, e_t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [\tau']_{V'}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$$

This further means that given H_{s1}, H_{t1} s.t $(n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$. Also given some $j < n, {}^s v_1$ s.t $e_s \Downarrow_j {}^s v_1$

And it suffices to prove that

$$\exists H'_{t1}, {}^t v_1 . (H_{t1}, e_t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [\tau']_{V'}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$$

Instantiating (Sub-E0) with the given H_{s1}, H_{t1} and $j < n, {}^s v_1$. We get

$$\exists H'_t, {}^t v . (H_{t1}, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v) \in [\tau]_{V'}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we have $({}^s \theta, n - j, {}^s v_1, {}^t v) \in [\tau]_{V'}^{\hat{\beta}}$ therefore from Statement(1) we get $({}^s \theta, n - j, {}^s v_1, {}^t v) \in [(\tau')]_{V'}^{\hat{\beta}}$

□

Theorem 1.70 (Deriving CG NI via compilation). $\forall e_s, {}^s v_1, {}^s v_2, {}^s v'_1, {}^s v'_2, n_1, n_2, H'_{s1}, H'_{s2}$.

let $\text{bool} = (\text{unit} + \text{unit})$.

$x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \wedge$

$\emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset \vdash {}^s v_2 : \text{Labeled } \top \text{ bool} \wedge$

$(\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1}^f (H'_{s1}, {}^s v'_1) \wedge$

$$\begin{aligned}
& (\emptyset, e_s[sv_2/x]) \Downarrow_{n_2}^f (H'_{s_2}, sv'_2) \\
& \implies \\
& sv'_1 = sv'_2
\end{aligned}$$

Proof. From the CG to FG translation we know that $\exists e_t$ s.t

$$x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \rightsquigarrow e_t$$

Similarly we also know that $\exists t v_1, t v_2$ s.t

$$\emptyset \vdash sv_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow t v_1 \text{ and } \emptyset \vdash sv_2 : \text{Labeled } \top \text{ bool} \rightsquigarrow t v_2 \quad (\text{NI-0})$$

From type preservation theorem we know that

$$\begin{aligned}
& x : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \vdash_\top e_t : (\text{unit} \xrightarrow{\perp} ((\text{unit} + \text{unit})^\perp + \text{unit})^\perp)^\perp \\
& \emptyset \vdash_\top t v_1 : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \\
& \emptyset \vdash_\top t v_2 : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \quad (\text{NI-1})
\end{aligned}$$

Since we have $\emptyset \vdash sv_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow t v_1$

And since sv_1 and $t v_1$ are closed terms (from given and NI-1)

Therefore from Theorem 1.68 we have (we choose n s.t $n > n_1$ and $n > n_2$)

$$(\emptyset, n, sv_1, t v_1) \in [\text{Labeled } \top \text{ bool}]_E^\emptyset \quad (\text{NI-2})$$

And therefore from Definition 1.64 and (NI-2) we have

$$(\emptyset, n, (x \mapsto sv_1), (x \mapsto t v_1)) \in [x \mapsto \text{Labeled } \top \text{ bool}]_V^\emptyset$$

From (NI-0) we know that $x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \rightsquigarrow e_t$

Therefore we can apply Theorem 1.68 to get

$$(\emptyset, n, e_s[sv_1/x], e_t[t v_1/x]) \in [\mathbb{C} \perp \perp \text{ bool}]_E^\emptyset \quad (\text{NI-3.1})$$

Applying Definition 1.61 on (NI-3.1) we get

$$\forall H_{s_2}, H_{t_2}. (n, H_{s_2}, H_{t_2}) \triangleright \hat{\beta} \emptyset \wedge \forall i < n. e_s[sv_1/x] \Downarrow_i sv \implies$$

$$\exists H'_{t_2}, t v. (H_{t_2}, e_t[t v_1/x]) \Downarrow (H'_{t_2}, t v) \wedge (\emptyset, n - i, sv, t v) \in [\mathbb{C} \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s_2}, H'_{t_2}) \triangleright \hat{\beta} \emptyset$$

Instantiating with \emptyset, \emptyset . From cg-val we know that $i = 0$ and $sv = e_s[sv_1/x]$.

Therefore we have

$$\exists H'_{t_2}, t v. (H_{t_2}, e_t[t v_1/x]) \Downarrow (H'_{t_2}, t v) \wedge (\emptyset, n, sv, t v) \in [\mathbb{C} \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n, H_{s_2}, H'_{t_2}) \triangleright \hat{\beta} \emptyset$$

From translation and from (NI-1) we know that $t v = e_t[t v_1/x] = \lambda_. e_{b_1}$ and therefore from fg-val we have $H'_{t_2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[sv_1/x], \lambda_. e_{b_1}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\emptyset$$

Expanding $(\emptyset, n, e_s[sv_1/x], \lambda_. e_{b_1}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\emptyset$ using Definition 1.60 we get

$$\forall^s \theta_e \sqsupseteq \emptyset, H_{s_3}, H_{t_3}, i, sv'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s_3}, H_{t_3}) \triangleright \hat{\beta}' (s\theta_e) \wedge (H_{s_3}, e_s[sv_1/x]) \Downarrow_i^f (H'_{s_1}, sv'_1) \wedge i < k \implies$$

$$\begin{aligned}
& \exists H''_{t_1}, t v'', (H_{t_3}, (\lambda_. e_{b_1})()) \Downarrow (H''_{t_1}, t v'') \wedge \exists s\theta' \sqsupseteq s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s_1}, H''_{t_1}) \triangleright \hat{\beta}'' s\theta' \wedge \exists t v''' . t v'' = \\
& \text{inl } t v''' \wedge (s\theta', k - i, sv'_1, t v''') \in [\text{bool}]_V^{\hat{\beta}''}
\end{aligned}$$

Instantiating with $\emptyset, \emptyset, \emptyset, n_1, sv'_1, n, \emptyset$ we get

$$\begin{aligned}
& \exists H''_{t_1}, t v'' . (\emptyset, (\lambda_. e_{b_1})()) \Downarrow (H''_{t_1}, t v'') \wedge \exists s\theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s_1}, H''_{t_1}) \triangleright \hat{\beta}'' s\theta' \wedge \exists t v''' . t v'' = \\
& \text{inl } t v''' \wedge (s\theta', n - n_1, sv'_1, t v''') \in [\text{bool}]_V^{\hat{\beta}''} \quad (\text{NI-3.2})
\end{aligned}$$

Since we have $\exists t v''' . t v'' = \text{inl } t v''' \wedge (s\theta', n - n_1, sv'_1, t v''') \in [(\text{unit} + \text{unit})]_V^{\hat{\beta}''}$, therefore from Definition 1.60 we know that 2 cases arise

- ${}^s v'_1 = \text{inl}^s v'_{i1}$ and ${}^t v'''_1 = \text{inl}^t v'_{i1}$:

And from Definition 1.60 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i1}, {}^t v'_{i1}) \in [\text{unit}]_{\mathcal{V}}^{\hat{\beta}''}$$

which means ${}^s v'_{i1} = {}^t v'_{i1} = ()$

- ${}^s v'_1 = \text{inr}^s v'_{i1}$ and ${}^t v'''_1 = \text{inr}^t v'_{i1}$:

Same reasoning as in the previous case

Thus no matter which case occurs we have ${}^s v'_1 = {}^t v'''_1$ (NI-3.3)

Similarly we can apply Theorem 1.68 with the other substitution to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\mathbb{C} \perp \perp \text{bool}]_E^{\emptyset} \quad (\text{NI-4.1})$$

Applying Definition 1.61 on (NI-4.1) we get

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} \emptyset \wedge \forall i < n, {}^s v_s.e_s[{}^s v_2/x] \Downarrow_i {}^s v_s \implies \exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow \\ (H'_{t2}, {}^t v_s) \wedge (\emptyset, n - i, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathcal{V}}^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset \end{aligned}$$

Instantiating with \emptyset, \emptyset . From cg-val we know that $i = 0$ and ${}^s v_s = e_s[{}^s v_2/x]$.

Therefore we have

$$\exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathcal{V}}^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset$$

Also from (NI-1) and from translation we know that ${}^t v = e_t[{}^t v_2/x] = \lambda_.e_{b2}$ and therefore from fg-val we know that $H'_{t2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[{}^s v_2/x], \lambda_.e_{b2}) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathcal{V}}^{\emptyset}$$

Expanding $(\emptyset, n, e_s[{}^s v_2/x], \lambda x.e_{b2}) \in [\mathbb{C} \perp \perp \text{bool}]_{\mathcal{V}}^{\emptyset}$ using Definition 1.60 we get

$$\forall {}^s \theta_e \sqsupseteq \emptyset, H_{s3}, H_{t3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, e_s[{}^s v_2/x]) \Downarrow_i^f (H'_{s2}, {}^s v''_2) \wedge i < k \implies$$

$$\begin{aligned} \exists H''_{t2}, {}^t v'', (H_{t3}, (\lambda_.e_{b2})) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_2. {}^t v''_2 = \\ \text{inl}^t v'''_2 \wedge ({}^s \theta', k - i, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_{\mathcal{V}}^{\hat{\beta}''} \end{aligned}$$

Instantiating with $\emptyset, \emptyset, \emptyset, n_2, {}^s v'_2, n, \emptyset$ we get

$$\begin{aligned} \exists H''_{t2}, {}^t v'' . (\emptyset, (\lambda_.e_{b2})) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_2. {}^t v''_2 = \\ \text{inl}^t v'''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_{\mathcal{V}}^{\hat{\beta}''} \quad (\text{NI-4.2}) \end{aligned}$$

Since we have $\exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_{\mathcal{V}}^{\hat{\beta}''}$, therefore from Definition 1.60 2 cases arise

- ${}^s v'_2 = \text{inl}^s v'_{i2}$ and ${}^t v'''_2 = \text{inl}^t v'_{i2}$:

And from Definition 1.60 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i2}, {}^t v'_{i2}) \in [\text{unit}]_{\mathcal{V}}^{\hat{\beta}''}$$

which means ${}^s v'_{i2} = {}^t v'_{i2} = ()$

- ${}^s v'_2 = \text{inr}^s v'_{i2}$ and ${}^t v'''_2 = \text{inr}^t v'_{i2}$:

Same reasoning as in the previous case

Thus no matter which case occurs we have ${}^s v_2' = {}^t v_2'''$ (NI-4.3)

From CG to FG translation we know that $\exists^t v_{i1}. {}^t v_1 = \text{inl } {}^t v_{i1}$ and similarly $\exists^t v_{i2}. {}^t v_2 = \text{inl } {}^t v_{i2}$

From (NI-1) since $\emptyset \vdash_{\top} {}^t v_1 : (\text{bool}^{\perp} + \text{unit})^{\top}$ therefore from CG-inl we know that $\emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\perp}$

And from CGsub-sum we know that $\emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\top}$
Therefore we also have $\emptyset \vdash_{\perp} {}^t v_{i1} : \text{bool}^{\top}$ (NI-5.1)

Similarly we also have $\emptyset \vdash_{\perp} {}^t v_{i2} : \text{bool}^{\top}$ (NI-5.2)

Next, let $e_T = (\lambda x : (\text{bool}^{\perp} + \text{unit})^{\top}. \text{case}(e_t(), y.y, z. {}^t v_b)) (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : \text{bool}^{\perp}$

where $\text{true} = \text{inl } ()$ and $\text{false} = \text{inr } ()$

We claim $u : \text{bool}^{\top} \vdash_{\perp} e_T : \text{bool}^{\perp}$

To show this we give its typing derivation

P2.3:

$$\frac{\frac{\frac{}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{false} : \text{bool}^{\perp}}{\text{FG-inl}}}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.2:

$$\frac{\frac{\frac{}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{true} : \text{bool}^{\perp}}{\text{FG-inl}}}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.1:

$$\frac{}{u : \text{bool}^{\top} \vdash_{\perp} u : \text{bool}^{\top}}$$

P2:

$$\frac{\text{P2.1} \quad \text{P2.2} \quad \text{P2.3} \quad \frac{}{\mathcal{L} \models (\text{bool}^{\perp} + \text{unit})^{\top} \searrow \perp}}{u : \text{bool}^{\top} \vdash_{\perp} (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : (\text{bool}^{\perp} + \text{unit})^{\top}}$$

P1.2:

$$\frac{\frac{\frac{}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t : (\text{unit} \xrightarrow{\perp} (\text{bool}^{\perp} + \text{unit})^{\perp})^{\perp}}{\text{NI-1}}}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} () : \text{unit}} \text{FG-unit}}{\frac{\mathcal{L} \models \perp \sqcup \perp \sqsubseteq \perp \quad \mathcal{L} \models (\text{bool}^{\perp} + \text{unit})^{\perp} \searrow \perp}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t() : (\text{bool}^{\perp} + \text{unit})^{\perp}} \text{FG-app}}$$

P1.1:

$$\frac{\text{P1.2} \quad \frac{}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, y : \text{bool}^{\perp} \vdash_{\perp} y : \text{bool}^{\perp}}{\text{FG-var}}}{\frac{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, z : \text{unit} \vdash_{\perp} \text{false} : \text{bool}^{\perp} \quad \mathcal{L} \models \text{bool}^{\perp} \searrow \perp}{u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} \text{case}(e_t(), y.y, z. {}^t v_b) : \text{bool}^{\perp}} \text{FG-case}}$$

P1:

$$\frac{\frac{P1.1}{u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top \vdash_\perp \text{case}(e_t(), y.y, z.^t v_b) : \text{bool}^\perp}}{u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top . \text{case}(e_t(), y.y, z.^t v_b)) : ((\text{bool}^\perp + \text{unit})^\top \xrightarrow{\perp} \text{bool}^\perp)^\perp}}$$

Main derivation:

$$\frac{\frac{P1 \quad P2 \quad \frac{}{\mathcal{L} \models \perp \sqcup \perp \sqsubseteq \perp} \quad \frac{}{\mathcal{L} \models \text{bool}^\perp \searrow \perp}}{u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top . \text{case}(e_t(), y.y, z.^t v_b)) (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : \text{bool}^\perp} \text{FG-app}}$$

Assuming $e_{b1}()$ reduces in n_{t1} steps in (NI-3.2) and $e_{b2}()$ reduces in n_{t2} steps in (NI-4.2).

We instantiate Theorem 1.87 with $e_T, ^t v_{i1}, ^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H''_{t1}, H''_{t2}$ and \perp and therefore from (NI-3.3) and (NI-4.3) we get $^t v_1''' = ^t v_2'''$ and thus $^s v_1' = ^s v_2'$

□

1.4 FG to CG translation

1.4.1 Type directed (direct) translation from FG to CG

Definition 1.71.

$$\begin{aligned}
(\mathbf{b}) &= \mathbf{b} \\
(\mathbf{unit}) &= \mathbf{unit} \\
(\tau_1 \xrightarrow{\ell_c} \tau_2) &= (\tau_1) \rightarrow \mathbb{C} \ell_c \perp (\tau_2) \\
(\tau_1 \times \tau_2) &= (\tau_1) \times (\tau_2) \\
(\tau_1 + \tau_2) &= (\tau_1) + (\tau_2) \\
(\mathbf{ref} \ A^\ell) &= \mathbf{ref} \ \ell \ (A) \\
(A^\ell) &= \mathbf{Labeled} \ (\ell) \ (A)
\end{aligned}$$

For $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$, define $(\Gamma) = x_1 : (\tau_1), \dots, x_n : (\tau_n)$.

We use a coercion function defined as follows:

$ \begin{aligned} \mathbf{coerce_taint} &: \mathbb{C} \ pc \ \ell_c \ \tau' \rightarrow \mathbb{C} \ pc \ \perp \ \tau' \quad \text{when } \tau' = \mathbf{Labeled} \ \ell'_c \ \tau \text{ and } \ell_c \sqsubseteq \ell'_c \\ \mathbf{coerce_taint} &\triangleq \lambda x. \mathbf{toLabeled}(\mathbf{bind}(x, y. \mathbf{unlabel}(y))) \end{aligned} $

$$\begin{aligned}
&\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \mathbf{ret} \ x} \text{FC-var} \\
&\frac{\Gamma, x : \tau_1 \vdash_{\ell_c} e : \tau_2 \rightsquigarrow e_{c1}}{\Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_c} \tau_2)^\perp \rightsquigarrow \mathbf{ret}(\mathbf{Lb}(\lambda x. e_{c1}))} \text{FC-lam} \\
&\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_c} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \ell \sqcup pc \sqsubseteq \ell_c \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} e_1 \ e_2 : \tau_2 \rightsquigarrow \mathbf{coerce_taint}(\mathbf{bind}(e_{c1}, a. \mathbf{bind}(e_{c2}, b. \mathbf{bind}(\mathbf{unlabel} \ a, c. (c \ b)))))} \text{FC-app} \\
&\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_{c1}, a. \mathbf{bind}(e_{c2}, b. \mathbf{ret}(\mathbf{Lb}(a, b))))} \text{FC-prod} \\
&\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \mathbf{fst}(e) : \tau_1 \rightsquigarrow \mathbf{coerce_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ (a), b. \mathbf{ret}(\mathbf{fst}(b)))))} \text{FC-fst} \\
&\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} \mathbf{snd}(e) : \tau_2 \rightsquigarrow \mathbf{coerce_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ (a), b. \mathbf{ret}(\mathbf{snd}(b)))))} \text{FC-snd} \\
&\frac{\Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \mathbf{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{ret}(\mathbf{Lbinl}(a)))} \text{FC-inl} \\
&\frac{\Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \mathbf{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{ret}(\mathbf{Lbinr}(a)))} \text{FC-inr} \\
&\frac{\Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \mathbf{case}(e, x. e_1, y. e_2) : \tau \rightsquigarrow \mathbf{coerce_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ a, b. \mathbf{case}(b, x. e_{c1}, y. e_{c2}))))} \text{FC-case} \\
&\frac{\Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \mathbf{new} \ (e) : (\mathbf{ref} \ \tau)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{new} \ (a), b. \mathbf{ret}(\mathbf{Lb}b)))} \text{FC-ref}
\end{aligned}$$

$$\frac{\Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{FC-deref}$$

$$\frac{\Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}()))} \text{FC-assign}$$

1.4.2 Type preservation for FG to CG translation

Theorem 1.72 (Type preservation: FG to CG). *If $\Gamma \vdash_{pc} e : \tau$ in FG then there exists e' such that $\Gamma \vdash_{pc} e : \tau \rightsquigarrow e'$ such that there is a derivation of $(\Gamma) \vdash e' : \mathbb{C} pc \perp (\tau)$ in CG.*

Proof. Proof by induction on the \rightsquigarrow relation

1. FC-var:

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var}$$

$$\frac{}{(\Gamma), x : (\tau) \vdash x : (\tau)} \text{CG-var}$$

$$\frac{}{(\Gamma), x : (\tau) \vdash \text{ret } x : \mathbb{C} pc \perp (\tau)} \text{CG-ret}$$

2. FC-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Gamma \vdash_{pc} \lambda x.e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\lambda x.e_{c1}))} \text{FC-lam}$$

$$T_0 = \mathbb{C} pc \perp ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) = \mathbb{C} pc \perp \text{Labeled} \perp ((\tau_1 \xrightarrow{\ell_e} \tau_2))$$

$$T_1 = \mathbb{C} pc \perp \text{Labeled} \perp (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.0} = \text{Labeled} \perp (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.1} = (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.2} = \mathbb{C} \ell_e \perp (\tau_2)$$

P1:

$$\frac{\frac{P2}{(\Gamma), x : (\tau_1) \vdash e_{c1} : T_{1.2}} \text{IH}}{(\Gamma) \vdash \lambda x.e_{c1} : T_{1.1}} \text{CG-lam}$$

Main derivation:

$$\frac{\frac{P1}{(\Gamma) \vdash (\text{Lb}(\lambda x.e_{c1})) : T_{1.0}} \text{CG-label}}{(\Gamma) \vdash \text{ret}(\text{Lb}(\lambda x.e_{c1})) : T_1} \text{CG-ret}$$

3. FC-app:

$$\frac{\Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))))} \text{FC-app}$$

$$T_0 = \mathbb{C} \text{ pc } \perp ((\tau_1 \xrightarrow{\ell_3} \tau_2)^\ell) = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell ((\tau_1 \xrightarrow{\ell_3} \tau_2))$$

$$T_1 = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.1} = \text{ Labeled } \ell (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.2} = \mathbb{C} \top \ell (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.3} = (\tau_1) \rightarrow \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.4} = \mathbb{C} \ell_e \perp (\tau_2)$$

$$T_{1.5} = \mathbb{C} \ell_e \ell (\tau_2)$$

$$T_{1.6} = \mathbb{C} \text{ pc } \ell (\mathbb{A}^{\ell_i})$$

$$T_{1.7} = \mathbb{C} \text{ pc } \ell \text{ Labeled } (\ell_i) (\mathbb{A})$$

$$T_{1.9} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell_i (\mathbb{A})$$

$$T_{1.10} = \mathbb{C} \text{ pc } \perp (\tau_2)$$

$$T_2 = \mathbb{C} \text{ pc } \perp (\tau_1)$$

$$T_{c4} = \text{ Labeled } \ell_i (\mathbb{A})$$

$$T_{c3} = \mathbb{C} \top \ell_i (\mathbb{A})$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i (\mathbb{A})$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell_i (\mathbb{A})$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{ Labeled } \ell_i (\mathbb{A})$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pc2:

$$\frac{\frac{}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-unlabel}}$$

Pc1:

$$\frac{}{(\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\frac{\frac{Pc1 \quad Pc2 \quad \frac{P0}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{\text{CG-bind}}}{(\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-tolabeled}}{(\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}}$$

Pc:

$$\frac{\frac{\frac{Pc0}{(\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{(\Gamma) \vdash \text{coerce_taint} : T_c} \text{From Definition of coerce_taint}}$$

P6:

$$\frac{}{(\Gamma), a : T_{1.1}, b : (\tau_1), c : T_{1.3} \vdash b : (\tau_1)} \text{CG-var}$$

P5:

$$\frac{}{(\Gamma), a : T_{1.1}, b : (\tau_1), c : T_{1.3} \vdash c : T_{1.3}} \text{CG-var}$$

P4:

$$\frac{\frac{P5 \quad P6}{(\Gamma), a : T_{1.1}, b : \langle \tau_2 \rangle, c : T_{1.3} \vdash c b : T_{1.4}} \text{CG-app}}{(\Gamma), a : T_{1.1}, b : \langle \tau_2 \rangle, c : T_{1.3} \vdash c b : T_{1.5}} \text{CGSub-monad}$$

P3:

$$\frac{}{(\Gamma), a : T_{1.1}, b : \langle \tau_1 \rangle \vdash a : T_{1.1}} \text{CG-var}$$

P2:

$$\frac{\frac{P3}{(\Gamma), a : T_{1.1}, b : \langle \tau_1 \rangle \vdash \text{unlabel } a : T_{1.2}} \text{CG-unlabel} \quad P4}{(\Gamma), a : T_{1.1}, b : \langle \tau_1 \rangle \vdash \text{bind}(\text{unlabel } a, c.(c b)) : T_{1.6}} \text{CG-bind}$$

P1:

$$\frac{\frac{}{(\Gamma), a : T_{1.1} \vdash e_{c2} : T_2} \text{IH2, Weakening} \quad P2}{(\Gamma), a : T_{1.1} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b))) : T_{1.6}} \text{CG-bind}$$

P0:

$$\frac{\frac{}{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{Given, } \tau_2 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{By inversion}$$

Main derivation:

$$\frac{Pc \quad \frac{\frac{}{(\Gamma) \vdash e_{c1} : T_1} \text{IH1} \quad P1}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))) : T_{1.7}} \text{CG-bind}}{(\Gamma) \vdash \text{coerce_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))) : T_{1.9}} \text{CG-app}}{(\Gamma) \vdash \text{coerce_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c b)))) : T_{1.10}} \text{Definition 1.71}$$

4. FC-prod:

$$\frac{\Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))))} \text{FC-prod}$$

$$T_1 = \mathbb{C} \text{ pc} \perp ((\tau_1 \times \tau_2)^\perp)$$

$$T_2 = \mathbb{C} \text{ pc} \perp \text{Labeled} \perp ((\tau_1 \times \tau_2))$$

$$T_3 = \mathbb{C} \text{ pc} \perp \text{Labeled} \perp (\tau_1) \times (\tau_2)$$

$$T_{3.1} = \text{Labeled} \perp (\tau_1) \times (\tau_2)$$

$$T_4 = \mathbb{C} \text{ pc} \perp (\tau_1)$$

$$T_5 = \mathbb{C} \text{ pc} \perp (\tau_2)$$

P4:

$$\frac{}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_1 \rangle \vdash a : \langle \tau_1 \rangle} \text{CG-var}$$

P3:

$$\frac{}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_1 \rangle \vdash b : \langle \tau_2 \rangle} \text{CG-var}$$

P2:

$$\frac{\frac{\frac{P3 \quad P4}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_1 \rangle \vdash (a, b) : \langle \tau_1 \rangle \times \langle \tau_2 \rangle} \text{CG-prod}}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_2 \rangle \vdash \text{Lb}(a, b) : T_{3.1}} \text{CG-label}}{(\Gamma), a : \langle \tau_1 \rangle, b : \langle \tau_2 \rangle \vdash \text{ret}(\text{Lb}(a, b)) : T_3} \text{CG-ret}$$

P1:

$$\frac{\frac{\text{IH2} \quad P2}{(\Gamma), a : \langle \tau_1 \rangle \vdash e_{c2} : T_5} \text{CG-bind}}{(\Gamma), a : \langle \tau_1 \rangle \vdash \text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))) : T_3}$$

Main derivation:

$$\frac{\frac{\frac{\text{IH1} \quad P1}{(\Gamma) \vdash e_{c1} : T_4} \text{CG-bind}}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_3} \text{Definition 1.71}}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_1}$$

5. FC-fst:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))} \text{FC-fst}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle \tau_1 \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \langle (\tau_1 \times \tau_2)^\ell \rangle$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\tau_1 \times \tau_2) \rangle$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.3} = \text{Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.4} = \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.5} = \mathbb{C} \top \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_3 = \mathbb{C} \top \ell \langle \tau_1 \rangle$$

$$T_{3.1} = \mathbb{C} \text{ pc } \ell \langle \tau_1 \rangle$$

$$T_{3.2} = \mathbb{C} \text{ pc } \ell \langle \mathbb{A}^{\ell_i} \rangle$$

$$T_{3.3} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_{3.5} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_{3.6} = \mathbb{C} \text{ pc } \perp \langle \mathbb{A}^{\ell_i} \rangle$$

$$T_{c4} = \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_{c3} = \mathbb{C} \top \ell_i \langle \mathbb{A} \rangle$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i \langle \mathbb{A} \rangle$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i \langle \mathbb{A} \rangle$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau_1 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\overline{(\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{ CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{Pg}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{(\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{ CG-bind}}{(\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{ CG-tolabeled}$$

Pc:

$$\frac{\overline{(\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{ CG-lam}}{(\Gamma) \vdash \text{coerce_taint} : T_c} \text{ From Definition of coerce_taint}$$

P2:

$$\frac{\overline{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}} \text{ CG-var}}{\frac{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{fst}(b) : \langle \tau_1 \rangle}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{fst}(b)) : T_3} \text{ CG-ret}} \text{ CG-fst}$$

P1:

$$\frac{\overline{(\Gamma), a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}} \text{ CG-unlabel} \quad P2}{(\Gamma), a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))) : T_{3.1}} \text{ CG-bind}$$

P0:

$$\frac{\overline{(\Gamma) \vdash e_c : T_{2.2}} \text{ IH} \quad P1}{(\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.1}} \text{ CG-bind}}{\frac{(\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.2}}{(\Gamma) \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.3}} \text{ Definition 1.71}}$$

Main derivation:

$$\frac{\overline{(\Gamma) \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.5}} \text{ CG-app}}{(\Gamma) \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.6}} \text{ Definition 1.71}}{(\Gamma) \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_1}$$

6. FC-snd:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau_2 \searrow \ell}{\Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b))))))} \text{FC-snd}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle \tau_2 \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \langle (\tau_1 \times \tau_2)^\ell \rangle$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell \langle (\tau_1 \times \tau_2) \rangle$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.3} = \text{ Labeled } \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.4} = \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_{2.5} = \mathbb{C} \top \ell \langle \tau_1 \rangle \times \langle \tau_2 \rangle$$

$$T_3 = \mathbb{C} \top \ell \langle \tau_2 \rangle$$

$$T_{3.1} = \mathbb{C} \text{ pc } \ell \langle \tau_2 \rangle$$

$$T_{3.2} = \mathbb{C} \text{ pc } \ell \langle A^{\ell_i} \rangle$$

$$T_{3.3} = \mathbb{C} \text{ pc } \ell \text{ Labeled } \ell_i \langle A \rangle$$

$$T_{3.5} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell_i \langle A \rangle$$

$$T_{3.6} = \mathbb{C} \text{ pc } \perp \langle A^{\ell_i} \rangle$$

$$T_{c4} = \text{ Labeled } \ell_i \langle A \rangle$$

$$T_{c3} = \mathbb{C} \top \ell_i \langle A \rangle$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i \langle A \rangle$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell_i \langle A \rangle$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{ Labeled } \ell_i \langle A \rangle$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau_2 = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{\langle \Gamma \rangle, x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{\langle \Gamma \rangle, x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\overline{\langle \Gamma \rangle, x : T_{c0} \vdash x : T_{c0}} \text{ CG-var}$$

Pc0:

$$\frac{\overline{\overline{\overline{Pc1} \quad Pc2} \quad \frac{Pg}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i}} \text{ CG-bind}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{ CG-tolabeled}} \langle \Gamma \rangle, x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}$$

Pc:

$$\frac{\frac{Pc0}{(\Gamma) \vdash \lambda x. \text{toLabeled}(\text{bind}(x, y. \text{unlabel}(y))) : T_c} \text{CG-lam}}{(\Gamma) \vdash \text{coerce_taint} : T_c} \text{From Definition of coerce_taint}$$

P2:

$$\frac{\frac{\frac{\frac{}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}} \text{CG-var}}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{snd}(b) : \langle \tau_2 \rangle} \text{CG-snd}}{(\Gamma), a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{snd}(b)) : T_3} \text{CG-ret}}$$

P1:

$$\frac{\frac{\frac{}{(\Gamma), a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}} \text{CG-unlabel} \quad P2}{(\Gamma), a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b))) : T_{3.1}} \text{CG-bind}}$$

P0:

$$\frac{\frac{\frac{\frac{}{(\Gamma) \vdash e_c : T_{2.2}} \text{IH} \quad P1}{(\Gamma) \vdash \text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.1}} \text{CG-bind}}{(\Gamma) \vdash \text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.2}}}{(\Gamma) \vdash \text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.3}} \text{Definition 1.71}$$

Main derivation:

$$\frac{\frac{\frac{\frac{Pc \quad P0}{(\Gamma) \vdash \text{coerce_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.5}} \text{CG-app}}{(\Gamma) \vdash \text{coerce_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_{3.6}} \text{Definition 1.71}}{(\Gamma) \vdash \text{coerce_taint}(\text{bind}(e_c, a. \text{bind}(\text{unlabel}(a), b. \text{ret}(\text{snd}(b)))) : T_1}}$$

7. FC-inl:

$$\frac{\Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a. \text{ret}(\text{Lbinl}(a)))} \text{FC-inl}$$

$$T_1 = \mathbb{C} \text{ pc} \perp ((\tau_1 + \tau_2)^\perp)$$

$$T_{1.1} = \mathbb{C} \text{ pc} \perp \text{Labeled} \perp ((\tau_1 + \tau_2))$$

$$T_{1.2} = \mathbb{C} \text{ pc} \perp \text{Labeled} \perp (\tau_1) + (\tau_2)$$

$$T_{1.3} = \text{Labeled} \perp (\tau_1) + (\tau_2)$$

$$T_2 = \mathbb{C} \text{ pc} \perp (\tau_1)$$

P1:

$$\frac{\frac{\frac{\frac{}{(\Gamma), a : (\tau_1) \vdash a : (\tau_1)} \text{CG-var}}{(\Gamma), a : (\tau_1) \vdash \text{inl}(a) : (\tau_1) + (\tau_2)} \text{CG-inl}}{(\Gamma), a : (\tau_1) \vdash \text{Lbinl}(a) : T_{1.3}} \text{CG-label}}{(\Gamma), a : (\tau_1) \vdash \text{ret}(\text{Lbinl}(a)) : T_{1.2}} \text{CG-ret}}$$

Main derivation:

$$\frac{\frac{\frac{}{(\Gamma) \vdash e_c : T_2} \text{IH} \quad P1}{(\Gamma) \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinl}(a))) : T_{1.2}} \text{CG-bind}}{(\Gamma) \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinl}(a))) : T_1} \text{Definition 1.71}}$$

8. FC-inr:

$$\frac{\Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Gamma \vdash_{pc} \mathbf{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinr}(a)))} \text{FC-inr}$$

$$T_1 = \mathbb{C} \text{ pc } \perp ((\tau_1 + \tau_2)^\perp)$$

$$T_{1.1} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \perp ((\tau_1 + \tau_2))$$

$$T_{1.2} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \perp (\tau_1) + (\tau_2)$$

$$T_{1.3} = \mathbf{Labeled} \perp (\tau_1) + (\tau_2)$$

$$T_2 = \mathbb{C} \text{ pc } \perp (\tau_2)$$

P1:

$$\frac{\frac{\frac{\frac{}{(\Gamma), a : (\tau_2) \vdash a : (\tau_2)} \text{CG-var}}{(\Gamma), a : (\tau_2) \vdash \mathbf{inr}(a) : (\tau_1) + (\tau_2)} \text{CG-inr}}{(\Gamma), a : (\tau_2) \vdash \mathbf{Lbinr}(a) : T_{1.3}} \text{CG-label}}{(\Gamma), a : (\tau_2) \vdash \mathbf{ret}(\mathbf{Lbinr}(a)) : T_{1.2}} \text{CG-ret}}$$

Main derivation:

$$\frac{\frac{\frac{}{(\Gamma) \vdash e_c : T_2} \text{IH} \quad P1}{(\Gamma) \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinr}(a))) : T_{1.2}} \text{CG-bind}}{(\Gamma) \vdash \mathbf{bind}(e_c, a.\mathbf{ret}(\mathbf{Lbinr}(a))) : T_1} \text{Definition 1.71}}$$

9. FC-case:

$$\frac{\Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \mathbf{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \mathbf{coerce_taint}(\mathbf{bind}(e_c, a.\mathbf{bind}(\mathbf{unlabel} \ a, b.\mathbf{case}(b, x.e_{c1}, y.e_{c2})))} \text{FC-case}}$$

$$T_1 = \mathbb{C} \text{ pc } \perp (\tau)$$

$$T_2 = \mathbb{C} \text{ pc } \perp ((\tau_1 + \tau_2)^\ell)$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \ \ell \ (\tau_1 + \tau_2)$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \mathbf{Labeled} \ \ell \ ((\tau_1) + (\tau_2))$$

$$T_{2.3} = \mathbf{Labeled} \ \ell \ ((\tau_1) + (\tau_2))$$

$$T_{2.4} = \mathbb{C} \top \ \ell \ ((\tau_1) + (\tau_2))$$

$$T_{2.5} = (\tau_1) + (\tau_2)$$

$$T_3 = \mathbb{C} (\text{pc } \sqcup \ell) \perp (\tau)$$

$$\begin{aligned}
T_4 &= \mathbb{C} (pc \sqcup \ell) \ell (\tau) \\
T_5 &= \mathbb{C} (pc) \ell (\mathbb{A}^{\ell_i}) \\
T_{5.1} &= \mathbb{C} (pc) \ell \text{ Labeled } \ell_i (\mathbb{A}) \\
T_{5.3} &= \mathbb{C} (pc) (\perp) \text{ Labeled } \ell_i (\mathbb{A}) \\
T_{5.4} &= \mathbb{C} (pc) (\perp) (\mathbb{A}^{\ell_i}) \\
T_{c4} &= \text{Labeled } \ell_i (\mathbb{A}) \\
T_{c3} &= \mathbb{C} \top \ell_i (\mathbb{A}) \\
T_{c2} &= \mathbb{C} pc \ell_i (\mathbb{A}) \\
T_{c1} &= \mathbb{C} pc \perp \text{ Labeled } \ell_i (\mathbb{A}) \\
T_{c0} &= \mathbb{C} pc \ell \text{ Labeled } \ell_i (\mathbb{A}) \\
T_c &= T_{c0} \rightarrow T_{c1}
\end{aligned}$$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\overline{(\Gamma), x : T_{c0} \vdash x : T_{c0}} \text{ CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{Pg}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i}}{\overline{(\Gamma), x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}}} \text{ CG-bind}}{(\Gamma), x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{ CG-tolabeled}$$

Pc:

$$\frac{\overline{Pc0} \text{ CG-lam}}{\overline{(\Gamma) \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c}} \text{ From Definition of } \text{coerce_taint} \\
(\Gamma) \vdash \text{coerce_taint} : T_c$$

P2:

$$\frac{\overline{(\Gamma), a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}} \text{ CG-var}}{(\Gamma), a : T_{2.3}, b : T_{2.5}, x : (\tau_1) \vdash e_{c1} : T_3} \text{ IH2, Weakening} \\
\frac{(\Gamma), a : T_{2.3}, b : T_{2.5}, y : (\tau_2) \vdash e_{c2} : T_3}{(\Gamma), a : T_{2.3}, b : T_{2.5} \vdash \text{case}(b, x.e_{c1}, y.e_{c2}) : T_3} \text{ IH3, Weakening} \\
\text{CG-case}$$

$$\begin{array}{c}
P1: \\
\frac{\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\text{CG-unlabel}} \quad P2}{\langle \Gamma \rangle, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})) : T_3} \text{CG-bind}}{\langle \Gamma \rangle, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})) : T_4} \text{CG-sub}
\end{array}$$

$$\begin{array}{c}
P0: \\
\frac{\frac{\frac{}{\langle \Gamma \rangle \vdash e_c : T_{2.2}}{\text{IH1}} \quad P1}{\langle \Gamma \rangle \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_5} \text{CG-bind}
\end{array}$$

$$\begin{array}{c}
P0.2: \\
\frac{P0}{\langle \Gamma \rangle \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_{5.1}} \text{Definition 1.71}
\end{array}$$

$$\begin{array}{c}
P0.1: \\
\frac{\frac{\frac{P_c \quad P0.2}{\langle \Gamma \rangle \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.3}}{\langle \Gamma \rangle \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.4}} \text{Definition 1.71}
\end{array}$$

Main derivation:

$$\frac{P0.1}{\langle \Gamma \rangle \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_1}$$

10. FC-ref:

$$\frac{\Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau \searrow_{pc}}{\Gamma \vdash_{pc} \text{new } (e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})))} \text{FC-ref}$$

$$\begin{aligned}
T_1 &= \mathbb{C} \text{ } pc \perp \langle (\text{ref } \tau)^\perp \rangle \\
T_{1.1} &= \mathbb{C} \text{ } pc \perp \langle (\text{ref } A^{\ell_i})^\perp \rangle \\
T_{1.2} &= \mathbb{C} \text{ } pc \perp \text{Labeled} \perp \langle (\text{ref } A^{\ell_i}) \rangle \\
T_{1.3} &= \mathbb{C} \text{ } pc \perp \text{Labeled} \perp \text{ref } \ell_i \langle A \rangle \\
T_2 &= \mathbb{C} \text{ } pc \perp \langle \tau \rangle \\
T_{2.1} &= \mathbb{C} \text{ } pc \perp \langle A^{\ell_i} \rangle \\
T_{2.2} &= \mathbb{C} \text{ } pc \perp \text{Labeled } \ell_i \langle A \rangle \\
T_{2.3} &= \text{Labeled } \ell_i \langle A \rangle \\
T_{2.4} &= \mathbb{C} \text{ } pc \perp \text{ref } \ell_i \langle A \rangle \\
T_{2.5} &= \text{ref } \ell_i \langle A \rangle \\
T_{2.51} &= \text{Labeled} \perp \text{ref } \ell_i \langle A \rangle
\end{aligned}$$

$$\begin{array}{c}
P2: \\
\frac{\frac{\frac{}{\langle \Gamma \rangle_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{CG-var}}}{\langle \Gamma \rangle_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash \text{Lbb} : T_{2.51}} \text{CG-label}}{\langle \Gamma \rangle_{\vec{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{1.3}} \text{CG-ret}
\end{array}$$

P1:

$$\frac{\frac{\overline{(\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{new } (a) : T_{2.4}} \text{ CG-new} \quad P2}{(\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})) : T_{1.3}} \text{ CG-bind}}{(\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})) : T_{1.3}} \text{ CG-bind}}$$

Main derivation:

$$\frac{\frac{\frac{\overline{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2}} \text{ IH} \quad P1}{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2}} \text{ CG-bind}}{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb}))) : T_{1.3}} \text{ CG-bind}}{(\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb}))) : T_1} \text{ Definition 1.71}}$$

11. FC-deref:

$$\frac{\Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{ FC-deref}$$

$$T_1 = \mathbb{C} \text{ pc } \perp \langle \tau' \rangle$$

$$T_{1.1} = \mathbb{C} \text{ pc } \perp \langle A'^{\ell'_i} \rangle$$

$$T_{1.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell'_i \langle A' \rangle$$

$$T_2 = \mathbb{C} \text{ pc } \perp \langle (\text{ref } \tau)^\ell \rangle$$

$$T_{2.1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\text{ref } \tau) \rangle$$

$$T_{2.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle (\text{ref } A^{\ell_i}) \rangle$$

$$T_{2.3} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \langle \text{ref } \ell_i \langle A \rangle \rangle$$

$$T_{2.4} = \text{Labeled } \ell \langle \text{ref } \ell_i \langle A \rangle \rangle$$

$$T_{2.5} = \mathbb{C} \top \ell \langle \text{ref } \ell_i \langle A \rangle \rangle$$

$$T_{2.6} = \text{ref } \ell_i \langle A \rangle$$

$$T_{2.7} = \mathbb{C} \top \perp \langle \text{Labeled } \ell_i \langle A \rangle \rangle$$

$$T_{2.8} = \mathbb{C} \top \ell \langle \text{Labeled } \ell'_i \langle A' \rangle \rangle$$

$$T_{2.9} = \mathbb{C} \text{ pc } \ell \langle \text{Labeled } \ell'_i \langle A' \rangle \rangle$$

$$T_{c4} = \text{Labeled } \ell_i \langle A \rangle$$

$$T_{c3} = \mathbb{C} \top \ell_i \langle A \rangle$$

$$T_{c2} = \mathbb{C} \text{ pc } \ell_i \langle A \rangle$$

$$T_{c1} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \langle A \rangle$$

$$T_{c0} = \mathbb{C} \text{ pc } \ell \text{Labeled } \ell_i \langle A \rangle$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pg:

$$\frac{\overline{\mathcal{L} \vdash A^{\ell_i} \searrow \ell} \text{ Given, } \tau' = A^{\ell_i}}{\mathcal{L} \vdash \ell \sqsubseteq \ell_i} \text{ By inversion}$$

Pc2:

$$\frac{\overline{(\Gamma), x : T_{c0}, y : T_{c4} \vdash y : T_{c4}} \text{ CG-var}}{(\Gamma), x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{ CG-unlabel}$$

Pc1:

$$\frac{}{\langle \Gamma \rangle, x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2 \quad \frac{Pg}{\mathcal{L} \models \ell \sqsubseteq \ell_i}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-bind}}{\langle \Gamma \rangle, x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{CG-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{\langle \Gamma \rangle \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{\langle \Gamma \rangle \vdash \text{coerce_taint} : T_c} \text{From Definition of coerce_taint}$$

P2:

$$\frac{\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{2.6} \vdash b : T_{2.6}} \text{CG-var}}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{2.6} \vdash !b : T_{2.7}} \text{CG-deref}}{\langle \Gamma \rangle, a : T_{2.4}, b : T_{2.6} \vdash !b : T_{2.8}} \text{CG-sub, Lemma 1.73}$$

P1:

$$\frac{\frac{}{\langle \Gamma \rangle, a : T_{2.4} \vdash \text{unlabel } a : T_{2.5}} \text{CG-unlabel} \quad P2}{\langle \Gamma \rangle, a : T_{2.4} \vdash \text{bind}(\text{unlabel } a, b.!b) : T_{2.8}} \text{CG-bind}$$

P0:

$$\frac{\frac{}{\langle \Gamma \rangle \vdash e_c : T_{2.3}} \quad P1}{\langle \Gamma \rangle \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.9}} \text{CG-bind}$$

Main derivation:

$$\frac{\frac{Pc \quad P0}{\langle \Gamma \rangle \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.2}} \text{CG-app}}{\langle \Gamma \rangle \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.1}} \text{Definition 1.71}$$

12. FC-assign:

$$\frac{\Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{FC-assign}$$

$$T_1 = \mathbb{C} \text{ } pc \perp (\text{unit})$$

$$T_{1.1} = \mathbb{C} \text{ } pc \perp \text{unit}$$

$$T_2 = \mathbb{C} \text{ } pc \perp ((\text{ref } \tau)^\ell)$$

$$T_{2.1} = \mathbb{C} \text{ } pc \perp \text{Labeled } \ell ((\text{ref } \tau))$$

$$T_{2.2} = \mathbb{C} \text{ } pc \perp \text{Labeled } \ell ((\text{ref } A^{\ell_i}))$$

$T_{2.3} = \mathbb{C} \text{ pc } \perp \text{ Labeled } \ell \text{ ref } \ell_i \text{ (A)}$

$T_{2.4} = \text{Labeled } \ell \text{ ref } \ell_i \text{ (A)}$

$T_{2.5} = \mathbb{C} \top (\ell) \text{ ref } \ell_i \text{ (A)}$

$T_{2.6} = \text{ref } \ell_i \text{ (A)}$

$T_{2.7} = \mathbb{C} (\text{pc } \sqcup \ell) \perp \text{unit}$

$T_{2.71} = \mathbb{C} (\text{pc } \sqcup \ell) \ell \text{unit}$

$T_{2.8} = \mathbb{C} \text{ pc } (\ell) \text{unit}$

$T_{2.9} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell \text{unit}$

$T_3 = \mathbb{C} \text{ pc } \perp (\tau)$

$T_{3.1} = \mathbb{C} \text{ pc } \perp (\text{A}^{\ell_i})$

$T_{3.2} = \mathbb{C} \text{ pc } \perp \text{Labeled } \ell_i \text{ (A)}$

$T_{3.3} = \text{Labeled } \ell_i \text{ (A)}$

P4:

$$\frac{}{(\Gamma), a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c : T_{2.6}} \text{CG-var}$$

P5:

$$\frac{}{(\Gamma), a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash b : T_{3.3}} \text{CG-var}$$

P3:

$$\frac{\begin{array}{c} \text{Given} \\ \mathcal{L} \vdash \tau \searrow (\text{pc } \sqcup \ell) \\ \text{By inversion} \\ \mathcal{L} \vdash (\text{pc } \sqcup \ell) \sqsubseteq \ell_i \end{array}}{(\Gamma), a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c := b : T_{2.7}} \text{CG-assign}$$

$$\frac{}{(\Gamma), a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c := b : T_{2.71}} \text{CGsub-monad}$$

P2:

$$\frac{\frac{}{(\Gamma), a : T_{2.4}, b : T_{3.3} \vdash \text{unlabel } a : T_{2.5}} \text{CG-unlabel} \quad P3}{(\Gamma), a : T_{2.4}, b : T_{3.3} \vdash \text{bind}(\text{unlabel } a, c.c := b) : T_{2.8}} \text{CG-bind}}$$

P1:

$$\frac{\frac{}{(\Gamma), a : T_{2.4} \vdash e_{c2} : T_{3.2}} \text{IH2} \quad P2}{(\Gamma), a : T_{2.4} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)) : T_{2.8}} \text{CG-bind}}$$

P0:

$$\frac{\frac{}{(\Gamma) \vdash e_{c1} : T_{2.3}} \text{IH1} \quad P1}{(\Gamma) \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))) : T_{2.8}} \text{CG-bind}}$$

P0.1:

$$\frac{P0}{(\Gamma) \vdash \text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) : T_{2.9}} \text{CG-toLabeled}}$$

Main derivation:

$$\frac{P0.1 \quad \frac{}{(\Gamma), d : \text{Labeled } \ell \text{ unit} \vdash \text{ret}() : T_{1.1}}{(\Gamma) \vdash \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) , d.\text{ret}()) : T_{1.1}} \text{CG-bind}}{(\Gamma) \vdash \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) , d.\text{ret}()) : T_{1.1}} \text{CG-bind}}$$

□

Lemma 1.73 (Subtyping - Type preservation). *The following holds:*

1. $\forall \tau, \tau'$.

$$\mathcal{L} \vdash \tau <: \tau' \implies \langle \tau \rangle <: \langle \tau' \rangle$$

2. $\forall A, A'$.

$$\mathcal{L} \vdash A <: A' \implies \mathcal{L} \vdash \langle A \rangle <: \langle A' \rangle$$

Proof. Proof by simultaneous induction on $\tau <: \tau$ and $A <: A$

Proof of statement (1)

Let $\tau = A_1^{\ell_1}$ and $\tau' = A_2^{\ell_2}$

P2:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\mathcal{L} \vdash A_1 <: A_2} \text{ By inversion} \quad P1}{\mathcal{L} \vdash \langle A_1 \rangle <: \langle A_2 \rangle} \text{ IH(2) on } A_1 <: A_2$$

P1:

$$\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\mathcal{L} \vdash \ell_1 \sqsubseteq \ell_2} \text{ By inversion}$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\mathcal{L} \vdash \text{Labeled } \ell_1 \langle A_1 \rangle <: \text{Labeled } \ell_2 \langle A_2 \rangle} \text{ CGsub-labeled}}{\mathcal{L} \vdash \langle A_1^{\ell_1} \rangle <: \langle A_2^{\ell_2} \rangle}$$

Proof of statement (2)

We proceed by cases on $A <: A$

1. FGsub-base:

$$\frac{\overline{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{ CG-refl}}{\mathcal{L} \vdash \langle \mathbf{b} \rangle <: \langle \mathbf{b} \rangle} \text{ Definition 1.71}$$

2. FGsub-ref:

$$\frac{\overline{\mathcal{L} \vdash \text{ref } \ell_i \langle A \rangle <: \text{ref } \ell_i \langle A \rangle} \text{ CG-refl}}{\mathcal{L} \vdash \langle \text{ref } A^{\ell_i} \rangle <: \langle \text{ref } A^{\ell_i} \rangle} \text{ Definition 1.71}$$

3. FGsub-prod:

P1:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\mathcal{L} \vdash \langle \tau_1 \rangle <: \langle \tau'_1 \rangle} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\mathcal{L} \vdash \langle \tau_2 \rangle <: \langle \tau'_2 \rangle} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\mathcal{L} \vdash \langle \tau_1 \rangle \times \langle \tau_2 \rangle <: \langle \tau'_1 \rangle \times \langle \tau'_2 \rangle} \text{ CGsub-prod}}{\mathcal{L} \vdash \langle \tau_1 \times \tau_2 \rangle <: \langle \tau'_1 \times \tau'_2 \rangle} \text{ Definition 1.71}$$

4. FGsub-sum:

P1:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\mathcal{L} \vdash \langle \tau_1 \rangle <: \langle \tau'_1 \rangle} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\mathcal{L} \vdash \langle \tau_2 \rangle <: \langle \tau'_2 \rangle} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\mathcal{L} \vdash \langle \tau_1 \rangle + \langle \tau_2 \rangle <: \langle \tau'_1 \rangle + \langle \tau'_2 \rangle} \text{ CGsub-prod}}{\mathcal{L} \vdash \langle \tau_1 + \tau_2 \rangle <: \langle \tau'_1 + \tau'_2 \rangle} \text{ Definition 1.71}$$

5. FGsub-arrow:

$$T_1 = \langle \tau_1 \rangle \rightarrow \mathbb{C} \ell_e \perp \langle \tau_2 \rangle$$

$$T_2 = \langle \tau'_1 \rangle \rightarrow \mathbb{C} \ell'_e \perp \langle \tau'_2 \rangle$$

P2:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau_2 <: \tau'_2} \text{ By inversion, Weakening}}{\frac{\overline{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{ Given}}{\mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e} \text{ By inversion, Weakening}}{\mathcal{L} \vdash \mathbb{C} \ell_e \perp \langle \tau_2 \rangle <: \mathbb{C} \ell'_e \perp \langle \tau'_2 \rangle} \text{ IH(1), CGsub-monad}$$

P1:

$$\frac{\frac{\overline{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{ Given}}{\mathcal{L} \vdash \tau'_1 <: \tau_1} \text{ By inversion, Weakening}}{\mathcal{L} \vdash \langle \tau'_1 \rangle <: \langle \tau_1 \rangle} \text{ IH(1)}$$

Main derivation:

$$\frac{P1 \quad P2}{\mathcal{L} \vdash (\tau_1 \xrightarrow{\ell_s} \tau_2) <: (\tau'_1 \xrightarrow{\ell'_s} \tau'_2)} \text{ Definition 1.71}$$

6. FGsub-unit:

$$\frac{\overline{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{ CGsub-unit}}{\mathcal{L} \vdash (\text{unit}) <: (\text{unit})} \text{ Definition 1.71}$$

□

1.4.3 Logical relation for FG to CG translation

Definition 1.74 (${}^s\theta_2$ extends ${}^s\theta_1$). ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$
 $\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$

Definition 1.75 ($\hat{\beta}_2$ extends $\hat{\beta}_1$). $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$
 $\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$

Definition 1.76 (Unary value relation).

$$\begin{aligned} [\mathbf{b}]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in [\mathbf{b}] \wedge {}^tv \in [\mathbf{b}] \wedge {}^sv = {}^tv\} \\ [\text{unit}]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in [\text{unit}] \wedge {}^tv \in [\text{unit}]\} \\ [\tau_1 \times \tau_2]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, ({}^sv_1, {}^sv_2), ({}^tv_1, {}^tv_2)) \mid \\ &\quad ({}^s\theta, m, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}\} \\ [\tau_1 + \tau_2]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \text{inl } {}^sv, \text{inl } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in [\tau_1]_V^{\hat{\beta}}\} \cup \\ &\quad \{({}^s\theta, m, \text{inr } {}^sv, \text{inr } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in [\tau_2]_V^{\hat{\beta}}\} \\ [\tau_1 \xrightarrow{\ell_s} \tau_2]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \lambda x. e_s, \lambda x. e_t) \mid \\ &\quad \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv, {}^tv, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^sv, {}^tv) \in [\tau_1]_V^{\hat{\beta}'} \implies \\ &\quad ({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in [\tau_2]_E^{\hat{\beta}'}\} \\ [\text{ref } \tau]_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, a_s, a_t) \mid {}^s\theta(a_s) = \tau \wedge ({}^sa, {}^ta) \in \hat{\beta}\} \\ [\mathbf{A}']_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, \text{Lb}({}^tv)) \mid ({}^s\theta, m, {}^sv, {}^tv) \in [\mathbf{A}]_V^{\hat{\beta}}\} \end{aligned}$$

Definition 1.77 (Unary expression relation).

$$\begin{aligned} [\tau]_E^{\hat{\beta}} &\triangleq \{({}^s\theta, n, e_s, e_t) \mid \\ &\quad \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^sv. (H_s, e_s) \Downarrow_i (H'_s, {}^sv) \implies \\ &\quad \exists H'_t, {}^tv. (H_t, e_t) \Downarrow^f (H'_t, {}^tv) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \\ &\quad \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in [\tau]_V^{\hat{\beta}'}\} \end{aligned}$$

Definition 1.78 (Unary heap well formedness).

$$\begin{aligned} (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta &\triangleq \text{dom}({}^s\theta) \subseteq \text{dom}(H_s) \wedge \\ &\quad \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \\ &\quad \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n - 1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a_1)]_V^{\hat{\beta}} \end{aligned}$$

Definition 1.79 (Value substitution). $\delta^s : Var \mapsto Val, \delta^t : Var \mapsto Val$

Definition 1.80 (Unary interpretation of Γ).

$$\llbracket \Gamma \rrbracket_V^{\hat{\beta}} \triangleq \{({}^s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\ \forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in \llbracket \Gamma(x) \rrbracket_V^{\hat{\beta}}\}$$

1.4.4 Soundness proof for FG to CG translation

Lemma 1.81 (Monotonicity). $\forall {}^s\theta, {}^s\theta', n, {}^s v, {}^t v, n', \beta, \beta'$.

1. $\forall A. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket A \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^s v, {}^t v) \in \llbracket A \rrbracket_V^{\hat{\beta}'}$
2. $\forall \tau. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket \tau \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^s v, {}^t v) \in \llbracket \tau \rrbracket_V^{\hat{\beta}'}$

Proof. Proof by simultaneous induction on A and τ

Proof of statement (1)

We case analyze A in the last step

1. Case b :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in \llbracket b \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in \llbracket b \rrbracket_V^{\hat{\beta}'}$$

Since $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket b \rrbracket_V^{\hat{\beta}}$ therefore from Definition 1.76 we know that ${}^s v \in \llbracket b \rrbracket \wedge {}^t v \in \llbracket b \rrbracket$ and ${}^s v = {}^t v$

Therefore from Definition 1.76 we get the desired

2. Case unit :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}'}$$

Since $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}}$ therefore from Definition 1.76 we know that ${}^s v \in \llbracket \text{unit} \rrbracket \wedge {}^t v \in \llbracket \text{unit} \rrbracket$

Therefore from Definition 1.76 we get the desired

3. Case $\tau_1 \times \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.76 we know that ${}^sv = ({}^sv_1, {}^sv_2)$ and ${}^tv = ({}^tv_1, {}^tv_2)$.

We also know that $({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}}$ and $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}$

IH1: $({}^s\theta', n', {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}'}$ (From Statement (2))

IH2: $({}^s\theta', n', {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}'}$ (From Statement (2))

Therefore from Definition 1.76, IH1 and IH2 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case $\tau_1 + \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.76 two cases arise

(a) ${}^sv = \text{inl}({}^sv')$ and ${}^tv = \text{inl}({}^tv')$:

IH: $({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$ (From Statement (2))

Therefore from Definition 1.76 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b) ${}^sv = \text{inr}({}^sv')$ and ${}^tv = \text{inr}({}^tv')$:

Symmetric reasoning as in the previous case

5. Case $\tau_1 \xrightarrow{\ell_e} \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}'}$$

From Definition 1.76 we know that

sv is of the form $\lambda x.e_s$ (for some e_s) and tv is of the form $\lambda x.e_t$ (for some e_t) s.t

$$\begin{aligned} \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta', j, {}^sv_1, {}^tv_1) \in [\tau_1]_{V}^{\hat{\beta}'_1} \implies \\ ({}^s\theta', j, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau_2]_{E}^{\hat{\beta}'_1} \quad (\text{A0}) \end{aligned}$$

Similarly from Definition 1.76 we are required to prove

$$\begin{aligned} \forall {}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}''_1. ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_{V}^{\hat{\beta}''_1} \implies \\ ({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_{E}^{\hat{\beta}''_1} \end{aligned}$$

This means we are given some

$${}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' \text{ s.t. } ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_{V}^{\hat{\beta}''}$$

and we are required to prove

$$({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_{E}^{\hat{\beta}''}$$

Instantiating (A0) with ${}^s\theta'', {}^sv_2, {}^tv_2, k, \hat{\beta}''$ since

$$\begin{aligned} {}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n \text{ and } \hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}'' \text{ therefore we get} \\ ({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_{E}^{\hat{\beta}''} \end{aligned}$$

6. Case ref τ :

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\text{ref } \tau]_{V}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \tau]_{V}^{\hat{\beta}'}$$

From Definition 1.76 we know that ${}^sv = a_s$ and ${}^tv = a_t$. We also know that

$${}^s\theta(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}$$

From Definition 1.76, Definition 1.74 and Definition 1.75 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \tau]_{V}^{\hat{\beta}'}$$

Proof of Statement (2)

Let $\tau = \mathbf{A}^{\ell''}$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{A}^{\ell''}]_{V}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

From Definition 1.76 we know that

$$\exists {}^tv_i. {}^tv = \text{Lb}({}^tv_i) \text{ and } ({}^s\theta, n, {}^sv, {}^tv_i) \in [\mathbf{A}]_{V}^{\hat{\beta}}$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbf{A}^{\ell''}]_{V}^{\hat{\beta}'}$$

This means from Definition 1.76 we need to prove

$$({}^s\theta', n', {}^sv, {}^tv_i) \in [\mathbf{A}]_{V}^{\hat{\beta}'}$$

$$\text{IH: } ({}^s\theta', n', {}^sv, {}^tv_i) \in [\mathbf{A}]_{V}^{\hat{\beta}'} \quad (\text{From Statement (1)})$$

Therefore we get the desired directly from IH. □

Lemma 1.82 (Unary monotonicity for Γ). $\forall \theta, \theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$.

$$(\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies (\theta', n', \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}'}^{\hat{\beta}'}$$

Proof. Given: $(\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove: $(\theta', n', \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}'}^{\hat{\beta}'}$

From Definition 1.80 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}}^{\hat{\beta}}$$

And again from Definition 1.80 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}'}^{\hat{\beta}'}$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t)$:

Given

- $\forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}'}^{\hat{\beta}'}$:

Since we know that $\forall x_i \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}}^{\hat{\beta}}$ (given)

Therefore from Lemma 1.81 we get

$$\forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_{\hat{\beta}'}^{\hat{\beta}'}$$

□

Lemma 1.83 (Unary monotonicity for H). $\forall {}^s\theta, H_s, H_t, n, n', \hat{\beta}$.

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Proof. Given: $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n$

To prove: $(n', H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$

From Definition 1.78 it is given that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$$

And again from Definition 1.78 we are required to prove that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$$

- $\text{dom}({}^s\theta) \subseteq \text{dom}(H_S)$:

Given

- $\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t))$:

Given

- $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$:

Since we know that $\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$ (given)

Therefore from Lemma 1.81 we get

$$\forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_{\hat{\beta}}^{\hat{\beta}}$$

□

Lemma 1.84 (Coercion lemma). $\forall H, e, v.$

$$(H, e) \Downarrow_{-}^f (H', \text{Lb } v) \implies (H, \text{coerce_taint } e) \Downarrow_{-}^f (H', \text{Lb } v)$$

Proof. Given: $(H, e) \Downarrow_{-}^f (H', \text{Lb } v)$

To prove: $(H, \text{coerce_taint } e) \Downarrow_{-}^f (H', \text{Lb } v)$

From Definition of `coerce_taint` and `cg-app` it suffices to prove that $(H, \text{toLabeled}(\text{bind}(e, y.\text{unlabel}(y)))) \Downarrow_{-}^f (H', \text{Lb } v)$

From `cg-tolabeled` it suffices to prove that $(H, \text{bind}(e, y.\text{unlabel}(y))) \Downarrow_{-}^f (H', v)$

From `cg-bind` it suffices to prove that

1. $(H, e) \Downarrow_{-}^f (H'_1, v_1)$:

We are given that $(H, e) \Downarrow_{-}^f (H', v)$ therefore we have $H'_1 = H'$ and $v'_1 = \text{Lb } v$

2. $(H'_1, \text{unlabel}(y)[v_1/y]) \Downarrow_{-}^f (H', v)$:

It suffices to prove that

$$(H', \text{unlabel}(\text{Lb } v)) \Downarrow_{-}^f (H', v)$$

We get this directly from `cg-unlabel`

□

Theorem 1.85 (Fundamental theorem). $\forall \Gamma, \tau, e_s, e_t, pc, \delta^s, \delta^t, {}^s\theta, n, \hat{\beta}.$

$$\begin{aligned} & \Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \wedge \\ & ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{V}^{\hat{\beta}} \\ & \implies \\ & ({}^s\theta, n, e_s, \delta^s, e_t, \delta^t) \in [\tau]_{E}^{\hat{\beta}} \end{aligned}$$

Proof. Proof by induction on the \rightsquigarrow relation

1. FC-var:

$$\frac{}{\Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\})]_{V}^{\hat{\beta}}$

To prove: $({}^s\theta, n, x, \delta^s, \text{ret}(x), \delta^t) \in [\tau]_{E}^{\hat{\beta}}$

From Definition 1.77 it suffices to prove that

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, x, \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(x), \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge \\ & ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_{V}^{\hat{\beta}'} \end{aligned}$$

This means given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that $i = 0, {}^s v = x \delta^s$. Also from cg-ret we know that ${}^t v = x \delta^t$ and $H'_t = H_t$

And we are required to prove

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'} \quad (\text{F-V0})$$

We choose ${}^s \theta'$ as ${}^s \theta$ and $\hat{\beta}'$ as $\hat{\beta}$

(a) $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$: Given

(b) $({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}}$:

Since we are given $({}^s \theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\})]_V^{\hat{\beta}}$, therefore from Definition 1.80 we get $({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}}$

2. FC-lam:

$$\frac{\Gamma, x : \tau_1 \vdash_{\ell_e} e_s : \tau_2 \rightsquigarrow e_t}{\Gamma \vdash_{pc} \lambda x. e_s : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb} \lambda x. e_t)} \text{FC-lam}$$

Also given is: $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s \theta, n, (\lambda x. e_s) \delta^s, \text{ret}(\text{Lb} \lambda x. e_t) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_E^{\hat{\beta}}$

From Definition 1.77 it suffices to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies$$

$$\exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\lambda x. e_t))) \delta^t \Downarrow^f (H'_t, {}^t v) \wedge$$

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_V^{\hat{\beta}'}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ and given some $i < n, {}^s v$ s.t $(H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that ${}^s v = (\lambda x. e_s) \delta^s, H'_s = H_s$ and $i = 0$. Also from cg-ret, cg-label and cg-FI we know that $H'_t = H_t$ and ${}^t v = (\text{Lb}(\lambda x. e_t)) \delta^t$

It suffices to prove that

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n, H_s, H_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v, {}^t v) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_V^{\hat{\beta}'}$$

We choose ${}^s \theta'$ as ${}^s \theta$ and $\hat{\beta}'$ as $\hat{\beta}$

(a) $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$: Given

(b) $({}^s\theta, n, \lambda x.e_s \delta^s, \text{Lb}(\lambda x.e_t) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp]_{\hat{\beta}}^V$:

From Definition 1.76 it suffices to prove that

$$({}^s\theta, n, \lambda x.e_s \delta^s, (\lambda x.e_t) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)]_{\hat{\beta}}^V$$

Again from Definition 1.76 it suffices to prove that

$$\begin{aligned} \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1]_{\hat{\beta}'}^V \implies \\ ({}^s\theta', j, e_s[{}^s v_d/x] \delta^s, e_t[{}^t v_d/x] \delta^t) \in [\tau_2]_{\hat{\beta}'}^E \end{aligned}$$

This further means that given ${}^s\theta' \sqsupseteq {}^s\theta, {}^s v_d, {}^t v_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t. $({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1]_{\hat{\beta}'}^V$

And we are required to prove

$$({}^s\theta', j, e_s[{}^s v_d/x] \delta^s, e_t[{}^t v_d/x] \delta^t) \in [\tau_2]_{\hat{\beta}'}^E \quad (\text{F-L0})$$

Since we are given $({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1]_{\hat{\beta}'}^V$, therefore from Definition 1.80 and Lemma 1.82 we have

$$({}^s\theta', j, \delta^s \cup \{x \mapsto {}^s v_d\}, \delta^t \cup \{x \mapsto {}^t v_d\}) \in [(\Gamma \cup \{x \mapsto \tau_1\})]_{\hat{\beta}'}^V.$$

Therefore from IH we get

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_d\}, e_t \delta^t \cup \{x \mapsto {}^t v_d\}) \in [\tau_2]_{\hat{\beta}'}^E$$

We get (F-L0) directly from IH

3. FC-app:

$$\frac{\Gamma \vdash_{pc} e_{s1} : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{t1} \quad \Gamma \vdash_{pc} e_{s2} : \tau_1 \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} e_{s1} e_{s2} : \tau_2 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b))))} \text{FC-app}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^V$

To prove:

$$({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \in [\tau]_{\hat{\beta}}^E$$

This means from Definition 1.77 it suffices to prove

$$\begin{aligned} \forall H_s, H_t, (n, H_s, H_t) \triangleright {}^s\theta \wedge \forall i < n, {}^s v.(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v.(H_t, \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2]_{\hat{\beta}'}^V \end{aligned}$$

This further means that given some H_s, H_t s.t. $(n, H_s, H_t) \triangleright {}^s\theta$ and given some $i < n, {}^s v$ s.t. $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} \exists H'_t, {}^t v.(H_t, \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2]_{\hat{\beta}'}^V \quad (\text{F-A0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{E}^{\hat{\beta}}$$

This means from Definition 1.77 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \\ ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{V}^{\hat{\beta}'_1} \end{aligned}$$

We instantiate with H_s, H_t . And since we know that $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t $(H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1)$.

This means we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - \\ j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{V}^{\hat{\beta}'_1} \quad (\text{F-A1.0}) \end{aligned}$$

Since we know that $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{V}^{\hat{\beta}'_1}$ therefore from Definition 1.76 we know that $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$ s.t

$$({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell]_{V}^{\hat{\beta}'_1} \quad (\text{F-A1.1})$$

From Definition 1.76 we know that ${}^s v_1 = \lambda x. e'_s$ and ${}^t v_i = \lambda x. e'_t$ s.t

$$\begin{aligned} \forall {}^s\theta''_1 \sqsupseteq {}^s\theta'_1, {}^s v', {}^t v', l < (n - j), \hat{\beta}''_1 \sqsubseteq \hat{\beta}'_1. \\ ({}^s\theta''_1, l, {}^s v', {}^t v') \in [\tau_1]_{V}^{\hat{\beta}''_1} \implies ({}^s\theta''_1, l, e'_s[{}^s v'/x], e'_t[{}^t v'/x]) \in [\tau_2]_{E}^{\hat{\beta}''_1} \quad (\text{F-A1}) \end{aligned}$$

IH2:

$$({}^s\theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1]_{E}^{\hat{\beta}'_1}$$

This means from Definition 1.77 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n - j, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s\theta \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2 \delta^t) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - \\ j - k, {}^s v_2, {}^t v_2) \in [\tau_1]_{V}^{\hat{\beta}'_2} \end{aligned}$$

We instantiate with H'_{s1}, H'_{t1} . And since we know that $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists k < i - j < n - j$ s.t $(H'_{s1}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$.

This means we have

$$\begin{aligned} \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - \\ j - k, {}^s v_2, {}^t v_2) \in [\tau_1]_{V}^{\hat{\beta}'_2} \quad (\text{F-A2}) \end{aligned}$$

We instantiate (F-A1) with θ''_1 as θ'_2 , ${}^s v'$ as ${}^s v_2$, ${}^t v'$ as ${}^t v_2$, l as $n - j - k$ and $\hat{\beta}''_1$ as $\hat{\beta}'_2$. Therefore we get

$$({}^s\theta'_2, n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2]_{E}^{\hat{\beta}'_2}$$

From Definition 1.77 we have

$$\begin{aligned} \forall H_s, H_t. (n - j - k, H_s, H_t) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge \forall a < n - j - k, {}^s v. (H_s, e'_s[{}^s v_2/x]) \Downarrow_i (H'_{s3}, {}^s v_3) \implies \\ \exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s\theta'_3 \sqsupseteq {}^s\theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2. \\ (n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'_3} {}^s\theta'_3 \wedge ({}^s\theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2]_{V}^{\hat{\beta}'_3} \end{aligned}$$

Instantiating with H'_{s2}, H'_{t2} . since we know that $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists a < i - j - k < n - j - k$ s.t $(H'_{s2}, e'_s[{}^s v/x] \delta^s) \Downarrow_a (H'_{s3}, {}^s v_3)$

Therefore we have

$$\begin{aligned} \exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s\theta'_3 \sqsupseteq {}^s\theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2. \\ (n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'_3} {}^s\theta'_3 \wedge ({}^s\theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2]_{V}^{\hat{\beta}'_3} \quad (\text{F-A3}) \end{aligned}$$

Let $\tau_2 = A_2^{\ell_i}$, since $\tau_2 \searrow \ell$ therefore $\ell \sqsubseteq \ell_i$ and

$$({}^s\theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2]_{V}^{\hat{\beta}'_3}$$

Therefore from Definition 1.76 we know that

$$({}^s\theta'_3, n - j - k - a, {}^s v_3, \text{Lb}^t v_{3i}) \in [\tau_2]_{V}^{\hat{\beta}'_3} \quad (\text{F-A3.1})$$

In order to prove (F-A0) we choose H'_t as H'_{t3} and ${}^t v$ as $\text{Lb}({}^t v_{3i})$. We need to prove:

$$(a) (H_t, \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b)))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$$

From Lemma 1.84 it suffices to prove that

$$(H_t, \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i}))$$

From cg-bind it further suffices to show that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1):$

We get this directly from (F-A1.0)

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c \ b))[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$

From cg-bind it suffices to prove that

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2):$

We get this directly from (F-A2)

- $(H'_{t2}, \text{bind}(\text{unlabel } a, c.c \ b)[{}^t v_1/a][{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i})):$

From cg-bind again it suffices to prove

- * $(H'_{t2}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t31}, {}^t v_{t2}):$

Since from (F-A1.1) we know that $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel and (F-A1) we know that $H'_{t31} = H'_{t2}$ and ${}^t v_{t2} = {}^t v_i = \lambda x. e'_t$

- * $((c \ b)[{}^t v_2/b][{}^t v_{t2}/c] \delta^t) \Downarrow {}^t v_{t21}:$

It suffices to prove that

$$((\lambda x. e'_t) {}^t v_2 \delta^t) \Downarrow {}^t v_{t21}$$

From cg-app we know that

$${}^t v_{t21} = e'_t[{}^t v_2/x] \delta^t$$

* $(H'_{t2}, {}^t v_{21}) \Downarrow^f (H'_{t3}, \text{Lb}({}^t v_{3i}))$:
 From (F-A3) and (F-A3.1) we get the desired

(b) $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau_2]_{V}^{\hat{\beta}'}$:

We choose ${}^s \theta'$ as ${}^s \theta'_3$ and $\hat{\beta}'$ as $\hat{\beta}'_3$. From fg-app we know that $i = j + k + a + 1$, ${}^s v = {}^s v_3$ and $H'_s = H'_{s3}$. Also from the termination proof (previous point) we know that $H'_t = H'_{t3}$ and ${}^t v = \text{Lb}({}^t v_3)$

We get $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta'$ from (F-A3) and Lemma 1.83

Since ${}^t v = \text{Lb}({}^t v_3)$ therefore from Definition 1.76 it suffices to prove that

$$({}^s \theta'_3, n - j - k - a - 1, {}^s v_3, {}^t v_3) \in [\tau_2]_{V}^{\hat{\beta}'_3}$$

We get this directly from (F-A3) and Lemma 1.81

4. FC-prod:

$$\frac{\Gamma \vdash_{pc} e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Gamma \vdash_{pc} e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Gamma \vdash_{pc} (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b))))} \text{prod}$$

Also given is: $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_{V}^{\hat{\beta}}$

To prove: $({}^s \theta, n, (e_{s1}, e_{s2}) \delta^s, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \in [(\tau_1 \times \tau_2)^\perp]_{E}^{\hat{\beta}}$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v_1, {}^s v_2. (H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2)) \implies \\ & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [(\tau_1 \times \tau_2)^\perp]_{V}^{\hat{\beta}'} \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$. Also given some $i < n, {}^s v_1, {}^s v_2$ s.t $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp]_{V}^{\hat{\beta}'} \quad (\text{F-P0}) \end{aligned}$$

IH1:

$$({}^s \theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1]_{E}^{\hat{\beta}}$$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_{V}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$ therefore $\exists j < i < n$ s.t $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_{V'}^{\hat{\beta}'_1} \quad (\text{F-P1}) \end{aligned}$$

IH2:

$$({}^s \theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_2]_E^{\hat{\beta}'_1}$$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s \theta'_1 \wedge \forall k < n - j, {}^s v_1. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_1) \implies \\ & \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_1) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2]_{V'}^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with H'_{s1}, H'_{t1} and since we know that $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$ therefore $\exists k < i - j < n - j$ s.t $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2]_{V'}^{\hat{\beta}'_2} \quad (\text{F-P2}) \end{aligned}$$

In order to prove (F-P0) we choose H_t as H'_{t2} and ${}^t v$ as $\text{Lb}({}^t v_1, {}^t v_2)$

(a) $(H_t, (\text{bind}(e_{t1}, a. \text{bind}(e_{t2}, b. \text{ret}(\text{Lb}(a, b)))))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$:

From cg-bind it suffices to prove that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{tb1}, {}^t v_{tb1})$:
From (F-P1) we know that $H'_{tb1} = H'_{t1}$ and ${}^t v_{tb1} = {}^t v_1$
- $(H'_{t1}, \text{bind}(e_{t2}, b. \text{ret}(\text{Lb}(a, b)))) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$:
From cg-bind it suffices to prove that

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{tb2}, {}^t v_{tb2})$:
From (F-P2) we know that $H'_{tb2} = H'_{t2}$ and ${}^t v_{tb2} = {}^t v_2$
- $(H'_{t2}, \text{ret}(\text{Lb}(a, b))) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$:
We get this from cg-ret, (F-P1) and (F-P2)

(b) $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp]_{V'}^{\hat{\beta}'_1}$:

We choose ${}^s \theta'$ as ${}^s \theta'_2$ and $\hat{\beta}'$ as $\hat{\beta}'_2$ and since from fg-prod $i = j + k + 1$ and $H'_s = H'_{s2}$. Therefore from (F-P2) and Lemma 1.83 we get

$$(n - i, H'_s, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'$$

In order to prove $({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp]_{V'}^{\hat{\beta}'_2}$

From Definition 1.76 it suffices to prove

$$\exists {}^t v_i. {}^t v = \text{Lb}({}^t v_i) \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v_i) \in [(\tau_1 \times \tau_2)]_{V'}^{\hat{\beta}'_2}$$

Since ${}^t v = \text{Lb}({}^t v_1, {}^t v_2)$ therefore we get the desired from (F-P1), (F-P2), Definition 1.76 and Lemma 1.81

5. FC-fst:

$$\frac{\Gamma \vdash_{pc} e_s : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_t \quad \mathcal{L} \vdash \tau_1 \searrow \ell}{\Gamma \vdash_{pc} \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))))} \text{fst}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \delta^t) \in [\tau_1]_{\hat{\beta}}^{\hat{\beta}}$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_{\hat{\beta}'}^{\hat{\beta}'} \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$

We need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_{\hat{\beta}'}^{\hat{\beta}'} \quad (\text{F-F0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}}^{\hat{\beta}}$$

This means from Definition 1.77 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-F1}) \end{aligned}$$

Since we know that $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1}$ therefore from Definition 1.76 we know that ${}^t v_1 = \text{Lb}({}^t v_i)$ s.t

$$({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \times \tau_2)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-F1.1})$$

From Definition 1.76 we know that ${}^s v_1 = ({}^s v_{i1}, {}^s v_{i2})$ and ${}^t v_i = ({}^t v_{i1}, {}^t v_{i2})$ s.t

$$({}^s\theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-F1.2})$$

Let $\tau_1 = \mathbf{A}_1^{\ell_i}$, since $\tau_1 \searrow \ell$ therefore $\ell \sqsubseteq \ell_i$ and

$$({}^s\theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\mathbf{A}_1^{\ell_i}]_V^{\hat{\beta}}$$

Therefore from Definition 1.76 we know that

$$({}^s\theta'_1, n - j, {}^s v_{i1}, \mathbf{Lb}^t v_{i11}) \in [\mathbf{A}_1]_V^{\hat{\beta}'_1} \quad (\text{F-F1.3})$$

In order to prove (F-F0) we choose H'_t as H'_{t1} and ${}^t v$ as ${}^t v_{i1} (= \mathbf{Lb}^t v_{i11})$ as we need to prove

$$(a) \quad (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_{t1}, \mathbf{Lb}^t v_{i11}):$$

From Lemma 1.84 it suffices to prove that

$$(H_t, \text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) \Downarrow^f (H'_{t1}, \mathbf{Lb}({}^t v_{i11}))$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{i11})$:

From (F-F1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{i11} = {}^t v_{i1} = \mathbf{Lb}({}^t v_{i1})$

- $(H'_{t1}, \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))[{}^t v_{i1}/a] \delta^t) \Downarrow^f (H'_{t1}, \mathbf{Lb}({}^t v_{i11}))$:

Again from cg-bind it suffices to prove that

- $(H'_{t1}, \text{unlabel}(a)[{}^t v_{i1}/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{i21})$:

Since ${}^t v_{i1} = \mathbf{Lb}({}^t v_{i1}, {}^t v_{i2})$ from (F-F1.1) and (F-F1.2) therefore we get the desired from cg-unlabel

So, $H'_{t21} = H'_{t1}$ and ${}^t v_{i21} = ({}^t v_{i1}, {}^t v_{i2})$

- $(H'_{t1}, \text{ret}(\text{fst}(b))[({}^t v_{i1}, {}^t v_{i2})/b] \delta^t) \Downarrow^f (H'_{t1}, \mathbf{Lb}({}^t v_{i11}))$:

We get the desired from cg-fst and cg-ret and (F-F1.3)

$$(b) \quad \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v_{i1}) \in [\tau_1]_V^{\hat{\beta}'_1}$$

We choose ${}^s\theta'$ as ${}^s\theta'_1$ and $\hat{\beta}'$ as $\hat{\beta}'_1$. And from fg-fst we know that $i = j + 1$ and $H'_s = H'_{s1}$ therefore from (F-F1) and Lemma 1.83 we get

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1$$

Since from fg-fst we know that ${}^s v = {}^s v_{i1}$ therefore from (F-F1.2) and Lemma 1.81 we get

$$({}^s\theta', n - i, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1]_V^{\hat{\beta}'_1}$$

6. FC-snd:

Symmetric reasoning as in the FC-fst case

7. FC-inl:

$$\frac{\Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_t}{\Gamma \vdash_{pc} \text{inl}(e_s) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{ret}(\mathbf{Lbinl}(a)))} \text{inl}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{inl}(e_s) \delta^s, \text{bind}(e_t, a.\text{ret}(\mathbf{Lbinl}(a))) \delta^t) \in [(\tau_1 + \tau_2)^\perp]_E^{\hat{\beta}}$

This means from Definition 1.77 we have

$$\begin{aligned}
& \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\
& \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\
& (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}
\end{aligned}$$

This means that we are given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright}^{\gamma, \hat{\beta}} s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned}
& \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\
& (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp]_V^{\hat{\beta}'} \quad (\text{F-IL0})
\end{aligned}$$

III:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1]_E^{\hat{\beta}}$$

This means from Definition 1.77 we need to prove

$$\begin{aligned}
& \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\
& \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\
& (n - j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} s\theta' \wedge ({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'_1}
\end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned}
& \exists H'_{t1}, {}^t v_1. (H_t, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\
& (n - j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} s\theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'_1} \quad (\text{F-IL1})
\end{aligned}$$

In order to prove (F-IL0) we choose H'_t as H'_{t1} and ${}^t v$ as $(\text{Lb inl}({}^t v_1))$ and we need to prove:

$$(a) (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1))):$$

From cg-bind it suffices to prove that

$$i. (H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11}):$$

From (F-IL1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t11} = {}^t v_1$

$$ii. (H'_{t1}, \text{ret}(\text{Lbinl}(a)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1))):$$

From cg-ret and (F-IL1)

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp]_V^{\hat{\beta}'}:$$

We choose ${}^s \theta'$ as ${}^s \theta'_1$ and $\hat{\beta}'$ as $\hat{\beta}'_1$. Since from fg-inl we know that $i = j + 1$ and $H'_s = H'_{s1}$ therefore from (F-IL1) and Lemma 1.83 we get

$$(n - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} s\theta'_1$$

Now we need to prove $({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp]_V^{\hat{\beta}'}$

Since ${}^s v = \text{inl } {}^s v_1$ and ${}^t v = \text{Lb}(\text{inl}({}^t v_1))$ therefore from Definition 1.76 it suffices to prove that

$$({}^s\theta', n - i, \text{inl } {}^s v_1, \text{inl } {}^t v_1) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}'}$$

Since from (F-IL1) we know that $({}^s\theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}'}$

Therefore from Lemma 1.81 and Definition 1.76 we get

$$({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)]_V^{\hat{\beta}'}$$

8. FC-inr:

Symmetric reasoning as in the FC-inl case

9. FC-case:

$$\frac{\Gamma \vdash_{pc} e_s : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_t \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \tau \searrow \ell}{\Gamma \vdash_{pc} \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2}))))} \text{case}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \in [\tau]_E^{\hat{\beta}}$$

This means from Definition 1.77 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge$$

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$$

This means we are given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge$$

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'} \quad (\text{F-C0})$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2)^\ell]_E^{\hat{\beta}}$$

This means from Definition 1.77 we have

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau]_V^{\hat{\beta}'_1}$$

Instantiating with H_s, H_t and since we know that $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \stackrel{\hat{\beta}'_1}{\triangleright} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell]_{V'}^{\hat{\beta}'_1} \quad (\text{F-C1}) \end{aligned}$$

Since from (F-C1) we have $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell]_{V'}^{\hat{\beta}'_1}$ therefore from Definition 1.76 we know that

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 + \tau_2)]_{V'}^{\hat{\beta}'_1} \quad (\text{F-C1.1})$$

2 cases arise

(a) ${}^s v_1 = \text{inl}({}^s v_{i1})$ and ${}^t v_i = \text{inl}({}^t v_{i1})$:

Also from Lemma 1.82 and Definition 1.80 we know that

$$({}^s \theta'_1, n - j, \delta^s \cup \{x \mapsto {}^s v_1\}, \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [(\Gamma, \{x \mapsto {}^s v_1\})]_{V'}^{\hat{\beta}'_1}$$

IH2:

$$({}^s \theta'_1, n - j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [\tau]_E^{\hat{\beta}'_1}$$

This means from Definition 1.77 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \stackrel{\hat{\beta}'_1}{\triangleright} {}^s \theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with H'_{s1}, H'_{t1} and since we know that $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists k < i - j < n - j$ s.t. $(H'_{s1}, e_{s1}) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2} \quad (\text{F-C2}) \end{aligned}$$

Let $\tau = A^{\ell_i}$ and since we know that $\tau \searrow \ell$ therefore we have $\ell \sqsubseteq \ell_i$

Since we have $({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2}$

Therefore from Definition 1.76 we have

$$({}^s \theta'_2, n - j - k, {}^s v_2, \text{Lb}({}^t v_{2i})) \in [A^{\ell_i}]_{V'}^{\hat{\beta}'_2} \quad (\text{F-C2.1})$$

In order to prove (F-C0) we choose H'_t as H'_{t2} and ${}^t v$ as ${}^t v_2 = \text{Lb}({}^t v_{2i})$

And we need to prove:

i. $(H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$:

From Lemma 1.84 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$:
From (F-C1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t11} = {}^t v_1$
- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})) [{}^t v_1 / a] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$:
From cg-bind it suffices to prove that

- $(H'_{t1}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$:
Since from (F-C1.1) we know that ${}^t v_1 = \text{Lb}({}^t v_i)$ therefore from cg-unlabel we know that
 $H'_{t21} = H'_{t1}$ and ${}^t v_{t21} = {}^t v_i$
- $(\text{case}(b, x.e_{t1}, y.e_{t2})[{}^t v_i/b] \delta^t) \Downarrow {}^t v_{t22}$:
Since we know that in this case ${}^t v_i = \text{inl}({}^t v_{i1})$
Therefore from cg-case we know that ${}^t v_{t22} = e_{t1}[{}^t v_{i1}/x] \delta^t$
- $(H'_{t1}, e_{t1}[{}^t v_{i1}/x] \delta^t) \Downarrow (H'_{t2}, \text{Lb}{}^t v_{2i})$:
From (F-C2) and (F-C2.1) we get the desired

ii. $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$:

We choose ${}^s \theta'$ as ${}^s \theta'_2$ and $\hat{\beta}'$ as $\hat{\beta}'_2$. Since from fg-case we know that $i = j + k + 1$ and $H'_s = H'_{s2}$ therefore from (F-C2) and Lemma 1.83 we get

$$(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$$

Now we need to prove $({}^s \theta'_2, n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'_2}$
Since ${}^s v = {}^s v_2$ and ${}^t v = {}^t v_2$ and since from (F-C2) we know that

$$({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}'_2}$$

Therefore from Lemma 1.81 and Definition 1.76 we get

$$({}^s \theta'_2, n - i, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}'_2}$$

(b) ${}^s v_1 = \text{inr}({}^s v_{i1})$ and ${}^t v_1 = \text{inr}({}^t v_{i1})$:

Symmetric reasoning as in the previous case

10. FC-ref:

$$\frac{\Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \quad \mathcal{L} \vdash \tau \searrow pc}{\Gamma \vdash_{pc} \text{new } (e_s) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b)))} \text{ref}$$

Also given is: $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{new } (e_s) \delta^s, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b)) \delta^t) \delta^t) \in [(\text{ref } \tau)^\perp]_E^{\hat{\beta}}$

This means from Definition 1.77 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$. Also given some $i < n, {}^s v$ s.t $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$.

And we are required to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp]_V^{\hat{\beta}'} \quad (\text{F-R0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau]_E^{\hat{\beta}}$$

This means from Definition 1.77 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \lfloor \tau \rfloor_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore we know that $\exists j < n$ s.t $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$.

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \lfloor \tau \rfloor_V^{\hat{\beta}'_1} \quad (\text{F-R1}) \end{aligned}$$

In order to prove (F-R0) we choose H'_t as $H'_1 \cup \{a_t \mapsto {}^t v_1\}$, ${}^t v = \text{Lb}(a_t)$, ${}^s \theta'$ as ${}^s \theta'_1 \cup \{a_s \mapsto \tau\}$ and $\hat{\beta}'$ as $\hat{\beta}'_1 \cup \{(a_s, a_t)\}$

And we need to prove:

$$(a) (H_t, \text{bind}(e_t, a. \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1}):$

From (F-R1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t1} = {}^t v_1$

- $(H'_{t1}, \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_t, {}^t v):$

From cg-bind it suffices to prove that

- i. $(H'_{t1}, \text{new } (a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_t, {}^t v_{t2}):$

From cg-new we know that $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$ and ${}^t v = a_t$

- ii. $(H'_1 \cup \{a_t \mapsto {}^t v_1\}, \text{ret}(\text{Lb } b)) [{}^t v_1/a] [a_t/b] \delta^t) \Downarrow^f (H'_t, {}^t v_{t2}):$

From cg-ret we know that $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$ and ${}^t v_{t2} = \text{Lb}(a_t)$

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V^{\hat{\beta}'_1}:$$

From (F-R1) we know that $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$ and since $H'_s = H'_{s1} \cup \{a_s \mapsto {}^s v_1\}$, $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$, ${}^s \theta' = {}^s \theta'_1 \cup \{a_s \mapsto \tau\}$

Therefore from Definition 1.78 and Lemma 1.83 we get $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta'$

To prove: $({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V^{\hat{\beta}'_1}$

Since we know that ${}^s v = a_s$ and ${}^t v = \text{Lb } a_t$ therefore we need to prove

$$({}^s \theta', n - i, a_s, \text{Lb}(a_t)) \in \lfloor (\text{ref } \tau)^\perp \rfloor_V^{\hat{\beta}'_1}$$

From Definition 1.76 it suffices to prove that

$$({}^s \theta', n - i, a_s, a_t) \in \lfloor (\text{ref } \tau) \rfloor_V^{\hat{\beta}'_1}$$

Again from Definition 1.76 it suffices to prove that

$${}^s \theta'(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1$$

We get this by construction

11. FC-deref:

$$\frac{\Gamma \vdash_{pc} e_s : (\text{ref } \tau)^\ell \rightsquigarrow e_t \quad \mathcal{L} \vdash \tau <: \tau' \quad \mathcal{L} \vdash \tau' \searrow \ell}{\Gamma \vdash_{pc} !e_s : \tau' \rightsquigarrow \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{deref}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_{\hat{\beta}}^{\hat{\beta}}$

To prove: $({}^s\theta, n, !e \delta^s, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))\delta^t) \in [\tau']_{\hat{\beta}}^{\hat{\beta}}$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, !e_s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau']_{\hat{\beta}'}^{\hat{\beta}'} \end{aligned}$$

This means that we are given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau']_{\hat{\beta}'}^{\hat{\beta}'} \quad (\text{F-DR0}) \end{aligned}$$

IIH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}}^{\hat{\beta}}$$

This means from Definition 1.77 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < n$ s.t $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-DR1}) \end{aligned}$$

From (F-DR1) we have $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1}$

From Definition 1.76 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau)^\ell]_{\hat{\beta}'_1}^{\hat{\beta}'_1} \quad (\text{F-DR1.1})$$

From Definition 1.76 we know that ${}^s v_1 = a_s$ and ${}^t v_i = a_t$

$${}^s\theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \quad (\text{F-DR1.2})$$

Since we are given that $(n, H_s, H_t) \hat{\triangleright}^{\hat{\beta}} s\theta$ therefore from Definition 1.78 we know that

$$({}^s\theta, n-1, H_s(a_s), H_t(a_t)) \in [{}^s\theta(a_s)]_V^{\hat{\beta}}$$

which means we have

$$({}^s\theta, n-1, H_s(a_s), H_t(a_t)) \in [\tau]_V^{\hat{\beta}}$$

From Lemma 1.86 we know that

$$({}^s\theta, n-1, H_s(a_s), H_t(a_t)) \in [\tau']_V^{\hat{\beta}}$$

Let $\tau' = A^{\ell_i}$ since $\tau' \searrow \ell$ therefore $\ell \sqsubseteq \ell_i$

Let $v_g = H_t(a_t)$ therefore from Definition 1.76 we have

$$({}^s\theta, n-1, H_s(a_s), \text{Lb } v_{gi}) \in [\tau']_V^{\hat{\beta}} \quad (\text{F-DR1.3})$$

In order to prove (F-DR0) we choose H'_t as H'_{t1} and ${}^t v$ as $H'_{t1}(a_t) = v_g = \text{Lb } v_{gi}$

$$(a) \ (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi}):$$

From Lemma 1.84 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi})$$

From cg-bind it suffices to prove

$$i. \ (H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1}):$$

From (F-DR1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t1} = {}^t v_1$

$$ii. \ (H'_{t1}, \text{bind}(\text{unlabel } a, b.!b)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi}):$$

From cg-bind it suffices to prove that

$$A. \ (H'_{t1}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$$

From (F-DR1.1) we know that ${}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel we know that $H'_{t21} = H'_{t1}$ and ${}^t v_{t21} = {}^t v_i$

$$B. \ (H'_{t1}, (!b)[{}^t v_1/a][{}^t v_i/b] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb } v_{gi}):$$

We get the desired from CG-deref, (F-DR1.2) and (F-DR1.3)

$$(b) \ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n-i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n-i, {}^s v, \text{Lb } v_{gi}) \in [\tau']_V^{\hat{\beta}'}:$$

We choose ${}^s\theta'$ as ${}^s\theta'_1$ and $\hat{\beta}'$ as $\hat{\beta}'_1$

Therefore from (F-DR1) we get $(n-j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$ and since $i = j+1$ therefore from Lemma 1.83 we get $(n-i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$

Since from (F-DR1.2) we know that $(a_s, a_t) \in \hat{\beta}'_1$ and ${}^s\theta'_1(a_s) = \tau$. Also from (F-

DR1) we have $(n-j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$. Therefore from Definition 1.77 we have $(n-j-1, H'_{s1}(a_s), H'_{t1}(a_t)) \in [{}^s\theta'_1(a_s)]_V^{\hat{\beta}'_1}$

Since $i = j+1$, ${}^s\theta'_1(a_s) = \tau$, $H'_{s1}(a_s) = {}^s v$ and $H'_{t1}(a_t) = {}^t v_g = \text{Lb } v_{gi}$

Therefore we get $({}^s\theta', n-i, {}^s v, {}^t v) \in [\tau']_V^{\hat{\beta}'}$

from (F-DR1.3) and Lemma 1.81

12. FC-assign:

$$\frac{\Gamma \vdash_{pc} e_{s1} : (\text{ref } \tau)^\ell \rightsquigarrow e_{t1} \quad \Gamma \vdash_{pc} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \mathcal{L} \vdash \tau \searrow (pc \sqcup \ell)}{\Gamma \vdash_{pc} e_{s1} := e_{s2} : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{assign}$$

Also given is: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \in [\text{unit}]_E^{\hat{\beta}}$$

This means from Definition 1.77 we are required to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some H_s, H_t s.t. $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t. $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}'} \quad (\text{F-AN0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{ref } \tau)^\ell]_E^{\hat{\beta}}$$

This means from Definition 1.77 we are required to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\gamma, \hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < n$ s.t. $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_V^{\hat{\beta}'_1} \quad (\text{F-AN1}) \end{aligned}$$

Since from (F-AN1) we know that $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell]_V^{\hat{\beta}'_1}$ therefore from Definition 1.76 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau)^\ell]_V^{\hat{\beta}'_1} \quad (\text{F-AN1.1})$$

From Definition 1.76 this further means that

$${}^s\theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \text{ where } {}^s v_1 = a_s \text{ and } {}^t v_1 = a_t \quad (\text{F-AN1.2})$$

IH2:

$$({}^s\theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau]_E^{\hat{\beta}'_1}$$

This means from Definition 1.77 we are required to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with H'_{s1}, H'_{t1} and since we know that $(H_s, (e_{s2} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists k < n - j$ s.t. $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_V^{\hat{\beta}'_2} \quad (\text{F-AN2}) \end{aligned}$$

In order to prove (F-AN0) we choose H'_t as $H'_{t2}[a_t \mapsto {}^s v_2]$, ${}^t v$ as ()

We need to prove

$$(a) (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

$$- (H_t, \text{toLabeled}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) \delta^t) \Downarrow^f (H'_T, {}^t v_T):$$

From cg-toLabeled it suffices to prove that

$$(H_t, \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))) \delta^t) \Downarrow^f (H'_T, {}^t v_{Ti})$$

where ${}^t v_T = \text{Lb} {}^t v_{Ti}$

From cg-bind it further suffices to prove that:

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11}):$

From (F-AN1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t11} = {}^t v_1$

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12}):$

From cg-bind it suffices to prove

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t13}, {}^t v_{t13}):$

From (F-AN2) we know that $H'_{t13} = H'_{t2}$ and ${}^t v_{t13} = {}^t v_2$

- $(H'_{t1}, \text{bind}(\text{unlabel } a, c.c := b) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_t, {}^t v_{t12}):$

From cg-bind it suffices to prove that

- * $(H'_{t1}, \text{unlabel } a [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$

From (F-AN1.1) we know that

$${}^t v_1 = \text{Lb} ({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau)]_V^{\hat{\beta}'_1}$$

Therefore from cg-unlabel we know that $H'_{t21} = H'_{t1}$ and ${}^t v_{t21} = {}^t v_i = a_t$

- * $(H'_{t1}, (c := b) [{}^t v_1/a] [{}^t v_2/b] [{}^t v_i/c] \delta^t) \Downarrow^f (H'_t, {}^t v):$

From cg-assign we know that $H'_t = H'_{t1}[a_t \mapsto {}^t v_2]$ and ${}^t v_{t12} = ()$

Since ${}^t v_{t12} = {}^t v_{T_i} = ()$ therefore ${}^t v_T = \text{Lb}()$

- $(H'_T, \text{ret}()[{}^t v_T/d]) \delta^t \Downarrow^f (H'_t, ()):$

From cg-ret and cg-val

(b) $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau]_{V'}^{\hat{\beta}'}$:

We choose ${}^s \theta'$ as ${}^s \theta'_2$ and $\hat{\beta}'$ as $\hat{\beta}'_2$

In order to prove $(n - i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$ it suffices to prove

- $\text{dom}({}^s \theta'_2) \subseteq \text{dom}(H'_s)$:

Since from (F-AN2) we know that $(n - j - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$ therefore from Definition 1.78 we get $\text{dom}({}^s \theta'_2) \subseteq \text{dom}(H'_s)$

- $\hat{\beta}'_2 \subseteq (\text{dom}({}^s \theta'_2) \times \text{dom}(H'_t))$:

Since from (F-AN2) we know that $(n - j - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$ therefore from Definition 1.78 we get

$$\hat{\beta}'_2 \subseteq (\text{dom}({}^s \theta'_2) \times \text{dom}(H'_t))$$

- $\forall (a_1, a_2) \in \hat{\beta}'_2. ({}^s \theta'_2, n - i - 1, H'_s(a_1), H'_t(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'_2}$:

$\forall (a_1, a_2) \in \hat{\beta}'_2.$

- $a_1 = a_s$ and $a_1 = a_t$:

Since from (F-AN2) we know that $({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2}$

Also from (F-AN1.2) and Definition 1.74 we know that ${}^s \theta'_2(a_1) = \tau$

Therefore from Lemma 1.81 we get

$$({}^s \theta'_2, n - i - 1, {}^s v_2, {}^t v_2) \in [\tau]_{V'}^{\hat{\beta}'_2}$$

- $a_1 \neq a_s$ and $a_1 \neq a_t$:

From (F-AN2) since we know that $(n - j - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'_2} {}^s \theta'_2$ therefore from Definition 1.78 we get

$$({}^s \theta'_2, n - j - k - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'_2}$$

Since $i = j + k + 1$ therefore from Lemma 1.81 we get

$$({}^s \theta'_2, n - i - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s \theta'_2(a_1)]_{V'}^{\hat{\beta}'_2}$$

- $a_1 = a_s$ and $a_1 \neq a_t$:

This case cannot arise

- $a_1 \neq a_s$ and $a_1 = a_t$:

This case cannot arise

And in order to prove $({}^s \theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'}$

Since we know that ${}^s v = ()$ and ${}^t v = ()$ therefore from Definition 1.76 we get $({}^s \theta', n - i, {}^s v, {}^t v) \in [\text{unit}]_{V'}^{\hat{\beta}'}$

□

Lemma 1.86 (Subtyping lemma). *The following holds:*

$\forall \mathcal{L}, \hat{\beta}.$

1. $\forall A, A'.$

$$(a) \mathcal{L} \vdash A <: A' \implies [(A)]_{V'}^{\hat{\beta}} \subseteq [(A')]_{V'}^{\hat{\beta}}$$

2. $\forall \tau, \tau'$.

$$(a) \mathcal{L} \vdash \tau <: \tau' \implies \llbracket (\tau) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau') \rrbracket_V^{\hat{\beta}}$$

$$(b) \mathcal{L} \vdash \tau <: \tau' \implies \llbracket (\tau) \rrbracket_E^{\hat{\beta}} \subseteq \llbracket (\tau') \rrbracket_E^{\hat{\beta}}$$

Proof. Proof by simultaneous induction on $A <: A'$ and $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of $A <: A'$ in the last step:

1. FGsub-arrow:

Given:

$$\frac{\mathcal{L} \vdash \tau'_1 <: \tau_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2 \quad \mathcal{L} \vdash \ell'_e \sqsubseteq \ell_e}{\mathcal{L} \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove: $\llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V^{\hat{\beta}}$

IH1: $\llbracket (\tau'_1) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau_1) \rrbracket_V^{\hat{\beta}}$ (Statement 2(a))

It suffices to prove: $\forall ({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V^{\hat{\beta}}$.

$({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V^{\hat{\beta}}$

This means that given some ${}^s\theta, m$ and $\lambda x.e_s, (\lambda x.e_t)$ s.t

$({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2)) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 1.76 we are given:

$$\begin{aligned} \forall {}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'_1} \implies \\ ({}^s\theta'_1, j, e_s[{}^s v_1/x] \delta^s, e_t[{}^t v_1/x] \delta^t) \in \llbracket \tau_2 \rrbracket_E^{\hat{\beta}'_1} \quad (\text{S-L0}) \end{aligned}$$

And it suffices to prove: $({}^s\theta, m, \lambda x.e_s, (\lambda x.e_t)) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2)) \rrbracket_V^{\hat{\beta}}$

Again from Definition 1.76, it suffices to prove:

$$\begin{aligned} \forall {}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}'_2} \implies \\ ({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in \llbracket \tau'_2 \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-L1}) \end{aligned}$$

This means that given ${}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2$ s.t $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}'_2}$

And we need to prove

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in \llbracket \tau'_2 \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-L2})$$

Instantiating (S-L0) with ${}^s\theta'_2, {}^s v_2, {}^t v_2, k, \hat{\beta}'_2$. Since we have $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau'_1 \rrbracket_V^{\hat{\beta}'_2}$ therefore from IH1 we also have

$$({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'_2}$$

Therefore we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau_2]_E^{\hat{\beta}'_2}$$

$$\text{IH2: } [(\tau_2)]_E^{\hat{\beta}} \subseteq [(\tau'_2)]_E^{\hat{\beta}} \text{ (Statement 2(b))}$$

Finally using IH2 we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2]_E^{\hat{\beta}'_2}$$

2. FGsub-prod:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

$$\text{To prove: } [((\tau_1 \times \tau_2))]_V^{\hat{\beta}} \subseteq [((\tau'_1 \times \tau'_2))]_V^{\hat{\beta}}$$

$$\text{IH1: } [(\tau_1)]_V^{\hat{\beta}} \subseteq [(\tau'_1)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

$$\text{IH2: } [(\tau_2)]_V^{\hat{\beta}} \subseteq [(\tau'_2)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

It suffices to prove:

$$\forall ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2))]_V^{\hat{\beta}}. ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2))]_V^{\hat{\beta}}$$

This means that given some ${}^s\theta, n$ and ${}^s v_1, {}^s v_2, {}^t v_1, {}^t v_2$ s.t

$$({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau_1 \times \tau_2))]_V^{\hat{\beta}}$$

Therefore from Definition 1.76 we are given:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}} \quad (\text{S-P0})$$

$$\text{And it suffices to prove: } ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in [((\tau'_1 \times \tau'_2))]_V^{\hat{\beta}}$$

Again from Definition 1.76, it suffices to prove:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau'_2]_V^{\hat{\beta}} \quad (\text{S-P1})$$

Since from (S-P0) we know that $({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$ therefore from IH1 we have

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}}$$

Similarly since we have $({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}$ from (S-P0) therefore from IH2 we have

$$({}^s\theta, m, {}^s v_2, {}^t v_2) \in [\tau'_2]_V^{\hat{\beta}}$$

3. FGsub-sum:

Given:

$$\frac{\mathcal{L} \vdash \tau_1 <: \tau'_1 \quad \mathcal{L} \vdash \tau_2 <: \tau'_2}{\mathcal{L} \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

$$\text{To prove: } [((\tau_1 + \tau_2))]_V^{\hat{\beta}} \subseteq [((\tau'_1 + \tau'_2))]_V^{\hat{\beta}}$$

$$\text{IH1: } [(\tau_1)]_V^{\hat{\beta}} \subseteq [(\tau'_1)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

IH2: $[(\tau_2)]_V^{\hat{\beta}} \subseteq [(\tau'_2)]_V^{\hat{\beta}}$ (Statement 2(a))

It suffices to prove: $\forall ({}^s\theta, n, {}^sv, {}^tv) \in [((\tau_1 + \tau_2))]_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, {}^tv) \in [((\tau'_1 + \tau'_2))]_V^{\hat{\beta}}$

This means that given: $({}^s\theta, n, {}^sv, {}^tv) \in [((\tau_1 + \tau_2))]_V^{\hat{\beta}}$

And it suffices to prove: $({}^s\theta, n, {}^sv, {}^tv) \in [((\tau'_1 + \tau'_2))]_V^{\hat{\beta}}$

2 cases arise

(a) ${}^sv = \text{inl } {}^sv_i$ and ${}^tv = \text{inl } {}^tv_i$:

From Definition 1.76 we are given:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau_1]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s\theta, n, {}^sv_i, {}^tv_i) \in [\tau'_1]_V^{\hat{\beta}}$$

From (S-S0) and IH1 get this

(b) ${}^sv = \text{inr } {}^sv_i$ and ${}^tv = \text{inr } {}^tv_i$:

Symmetric reasoning as in the previous case

4. FGsub-ref:

Given:

$$\frac{}{\mathcal{L} \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove: $[((\text{ref } \tau))]_V^{\hat{\beta}} \subseteq [((\text{ref } \tau))]_V^{\hat{\beta}}$

It suffices to prove: $\forall ({}^s\theta, n, a_s, a_t) \in [((\text{ref } \tau))]_V^{\hat{\beta}}. ({}^s\theta, n, a_s, a_t) \in [((\text{ref } \tau))]_V^{\hat{\beta}}$

We get this directly from Definition 1.76

5. FGsub-base:

Given:

$$\frac{}{\mathcal{L} \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove: $[((\mathbf{b}))]_V^{\hat{\beta}} \subseteq [((\mathbf{b}))]_V^{\hat{\beta}}$

Directly from Definition 1.76

6. FGsub-unit:

Given:

$$\frac{}{\mathcal{L} \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove: $[((\text{unit}))]_V^{\hat{\beta}} \subseteq [((\text{unit}))]_V^{\hat{\beta}}$

Directly from Definition 1.76

Proof of statement 2(a)

Given:

$$\frac{\mathcal{L} \vdash \ell' \sqsubseteq \ell'' \quad \mathcal{L} \vdash A <: A'}{\mathcal{L} \vdash A^{\ell'} <: A^{\ell''}} \text{FGsub-label}$$

To prove: $[(A^{\ell'})]_V^{\hat{\beta}} \subseteq [(A^{\ell''})]_V^{\hat{\beta}}$

This means from Definition 1.76 we need to prove

$$\forall ({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in [A^{\ell'}]_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in [A^{\ell''}]_V^{\hat{\beta}}$$

This means that given $({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in [A^{\ell'}]_V^{\hat{\beta}}$

From Definition 1.76 it further means that we are given

$$({}^s\theta, n, {}^s v, {}^t v_i) \in [A]_V^{\hat{\beta}} \quad (\text{S-LB0})$$

And we need to prove

$$({}^s\theta, n, {}^s v, \text{Lb}({}^t v_i)) \in [A^{\ell''}]_V^{\hat{\beta}}$$

Again from Definition 1.76 it suffices to prove that

$$({}^s\theta, n, {}^s v, {}^t v_i) \in [A']_V^{\hat{\beta}}$$

Since $\ell' \sqsubseteq \ell''$ and $A' <: A''$ therefore from IH (Statement 1(a)) and (S-LB0) we get the desired

Proof of statement 2(b)

Given: $\mathcal{L} \vdash \tau <: \tau'$

To prove: $[(\tau)]_E^{\hat{\beta}} \subseteq [(\tau')]_E^{\hat{\beta}}$

This means we need to prove that

$$\forall (\theta, n, e_s, e_t) \in [(\tau)]_E^{\hat{\beta}}. (\theta, n, e_s, e_t) \in [(\tau')]_E^{\hat{\beta}}$$

This means given $(\theta, n, e_s, e_t) \in [(\tau)]_E^{\hat{\beta}}$

This means from Definition 1.77 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, e_s) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'_1} \quad (\text{S-E0})$$

And it suffices to prove that $({}^s\theta, n, e_s, e_t) \in [(\tau')]_E^{\hat{\beta}}$

Again from Definition 1.77 it means we need to prove

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau']_V^{\hat{\beta}'_1}$$

This means that given some H_{s1}, H_{t1} s.t $(n, H_{s1}, H_{t1}) \triangleright^{\ell_2, \hat{\beta}} {}^s\theta$. Also given some $j < n, {}^s v_1$ s.t $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

And we need to prove

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau']_V^{\hat{\beta}'_1} \quad (\text{S-E1})$$

Instantiating (S-E0) with H_{s1}, H_{t1} and with $j, {}^s v_1$. Then we get

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_t) & \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau]_V^{\hat{\beta}'_1} \end{aligned}$$

Since we have $\tau <: \tau'$. Therefore from IH (Statement 2(a)) we get

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) & \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau']_V^{\hat{\beta}'_1} \end{aligned}$$

□

Theorem 1.87 (Deriving FG NI via compilation). $\forall e_s, {}^s v_1, {}^s v_2, n_1, n_2, H'_{s1}, H'_{s2}, \perp$.

$$\begin{aligned} & \text{Let } \mathbf{bool} = (\mathbf{unit} + \mathbf{unit}) \\ & x : \mathbf{bool}^\top \vdash_\perp e_s : \mathbf{bool}^\perp \wedge \\ & \emptyset \vdash_\perp {}^s v_1 : \mathbf{bool}^\top \wedge \emptyset \vdash_\perp {}^s v_2 : \mathbf{bool}^\top \wedge \\ & (\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1} (H'_{s1}, {}^s v'_1) \wedge \\ & (\emptyset, e_s[{}^s v_2/x]) \Downarrow_{n_2} (H'_{s2}, {}^s v'_2) \wedge \\ & \implies \\ & {}^s v'_1 = {}^s v'_2 \end{aligned}$$

Proof. From the FG to CG translation we know that $\exists e_t$ s.t

$$\begin{aligned} & x : \mathbf{bool}^\top \vdash e_s : \mathbf{bool}^\perp \rightsquigarrow e_t \\ & \text{Similarly we also know that } \exists {}^t v_1, {}^t v_2 \text{ s.t} \\ & \emptyset \vdash {}^s v_1 : \mathbf{bool}^\top \rightsquigarrow {}^t v_1 \text{ and } \emptyset \vdash {}^s v_2 : \mathbf{bool}^\top \rightsquigarrow {}^t v_2 \quad (\text{NI-0}) \end{aligned}$$

From type preservation theorem (choosing $\alpha = \bar{\beta} = \perp$) we know that

$$\begin{aligned} & x : \text{Labeled } \top \mathbf{bool} \vdash e_t : \mathbb{C} \perp \perp \text{Labeled } \perp \mathbf{bool} \\ & \emptyset \vdash {}^t v_1 : \mathbb{C} \perp \perp \text{Labeled } \top \mathbf{bool} \\ & \emptyset \vdash {}^t v_2 : \mathbb{C} \perp \perp \text{Labeled } \top \mathbf{bool} \quad (\text{NI-1}) \end{aligned}$$

Since we have $\emptyset \vdash {}^s v_1 : \mathbf{bool}^\top \rightsquigarrow {}^t v_1$

And since ${}^s v_1$ and ${}^t v_1$ are closed terms (from given and NI-1)

Therefore from Theorem 1.85 we have (we choose $n > n_1$ and $n > n_2$)

$$(\emptyset, n, {}^s v_1, {}^t v_1) \in [\mathbf{bool}^\top]_E^\emptyset \quad (\text{NI-2})$$

Therefore from Definition 1.77 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_1) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ (n - i, H'_s, H'_t) & \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{11}) \in [\mathbf{bool}^\top]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with \emptyset, \emptyset and from fg-val we know that $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$. Therefore we have

$$\begin{aligned} & \exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ (n, H'_s, H'_t) & \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{11}) \in [\mathbf{bool}^\top]_V^{\hat{\beta}'_1} \quad (\text{NI-2.1}) \end{aligned}$$

From Definition 1.76 we know that

$${}^t v_{11} = \mathbf{Lb}({}^t v_{i11}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{i11}) \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}'_1}$$

Again from Definition 1.76 we know that

Either a) ${}^s v_1 = \mathbf{inl}()$ and ${}^t v_{i11} = \mathbf{inl}()$ or b) ${}^s v_1 = \mathbf{inr}()$ and ${}^t v_{i11} = \mathbf{inr}()$

But in either case we have that $\emptyset \vdash {}^t v_{i11} : (\mathbf{unit} + \mathbf{unit})$ (NI-2.2)

As a result we have $\emptyset \vdash {}^t v_{11} : \text{Labeled } \top (\mathbf{unit} + \mathbf{unit})$ (NI-2.3)

We give it typing derivation

$$\frac{\overline{\emptyset \vdash {}^t v_{i11} : (\text{unit} + \text{unit})} \text{ (NI-2.2)}}{\emptyset \vdash \text{Lb}({}^t v_{i11}) : \text{Labeled } \top \text{ (unit} + \text{unit)}}$$

From Definition 1.80 and (NI-2.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto {}^t v_{11})) \in \llbracket x \mapsto \text{bool}^\top \rrbracket_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 1.85 to get

$$(\emptyset, n, e_s[{}^s v_1/x], e_t[{}^t v_{11}/x]) \in \llbracket \text{bool}^\perp \rrbracket_E^{\hat{\beta}'} \quad (\text{NI-2.4})$$

From Definition 1.77 we get

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}'} \emptyset \wedge \forall i < n, {}^s v_1'' \cdot (H_s, e_s[{}^s v_1/x]) \Downarrow_i (H'_{s1}, {}^s v_1'') \implies \exists H'_{t1}, {}^t v_1'' \cdot (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v_1'') \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v_1'', {}^t v_1'') \in \llbracket \text{bool}^\perp \rrbracket_V^{\hat{\beta}''}$$

Instantiating with $\emptyset, \emptyset, n_1, {}^s v_1'$ we get

$$\exists H'_{t1}, {}^t v_1'' \cdot (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v_1'') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}'.$$

$$(n - n_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v_1', {}^t v_1'') \in \llbracket \text{bool}^\perp \rrbracket_V^{\hat{\beta}''} \quad (\text{NI-2.5})$$

Since we have $({}^s \theta', n - n_1, {}^s v_1', {}^t v_1'') \in \llbracket \text{bool}^\perp \rrbracket_V^{\hat{\beta}''}$ therefore from Definition 1.76 we have

$$\exists {}^t v_{i1}. {}^t v'' = \text{Lb}({}^t v_{i1}) \wedge ({}^s \theta', n - n_1, {}^s v_1', {}^t v_{i1}) \in \llbracket \text{bool} \rrbracket_V^{\hat{\beta}''}$$

Since $({}^s \theta', n - n_1, {}^s v_1', {}^t v_{i1}) \in \llbracket (\text{unit} + \text{unit}) \rrbracket_V^{\hat{\beta}''}$ therefore from Definition 1.76 two cases arise

- ${}^s v_1' = \text{inl } {}^s v_{i11}$ and ${}^t v_{i1} = \text{inl } {}^t v_{i11}$:

From Definition 1.76 we have

$$({}^s \theta', n - n_1, {}^s v_{i11}, {}^t v_{i11}) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}''}$$

which means we have ${}^s v_{i11} = {}^t v_{i11}$

- ${}^s v_1' = \text{inr } {}^s v_{i11}$ and ${}^t v_{i1} = \text{inr } {}^t v_{i11}$:

Symmetric reasoning as in the previous case

So no matter which case arise we have ${}^s v_1' = {}^t v_{i1}$

$$\text{Similarly with other substitution we have } (\emptyset, n, {}^s v_2, {}^t v_2) \in \llbracket \text{bool}^\top \rrbracket_E^\emptyset \quad (\text{NI-3})$$

Therefore from Definition 1.77 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^\emptyset \emptyset \wedge \forall i < n, {}^s v \cdot (H_s, {}^s v_2) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v_{22} \cdot (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{22}) \in \llbracket \text{bool}^\top \rrbracket_V^{\hat{\beta}'}$$

Instantiating with \emptyset, \emptyset and from fg-val we know that $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$. Therefore we have

$$\exists H'_t, {}^t v_{22} \cdot (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in \llbracket \text{bool}^\top \rrbracket_V^{\hat{\beta}'} \quad (\text{NI-3.1})$$

From Definition 1.76 we know that

$${}^t v_2 = \text{Lb}({}^t v_{22}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in \llbracket (\text{unit} + \text{unit}) \rrbracket_V^{\hat{\beta}'}$$

Again from Definition 1.76 we know that

Either a) ${}^s v_2 = \text{inl}()$ and ${}^t v_{i22} = \text{inl}()$ or b) ${}^s v_2 = \text{inr}()$ and ${}^t v_{i22} = \text{inr}()$
 But in either case we have that $\emptyset \vdash {}^t v_{i22} : (\text{unit} + \text{unit})$ (NI-3.2)

As a result we have $\emptyset \vdash {}^t v_{22} : \text{Labeled } \top (\text{unit} + \text{unit})$ (NI-3.3)

We give it typing derivation

$$\frac{\overline{\emptyset \vdash {}^t v_{i22} : (\text{unit} + \text{unit})} \text{ (NI-3.2)}}{\emptyset \vdash \text{Lb}({}^t v_{i22}) : \text{Labeled } \top (\text{unit} + \text{unit})}$$

From Definition 1.80 and (NI-3.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_2), (x \mapsto {}^t v_{22})) \in [x \mapsto \text{bool}^\top]_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 1.85 to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_{22}/x]) \in [\text{bool}^\perp]_E^{\hat{\beta}'}$$
 (NI-3.4)

From Definition 1.77 we get

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright \emptyset \wedge \forall i < n, {}^s v_2''.(H_s, e_s[{}^s v_2/x]) \Downarrow_i (H_{s2}', {}^s v_2'') \implies \\ & \exists H_{t2}', {}^t v_2''.(H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H_{t2}', {}^t v_2'') \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - i, H_{s2}', H_{t2}') \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v_2'', {}^t v_2'') \in [\text{bool}^\perp]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating with $\emptyset, \emptyset, n_2, {}^s v_2'$ we get

$$\begin{aligned} & \exists H_{t2}', {}^t v_2''.(H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H_{t2}', {}^t v_2'') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - n_1, H_s', H_{t2}') \triangleright^{\hat{\beta}''} {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v_2', {}^t v_2'') \in [\text{bool}^\perp]_V^{\hat{\beta}''} \end{aligned}$$
 (NI-3.5)

Since we have $({}^s \theta', n - n_2, {}^s v_2', {}^t v_2'') \in [\text{bool}^\perp]_V^{\hat{\beta}''}$ therefore from Definition 1.76 we have

$$\exists {}^t v_{i2}. {}^t v_2'' = \text{Lb}({}^t v_{i2}) \wedge ({}^s \theta', n - n_2, {}^s v_2', {}^t v_{i2}) \in [\text{bool}]_V^{\hat{\beta}''}$$

Since $({}^s \theta', n - n_2, {}^s v_2', {}^t v_{i2}) \in [(\text{unit} + \text{unit})]_V^{\hat{\beta}''}$ therefore from Definition 1.76 two cases arise

- ${}^s v_2' = \text{inl } {}^s v_{i22}$ and ${}^t v_{i2} = \text{inl } {}^t v_{i22}$:

From Definition 1.76 we have

$$({}^s \theta', n - n_2, {}^s v_{i22}, {}^t v_{i22}) \in [\text{unit}]_V^{\hat{\beta}''}$$

which means we have ${}^s v_{i22} = {}^t v_{i22}$

- ${}^s v_2' = \text{inr } {}^s v_{i22}$ and ${}^t v_{i2} = \text{inr } {}^t v_{i22}$:

Symmetric reasoning as in the previous case

So no matter which case arise we have ${}^s v_2' = {}^t v_{i2}$

We know that $\emptyset \vdash {}^t v_{11} : \text{Labeled } \top \text{ bool}$ (NI-2.3)

Also we have $\emptyset \vdash {}^t v_{22} : \text{Labeled } \top \text{ bool}$ (NI-3.3)

Let $e_T = \text{bind}(e_t, y.\text{unlabel}(y))$

We show that $x : \text{Labeled } \top \text{ bool} \vdash e_T : \mathbb{C} \perp \perp \text{ bool}$ by giving a typing derivation

P2:

$$\frac{\frac{}{x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash y : \text{Labeled } \perp \text{ bool}}{\text{CG-var}}}{x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash \text{unlabel}(y) : \mathbb{C} \perp \perp \text{ bool}} \text{CG-unlabel}$$

P1:

$$\frac{}{x : \text{Labeled } \top \text{ bool} \vdash e_t : \mathbb{C} \perp \perp \text{Labeled } \perp \text{ bool}} \text{From (NI-1)}$$

Main derivation:

$$\frac{P1 \quad P2}{x : \text{Labeled } \top \text{ bool} \vdash \text{bind}(e_t, y.\text{unlabel}(y)) : \mathbb{C} \perp \perp \text{ bool}}$$

Say $e_t[{}^t v_{11}/x]$ reduces in n_{t1} steps in (NI-2.5) and $e_t[{}^t v_{22}/x]$ reduces in n_{t2} steps in (NI-3.5)

We instantiate Theorem 1.70 with $e_T, {}^t v_{11}, {}^t v_{22}, {}^t v_{i1}, {}^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H'_{t1}, H'_{t2}$ and from (NI-2.5) and (NI-3.5) we have ${}^t v_{i1} = {}^t v_{i2}$ and thus ${}^s v'_1 = {}^s v'_2$

□

2 Part II: Alternate development with original HLIO in place of CG

2.1 Fine-grained IFC enforcement (FG)

2.1.1 FG type system

Syntax, types, constraints:

Expressions	e	$::=$	$x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, x.e) \mid \text{new } e \mid !e \mid e := e \mid \Lambda e \mid e [] \mid \nu e \mid e \bullet$
Labels	ℓ, pc	$::=$	$l \mid \alpha \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
(Labeled) Types	τ	$::=$	A^ℓ
Unlabeled types	A	$::=$	$b \mid \tau \xrightarrow{\ell_c} \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \tau \mid \text{unit} \mid \forall \alpha. (\ell_e, \tau) \mid c \xrightarrow{\ell_c} \tau$
Constraints	c	$::=$	$\ell \sqsubseteq \ell \mid (c, c)$

Lemma 2.1 (FG: Reflexivity of subtyping). *The following hold:*

1. For all $\Sigma, \Psi, \tau: \Sigma; \Psi \vdash \tau <: \tau$
2. For all $\Sigma, \Psi, A: \Sigma; \Psi \vdash A <: A$

Proof. Proof by simultaneous induction on τ and A .

Proof of statement (1)

Let $\tau = A^\ell$. Then, we have:

$$\frac{\frac{}{\Sigma; \Psi \vdash A <: A} \text{IH(2)} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell}{\Sigma; \Psi \vdash A^\ell <: A^\ell} \text{FGsub-label}$$

Proof of statement (2)

We proceed by cases on A .

1. $A = b$:

$$\frac{}{\Sigma; \Psi \vdash b <: b} \text{FGsub-base}$$

2. $A = \text{ref } \tau$:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

3. $A = \tau_1 \times \tau_2$:

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau_1 \times \tau_2}$$

Type system: $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau}$

$$\begin{array}{c}
\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var} \qquad \frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp} \text{FG-lam} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2} \text{FG-app} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp} \text{FG-prod} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1} \text{FG-fst} \quad \frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp} \text{FG-inl} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau} \text{FG-case} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc'} e : \tau' \quad \Sigma; \Psi \vdash pc \sqsubseteq pc' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau} \text{FG-sub} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^\perp} \text{FG-ref} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau'} \text{FG-deref} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}} \text{FG-assign} \\
\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} () : \text{unit}^\perp} \text{FG-unitI} \qquad \frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha. (\ell_e, \tau))^\perp} \text{FG-FI} \\
\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell'/\alpha]} \text{FG-FE} \\
\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp} \text{FG-CI} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau} \text{FG-CE}
\end{array}$$

Figure 8: Type system for FG

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash A <: A'}{\Sigma; \Psi \vdash A^\ell <: A'^{\ell'}} \text{FGsub-label} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref} \qquad \frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2} \text{FGsub-arrow} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2} \text{FGsub-forall} \\
\\
\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2} \text{FGsub-constraint}
\end{array}$$

Figure 9: FG subtyping

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash A \text{ WF} \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi \vdash A^\ell \text{ WF}} \text{FG-wff-label} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{b} \text{ WF}} \text{FG-wff-base} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{unit} \text{ WF}} \text{FG-wff-unit} \qquad \frac{\Sigma; \Psi \vdash \tau_1 \text{ WF} \quad \Sigma; \Psi \vdash \tau_2 \text{ WF} \quad \text{FV}(\ell_e) \in \Sigma}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 \text{ WF}} \text{FG-wff-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ WF} \quad \Sigma; \Psi \vdash \tau_2 \text{ WF}}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 \text{ WF}} \text{FG-wff-prod} \qquad \frac{\Sigma; \Psi \vdash \tau_1 \text{ WF} \quad \Sigma; \Psi \vdash \tau_2 \text{ WF}}{\Sigma; \Psi \vdash \tau_1 + \tau_2 \text{ WF}} \text{FG-wff-sum} \\
\\
\frac{\text{FV}(\tau) = \emptyset}{\Sigma; \Psi \vdash (\text{ref } \tau) \text{ WF}} \text{FG-wff-ref} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau \text{ WF} \quad \text{FV}(\ell_e) \in \Sigma}{\Sigma; \Psi \vdash (\forall \alpha. (\ell_e, \tau)) \text{ WF}} \text{FG-wff-forall} \\
\\
\frac{\Sigma; \Psi \vdash \tau \text{ WF} \quad \text{FV}(c) \in \Sigma \quad \text{FV}(\ell_e) \in \Sigma}{\Sigma; \Psi \vdash (c \xrightarrow{\ell_\xi} \tau) \text{ WF}} \text{FG-wff-constraint}
\end{array}$$

Figure 10: Well-formedness relation for FG

4. $A = \tau_1 + \tau_2$:

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau_1 + \tau_2}$$

5. $A = \tau_1 \xrightarrow{\ell_e} \tau_2$:

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(2) on } \tau_2 \quad \frac{}{\Sigma; \Psi \vdash \ell_e \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau_1 \xrightarrow{\ell_e} \tau_2}$$

6. $A = \text{unit}$:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}$$

7. $A = \forall \alpha. \tau_i$:

$$\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash \forall \alpha. \tau_i <: \forall \alpha. \tau_i}$$

8. $A = c \Rightarrow \tau_i$:

$$\frac{\frac{}{\Sigma; \Psi \vdash c \Rightarrow c} \quad \frac{}{\Sigma; \Psi, c \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash c \Rightarrow \tau <: c \Rightarrow \tau_i}$$

□

2.1.2 FG semantics

Judgement: $(H, e) \Downarrow_i (H', v)$

The semantics are described in Figure 11

2.1.3 Logical relation for FG

$W : ((Loc \mapsto Type) \times (Loc \mapsto Type) \times (Loc \leftrightarrow Loc))$

Definition 2.2 (FG: θ_2 extends θ_1). $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

Definition 2.3 (FG: W_2 extends W_1). $W_1 \sqsubseteq W_2 \triangleq$

1. $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2. $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

$$\begin{array}{c}
\frac{(H, e_1) \Downarrow_i (H', \lambda x. e_i) \quad (H', e_2) \Downarrow_j (H'', v_2) \quad (H'', e_i[v_2/x]) \Downarrow_k (H''', v_3)}{(H, e_1 e_2) \Downarrow_{i+j+k+1} (H''', v_3)} \text{fg-app} \\
\\
\frac{(H, e_1) \Downarrow_i (H', v_1) \quad (H', e_2) \Downarrow_j (H'', v_2)}{(H, (e_1, e_2)) \Downarrow_{i+j+1} (H'', (v_1, v_2))} \text{fg-prod} \quad \frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{fst}(e)) \Downarrow_{i+1} (H', v_1)} \text{fg-fst} \\
\\
\frac{(H, e) \Downarrow_i (H', (v_1, v_2))}{(H, \text{snd}(e)) \Downarrow_{i+1} (H', v_2)} \text{fg-snd} \quad \frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inl}(e)) \Downarrow_{i+1} (H', \text{inl}(v))} \text{fg-inl} \\
\\
\frac{(H, e) \Downarrow_i (H', v)}{(H, \text{inr}(e)) \Downarrow_{i+1} (H', \text{inr}(v))} \text{fg-inr} \quad \frac{(H, e) \Downarrow_i (H', \text{inl } v) \quad (H', e_1[v/x]) \Downarrow_j (H'', v_1)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_1)} \text{fg-case1} \\
\\
\frac{(H, e) \Downarrow_i (H', \text{inr } v) \quad (H', e_2[v/x]) \Downarrow_j (H'', v_2)}{(H, \text{case}(e, x.e_1, y.e_2)) \Downarrow_{i+j+1} (H'', v_2)} \text{fg-case2} \\
\\
\frac{(H, e) \Downarrow_i (H', \Lambda e_i) \quad (H', e_i) \Downarrow_j (H'', v)}{(H, e[]) \Downarrow_{i+j+1} (H'', v)} \text{fg-FE} \\
\\
\frac{(H, e) \Downarrow_i (H', \nu e_i) \quad (H', e_i) \Downarrow_j (H'', v)}{(H, e\bullet) \Downarrow_{i+j+1} (H'', v)} \text{fg-CE} \\
\\
\frac{(H, e) \Downarrow_i (H', v) \quad a \notin \text{dom}(H)}{(H, \text{new } (e)) \Downarrow_{i+1} (H'[a \mapsto v], a)} \text{fg-ref} \quad \frac{(H, e) \Downarrow_i (H', a)}{(H, !e) \Downarrow_{i+1} (H', H(a))} \text{fg-deref} \\
\\
\frac{(H, e_1) \Downarrow_i (H', a) \quad (H', e_2) \Downarrow_j (H'', v)}{(H, e_1 := e_2) \Downarrow_{i+j+1} (H''[a \mapsto v], ())} \text{fg-assign} \quad \frac{e \in \{x, \lambda y. -, \Lambda -, \nu -\}}{(H, e) \Downarrow_0 (H, e)} \text{fg-val}
\end{array}$$

Figure 11: FG semantics

Definition 2.4 (FG: Binary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^A &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V^A &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{b} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^A &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A\} \\
[\tau_1 + \tau_2]_V^A &\triangleq \{(W, n, \text{inl } v, \text{inl } v') \mid (W, n, v, v') \in [\tau_1]_V^A\} \cup \\
&\quad \{(W, n, \text{inr } v, \text{inr } v') \mid (W, n, v, v') \in [\tau_2]_V^A\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\
&\quad ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e}) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j, v_c. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})\} \\
[\forall \alpha. (\ell_e, \tau)]_V^A &\triangleq \{(W, n, \Lambda e_1, \Lambda e_2) \mid \\
&\quad \forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. \\
&\quad ((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j, \ell'' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j, \ell'' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})\} \\
[c \xrightarrow{\ell_e} \tau]_V^A &\triangleq \{(W, n, \nu e_1, \nu e_2) \mid \\
&\quad \forall W' \sqsupseteq W, n' < n. \\
&\quad \mathcal{L} \models c \implies (W', n', e_1, e_2) \in [\tau]_E^A \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E^{\ell_e} \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E^{\ell_e}\} \\
[\text{ref } \tau]_V^A &\triangleq \{(W, n, a_1, a_2) \mid \\
&\quad (a_1, a_2) \in W. \hat{\beta} \wedge W. \theta_1(a_1) = W. \theta_2(a_2) = \tau\}
\end{aligned}$$

$$[\mathbf{A}^{\ell'}]_V^A \triangleq \begin{cases} \{(W, n, v_1, v_2) \mid (W, n, v_1, v_2) \in [\mathbf{A}]_V^A\} & \ell' \sqsubseteq \mathbf{A} \\ \{(W, n, v_1, v_2) \mid \forall i \in \{1, 2\}. \forall m. (W(n). \theta_i, m, v_i) \in [\mathbf{A}]_V\} & \ell' \not\sqsubseteq \mathbf{A} \end{cases}$$

Definition 2.5 (FG: Binary expression relation).

$$\begin{aligned}
[\tau]_E^A &\triangleq \{(W, n, e_1, e_2) \mid \\
&\quad \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge \\
&\quad (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\
&\quad \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_V^A\}
\end{aligned}$$

Definition 2.6 (FG: Unary value relation).

$$\begin{aligned}
[\mathbf{b}]_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in [\tau_1]_V \wedge (\theta, m, v_2) \in [\tau_2]_V\} \\
[\tau_1 + \tau_2]_V &\triangleq \{(\theta, m, \text{inl } v) \mid (\theta, m, v) \in [\tau_1]_V\} \cup \{(\theta, m, \text{inr } v) \mid (\theta, m, v) \in [\tau_2]_V\} \\
[\tau_1 \xrightarrow{\ell_e} \tau_2]_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies \\
&\quad (\theta', j, e[v/x]) \in [\tau_2]_E^{\ell_e}\} \\
[\forall \alpha. (\ell_e, \tau)]_V &\triangleq \{(\theta, m, \Lambda e) \mid \forall \theta'. \theta \sqsubseteq \theta'. \forall m' < m. \forall \ell' \in \mathcal{L}. (\theta', m', e) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}\} \\
[c \xrightarrow{\ell_e} \tau]_V &\triangleq \{(\theta, m, \nu e) \mid \forall \theta'. \theta \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e) \in [\tau]_E^{\ell_e}\} \\
[\text{ref } \tau]_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \tau\}
\end{aligned}$$

$$[A^{\ell'}]_V \triangleq [A]_V$$

Definition 2.7 (FG: Unary expression relation).

$$\begin{aligned} [\tau]_E^{pc} \triangleq & \{(\theta, n, e) \mid \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)\} \end{aligned}$$

Definition 2.8 (FG: Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in [\theta(a)]_V$$

Definition 2.9 (FG: Binary heap well formedness).

$$\begin{aligned} (n, H_1, H_2) \hat{\triangleright}^A W \triangleq & \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ & (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ & \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ & (W, n - 1, H_1(a_1), H_2(a_2)) \in [W.\theta_1(a_1)]_V^A \wedge \\ & \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in [W.\theta_i(a_i)]_V \end{aligned}$$

Definition 2.10 (FG: Label substitution). $\sigma : Lvar \mapsto Label$

Definition 2.11 (FG: Value substitution to value pairs). $\gamma : Var \mapsto (Val, Val)$

Definition 2.12 (FG: Value substitution to values). $\delta : Var \mapsto Val$

Definition 2.13 (FG: Unary interpretation of Γ).

$$[\Gamma]_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V\}$$

Definition 2.14 (FG: Binary interpretation of Γ).

$$[\Gamma]_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A\}$$

2.1.4 Soundness proof for FG

Lemma 2.15 (FG: Binary value relation subsumes unary value relation). $\forall W, v_1, v_2, \mathcal{A}, n.$

The following holds:

1. $\forall \mathbf{A}.$

$$(W, n, v_1, v_2) \in [A]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [A]_V$$

2. $\forall \tau.$

$$(W, n, v_1, v_2) \in [\tau]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$$

Proof. Proof by simultaneous induction on \mathbf{A} and τ

Proof of statement (1)

We analyze the various cases of \mathbf{A} in the last step:

1. Case **b**:

From Definition 2.6

2. Case $\tau_1 \times \tau_2$:

Given: $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$ (P01)

and

$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$ (P02)

From Definition 2.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$ (P1)

IH1a: $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$ and

IH1b: $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a: $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$ and

IH2b: $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some m we need to prove

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly from (P02) we know that given some m we need to prove

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

We instantiate IH1a and IH2a with the given m from (P01) to get

$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$ and $(W.\theta_1, m, v_{i2}) \in [\tau_2]_V$

Then from Definition 2.6, we get

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly we instantiate IH1b and IH2b with the given m from (P02) to get

$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$ and $(W.\theta_2, m, v_{j2}) \in [\tau_2]_V$

Then from Definition 2.6, we get

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

3. Case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v_1 = \text{inl}(v_{i1})$ and $v_2 = \text{inl}(v_{j1})$

Given: $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$ (S01)

and

$\forall m. (W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$ (S02)

From Definition 2.4 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A$ (S0)

IH1: $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$ and

IH2: $\forall m_2. (W.\theta_2, m_2, v_{j1}) \in \lfloor \tau_1 \rfloor_V$

From (S01) we know that given some m and we are required to prove:

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

Also from (S02) we know that given some m and we are required to prove:

$$(W.\theta_2, m, \text{inl}(v_{i2})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

We instantiate IH1 with m from (S01) to get

$$(W.\theta_1, m, v_{i1}) \in \lfloor \tau_1 \rfloor_V$$

Therefore from Definition 2.6, we get

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

We instantiate IH2 with m from (S02) to get

$$(W.\theta_2, m, v_{j1}) \in \lfloor \tau_1 \rfloor_V$$

Therefore from Definition 2.6, we get

$$(W.\theta_2, m, \text{inl}(v_{j1})) \in \lfloor \tau_1 + \tau_2 \rfloor_V$$

(b) $v_1 = \text{inr}(v_{i2})$ and $v_2 = \text{inr}(v_{j2})$

Symmetric case as (a)

4. Case $\tau_1 \xrightarrow{\ell_e} \tau_2$:

$$\text{Given: } (W, n, \lambda x.e_1, \lambda x.e_2) \in \lceil \tau_1 \xrightarrow{\ell_e} \tau_2 \rceil_V^A$$

This means from Definition 2.4 we know that

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in \lceil \tau_1 \rceil_V^A &\implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in \lceil \tau_2 \rceil_E^A) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in \lfloor \tau_1 \rfloor_V &\implies (\theta_l, i, e_1[v_c/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in \lfloor \tau_1 \rfloor_V &\implies (\theta_l, k, e_2[v_c/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}) \end{aligned} \quad (\text{L0})$$

To prove:

$$(a) \forall m. (W.\theta_1, m, \lambda x.e_1) \in \lfloor \tau_1 \xrightarrow{\ell_e} \tau_2 \rfloor_V:$$

This means from Definition 2.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

This further means that we have some θ', j and v s.t

$$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in \lfloor \tau_1 \rfloor_V$$

$$\text{And we need to prove: } (\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

Instantiating θ_l, i and v_c in the second conjunct of L0 with θ', j and v respectively and since we know that $W.\theta_1 \sqsubseteq \theta'$ and $(\theta', j, v) \in \lfloor \tau_1 \rfloor_V$

$$\text{Therefore we get } (\theta', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E^{\ell_e}$$

$$(b) \forall m. (W.\theta_2, m, \lambda x.e_2) \in \lfloor \tau_1 \xrightarrow{\ell_e} \tau_2 \rfloor_V:$$

Similar reasoning with e_2

5. Case $\forall \alpha. (\ell_e, \tau)$:

$$\text{Given: } (W, n, \Lambda e_1, \Lambda e_2) \in \lceil \forall \alpha. (\ell_e, \tau) \rceil_V^A$$

This means from Definition 2.4 we know that

$$\begin{aligned}
& \forall W_b \sqsupseteq W, n_b < n, \ell' \in \mathcal{L}. ((W_b, n_b, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_2, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}) \quad (\text{F0})
\end{aligned}$$

To prove:

$$(a) \forall m. (W.\theta_1, m, \Lambda e_1) \in [\forall \alpha. (\ell_e, \tau)]_V:$$

This means from Definition 2.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \forall \ell_u \in \mathcal{L}. (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E^{\ell_e[\ell_u/\alpha]}$$

This further means that we are given some θ' , m' and ℓ_u s.t $W.\theta_1 \sqsubseteq \theta'$, $m' < m$ and $\ell_u \in \mathcal{L}$

$$\text{And we need to prove: } (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E^{\ell_e[\ell_u/\alpha]}$$

Instantiating θ_l, i and ℓ'' in the second conjunct of F0 with θ' , m' and ℓ_u respectively and since we know that $W.\theta_1 \sqsubseteq \theta'$ and $\ell_u \in \mathcal{L}$

$$\text{Therefore we get } (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E^{\ell_e[\ell_u/\alpha]}$$

$$(b) \forall m. (W.\theta_2, m, \Lambda e_2) \in [\forall \alpha. (\ell_e, \tau)]_V:$$

Symmetric reasoning for e_2

6. Case $c \xrightarrow{\ell_e} \tau$:

$$\text{Given: } (W, n, \nu e_1, \nu e_2) \in [c \xrightarrow{\ell_e} \tau]_V^A$$

This means from Definition 2.4 we know that

$$\begin{aligned}
& \forall W_b \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W_b, n', e_1, e_2) \in [\tau]_E^A \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E^{\ell_e} \\
& \wedge \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E^{\ell_e} \quad (\text{C0})
\end{aligned}$$

To prove:

$$(a) \forall m. (W.\theta_1, m, \nu e_1) \in [c \xrightarrow{\ell_e} \tau]_V:$$

This means from Definition 2.6 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e_1) \in [\tau]_E^{\ell_e}$$

This further means that we are given some θ' and m' s.t $W.\theta_1 \sqsubseteq \theta'$, $m' < m$ and $\mathcal{L} \models c$

$$\text{And we need to prove: } (\theta', m', e_1) \in [\tau]_E^{\ell_e}$$

Instantiating θ_l, j in the second conjunct of C0 with θ' , m' respectively and since we know that $W.\theta_1 \sqsubseteq \theta'$ and $\mathcal{L} \models c$

$$\text{Therefore we get } (\theta', m', e_1) \in [\tau]_E^{\ell_e}$$

$$(b) \forall m. (W.\theta_2, m, \nu e_2) \in [c \xrightarrow{\ell_e} \tau]_V:$$

Symmetric reasoning for e_2

7. Case ref τ :

From Definition 2.4 and 2.6

Proof of statement (2)

Let $\tau = A^\ell$

2 cases arise:

1. $\ell \sqsubseteq \mathcal{A}$:

From IH (statement(1))

2. $\ell \not\sqsubseteq \mathcal{A}$:

Directly from Definition 2.4

□

Lemma 2.16 (FG: Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m'$.

1. $\forall \mathbf{A}. (\theta, m, v) \in \llbracket \mathbf{A} \rrbracket_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \llbracket \mathbf{A} \rrbracket_V$

2. $\forall \tau. (\theta, m, v) \in \llbracket \tau \rrbracket_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \llbracket \tau \rrbracket_V$

Proof. Proof by simultaneous induction on \mathbf{A} and τ

Proof of statement (1)

We analyze the various cases of \mathbf{A} in the last step:

1. case \mathbf{b} :

Directly from Definition 2.6

2. case $\tau_1 \times \tau_2$:

Given: $(\theta, m, (v_1, v_2)) \in \llbracket \tau_1 \times \tau_2 \rrbracket_V$

To prove: $(\theta', m', (v_1, v_2)) \in \llbracket \tau_1 \times \tau_2 \rrbracket_V$

This means from Definition 2.6 we know that

$(\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V \wedge (\theta, m, v_2) \in \llbracket \tau_2 \rrbracket_V$

IH1 : $(\theta', m', v_1) \in \llbracket \tau_1 \rrbracket_V$

IH2 : $(\theta', m', v_2) \in \llbracket \tau_2 \rrbracket_V$

We get the desired from IH1, IH2 and Definition 2.6

3. case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v = \text{inl}(v_1)$:

Given: $(\theta, m, (\text{inl } v_1)) \in \llbracket \tau_1 + \tau_2 \rrbracket_V$

To prove: $(\theta', m', \text{inl } v_1) \in \llbracket \tau_1 + \tau_2 \rrbracket_V$

This means from Definition 2.6 we know that

$(\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V$

IH : $(\theta', m', v_1) \in \llbracket \tau_1 \rrbracket_V$

Therefore from IH and Definition 2.6 we get the desired

(b) $v = \text{inr}(v_2)$

Symmetric case

4. case $\tau_1 \xrightarrow{\ell_e} \tau_2$:

Given: $(\theta, m, (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

To prove: $(\theta', m', (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V$

This means from Definition 2.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in [\tau_1]_V \implies (\theta'', j, e_1[v/x]) \in [\tau_2]_E^{\ell_e} \quad (69)$$

Similarly from Definition 2.6 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall v_1. (\theta''', k, v_1) \in [\tau_1]_V \implies (\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$$

This means that given some θ''', k and v_1 such that $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in [\tau_1]_V$

And we are required to prove $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating Equation 143 with θ''', k and v_1 and since we know that $\theta' \sqsubseteq \theta'''$ and $\theta \sqsubseteq \theta'$ therefore we have $\theta \sqsubseteq \theta'''$. Also, we know that $k < m' < m$ and $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E^{\ell_e}$

5. case ref τ :

From Definition 2.6 and Definition 2.2

6. case $\forall \alpha. (\ell_e, \tau)$:

Given: $(\theta, m, (\Lambda e_1)) \in [\forall \alpha. (\ell_e, \tau)]_V$

To prove: $(\theta', m', (\Lambda e_1)) \in [\forall \alpha. (\ell_e, \tau)]_V$

This means from Definition 2.6 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall \ell_i \in \mathcal{L}. (\theta'', j, e_1) \in [\tau[\ell_i/\alpha]]_E^{\ell_e[\ell_i/\alpha]} \quad (70)$$

Similarly from Definition 2.6 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall \ell_j \in \mathcal{L}. (\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E^{\ell_e[\ell_j/\alpha]}$$

This means that given some θ''', k and ℓ_j such that $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E^{\ell_e[\ell_j/\alpha]}$

Instantiating Equation 70 with θ''', k and ℓ_j and since we know that $\theta' \sqsubseteq \theta'''$ and $\theta \sqsubseteq \theta'$ therefore we have $\theta \sqsubseteq \theta'''$. Also, we know that $k < m' < m$ and $\ell_j \in \mathcal{L}$

Therefore we get $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E^{\ell_e[\ell_j/\alpha]}$

7. case $c \xrightarrow{\ell_e} \tau$:

Given: $(\theta, m, (\nu e_1)) \in [c \xrightarrow{\ell_e} \tau]_V$

To prove: $(\theta', m', (\nu e_1)) \in [c \xrightarrow{\ell_e} \tau]_V$

This means from Definition 2.6 we know that

$$\forall \theta'' . \theta \sqsubseteq \theta'' \wedge \forall j < m . \mathcal{L} \models c \implies (\theta'' , j, e_1) \in [\tau]_E^{\ell_e} \quad (71)$$

Similarly from Definition 2.6 we know that we are required to prove

$$\forall \theta''' . \theta' \sqsubseteq \theta''' \wedge \forall k < m' . \mathcal{L} \models c \implies (\theta''' , k, e_1) \in [\tau]_E^{\ell_e}$$

This means that given some θ''' , k and ℓ_j such that $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove $(\theta''' , k, e_1) \in [\tau]_E^{\ell_e}$

Instantiating Equation 71 with θ''' , k and since we know that $\theta' \sqsubseteq \theta'''$ and $\theta \sqsubseteq \theta'$ therefore we have $\theta \sqsubseteq \theta'''$. Also, we know that $k < m' < m$ and $\mathcal{L} \models c$

Therefore we get $(\theta''' , k, e_1) \in [\tau]_E^{\ell_e}$

Proof of statement (2)

Let $\tau = A^\ell$

Since $[A^\ell]_V = [A]_V$, therefore from IH (statement 1) □

Lemma 2.17 (FG: Monotonicity binary). *The following holds:*

$\forall W, W', v_1, v_2, \mathcal{A}, n, n'$.

$$1. \forall A. (W, n, v_1, v_2) \in [A]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [A]_V^A$$

$$2. \forall \tau. (W, n, v_1, v_2) \in [\tau]_V^A \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in [\tau]_V^A$$

Proof. Proof by simultaneous induction on A and τ

Proof of statement (1)

We analyze the different cases of A in the last step:

1. Case b :

From Definition 2.4

2. Case $\tau_1 \times \tau_2$:

Given: $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove: $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

From Definition 2.4 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$$

$$\text{IH1 : } (W', n', v_{i1}, v_{j1}) \in [\tau_1]_V^A$$

$$\text{IH2 : } (W', n', v_{i2}, v_{j2}) \in [\tau_2]_V^A$$

From IH1, IH2 and Definition 2.4 we get the desired.

3. Case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v_1 = \text{inl } v_{i1}$ and $v_2 = \text{inl } v_{i2}$:

Given: $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

To prove: $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

From Definition 2.4 we know that we are given

$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

IH : $(W', n', v_{i1}, v_{i2}) \in [\tau_1]_V^A$

Therefore from Definition 2.4 we get

$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_V^A$

(b) $v_1 = \text{inr}(v_{i2})$ and $v_2 = \text{inr}(v_{i2})$:

Symmetric case

4. Case $\tau_1 \xrightarrow{\ell_e} \tau_2$:

Given: $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

To prove: $(\theta', n', (\lambda x.e_1), (\lambda x.e_1)) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^A$

This means from Definition 2.4 we know that the following holds

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A)$
(BM-A0)

$\forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E^{\ell_e})$ (BM-A1)

$\forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E^{\ell_e})$ (BM-A2)

Similarly from Definition 2.4 we know that we are required to prove

(a) $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$:

This means that we are given some $W'' \sqsupseteq W', k < n'$ and v'_1, v'_2 s.t

$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

And we a required to prove: $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with W'', k and v'_1, v'_2 we get

$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

(b) $\forall \theta'_l \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e})$:

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_1, k$ and v'_c s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we a required to prove: $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A1 with θ'_l, k and v'_c we get

$(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

(c) $\forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e})$:

This means that we are given some $\theta'_l \sqsupseteq W.\theta_2, k$ and v'_c s.t

$(\theta'_l, k, v'_c) \in [\tau_1]_V$

And we a required to prove: $(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

Instantiating BM-A1 with θ'_l, k and v'_c we get

$(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E^{\ell_e}$

5. Case ref τ :

From Definition 2.4 and Definition 2.3

6. Case $\forall\alpha.(\ell_e, \tau)$:

Given: $(W, n, (\Lambda e_1), (\Lambda e_2)) \in [\forall\alpha.(\ell_e, \tau)]_V^A$

To prove: $(\theta', n', (\Lambda e_1), (\Lambda e_1)) \in [\forall\alpha.(\ell_e, \tau)]_V^A$

This means from Definition 2.4 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \quad (\text{BM-F0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}) \quad (\text{BM-F1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}) \quad (\text{BM-F2})$$

Similarly from Definition 2.4 we know that we are required to prove

$$(a) \quad \forall W'' \sqsupseteq W', n'' < n', \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A):$$

This means that we are given some $W'' \sqsupseteq W', n'' < n'$ and $\ell'' \in \mathcal{L}$

And we are required to prove: $((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$

Instantiating BM-F0 with W'', n'' and ℓ'' . And since $W'' \sqsupseteq W'$ and $W' \sqsupseteq W$ therefore $W'' \sqsupseteq W$. Also since $n'' < n'$ and $n' < n$ therefore $n'' < n$. And finally since $\ell'' \in \mathcal{L}$ therefore we get

$$((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$$

$$(b) \quad \forall \theta'_l \sqsupseteq W'.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}):$$

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_1, k$ and $\ell'' \in \mathcal{L}$

And we are required to prove: $((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$

Instantiating BM-F1 with θ'_l, k and ℓ'' . And since $\theta'_l \sqsupseteq W'.\theta_1$ and $W' \sqsupseteq W$ therefore $\theta'_l \sqsupseteq W.\theta_1$. And since $\ell'' \in \mathcal{L}$ therefore we get

$$((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$$

$$(c) \quad \forall \theta_l \sqsupseteq W.\theta_2, j, \ell'' \in \mathcal{L}. ((\theta_l, k, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]}):$$

This means that we are given some $\theta_l \sqsupseteq W.\theta_2, k$ and $\ell'' \in \mathcal{L}$

And we are required to prove: $((\theta_l, k, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$

Instantiating BM-F2 with θ_l, k and ℓ'' . And since $\theta_l \sqsupseteq W.\theta_2$ and $W' \sqsupseteq W$ therefore $\theta_l \sqsupseteq W.\theta_2$. And since $\ell'' \in \mathcal{L}$ therefore we get

$$((\theta_l, k, e_2) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]})$$

7. Case $c \stackrel{\ell_e}{\Rightarrow} \tau$:

Given: $(W, n, (\nu e_1), (\nu e_2)) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^A$

To prove: $(\theta', n', (\nu e_1), (\nu e_1)) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^A$

This means from Definition 2.4 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W', n', e_1, e_2) \in [\tau]_E^A \quad (\text{BM-C0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E^{\ell_e} \quad (\text{BM-C1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E^{\ell_e} \quad (\text{BM-C2})$$

Similarly from Definition 2.4 we know that we are required to prove

(a) $\forall W'' \sqsupseteq W', n'' < n. \mathcal{L} \models c \implies (W'', n'', e_1, e_2) \in [\tau]_E^A$:

This means that we are given some $W'' \sqsupseteq W', n'' < n'$ and $\mathcal{L} \models c$

And we are required to prove: $(W'', n'', e_1, e_2) \in [\tau]_E^A$

Instantiating BM-C0 with W'', n'' . And since $W'' \sqsupseteq W'$ and $W' \sqsupseteq W$ therefore $W'' \sqsupseteq W$. And since $\mathcal{L} \models c$ therefore we get

$(W'', n'', e_1, e_2) \in [\tau]_E^A$

(b) $\forall \theta'_i \sqsupseteq W'.\theta_1, k. \mathcal{L} \models c \implies (\theta'_i, k, e_1) \in [\tau]_E^{\ell_e}$:

This means that we are given some $\theta'_i \sqsupseteq W'.\theta_1, k$ and $\mathcal{L} \models c$

And we are required to prove: $(\theta'_i, k, e_1) \in [\tau]_E^{\ell_e}$

Instantiating BM-F1 with θ'_i, k . And since $\theta'_i \sqsupseteq W'.\theta_1$ and $W' \sqsupseteq W$ therefore $\theta'_i \sqsupseteq W.\theta_1$. And since $\mathcal{L} \models c$ therefore we get

$(\theta'_i, k, e_1) \in [\tau]_E^{\ell_e}$

(c) $\forall \theta'_i \sqsupseteq W'.\theta_2, k. \mathcal{L} \models c \implies (\theta'_i, k, e_2) \in [\tau]_E^{\ell_e}$:

This means that we are given some $\theta'_i \sqsupseteq W'.\theta_2, k$ and $\mathcal{L} \models c$

And we are required to prove: $(\theta'_i, k, e_2) \in [\tau]_E^{\ell_e}$

Instantiating BM-F1 with θ'_i, k . And since $\theta'_i \sqsupseteq W'.\theta_2$ and $W' \sqsupseteq W$ therefore $\theta'_i \sqsupseteq W.\theta_2$. And since $\mathcal{L} \models c$ therefore we get

$(\theta'_i, k, e_2) \in [\tau]_E^{\ell_e}$

Proof of statement (2)

Let $\tau = A^\ell$

2 cases arise:

1. $\ell \sqsubseteq \mathcal{A}$:

From IH (statement 1)

2. $\ell \not\sqsubseteq \mathcal{A}$:

From Lemma 2.16 and Definition 2.4

□

Lemma 2.18 (FG: Unary monotonicity for Γ). $\forall \theta, \theta', \delta, \Gamma, n, n'$.

$(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in [\Gamma]_V$

Proof. Given: $(\theta, n, \delta) \in [\Gamma]_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$

To prove: $(\theta', n', \delta) \in [\Gamma]_V$

From Definition 2.13 it is given that

$dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in [\Gamma(x)]_V$

And again from Definition 2.13 we are required to prove that

$dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in [\Gamma(x)]_V$

• $dom(\Gamma) \subseteq dom(\delta)$:

Given

- $\forall x \in \text{dom}(\Gamma).(\theta', n', \delta(x)) \in \lfloor \Gamma(x) \rfloor_V$:
 Since we know that $\forall x \in \text{dom}(\Gamma).(\theta, n, \delta(x)) \in \lfloor \Gamma(x) \rfloor_V$ (given)
 Therefore from Lemma 2.16 we get
 $\forall x \in \text{dom}(\Gamma).(\theta', n', \delta(x)) \in \lfloor \Gamma(x) \rfloor_V$

□

Lemma 2.19 (FG: Binary monotonicity for Γ). $\forall W, W', \delta, \Gamma, n, n'$.
 $(W, n, \gamma) \in \lfloor \Gamma \rfloor_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in \lfloor \Gamma \rfloor_V$

Proof. Given: $(W, n, \gamma) \in \lfloor \Gamma \rfloor_V \wedge n' < n \wedge W \sqsubseteq W'$
 To prove: $(W', n', \gamma) \in \lfloor \Gamma \rfloor_V$

From Definition 2.14 it is given that
 $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$

And again from Definition 2.13 we are required to prove that
 $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma)$:
 Given
- $\forall x \in \text{dom}(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$:
 Since we know that $\forall x \in \text{dom}(\Gamma).(W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$ (given)
 Therefore from Lemma 2.17 we get
 $\forall x \in \text{dom}(\Gamma).(W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \lceil \Gamma(x) \rceil_V^A$

□

Lemma 2.20 (FG: Unary monotonicity for H). $\forall \theta, H, n, n'$.
 $(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$

Proof. Given: $(n, H) \triangleright \theta \wedge n' < n$
 To prove: $(n', H) \triangleright \theta$

From Definition 2.8 it is given that
 $\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$

And again from Definition 2.13 we are required to prove that
 $\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$

- $\text{dom}(\theta) \subseteq \text{dom}(H)$:
 Given
- $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$:
 Since we know that $\forall a \in \text{dom}(\theta).(\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$ (given)
 Therefore from Lemma 2.16 we get
 $\forall a \in \text{dom}(\theta).(\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$

□

Lemma 2.21 (FG: Binary monotonicity for heaps). $\forall W, H_1, H_2, n, n'$.
 $(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$

Proof. Given: $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$
 To prove: $(n', H_1, H_2) \triangleright W$

From Definition 2.9 it is given that

$$\begin{aligned} \text{dom}(W.\theta_1) &\subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ (W.\hat{\beta}) &\subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ \forall(a_1, a_2) \in (W.\hat{\beta}). &(W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ (W, n-1, H_1(a_1), H_2(a_2)) &\in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). &(W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

And again from Definition 2.9 we are required to prove:

- $\text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2)$:
 Given
- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$:
 Given
- $\forall(a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2))$ and $(W, n'-1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$:
 $\forall(a_1, a_2) \in (W.\hat{\beta})$.
 - $(W.\theta_1(a_1) = W.\theta_2(a_2))$: Given
 - $(W, n'-1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$:
 Given and from Lemma 2.17
- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$:
 Given

□

Theorem 2.22 (FG: Fundamental theorem unary). $\forall \Sigma, \Psi, \Gamma, pc, \theta, \mathcal{L}, e, \tau, \sigma, \delta, n$.

$$\begin{aligned} \Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \wedge \\ \mathcal{L} \models \Psi \sigma \wedge \\ (\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V \implies \\ (\theta, n, e \delta) \in \lfloor \tau \sigma \rfloor_E^{pc} \end{aligned}$$

Proof. Proof by induction on *FG* typing derivation

1. FG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

To prove: $(\theta, n, x \delta) \in \lfloor \tau \sigma \rfloor_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap H and $j < n$ s.t $(n, H) \triangleright \theta \wedge (H, x \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-V0})
\end{aligned}$$

In order to prove FU-V0 we instantiate θ' with θ . From reduction relation we know that $H' = H$, $v' = \delta(x)$ and $j = 1$

We need to prove the following:

- (a) $\theta \sqsubseteq \theta \wedge (n - 1, H) \triangleright \theta \wedge (\theta, n - 1, v') \in \lfloor \tau \sigma \rfloor_V$:
- $\theta \sqsubseteq \theta$: From Definition 2.2
 - $(n - 1, H) \triangleright \theta$: From Lemma 2.20
 - $(\theta, n - 1, v') \in \lfloor \tau \sigma \rfloor_V$:
Since we are given that $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$ and $v' = \delta(x)$
Therefore $(\theta, n, v') \in \lfloor \Gamma(x) \sigma \rfloor_V$, where $\Gamma(x) = \tau$
And finally from Lemma 2.16 we get $(\theta, n - 1, v') \in \lfloor \tau \sigma \rfloor_V$
- (b) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$:
Since $H' = H$, so we are done
- (c) $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$:
Since $\theta' = \theta$, so we are done

2. FG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove: $(\theta, \lambda x. e_i \delta) \in \lfloor ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \sigma \rfloor_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, (\lambda x. e_i) \delta) \Downarrow_j (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap H and $j < n$ s.t $(n, H) \triangleright \theta \wedge (H, (\lambda x. e_i) \delta) \Downarrow_j (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma \rfloor_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-L0})
\end{aligned}$$

IH1:

$\forall \theta_i, v_x, n. (\theta_i, n, e_i \delta \cup \{x \mapsto v_x\}) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$, s.t $(\theta_i, n, v_x) \in [\tau_1 \sigma]_V$

In order to prove FU-L0 we instantiate θ' with θ . From reduction relation we know that $H' = H$, $j = 0$ and $v' = \lambda x. e_i \delta$

(a) $\theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in [((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \sigma]_V$:

- $\theta \sqsubseteq \theta$: From Definition 2.2
- $(n, H) \triangleright \theta$: Given
- $(\theta, n, (\lambda x. e_i) \delta) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp) \sigma]_V$:

From Definition 2.6 it suffices to prove that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < n. \forall v. (\theta'', j, v) \in [\tau_1 \sigma]_V \implies (\theta'', j, e_i[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

This means given some θ'' , j and v such that $\theta \sqsubseteq \theta''$, $j < n$ and $(\theta'', j, v) \in [\tau_1 \sigma]_V$
It suffices to prove that $(\theta'', j, e_i[v/x] \delta) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

Since $(\theta, n, \delta) \in [\Gamma \sigma]_V$ and $j < n$ therefore from Lemma 2.18 we have
 $(\theta, j, \delta) \in [\Gamma \sigma]_V$

So we can apply IH1 instantiated with θ'' , v and j we get
 $(\theta'', j, e_i[v/x] \delta) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

(b) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sqsubseteq \ell')$:

Since $H' = H$ so we are done

(c) $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc)$:

Since $\theta' = \theta$ so we are done

3. FG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove: $(\theta, n, (e_1 e_2) \delta) \in [\tau_2 \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2 \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap H s.t $(n, H) \triangleright \theta \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_2 \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = A^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad \text{(FU-P0)} \end{aligned}$$

IH1:

$$\begin{aligned} & \forall n_1, H_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V \wedge \end{aligned}$$

$$\begin{aligned}
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH1 with n, H and since we know that $(n, H) \triangleright \theta \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad \text{(FU-P1)}
\end{aligned}$$

From evaluation rule we know that $v'_1 = \lambda x. e_i$. Since from FU-P1 we know that

$$(\theta'_1, n - i, \lambda x. e_i) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_V$$

This means from Definition 2.6 we have

$$\forall \theta''. \theta'_1 \sqsubseteq \theta'' \wedge \forall j < (n - i). \forall v. (\theta'', j, v) \in [\tau_1 \sigma]_V \implies (\theta'', j, e_i[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma} \quad (72)$$

IH2:

$$\begin{aligned}
& \forall n_2, \forall H_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall k < n_2. (H_2, (e_2) \delta) \Downarrow_k (H'_2, v'_2) \implies \\
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - k, v'_2) \in [(\tau_1) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH2 with $n - i, H'_1$ and since we know that $(n - i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1 e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - k, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - k, v'_2) \in [(\tau_1) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \sigma) \quad \text{(FU-P2)}
\end{aligned}$$

Instantiating θ'', j and v in Equation 72 with $\theta'_2, n - i - k$ and v'_2 from FU-P2 respectively, we get

$$(\theta'_2, n - i - k, e_i[v'_2/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

This means from Definition 2.7 we have

$$\begin{aligned}
& \forall H_3. (n - i - k, H_3) \triangleright \theta'_2 \wedge \forall l < (n - i - k). (H_3, e_i[v'_2/x]) \Downarrow_l (H'_3, v'_3) \implies \\
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n - i - k - l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n - i - k - l), v'_3) \in [\tau_2 \sigma]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e \sigma)
\end{aligned}$$

Instantiating H_3 with H'_2 from FU-P2 and since we know that $((n - i - k), H'_2) \triangleright \theta'_2$ and since the reduction happens therefore we have

$$\begin{aligned}
& \exists \theta'_3. \theta'_2 \sqsubseteq \theta'_3 \wedge ((n - i - k - l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n - i - k - l), v'_3) \in [\tau_2 \sigma]_V \wedge \\
& (\forall a. H_3(a) \neq H'_3(a) \implies \exists \ell'. \theta'_2(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_2). \theta'_3(a) \searrow \ell_e \sigma) \quad \text{(FU-P3)}
\end{aligned}$$

In order to prove FU-P0 we choose θ' as θ'_3 from FU-P3. Also we know that $H' = H'_3$, $v' = v'_3$ and $n' = i + k + l$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_3 \wedge ((n - i - k - l), H'_3) \triangleright \theta'_3 \wedge (\theta'_3, (n - i - k - l), v'_3) \in \llbracket \tau_2 \sigma \rrbracket_V$:
- $\theta \sqsubseteq \theta'_3$:
Since $\theta \sqsubseteq \theta'_1$ from FU-P1, $\theta'_1 \sqsubseteq \theta'_2$ from FU-P2 and $\theta'_2 \sqsubseteq \theta'_3$ from FU-P3 therefore from Definition 2.2 we get $\theta \sqsubseteq \theta'_3$
 - $((n - i - k - l), H'_3) \triangleright \theta'_3$:
From FU-P3 we get $((n - i - k - l), H'_3) \triangleright \theta'_3$
 - $(\theta'_3, (n - i - k - l), v'_3) \in \llbracket \tau_2 \sigma \rrbracket_V$:
From FU-P3 we get $(\theta'_3, (n - i - k - l), v'_3) \in \llbracket \tau_2 \sigma \rrbracket_V$
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'_3(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$
Since $pc \sigma \sqsubseteq \ell_e \sigma$ therefore we get the desired from FU-P1, FU-P2 and FU-P3
- (c) $(\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta). \theta'_3(a) \searrow pc \sigma)$
Since $pc \sigma \sqsubseteq \ell_e \sigma$ therefore we get the desired from FU-P1, FU-P2 and FU-P3

4. FG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove: $(\theta, n, (e_1, e_2) \delta) \in \llbracket (\tau_1 \times \tau_2)^\perp \sigma \rrbracket_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\tau_1 \times \tau_2)^\perp \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap H s.t $H \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\tau_1 \times \tau_2)^\perp \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad \text{(FU-PA0)} \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in \llbracket \tau_1 \sigma \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \end{aligned}$$

We instantiate IH1 with H and n . And since we know that $(n, H) \triangleright \theta \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in \llbracket \tau_1 \sigma \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad \text{(FU-PA1)} \end{aligned}$$

IH2:

$$\begin{aligned} & \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in \llbracket (\tau_2) \sigma \rrbracket_V \wedge \end{aligned}$$

$$\begin{aligned}
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \ \sigma)
\end{aligned}$$

We instantiate IH2 with H'_1 and $n - i$. And since we know that $(n - i, H'_1) \triangleright \theta'_1 \wedge (H, (e_1, e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau_2) \ \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc \ \sigma) \quad (\text{FU-PA2})
\end{aligned}$$

In order to prove FU-PA0 we choose θ' as θ'_2 from FU-PA2. Also we know from the evaluation rule, that let $v' = (v'_1, v'_2)$, $H' = H'_2$ and $n' = i + j + 1$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v') \in [(\tau_1 \times \tau_2)^\perp]_V$:
- $\theta \sqsubseteq \theta'_2$:
Since $\theta \sqsubseteq \theta'_1$ from FU-PA1 and $\theta'_1 \sqsubseteq \theta'_2$ from FU-PA2 therefore from Definition 2.2 we get $\theta \sqsubseteq \theta'_2$
 - $(n - i - j - 1, H'_2) \triangleright \theta'_2$:
From FU-PA2 we get $(n - i - j, H'_2) \triangleright \theta'_2$ therefore from Lemma 2.20 we get $(n - i - j - 1, H'_2) \triangleright \theta'_2$
 - $(\theta'_2, n - i - j, v') \in [(\tau_1 \times \tau_2)^\perp]_V$:
From Definition 2.6 it suffices to show
 - i. $(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1) \ \sigma]_V$:
Since from FU-PA1 we know that $(\theta'_1, n - i, v'_1) \in [(\tau_1) \ \sigma]_V$ and since $\theta'_1 \sqsubseteq \theta'_2$ (from FU-PA2) therefore from Lemma 2.16 we get $(\theta'_2, n - i - j - 1, v'_1) \in [(\tau_1) \ \sigma]_V$
 - ii. $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2) \ \sigma]_V$:
From FU-PA2 we know that $(\theta'_2, n - i - j, v'_2) \in [(\tau_2) \ \sigma]_V$ therefore from Lemma 2.16 we get $(\theta'_2, n - i - j - 1, v'_2) \in [(\tau_2) \ \sigma]_V$
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell')$
From FU-PA1 and FU-PA2
- (c) $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \ \sigma)$
From FU-PA1 and FU-PA2

5. FG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$

To prove: $(\theta, n, \text{fst}(e_i) \delta) \in [\tau_1 \ \sigma]_E^{pc \ \sigma}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{fst}(e_i) \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1 \ \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \ \sigma)
\end{aligned}$$

This means that given some heap H s.t $(n, H) \triangleright \theta \wedge (H, \text{fst}(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau_1 \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-F0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \end{aligned}$$

Instantiating IH1 with H and n . Since we know that $H \triangleright \theta \wedge (H, \text{fst}(e_i) \delta) \Downarrow (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad (\text{FU-F1}) \end{aligned}$$

From evaluation rule we know that $v'_1 = (v''_1, v''_2)$

In order to prove FU-F0 we choose θ' as θ'_1 from FU-F1. Also we know that $H' = H'_1$ and $v' = v''_1$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v'_1) \in [\tau_1 \sigma]_V$:
 - $\theta \sqsubseteq \theta'_1$:
From FU-F1
 - $(n - i - 1, H'_1) \triangleright \theta'_1$:
From FU-F1 we know $(n - i, H'_1) \triangleright \theta'_1$ therefore from Lemma 2.20 we get $(n - i - 1, H'_1) \triangleright \theta'_1$
 - $(\theta'_1, n - i, v''_1) \in [\tau_1 \sigma]_V$:
Since from FU-F1 we know that $(\theta'_1, n - i, (v''_1, v''_2)) \in [(\tau_1 \times \tau_2) \sigma]_V$
Therefore from Definition 2.6 we know that $(\theta'_1, n - i, v''_1) \in [\tau_1 \sigma]_V$
From Lemma 2.16 we get $(\theta'_1, n - i - 1, v''_1) \in [\tau_1 \sigma]_V$
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$
From FU-F1
- (c) $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)$
From FU-F1

6. FG-snd:

Symmetric case to FG-fst

7. FG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove: $(\theta, n, \text{inl}(e_i) \delta) \in [(\tau_1 + \tau_2)^\perp \sigma]_E^{pc \sigma}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 + \tau_2)^\perp]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma)
\end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\tau_1 + \tau_2)^\perp]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad \text{(FU-LE0)}
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau_1 \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH1 with H and n . Since we know that $(n, H) \triangleright \theta \wedge (H, \text{inl}(e_i) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau_1 \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad \text{(FU-LE1)}
\end{aligned}$$

In order to prove FU-LE0 we choose θ' as θ'_1 from FU-LE1. Also we know from the evaluation rule, that let $v' = \text{inl}(v'_1)$, $H' = H'_1$ and $n' = i + 1$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2)]_V$:
- $\theta \sqsubseteq \theta'_1$:
From FU-LE1
 - $(n - i - 1, H') \triangleright \theta'_1$:
From FU-LE1 we know that $(n - i, H') \triangleright \theta'_1$ therefore from Lemma 2.20 we get $(n - i - 1, H') \triangleright \theta'_1$
 - $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2) \sigma]_V$:
Since $v' = \text{inl}(v'_1)$ and from FU-LE1 we know that $(\theta'_1, n - i, v'_1) \in [\tau_1 \sigma]_V$
Therefore from Definition 2.6 we get $(\theta'_1, n - i, v') \in [(\tau_1 + \tau_2) \sigma]_V$
From Lemma 2.16 we get $(\theta'_1, n - i - 1, v') \in [(\tau_1 + \tau_2) \sigma]_V$
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$
From FU-LE1
- (c) $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)$
From FU-LE1

8. FG-inr:

Symmetric case to FG-inl

9. FG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau}$$

To prove: $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in [\tau \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap H and n s.t. $(n, H) \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-C0}) \end{aligned}$$

IH1:

$$\begin{aligned} & \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_c) \delta) \Downarrow_i (H'_1, v'_c) \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_c) \in [(\tau_1 + \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \end{aligned}$$

Instantiating IH1 with H and n . Since we know that $H \triangleright \theta \wedge (H, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_c) \in [(\tau_1 + \tau_2)^\ell \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad (\text{FU-C1}) \end{aligned}$$

2 cases arise:

(a) $v'_c = \text{inl}(v_{ci})$:

IH2:

$$\begin{aligned} & \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_1) \delta \cup \{x \mapsto v_{ci}\}) \Downarrow_j (H'_2, v'_2) \implies \\ & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n_2 - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in [(\tau) \sigma]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

Instantiating IH2 with H'_1 and $n-i$ since we know that $H'_1 \triangleright \theta'_1 \wedge (H'_1, (\text{case } e_c, x.e_1, y.e_2) \delta) \Downarrow (H', v')$ therefore we have

$$\begin{aligned} & \exists \theta'_2. \theta'_1 \sqsubseteq \theta'_2 \wedge (n - i - j, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau) \sigma]_V \wedge \\ & (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma) \quad (\text{FU-C2}) \end{aligned}$$

In order to prove FU-C0 we choose θ' as θ'_2 from FU-C2. Also we know that $H' = H'_2$, $v' = v'_2$ and $n' = i + j + 1$. Now we are required to show

- i. $\theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, v'_2) \in \lfloor \tau \sigma \rfloor_V$:
- $\theta \sqsubseteq \theta'_2$:
Since $\theta \sqsubseteq \theta'_1$ from FU-C1 and $\theta'_1 \sqsubseteq \theta'_2$ from FU-C2 therefore from Definition 2.2 we get $\theta \sqsubseteq \theta'_2$
 - $(n - i - j - 1, H'_2) \triangleright \theta'_2$:
From FU-C2 we know that $(n - i - j, H'_2) \triangleright \theta'_2$ therefore from Lemma 2.20 we get $(n - i - j - 1, H'_2) \triangleright \theta'_2$
 - $(\theta'_2, n - i - j - 1, v'_2) \in \lfloor \tau \sigma \rfloor_V$:
From FU-C2 we know that $(\theta'_2, n - i - j, v'_2) \in \lfloor \tau \sigma \rfloor_V$ therefore from Lemma 2.16 we get $(\theta'_2, n - i - j - 1, v'_2) \in \lfloor \tau \sigma \rfloor_V$
- ii. $(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$:
Since from FU-C2 we know that
 $(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell')$
therefore we also have
 $(\forall a. H'_1(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge (pc) \sigma \sqsubseteq \ell')$
and from FU-C1 we know that
 $(\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge (pc) \sigma \sqsubseteq \ell')$
Combining the two we get
 $(\forall a \in \text{dom}(H). H(a) \neq H'_2(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$
- iii. $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \sigma)$:
Since from FU-C2 we know that
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma)$
therefore we also have
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow (pc) \sigma)$
and from FU-C1 we know that
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow (pc \sqcup \ell) \sigma)$
Combining the two we get
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \sigma)$
- (b) $v'_c = \text{inr}(v_{ci})$:
Symmetric case as $v'_c = \text{inl}(v_{ci})$

10. FG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove: $(\theta, n, \text{new } (e_i) \delta) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \lfloor (\text{ref } \tau)^\perp \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, \text{new } (e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(\mathbf{ref} \ \tau)^\perp]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \ \sigma) \quad (\text{FU-R0})
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \ \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [\tau \ \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \ \sigma)
\end{aligned}$$

Instantiating IH1 with H and n . Since we know that $(n, H) \triangleright \theta \wedge (H, \mathbf{new} \ (e_i) \ \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [\tau \ \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \ \sigma) \quad (\text{FU-R1})
\end{aligned}$$

From the evaluation rule we know that $H' = H'_1[a \mapsto v'_1]$ where $a \notin \text{dom}(H'_1)$, $v' = a$ and $n' = i + 1$. In order to prove FU-R0 we choose θ' as $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau \ \sigma\})$. Now we are required to show

$$(a) \ \theta \sqsubseteq \theta'_2 \wedge (n - i - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - 1, v') \in [(\mathbf{ref} \ \tau)^\perp \ \sigma]_V:$$

- $\theta \sqsubseteq \theta'_2$:

From FU-R1 we know that $\theta \sqsubseteq \theta'_1$ therefore from Definition 2.2 $\theta \sqsubseteq \theta'_2$

- $(n - i - 1, H') \triangleright \theta'_2$:

From FU-R1 we know that $(n - i, H'_1) \triangleright \theta'_1$. Therefore from Lemma 2.20 we get $(n - i - 1, H'_1) \triangleright \theta'_1$.

We also know that $(\theta'_1, n - i, v'_1) \in [\tau \ \sigma]_V$ (from FU-R1) therefore from Lemma 2.16 we get $(\theta'_1, n - i - 1, v'_1) \in [\tau \ \sigma]_V$

Since $H' = H'_1[a \mapsto v'_1]$ and $\theta'_2 = (\theta'_1 \cup \{a \mapsto \tau \ \sigma\})$ therefore from Definition 2.8 we get $(n - i - 1, H') \triangleright \theta'_2$

- $(\theta'_2, n - i - 1, a) \in [(\mathbf{ref} \ \tau)^\perp \ \sigma]_V$:

Since $\theta'_2(a) = \tau \ \sigma$ therefore from Definition 2.6 we get $(\theta'_2, n - i - 1, a) \in [(\mathbf{ref} \ \tau)^\perp \ \sigma]_V$

$$(b) \ (\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \ \sigma \sqsubseteq \ell')$$

From FU-R1

$$(c) \ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc \ \sigma):$$

We get this from FU-R1 and $\tau \ \sigma \searrow pc \ \sigma$ (given)

11. FG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\mathbf{ref} \ \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e_i : \tau'}$$

To prove: $(\theta, n, (!e_i) \ \delta) \in [\tau' \ \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (!e_i) \ \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau' \ \sigma]_V \wedge
\end{aligned}$$

$$\begin{aligned}
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap H and n s.t. $(n, H) \triangleright \theta \wedge (H, !(e_i) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau' \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-D0})
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\text{ref } \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)
\end{aligned}$$

Instantiating IH1 with H and n . Since we know that $(n, H) \triangleright \theta \wedge (H, !(e_i) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\text{ref } \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-D1})
\end{aligned}$$

In order to prove FU-D0 we choose θ' as θ'_1 from FU-D1. Also we know from the evaluation rule, that $H' = H'_1$, $v' = H'_1(a)$, $v'_1 = a$ and $n' = i + 1$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_1 \wedge (n - i - 1, H') \triangleright \theta'_1 \wedge (\theta'_1, n - i - 1, v') \in [\tau \sigma]_V$:
- $\theta \sqsubseteq \theta'_1$:
From FU-D1
 - $(n - i - 1, H') \triangleright \theta'_1$:
From FU-D1 we know that $(n - i, H') \triangleright \theta'_1$ therefore from Lemma 2.20 we get $(n - i - 1, H') \triangleright \theta'_1$
 - $(\theta'_1, n - i - 1, v') \in [\tau \sigma]_V$:
Since from FU-D1 we know that $(n - i, H'_1) \triangleright \theta'_1$ therefore from the Definition 2.8 we get $(\theta'_1, n - i, H'_1(a)) \in [\tau \sigma]_V$
From Lemma 2.16 we get $(\theta'_1, n - i - 1, H'_1(a)) \in [\tau \sigma]_V$
Since $\tau \sigma <: \tau' \sigma$ therefore from Lemma 2.24 we get $(\theta'_1, n - i - 1, H'_1(a)) \in [\tau' \sigma]_V$
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$
From FU-D1
- (c) $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)$
From FU-D1

12. FG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}}$$

To prove: $(\theta, n, (e_1 := e_2) \delta) \in [\text{unit } \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned}
& \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v') \implies \\
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\mathbf{unit}]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)
\end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned}
& \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\mathbf{unit}]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \quad (\text{FU-A0})
\end{aligned}$$

IH1:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_1) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\mathbf{ref} \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc)
\end{aligned}$$

Instantiating IH1 with H and n . Since we know that $(n, H) \triangleright \theta \wedge (H, (e_1 := e_2) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\mathbf{ref} \tau)]^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \quad (\text{FU-A1})
\end{aligned}$$

IH2:

$$\begin{aligned}
& \forall H_2, n_2. (n_2, H_2) \triangleright \theta'_1 \wedge \forall j < n_2. (H_2, (e_2) \delta) \Downarrow_j (H'_2, v'_2) \implies \\
& \exists \theta'_2. \theta'_1 \sqsubseteq (n_2 - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n_2 - j, v'_2) \in [(\tau) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc)
\end{aligned}$$

Instantiating IH2 with H'_1 and since we know that $H'_1 \triangleright \theta'_1 \wedge (H, (e_1 := e_2) \delta) \Downarrow (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_2. \theta'_1 \sqsubseteq (n - i - j, \theta'_2) \wedge H'_2 \triangleright \theta'_2 \wedge (\theta'_2, n - i - j, v'_2) \in [(\tau) \sigma]_V \wedge \\
& (\forall a. H_2(a) \neq H'_2(a) \implies \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta'_1). \theta'_2(a) \searrow pc) \quad (\text{FU-A2})
\end{aligned}$$

In order to prove FU-A0 we choose θ' as θ'_2 from FU-A2. Also we know from the evaluation rule, assign, that let $v'_1 = a_1$, $H' = H'_2[a_1 \mapsto v'_2]$, $v' = ()$ and $n' = i + j + 1$. Now we are required to show

$$(a) \theta \sqsubseteq \theta'_2 \wedge (n - i - j - 1, H') \triangleright \theta'_2 \wedge (\theta'_2, n - i - j - 1, ()) \in [\mathbf{unit}]_V:$$

- $\theta \sqsubseteq \theta'_2$:
Since $\theta \sqsubseteq \theta'_1$ from FU-A1 and $\theta'_1 \sqsubseteq \theta'_2$ from FU-A2 therefore from Definition 2.2 we get $\theta \sqsubseteq \theta'_2$
- $(n - i - j - 1, H') \triangleright \theta'_2$:
From Definition 2.8 it suffices to prove that
 - i. $\text{dom}(\theta'_2) \subseteq \text{dom}(H')$: From FU-A2
 - ii. $\forall a \in \text{dom}(\theta'_2). (\theta'_2, n - i - j - 1, H'(a)) \in [\theta'_2(a)]_V$:
 $\forall a \in \text{dom}(\theta'_2).$

- $a = a_1$:
From FU-A2 (since we know that $(\theta'_2, n - i - j, v'_2) \in \llbracket (\tau) \sigma \rrbracket_V$)
Therefore from Lemma 2.16 we get $(\theta'_2, n - i - j - 1, v'_2) \in \llbracket (\tau) \sigma \rrbracket_V$
 - $a \neq a_1$:
From FU-A2 (since we know that $(n - i - j, H'_2) \triangleright \theta'_2$ therefore from Lemma 2.20
we get $(n - i - j - 1, H'_2) \triangleright \theta'_2$)
 - $(\theta'_2, n - i - j - 1, ()) \in \llbracket \text{unit} \rrbracket_V$:
From Definition 2.6
- (b) $(\forall a \in \text{dom}(H). H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$
 $\forall a \in \text{dom}(H)$.
- $a = a_1$:
Since we know that $H(a_1) \neq H'(a_1)$ and $\theta(a_1) = \tau = \mathbf{A}^{\ell_i}$ (given)
It is given that $\tau \sigma \searrow pc \sigma$ therefore $pc \sigma \sqsubseteq \ell_i \sigma$
 - $a \neq a_1$:
From FU-A2
- (c) $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta). \theta'_2(a) \searrow pc)$
From FU-A2

13. FG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e_i : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_i : (\forall \alpha. (\ell_e, \tau))^\perp}$$

To prove: $(\theta, n, (\Lambda e_i) \delta) \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (\Lambda e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, (\Lambda e_i) \delta) \Downarrow (H', v')$

It suffices to prove

$$\begin{aligned} & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-FI0}) \end{aligned}$$

IH1:

$$\forall n_1, \theta_i, \ell' \in \mathcal{L}. (\theta_i, n_1, e_i \delta) \in \llbracket \tau \sigma \cup \{\alpha \mapsto \ell''\} \rrbracket_E^{\ell_e \sigma \cup \{\alpha \mapsto \ell''\}}$$

In order to prove FU-FI0 we choose θ' as θ . Also we know from the evaluation rule, that $H' = H$ and $n' = 0$. Now we are required to show

- (a) $\theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in \llbracket (\forall \alpha. (\ell_e, \tau))^\perp \sigma \rrbracket_V$:
- $\theta \sqsubseteq \theta$: From Definition 2.2
 - $(n, H) \triangleright \theta$: Given

- $(\theta, n, (\Lambda e_i)\delta) \in [(\forall\alpha.(\ell_e, \tau))^\perp]_V \sigma$:

From Definition 2.6 it suffices to prove that

$$\forall\theta''.\theta \sqsubseteq \theta'' \wedge \forall j < n.\forall\ell_d \in \mathcal{L} \implies (\theta'', j, e_i) \in [\tau[\ell_d/\alpha] \sigma]_E^{\ell_e[\ell_d/\alpha]} \sigma$$

This means given some θ'', j and ℓ_d such that $\theta \sqsubseteq \theta'', j < n$ and $\ell_d \in \mathcal{L}$

It suffices to prove that $(\theta'', j, e_i) \in [\tau[\ell_d/\alpha] \sigma]_E^{\ell_e[\ell_d/\alpha]} \sigma$

Instantiating IH1 with j, θ'' and ℓ_d we get $(\theta_i, j, e_i \delta) \in [\tau \sigma \cup \{\alpha \mapsto \ell_d\}]_E^{\ell_e} \sigma \cup \{\alpha \mapsto \ell_d\}$

- (b) $(\forall a.H(a) \neq H'(a) \implies \exists\ell'.\theta.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$:

Since $H' = H$ so we are done

- (c) $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta).\theta(a) \searrow pc)$:

Since $\theta' = \theta$ so we are done

14. FG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\forall\alpha.(\ell_e, \tau))^\ell \quad \ell'' \in \text{FV}(\Sigma) \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell''/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell''/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_i [] : \tau[\ell''/\alpha]}$$

To prove: $(\theta, n, (e_i[]) \delta) \in [\tau[\ell''/\alpha] \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n.(n, H) \triangleright \theta \wedge \forall n' < n.(H, (e_i[]) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists\theta'.\theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a.H(a) \neq H'(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta).\theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, (e_i[]) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} & \exists\theta'.\theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a.H(a) \neq H'(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta).\theta'(a) \searrow pc \sigma) \quad (\text{FU-FE0}) \end{aligned}$$

IH:

$$\begin{aligned} & \forall H_1, n_1.(n_1, H_1) \triangleright \theta \wedge \forall i < n_1.(H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\ & \exists\theta'_1.\theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(\forall\alpha.(\ell_e, \tau))^\ell \sigma]_V \wedge \\ & (\forall a.H_1(a) \neq H'_1(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta).\theta'_1(a) \searrow pc \sigma) \end{aligned}$$

Instantiating IH with H and n . Since we know that $(n, H) \triangleright \theta \wedge (H, (e_i[]) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned} & \exists\theta'_1.\theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(\forall\alpha.(\ell_e, \tau))^\ell \sigma]_V \wedge \\ & (\forall a.H_1(a) \neq H'_1(a) \implies \exists\ell'.\theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta).\theta'_1(a) \searrow pc \sigma) \quad (\text{FU-FE1}) \end{aligned}$$

From evaluation rule we know that $v'_1 = \Lambda e_{i1}$. Since from FU-FE1 we know that

$$(\theta'_1, n - i, \Lambda e_{i1}) \in [(\forall\alpha.(\ell_e, \tau))^\ell \sigma]_V$$

This means from Definition 2.6 we have

$$\forall \theta'' . \theta'_1 \sqsubseteq \theta'' \wedge \forall j < n - i . \forall \ell_g \in \mathcal{L} \implies (\theta'' , j , e_{i1}) \in [\tau[\ell_g/\alpha] \sigma]_E^{\ell_e[\ell_g/\alpha]} \sigma \quad (73)$$

Instantiating Equation 73 with θ'_1 , $n - i - 1$ and ℓ'' we get

$$(\theta'_1, n - i - 1, e_{i1}) \in [\tau[\ell''/\alpha] \sigma]_E^{\ell_e[\ell''/\alpha]} \sigma$$

This means from Definition 2.7 we have

$$\begin{aligned} & \forall H_3 . (n - i - 1, H_3) \triangleright \theta'_1 \wedge \forall k < n - i - 1 . (H_3, e_{i1}) \Downarrow_k (H'_3, v'_3) \implies \\ & \exists \theta'_3 . \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a . H_3(a) \neq H'_3(a) \implies \exists \ell' . \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1) . \theta'_3(a) \searrow \ell_e \sigma) \end{aligned}$$

Instantiating H_3 with H'_1 from FU-FE1 and since we know that $(n - i - 1, H'_1) \triangleright \theta'_1$ (Lemma 2.20) and since we know that $e_i \Downarrow \gamma \downarrow_1$ reduces in n' steps where $n' = i + k + 1$ and since $n' < n$ therefore we have $k < n - i - 1$ s.t. $(H'_1, e_{i1}) \Downarrow_k (H'_3, v'_3)$. Therefore we get

$$\begin{aligned} & \exists \theta'_3 . \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a . H_3(a) \neq H'_3(a) \implies \exists \ell' . \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1) . \theta'_3(a) \searrow \ell_e \sigma) \quad (\text{FU-FE2}) \end{aligned}$$

In order to prove FU-FE0 we choose θ' as θ'_3 from FU-FE2. Also we know that $H' = H'_3$, $v' = v'_3$ and $n' = i + k + 1$. Now we are required to show

- (a) $\theta \sqsubseteq \theta'_3 \wedge (n - i - k - 1, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$:
- $\theta \sqsubseteq \theta'_3$:
Since $\theta \sqsubseteq \theta'_1$ from FU-FE1 and $\theta'_1 \sqsubseteq \theta'_3$ from FU-FE2 therefore from Definition 2.2 we get $\theta \sqsubseteq \theta'_3$
 - $(n - i - k - 1, H'_3) \triangleright \theta'_3$:
From FU-FE2 we know that $(n - i - k - 1, H'_3) \triangleright \theta'_3$
 - $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$:
From FU-FE2 we know that $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$
- (b) $(\forall a \in \text{dom}(H) . H(a) \neq H'_3(a) \implies \exists \ell' . \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$
Since $pc \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$ therefore we get the desired from FU-FE1 and FU-FE2
- (c) $(\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta) . \theta'_3(a) \searrow pc \sigma)$
Since $pc \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$ therefore we get the desired from FU-FE1 and FU-FE2

15. FG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e_i : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e_i : (c \xrightarrow{\ell_e} \tau)^\perp}$$

To prove: $(\theta, n, (\nu e_i) \delta) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} & \forall H, n . (n, H) \triangleright \theta \wedge \forall n' < n . (H, (\nu e_i) \delta) \Downarrow_{n'} (H', v') \implies \\ & \exists \theta' . \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V \wedge \\ & (\forall a . H(a) \neq H'(a) \implies \exists \ell' . \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta) . \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, (\nu e_i) \delta) \Downarrow (H', v')$

It suffices to prove

$$\begin{aligned} \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-CI0}) \end{aligned}$$

IH1:

$$\forall \theta_i, n_1. (\theta_i, n_1, e_i \delta) \in [\tau \sigma]_E^{\ell_e} \sigma \text{ such that } \mathcal{L} \models c \sigma$$

In order to prove FU-FI0 we choose θ' as θ . Also we know from the evaluation rule, that $H' = H$, $v' = \nu e_i \delta$ and $n' = 0$. Now we are required to show

$$(a) \theta \sqsubseteq \theta \wedge (n, H) \triangleright \theta \wedge (\theta, n, v') \in [(c \xrightarrow{\ell_e} \tau)^\perp]_V \sigma:$$

• $\theta \sqsubseteq \theta$: From Definition 2.2

• $(n, H) \triangleright \theta$: Given

• $(\theta, n, (\nu e_i) \delta) \in [(c \xrightarrow{\ell_e} \tau)^\perp]_V \sigma$:

From Definition 2.6 it suffices to prove that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < n. \mathcal{L} \models c \sigma \implies (\theta'', j, e_i \delta) \in [\tau \sigma]_E^{\ell_e} \sigma$$

This means given some θ'' such that $\theta \sqsubseteq \theta''$, $j < n$ and $\mathcal{L} \models c$

It suffices to prove that $(\theta'', j, e_i \delta) \in [\tau \sigma]_E^{\ell_e} \sigma$

Instantiating IH1 with θ'' and j we get $(\theta'', j, e_i \delta) \in [\tau \sigma]_E^{\ell_e} \sigma$

$$(b) (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell'):$$

Since $H' = H$ so we are done

$$(c) (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta(a) \searrow pc):$$

Since $\theta' = \theta$ so we are done

16. FG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_i \bullet : \tau}$$

To prove: $(\theta, n, (e_i \bullet) \delta) \in [\tau \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\begin{aligned} \forall H, n. (n, H) \triangleright \theta \wedge \forall n' < n. (H, (e_i \bullet) \delta) \Downarrow_{n'} (H', v') \implies \\ \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \end{aligned}$$

This means that given some heap H and n s.t $(n, H) \triangleright \theta \wedge (H, (e_i \bullet) \delta) \Downarrow_{n'} (H', v')$

It suffices to prove

$$\begin{aligned} \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - n', H') \triangleright \theta' \wedge (\theta', n - n', v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc \sigma) \quad (\text{FU-CE0}) \end{aligned}$$

IH:

$$\begin{aligned}
& \forall H_1, n_1. (n_1, H_1) \triangleright \theta \wedge \forall i < n_1. (H_1, (e_i) \delta) \Downarrow_i (H'_1, v'_1) \implies \\
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n_1 - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n_1 - i, v'_1) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma)
\end{aligned}$$

Instantiating IH with H and n . And since we know that $(n, H) \triangleright \theta \wedge (H, (e_i) \delta) \Downarrow_{n'} (H', v')$ therefore we have

$$\begin{aligned}
& \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - i, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - i, v'_1) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V \wedge \\
& (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc \sigma) \quad (\text{FU-CE1})
\end{aligned}$$

From evaluation rule we know that $v'_1 = \nu e_{i1}$. Since from FU-CE1 we know that

$$(\theta'_1, n - i, \nu e_{i1}) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_V$$

This means from Definition 2.6 we have

$$\forall \theta''. \theta'_1 \sqsubseteq \theta'' \wedge \forall j < n - i. \mathcal{L} \models c \sigma \implies (\theta'', j, e_{i1}) \in [\tau \sigma]_E^{\ell_e \sigma} \quad (74)$$

Instantiating Equation 74 with θ'_1 and $n - i - 1$ since we know that $\mathcal{L} \models c \sigma$ therefore we get

$$(\theta'_1, n - i - 1, e_{i1}) \in [\tau \sigma]_E^{\ell_e \sigma}$$

This means from Definition 2.7 we have

$$\begin{aligned}
& \forall H_3. (n - i - 1, H_3) \triangleright \theta'_1 \wedge \forall k < n - i - 1. (H_3, e_{i1}) \Downarrow_k (H'_3, v'_3) \implies \\
& \exists \theta'_3. \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau \sigma]_V \wedge (\forall a. H_3(a) \neq H'_3(a) \implies \\
& \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1). \theta'_3(a) \searrow \ell_e \sigma)
\end{aligned}$$

Instantiating H_3 with H'_1 from FU-CE1 and since we know that $(n - i - 1, H'_1) \triangleright \theta'_1$ (Lemma 2.20) and since we know that $e_i \bullet \gamma \Downarrow_1$ reduces in n' steps where $n' = i + k + 1$ and since $n' < n$ therefore we have $k < n - i - 1$ s.t. $(H'_1, e_{i1}) \Downarrow_k (H'_3, v'_3)$. Therefore we get

$$\begin{aligned}
& \exists \theta'_3. \theta'_1 \sqsubseteq \theta'_3 \wedge (n - i - 1 - k, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - 1 - k, v'_3) \in [\tau \sigma]_V \wedge (\forall a. H_3(a) \neq H'_3(a) \implies \\
& \exists \ell'. \theta'_1(a) = \mathbf{A}^{\ell'} \wedge \ell_e \sigma \sqsubseteq \ell') \wedge (\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta'_1). \theta'_3(a) \searrow \ell_e \sigma) \quad (\text{FU-CE2})
\end{aligned}$$

In order to prove FU-CE0 we choose θ' as θ'_3 from FU-CE2. Also we know that $H' = H'_3$, $v' = v'_3$ and $n' = i + k + 1$. Now we are required to show

$$(a) \theta \sqsubseteq \theta'_3 \wedge (n - i - k - 1, H'_3) \triangleright \theta'_3 \wedge (\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V:$$

$$\bullet \theta \sqsubseteq \theta'_3:$$

Since $\theta \sqsubseteq \theta'_1$ from FU-CE1 and $\theta'_1 \sqsubseteq \theta'_3$ from FU-CE2 therefore from Definition 2.2 we get $\theta \sqsubseteq \theta'_3$

$$\bullet (n - i - k - 1, H'_3) \triangleright \theta'_3:$$

From FU-CE3 we know that $(n - i - k - 1, H'_3) \triangleright \theta'_3$

$$\bullet (\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V:$$

From FU-CE3 we know that $(\theta'_3, n - i - k - 1, v'_3) \in [\tau[\ell''/\alpha] \sigma]_V$

$$(b) (\forall a \in \text{dom}(H). H(a) \neq H'_3(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sigma \sqsubseteq \ell')$$

Since $pc \sigma \sqsubseteq \ell_e \sigma$ therefore we get the desired from FU-CE1 and FU-CE2

(c) $(\forall a \in \text{dom}(\theta'_3) \setminus \text{dom}(\theta)).\theta'_3(a) \searrow pc \sigma$

Since $pc \sigma \sqsubseteq \ell_e \sigma$ therefore we get the desired from FU-CE1 and FU-CE2

□

Lemma 2.23 (FG: Expression subtyping with closed labels and types). $\forall pc, pc', \tau$.

$$\mathcal{L} \models pc \sqsubseteq pc' \implies \lfloor \tau \rfloor_E^{pc'} \subseteq \lfloor \tau \rfloor_E^{pc}$$

Proof. Given: $\mathcal{L} \models pc \sqsubseteq pc'$

$$\text{To prove: } \lfloor (\tau) \rfloor_E^{pc'} \subseteq \lfloor (\tau) \rfloor_E^{pc}$$

This means we need to prove that

$$\forall (\theta, n, e) \in \lfloor (\tau) \rfloor_E^{pc'} . (\theta, n, e) \in \lfloor (\tau) \rfloor_E^{pc}$$

This means given $\forall (\theta, n, e) \in \lfloor (\tau) \rfloor_E^{pc'}$

It suffices to prove that $(\theta, n, e) \in \lfloor (\tau) \rfloor_E^{pc}$

From Definition 2.7 for the chosen θ, n, e we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall j < n.(H, e) \Downarrow_j (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - j, H') \triangleright \theta' \wedge (\theta', n - j, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta)).\theta'(a) \searrow pc' \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall k < n.(H_1, e) \Downarrow_k (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta)).\theta'_1(a) \searrow pc \end{aligned}$$

This means that we are given some H_1 and k such that $(n, H_1) \triangleright \theta$, $k < n$ and $(H_1, e) \Downarrow_k (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - k, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - k, v') \in \lfloor \tau \rfloor_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta)).\theta'_1(a) \searrow pc \end{aligned}$$

Instantiate H in (A) with H_1 and then we choose θ'_1 as θ'

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n - k, H'_1) \triangleright \theta' \wedge (\theta', n - k, v') \in \lfloor \tau \rfloor_V$:

Given

- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$:

Since $pc \sqsubseteq pc'$ and we are given

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc' \sqsubseteq \ell')$$

Therefore

$$(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$$

- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta)).\theta'(a) \searrow pc$:

We are given

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta)).\theta'(a) \searrow pc'$$

and since $pc \sqsubseteq pc'$ Therefore

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta)).\theta'(a) \searrow pc$$

□

Lemma 2.24 (FG: Subtyping unary). *The following holds:*

$\forall \Sigma, \Psi, \sigma.$

1. $\forall A, A'.$

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies [(A \sigma)]_V \subseteq [(A' \sigma)]_V$$

2. $\forall \tau, \tau'.$

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$$

$$(b) \forall pc. \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^{pc} \subseteq [(\tau' \sigma)]_E^{pc}$$

Proof. Proof by simultaneous induction on $A <: A'$ and $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of $A <: A'$ in the last step:

1. FGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

$$\text{To prove: } [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V \subseteq [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V$$

$$\text{IH1: } [(\tau'_1 \sigma)]_V \subseteq [(\tau_1 \sigma)]_V \text{ (Statement 2(a))}$$

$$\text{IH2: } \forall pc. [(\tau_2 \sigma)]_E^{pc} \subseteq [(\tau'_2 \sigma)]_E^{pc} \text{ (Statement 2(b))}$$

$$\text{It suffices to prove: } \forall (\theta, n, \lambda x.e_i) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V. (\theta, n, \lambda x.e_i) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V$$

This means that given some θ, n and $\lambda x.e_i$ s.t $(\theta, n, \lambda x.e_i) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V$

Therefore from Definition 2.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in [\tau_1 \sigma]_V \implies (\theta_1, i, e_i[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma} \quad (75)$$

$$\text{And it suffices to prove: } (\theta, n, \lambda x.e_i) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V$$

Again from Definition 2.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in [\tau'_1 \sigma]_V \implies (\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$$

This means that given some $\theta_2, j < n, v$ s.t $\theta \sqsubseteq \theta_2$ and $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$

$$\text{And we are required to prove: } (\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$$

Since $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$ therefore from IH1 we know that $(\theta_2, j, v) \in [\tau_1 \sigma]_V$

As a result from Equation 75 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

From IH2, we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E^{\ell_e \sigma}$$

Since $\mathcal{L} \models \ell'_e \sigma \sqsubseteq \ell_e \sigma$ therefore from Lemma 2.23 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$$

2. FGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove: $[\lceil((\tau_1 \times \tau_2) \sigma)\rceil]_V \subseteq [\lceil((\tau'_1 \times \tau'_2) \sigma)\rceil]_V$

IH1: $[\lceil(\tau_1 \sigma)\rceil]_V \subseteq [\lceil(\tau'_1 \sigma)\rceil]_V$ (Statement 2(a))

IH2: $[\lceil(\tau_2 \sigma)\rceil]_V \subseteq [\lceil(\tau'_2 \sigma)\rceil]_V$ (Statement 2(a))

It suffices to prove: $\forall(\theta, n, (v_1, v_2)) \in [\lceil((\tau_1 \times \tau_2) \sigma)\rceil]_V. (\theta, n, (v_1, v_2)) \in [\lceil((\tau'_1 \times \tau'_2) \sigma)\rceil]_V$

This means that given some θ, n and (v_1, v_2) $(\theta, (v_1, v_2)) \in [\lceil((\tau_1 \times \tau_2) \sigma)\rceil]_V$

Therefore from Definition 2.6 we are given:

$$(\theta, n, v_1) \in [\tau_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau_2 \sigma]_V \quad (76)$$

And it suffices to prove: $(\theta, (v_1, v_2)) \in [\lceil((\tau'_1 \times \tau'_2) \sigma)\rceil]_V$

Again from Definition 2.6, it suffices to prove:

$$(\theta, n, v_1) \in [\tau'_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau'_2 \sigma]_V$$

Since from Equation 76 we know that $(\theta, n, v_1) \in [\tau_1 \sigma]_V$ therefore from IH1 we have $(\theta, n, v_1) \in [\tau'_1 \sigma]_V$

Similarly since $(\theta, n, v_2) \in [\tau_2 \sigma]_V$ from Equation 76 therefore from IH2 we have $(\theta, n, v_2) \in [\tau'_2 \sigma]_V$

3. FGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove: $[\lceil((\tau_1 + \tau_2) \sigma)\rceil]_V \subseteq [\lceil((\tau'_1 + \tau'_2) \sigma)\rceil]_V$

IH1: $[\lceil(\tau_1 \sigma)\rceil]_V \subseteq [\lceil(\tau'_1 \sigma)\rceil]_V$ (Statement 2(a))

IH2: $[\lceil(\tau_2 \sigma)\rceil]_V \subseteq [\lceil(\tau'_2 \sigma)\rceil]_V$ (Statement 2(a))

It suffices to prove: $\forall(\theta, n, v_s) \in [\lceil((\tau_1 + \tau_2) \sigma)\rceil]_V. (\theta, v_s) \in [\lceil((\tau'_1 + \tau'_2) \sigma)\rceil]_V$

This means that given: $(\theta, n, v_s) \in [\lceil((\tau_1 + \tau_2) \sigma)\rceil]_V$

And it suffices to prove: $(\theta, n, v_s) \in [\lceil((\tau'_1 + \tau'_2) \sigma)\rceil]_V$

2 cases arise

(a) $v_s = \text{inl } v_i$:

From Definition 2.6 we are given:

$$(\theta, n, v_i) \in \lfloor \tau_1 \sigma \rfloor_V \quad (77)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \lfloor \tau'_1 \sigma \rfloor_V$$

From Equation 77 and IH1 we know that

$$(\theta, n, v_i) \in \lfloor \tau'_1 \sigma \rfloor_V$$

(b) $v_s = \text{inr } v_i$:

From Definition 2.6 we are given:

$$(\theta, n, v_i) \in \lfloor \tau_2 \sigma \rfloor_V \quad (78)$$

And we are required to prove that:

$$(\theta, n, v_i) \in \lfloor \tau'_2 \sigma \rfloor_V$$

From Equation 78 and IH2 we know that

$$(\theta, n, v_i) \in \lfloor \tau'_2 \sigma \rfloor_V$$

4. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

To prove: $\lfloor ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rfloor_V \subseteq \lfloor (\forall \alpha. (\ell'_e, \tau_2)) \sigma \rfloor_V$

IH1: $\forall pc. \lfloor (\tau_1 \sigma) \rfloor_E^{pc} \subseteq \lfloor (\tau_2 \sigma) \rfloor_E^{pc}$ (Statement 2(b))

It suffices to prove: $\forall (\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rfloor_V. (\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rfloor_V$

This means that given: $(\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rfloor_V$

Therefore from Definition 2.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall \ell' \in \mathcal{L} \implies (\theta_1, i, e_i) \in \lfloor \tau_1 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])} \quad (79)$$

And it suffices to prove: $(\theta, n, \Lambda e_i) \in \lfloor ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rfloor_V$

Again from Definition 2.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall \ell' \in \mathcal{L} \implies (\theta_2, j, e_i) \in \lfloor \tau_2 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$$

This means that given some $\theta_2, j < n, \ell' \in \mathcal{L}$ s.t $\theta \sqsubseteq \theta_2$

And we are required to prove: $(\theta_2, j, e_i) \in \lfloor \tau_2 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$

Since we are given $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \ell' \in \mathcal{L}$ therefore from Equation 79 we have

$$(\theta_2, j, e_i) \in \lfloor \tau_1 (\sigma \cup [\alpha \mapsto \ell']) \rfloor_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$$

From IH1, we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E^{\ell_e (\sigma \cup [\alpha \mapsto \ell'])}$$

Since $\mathcal{L} \models \ell'_e (\sigma \cup [\alpha \mapsto \ell']) \sqsubseteq \ell_e (\sigma \cup [\alpha \mapsto \ell'])$ therefore from Lemma 2.23 we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E^{\ell'_e (\sigma \cup [\alpha \mapsto \ell'])}$$

5. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell_e} \tau_1 <: c_2 \xrightarrow{\ell'_e} \tau_2} \text{FGsub-constraint}$$

$$\text{To prove: } [((c_1 \xrightarrow{\ell_e} \tau_1) \sigma)]_V \subseteq [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V$$

$$\text{IH1: } \forall pc. [(\tau_1 \sigma)]_E^{pc} \subseteq [(\tau_2 \sigma)]_E^{pc} \text{ (Statement 2(b))}$$

$$\text{It suffices to prove: } \forall (\theta, n, \nu e_i) \in [((c_1 \xrightarrow{\ell_e} \tau_1) \sigma)]_V. (\theta, n, \nu e_i) \in [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V$$

$$\text{This means that given: } (\theta, n, \nu e_i) \in [((c_1 \xrightarrow{\ell_e} \tau_1) \sigma)]_V$$

Therefore from Definition 2.6 we are given:

$$\forall \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \mathcal{L} \models c_1 \sigma \implies (\theta_1, i, e_i) \in [\tau_1 (\sigma)]_E^{\ell_e \sigma} \quad (80)$$

$$\text{And it suffices to prove: } (\theta, n, \nu e_i) \in [((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma)]_V$$

Again from Definition 2.6, it suffices to prove:

$$\forall \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \mathcal{L} \models c_2 \sigma \implies (\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell'_e \sigma}$$

This means that given some θ_2, j s.t $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$

$$\text{And we are required to prove: } (\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell'_e \sigma}$$

Since we are given $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$ therefore from Equation 80 we have

$$(\theta_2, j, e_i) \in [\tau_1 (\sigma)]_E^{\ell_e \sigma}$$

From IH1, we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell_e \sigma}$$

Since $\mathcal{L} \models \ell'_e \sigma \sqsubseteq \ell_e \sigma$ therefore from Lemma 2.23 we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E^{\ell'_e \sigma}$$

6. FGsub-ref:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

$$\text{To prove: } [((\text{ref } \tau) \sigma)]_V \subseteq [((\text{ref } \tau) \sigma)]_V$$

$$\text{It suffices to prove: } \forall (\theta, n, a) \in [((\text{ref } \tau) \sigma)]_V. (\theta, n, a) \in [((\text{ref } \tau) \sigma)]_V$$

Trivial

7. FGsub-base:

Given:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove: $\llbracket ((\mathbf{b}) \sigma) \rrbracket_V \subseteq \llbracket ((\mathbf{b}) \sigma) \rrbracket_V$

Directly from Definition 2.6

8. FGsub-unit:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove: $\llbracket ((\text{unit}) \sigma) \rrbracket_V \subseteq \llbracket ((\text{unit}) \sigma) \rrbracket_V$

Directly from Definition 2.6

Proof of statement 2(a)

Given:

$$\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash \mathbf{A} <: \mathbf{A}'}{\Sigma; \Psi \vdash \mathbf{A}^\ell <: \mathbf{A}^{\ell'}} \text{FGsub-label}$$

To prove: $\llbracket ((\mathbf{A}^\ell) \sigma) \rrbracket_V \subseteq \llbracket ((\mathbf{A}^{\ell'}) \sigma) \rrbracket_V$

From Definition 2.6 it suffices to prove: $\llbracket ((\mathbf{A}) \sigma) \rrbracket_V \subseteq \llbracket ((\mathbf{A}') \sigma) \rrbracket_V$

This we get directly from IH (Statement 1(a))

Proof of statement 2(b)

Given: $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma$

To prove: $\llbracket (\tau \sigma) \rrbracket_E^{pc} \subseteq \llbracket (\tau' \sigma) \rrbracket_E^{pc}$

This means we need to prove that

$$\forall (\theta, n, e) \in \llbracket (\tau \sigma) \rrbracket_E^{pc}. (\theta, n, e) \in \llbracket (\tau' \sigma) \rrbracket_E^{pc}$$

This means given $(\theta, n, e) \in \llbracket (\tau \sigma) \rrbracket_E^{pc}$

It suffices to prove that $(\theta, n, e) \in \llbracket (\tau' \sigma) \rrbracket_E^{pc}$

From Definition 2.7 we know we are given:

$$\begin{aligned} & \forall H.(n, H) \triangleright \theta \wedge \forall i < n.(H, e) \Downarrow_i (H', v') \implies \\ & \exists \theta'. \theta \sqsubseteq \theta' \wedge (n - i, H') \triangleright \theta' \wedge (\theta', n - i, v') \in \llbracket \tau \sigma \rrbracket_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc) \end{aligned} \quad (\text{A})$$

And we need prove that

$$\begin{aligned} & \forall H_1.(n, H_1) \triangleright \theta \wedge \forall j < n.(H_1, e) \Downarrow_j (H'_1, v') \implies \\ & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n - j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n - j, v') \in \llbracket \tau' \sigma \rrbracket_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

This means that we are given some H_1 and $j < n$ s.t $(n, H_1) \triangleright \theta \wedge (H_1, e) \Downarrow_j (H'_1, v')$

It suffices to prove:

$$\begin{aligned} & \exists \theta'_1. \theta \sqsubseteq \theta'_1 \wedge (n-j, H'_1) \triangleright \theta'_1 \wedge (\theta'_1, n-j, v') \in [\tau' \sigma]_V \wedge \\ & (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta). \theta'_1(a) \searrow pc) \end{aligned}$$

Instantiate H in (A) with H_1 and i with j then we choose θ'_1 as θ'
Also we have IH1 as $[\tau \sigma]_V \subseteq [\tau' \sigma]_V$ (Statement 2(a))

- $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H'_1) \triangleright \theta' \wedge (\theta', n-j, v') \in [\tau' \sigma]_V$:
We are given $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H'_1) \triangleright \theta' \wedge (\theta', n-j, v') \in [\tau \sigma]_V$
From IH1 we know that $[\tau \sigma]_V \subseteq [\tau' \sigma]_V$
Therefore, $\exists \theta'. \theta \sqsubseteq \theta' \wedge (n-j, H'_1) \triangleright \theta' \wedge (\theta', n-j, v') \in [\tau' \sigma]_V$
- $(\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta(a) = \mathbf{A}^{\ell'} \wedge pc \sqsubseteq \ell')$:
Given
- $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta). \theta'(a) \searrow pc)$:
Given

□

Lemma 2.25 (FG: Binary interpretation of Γ implies Unary interpretation of Γ). $\forall W, \gamma, \Gamma, n.$
 $(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

Proof. Given: $(W, n, \gamma) \in [\Gamma]_V^A$

To prove: $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 2.14 we know that we are given:

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And we are required to prove:

$$\forall i \in \{1, 2\}. \forall m.$$

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i) \wedge \forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$$

Case $i = 1$

Given some m we need to show:

- $\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i)$:
 $\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_i)$
Therefore, $\text{dom}(\Gamma) \subseteq (\text{dom}(\gamma) = \text{dom}(\gamma \downarrow_i))$ (Given)
- $\forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$:
We are given: $\forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$
Therefore from Lemma 2.15 we know that
 $\forall m'. (W.\theta_i, m', \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$
Instantiating m' with m we get
 $(W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Case $i = 2$

Symmetric case as $i = 1$

□

Theorem 2.26 (FG: Fundamental theorem binary). $\forall \Sigma, \Psi, \Gamma, pc, W, \mathcal{A}, \mathcal{L}, e, \tau, \sigma, \gamma, n.$

$$\begin{aligned} \Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \wedge \mathcal{L} \models \Psi \sigma \wedge (W, n, \gamma) \in [\Gamma]_{\mathcal{V}}^{\mathcal{A}} &\implies \\ (W, n, e (\gamma \downarrow_1), e (\gamma \downarrow_2)) \in [\tau \sigma]_E^{\mathcal{A}} & \end{aligned}$$

Proof. Proof by induction on the typing derivation

1. FG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau} \text{FG-var}$$

To prove: $(W, n, x (\gamma \downarrow_1), x (\gamma \downarrow_2)) \in [\tau \sigma]_E^{\mathcal{A}}$

Say $e_1 = x (\gamma \downarrow_1)$ and $e_2 = x (\gamma \downarrow_2)$

From Definition of $[\tau]_E^{\mathcal{A}}$ it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^{\mathcal{A}} W \wedge \forall j < n. (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) &\implies \\ \exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}} & \end{aligned}$$

This means given some H_1, H_2 and j s.t $(n, H_1, H_2) \triangleright^{\mathcal{A}} W \wedge (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$

We are required to prove: $\exists W' \sqsubseteq W. (n - j, H'_1, H'_2) \triangleright^{\mathcal{A}} W' \wedge (W', n - j, v'_1, v'_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Here

- $H'_1 = H_1$ and $H'_2 = H_2$
- $e_1 = v'_1 = \gamma(x) \downarrow_1$
- $e_2 = v'_2 = \gamma(x) \downarrow_2$
- $j = 1$

We choose $W' = W$.

- $W \sqsubseteq W$: From Definition 2.3

- $(n - 1, H_1, H_2) \triangleright^{\mathcal{A}} W$:

Since we know that $(n, H_1, H_2) \triangleright^{\mathcal{A}} W$ therefore from Lemma 2.21 we get

$$(n - 1, H_1, H_2) \triangleright^{\mathcal{A}} W$$

- $(W, n - 1, \gamma(x) \downarrow_1, \gamma(x) \downarrow_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$:

We are given that $(W, n, \gamma) \in [\Gamma]_{\mathcal{V}}^{\mathcal{A}}$ therefore from Lemma 2.19 we get

$$(W, n - 1, \gamma) \in [\Gamma]_{\mathcal{V}}^{\mathcal{A}}$$

which means from Definition 2.14 we have

$$(W, n - 1, \gamma(x) \downarrow_1, \gamma(x) \downarrow_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$$

2. FG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e_i : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e_i : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp}$$

To prove: $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_E^A$

Say $e_1 = \lambda x. e (\gamma \downarrow_1)$ and $e_2 = \lambda x. e (\gamma \downarrow_2)$

From Definition of $[(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_E^A$ it suffices to prove that

$$\begin{aligned} \forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^A \end{aligned}$$

This means that given H_1, H_2 and j s.t $(n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$

It suffices to prove:

$$\exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^A \quad (\text{FB-L0})$$

IH1:

$$\forall W, n. (W, n, e (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2 \sigma]_E^A$$

s.t

$$(W, n, (v_1, v_2)) \in [\tau_1 \sigma]_V^A$$

We know from the evaluation rules that $H'_1 = H_1$, $H'_2 = H_2$, $v'_1 = e_1 = \lambda x. e (\gamma \downarrow_1)$, $v'_2 = e_2 = \lambda x. e (\gamma \downarrow_2)$ and $j = 0$. In order to prove FB-L0 we choose $W' = W$ and we need to prove the following:

- $W \sqsubseteq W$: From Definition 2.3
- $(n, H_1, H_2) \triangleright^A W$: Given
- $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^A$

From Definition 2.4 it suffices to prove that:

$$\forall W'' \sqsupseteq W, k < n, v_1, v_2.$$

$$((W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, k, v_c.$$

$$((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, k, e (\gamma \downarrow_1)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}) \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, v_c.$$

$$((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, k, e (\gamma \downarrow_2)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$$

This means that we need to prove the following:

$$\begin{aligned} - \forall W'' \sqsupseteq W, k < n, v_1, v_2. ((W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies \\ (W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A): \end{aligned}$$

This means given $W'' \sqsupseteq W, k < n, v_1, v_2$ s.t $((W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A$

We need to prove: $(W'', k, e (\gamma \downarrow_1)[v_1/x], e (\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A$

We instantiate IH1 with W'' and k
 And since $(W'', k, v_1, v_2) \in [\tau_1 \sigma]_V^A$ therefore we get
 $(W'', k, e(\gamma \downarrow_1)[v_1/x], e(\gamma \downarrow_2)[v_2/x]) \in [\tau_2 \sigma]_E^A$

– $\forall \theta_l \sqsupseteq W.\theta_1, k, v_c. ((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies$
 $(\theta_l, k, e(\gamma \downarrow_1)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$:

This means that we are given θ_l, k and v_c s.t

$\theta_l \sqsupseteq W.\theta_1$ and $(\theta_l, k, v_c) \in [\tau_1 \sigma]_V$

And we are required to prove:

$(\theta_l, k, e(\gamma \downarrow_1)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

It is given to us that

$\forall v_1, v_2. (W, n, \gamma \in [\Gamma]_V^A$

Therefore from Lemma 2.25 we know that

$\forall m. (W.\theta_1, m, (\gamma \downarrow_1) \in [\Gamma]_V$

Therefore, we can apply Theorem 2.22 to obtain

$\forall m. (W.\theta_1, m, \lambda x. e \gamma \downarrow_1) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V$

From Definition 2.6 it means that we have

$\forall m. \forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1 \sigma]_V \implies$

$(\theta', j, e[v/x]\gamma \downarrow_1) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

We instantiate m with some $l > k$, θ' with θ_l , j with k and v with v_c to get

$W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

Since we show that $W.\theta_1 \sqsubseteq \theta_l \wedge k < l \wedge (\theta_l, k, v_c) \in [\tau_1 \sigma]_V$ therefore we get

$(\theta_l, k, e[v_c/x]\gamma \downarrow_1) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

– $\forall \theta_l \sqsupseteq W.\theta_2, v_c. ((\theta_l, k, v_c) \in [\tau_1 \sigma]_V \implies$

$(\theta_l, k, e(\gamma \downarrow_2)[v_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$:

Symmetric case as above

3. FG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2}$$

To prove: $(W, n, (e_1 e_2)(\gamma \downarrow_1), (e_1 e_2)(\gamma \downarrow_2)) \in [(\tau_2) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$\forall H_1, H_2, n' < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow$
 $(H'_2, v'_2) \implies$

$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2) \sigma]_V^A$

This further means that given $H_1, H_2, n' < n$ s.t

$(n, H_1, H_2) \triangleright^A W \wedge (H_1, (e_1 e_2)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_1 e_2)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$

It suffices to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2) \sigma]_{\mathcal{V}}^A \quad (\text{FB-A0})$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}, i < n. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge (H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies \exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{\mathcal{V}}^A$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $(e_1 e_2)$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps. Therefore $\exists i < n' < n$ s.t $(H_{i1}, e_1 (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1)$. $(H_{i2}, e_1 (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$ is known because $(e_1 e_2)$ reduces to value with $\gamma \downarrow_2$. Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{\mathcal{V}}^A \quad (81)$$

$$\underline{\text{IH2}}: (W'_1, n - i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [(\tau_1) \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{j1}, H_{j2}, j < (n - i). (n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge (H_{j1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies \exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1) \sigma]_{\mathcal{V}}^A$$

Instantiating H_{j1} with H'_1 and H_{j2} with H'_2 in IH2. Since the $(e_1 e_2)$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps. Also, e_1 reduces to value $\gamma \downarrow_1$ in $i < n'$ steps. Therefore $\exists j < n' - i < n - i$ s.t $(H_{i1}, e_2 (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1})$. $(H_{i2}, e_2 (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$ is known because $(e_1 e_2)$ reduces to value with $\gamma \downarrow_2$. Hence we get

$$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau_1) \sigma]_{\mathcal{V}}^A \quad (82)$$

We case analyze on $(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \sigma]_{\mathcal{V}}^A$ from Equation 81

- Case $\ell \sigma \sqsubseteq \mathcal{A}$:

From Definition 2.4 we know that this would mean that

$$(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_{\mathcal{V}}^A$$

This means

$$(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 \sigma \xrightarrow{\ell_e \sigma} \tau_2 \sigma)]_{\mathcal{V}}^A$$

$$\text{Let } v'_1 = \lambda x.e_{h1} \text{ and } v'_2 = \lambda x.e_{h2}$$

Again from Definition 2.4 it means that

$$\forall W'_{h1} \sqsupseteq W'_1, j_1 < (n - i), v_1, v_2.$$

$$((W'_{h1}, j_1, v_1, v_2) \in [\tau_1 \sigma]_{\mathcal{V}}^A \implies (W'_{h1}, j_1, e_{h1}[v_1/x], e_{h2}[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_1, m_1, v_c.$$

$$\wedge ((\theta_{l1}, m_1, v_1) \in [\tau_1 \sigma]_{\mathcal{V}} \implies (W'_{h1}.\theta_1, e_{h1}[v_1/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}) \wedge$$

$$\forall \theta_{l1} \sqsupseteq W'_1.\theta_2, m_1, v_c.$$

$$\wedge (\theta_{l1}, m_1, v_2) \in [\tau_1 \sigma]_{\mathcal{V}} \implies (W'_{h1}.\theta_2, e_{h2}[v_2/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma})$$

We instantiate W'_{h1} with W'_2 obtained from Equation 82. Similarly we also instantiate v_1 and v_2 with v'_{j1} and v'_{j2} respectively from Equation 82, and j_1 with $n - i - j$. And we get

$$(W'_2, n - i - j, e_{h1}[v'_{j1}/x], e_{h2}[v'_{j2}/x]) \in [\tau_2 \sigma]_E^A$$

From Definition 2.5 we get

$$\begin{aligned} & \forall H_1, H_2, k_e < (n - i - j). (n - i - j, H_1, H_2) \triangleright^A W'_2 \wedge \\ & (H_1, e_{h1}[v'_{j1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1}) \wedge (H_2, e_{h2}[v'_{j2}/x]) \Downarrow (H'_{f2}, v_{f2}) \implies \\ & \exists W' \sqsupseteq W'_2. (n - i - j - k_e, H'_{f1}, H'_{f2}) \triangleright^A W' \wedge (W', n - i - j - k_e, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A \end{aligned}$$

Instantiating H_1 with H'_{j1} and H_2 with H'_{j2} obtained from Equation 82. And since we know that $e_1 e_2$ reduces with $\gamma \downarrow_1$ in $n' < n$ steps. And e_2 reduces to value $\gamma \downarrow_1$ in $j < n' - 1 < n - i$ steps. Therefore $\exists k_e = n' - i - j < n - i - j$ s.t $(H_1, e_{h1}[v'_{j1}/x]) \Downarrow_{k_e} (H'_{f1}, v_{f1})$. $(H_2, e_{h2}[v'_{j2}/x]) \Downarrow (H'_{f2}, v_{f2})$ is known because $(e_1 e_2)$ reduces to value with $\gamma \downarrow_2$. Hence we get

$$\exists W' \sqsupseteq W'_2. ((n - i - j - k_e), H'_{f1}, H'_{f2}) \triangleright^A W' \wedge (W', (n - i - j - k_e), v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A \quad (83)$$

This concludes the proof in this case.

- Case $\ell \sigma \not\sqsubseteq \mathcal{A}$:

From FB-A0 we know that we need to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_2) \sigma]_V^A$$

In this case since we know that $\ell \sigma \not\sqsubseteq \mathcal{A}$. Let $\tau_2 \sigma = A^{\ell_i}$ and since $\tau_2 \sigma \searrow \ell \sigma$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

Therefore from Definition 2.4 it will suffice to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau_2) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau_2) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau_2) \sigma]_V)$$

This means given m_1 and m_2 it suffices to prove:

$$(\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau_2) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau_2) \sigma]_V \quad (84)$$

In this case from Definition 2.6 we know that

$$\forall m. (W'_1.\theta_1, m, \lambda x. e_{h1}) \in [(\tau_1 \sigma \xrightarrow{\ell_e \sigma} \tau_2 \sigma)]_V \quad (85)$$

$$\forall m. (W'_1.\theta_2, m, \lambda x. e_{h2}) \in [(\tau_1 \sigma \xrightarrow{\ell_e \sigma} \tau_2 \sigma)]_V \quad (86)$$

Applying Definition 2.6 on Equation 85 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall v. (\theta', j_1, v) \in [\tau_1 \sigma]_V \implies (\theta', j_1, e_{h1}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

where $\theta = W'_1.\theta_1$

We instantiate m with $m_1 + 2 + t_1$ where t_1 is the number of steps in which e_{h1} reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 1 + t_1). \forall v. (\theta', j_1, v) \in [\tau_1 \sigma]_V \implies$$

$$(\theta', j_1, e_{h1}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma} \quad (\text{FB-AC1})$$

Since from Equation 82 we have

$$(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau_1) \sigma]_V^A$$

Therefore from Lemma 2.15 we get

$$\forall m. (W'_2.\theta_1, m, v'_{j_1}) \in [\tau_1 \sigma]_V$$

Instantiating m with $m_1 + 1 + t_1$ we get

$$(W'_2.\theta_1, m_1 + 1 + t_1, v'_{j_1}) \in [\tau_1 \sigma]_V$$

Instantiating θ' with $W'_2.\theta_1$, j_1 with $m_1 + t_1$ and v with v'_{j_1} from Equation 82.

$$\text{Therefore we get } (W'_2.\theta_1, m_1 + 1 + t_1, e_{h1}[v'_{j_1}/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

From Definition 2.7, we get

$$\forall H. (m_1 + 1 + t_1, H) \triangleright W'_2.\theta_1 \wedge \forall k_c < (m_1 + 1 + t_1). (H, e_{h1}[v'_{j_1}/x]) \Downarrow_{k_c} (H'_1, v'_1) \implies$$

$$\exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1 + t_1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1 + t_1 - k_c), v'_1) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma))$$

Since from Equation 82 we have $(n - i - j, H'_{j_1}, H'_{j_1}) \triangleright W'_2$

Therefore from Lemma 2.27 we get $\forall m. (m, H'_{j_1}) \triangleright W'_2.\theta_1$

Instantiating m with $m_1 + 1 + t_1$ we get $(m_1 + 1 + t_1, H'_{j_1}) \triangleright W'_2.\theta_1$

Now instantiating H with H'_{j_1} from Equation 82 and k_c with t_1 we get

$$\exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \quad (\text{R1})$$

Similarly we can apply Definition 2.6 on Equation 86 to get

$$\forall m. \forall \theta'_2. (m, W'_1.\theta_2) \sqsubseteq \theta'_2 \wedge \forall j_2 < m. \forall v. (\theta'_2, j_2, v) \in [\tau_1 \sigma]_V \implies$$

$$(\theta'_2, j_2, e_{h2}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$$

We instantiate m with $m_2 + 2 + t_2$ where t_2 is the number of steps in which e_{h2} reduces

$$\forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall v. (\theta', j_1, v) \in [\tau_1 \sigma]_V \implies$$

$$(\theta', j_1, e_{h2}[v/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma} \quad (\text{FB-AC2})$$

Since from Equation 82 we have

$$(W'_2, n - i - j, v'_{j_1}, v'_{j_2}) \in [(\tau_1) \sigma]_V^A$$

Therefore from Lemma 2.15 we get

$$\forall m. (W'_2.\theta_2, m, v'_{j_2}) \in [\tau_1 \sigma]_V$$

Instantiating m with $m_2 + 1 + t_2$ we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, v'_{j_2}) \in [\tau_1 \sigma]_V$$

Instantiating θ' with $W'_2.\theta_2$, j_1 with $m_2 + 1 + t_2$ and v with v'_{j_2} from Equation 82 in FB-AC2 we get

$$(W'_2.\theta_2, m_2 + 1 + t_2, e_{h_2}[v'_{j_2}/x]) \in [\tau_2 \sigma]_E^{\ell_e} \sigma$$

From Definition 2.7, we get

$$\begin{aligned} \forall H.(m_2 + 1 + t_2, H) \triangleright W'_2.\theta_2 \wedge \forall k_c < (m_2 + 1 + t_2).(H, e_{h_2}[v'_{j_1}/x]) \Downarrow_{k_c} (H'_2, v'_2) \implies \\ \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1 + t_2 - k_c), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1 + t_2 - k_c)v'_2) \in [\tau_2 \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since from Equation 82 we have $(n - i - j, H'_{j_1}, H'_{j_1}) \triangleright W'_2$

Therefore from Lemma 2.27 we get $\forall m.(m, H'_{j_2}) \triangleright W'_2.\theta_2$

Instantiating m with $m_2 + 1 + t_2$ we get $(m_2 + 1 + t_2, H'_{j_2}) \triangleright W'_2.\theta_2$

Now Instantiating H with H'_{j_2} from Equation 82 and and k_c with t_2 .

$$\begin{aligned} \exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2 \sigma]_V \wedge \\ (\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \quad (\text{R2}) \end{aligned}$$

In order to prove FB-A0 we choose W' to be $(\theta'_1, \theta'_2, W'_2.\beta)$. Now we need to show two things:

(a) $(n - n', H'_1, H'_2) \triangleright W'$:

From Definition 2.9 it suffices to show that

– $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$:

From R1 we know that $(m_1 + 1, H'_1) \triangleright \theta'_1$, therefore from Definition 2.8 we get $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from R2 we know that $(m_2 + 1, H'_2) \triangleright \theta'_2$, therefore from Definition 2.8 we get $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

– $(W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2))$:

Since from Equation 82 we know that $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$ therefore from

Definition 2.9 we know that $(W'_2.\hat{\beta}) \subseteq (\text{dom}(W'_2.\theta_1) \times \text{dom}(W'_2.\theta_2))$

From R1 and R2 we know that $W'_2.\theta_1 \sqsubseteq \theta'_1$ and $W'_2.\theta_2 \sqsubseteq \theta'_2$ therefore

$(W'_2.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

– $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge$

$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

4 cases arise for each $(a_1, a_2) \in W'_2.\hat{\beta}$

i. $H'_{j_1}(a_1) = H'_1(a_1) \wedge H'_{j_2}(a_2) = H'_2(a_2)$:

* $W'.\theta_1(a_1) = W'.\theta_2(a_2)$:

We know from Equation 82 that $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright W'_2$

Therefore from Definition 2.9 we have

$\forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

Since $W'.\hat{\beta} = W'_2.\hat{\beta}$ by construction therefore

$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$

From R1 and R2 we know that $W'_2.\theta_1 \sqsubseteq \theta'_1$ and $W'_2.\theta_2 \sqsubseteq \theta'_2$ respectively.

Therefore from Definition 2.2

$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$

* $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

From Equation 82 we know that $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$

This means from Definition 2.9 that

$$\forall (a_{i_1}, a_{i_2}) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$$

Instantiating with a_1 and a_2 and since $W'_2 \sqsubseteq W'$ and $n - n' - 1 < n - i - j - 1$ (since $n' = i + j + t_1$ where t_1 is the number of steps taken by e_{h_1} , i is the number of steps taken by $e_1 \gamma \downarrow_1$ to reduce and j is the number of steps taken by $e_2 \gamma \downarrow_1$ to reduce) therefore from Lemma 2.17 we get

$$(W', n - n' - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

ii. $H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$:

* $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

Same reasoning as in the previous case

* $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

From R1 and R2 we know that

$$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_2.\theta_1(a_1) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_2.\theta_2(a_2) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell'$$

Since $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$ (given) and $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_e \sigma \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Also from R1 and R2, $(m_1 + 1, H'_1) \triangleright \theta'_1$ and $(m_2 + 1, H'_2) \triangleright \theta'_2$. Therefore from Definition 2.8 we have

$$(\theta'_1, m_1, H'_1(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V$$
 and

$$(\theta'_2, m_2, H'_2(a_1)) \in \lfloor \theta'_2(a_2) \rfloor_V$$

Since m_1 and m_2 are arbitrary indices therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

iii. $H'_{j_1}(a_1) = H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$:

* $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

Same reasoning as in the previous case

* $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

From R2 we know that

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means that a_2 was protected at $\ell_e \sigma$ in the world before the modification. Since $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$ (given) and $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_e \sigma \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 82 we know that $(n - i - j, H'_{j_1}, H'_{j_2}) \triangleright^A W'_2$ that means from Definition 2.9 that $(W'_2, n - i - j - 1, H'_{j_1}(a_1), H'_{j_2}(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$. Since $(\ell_e \sigma) \sqsubseteq \ell'$ therefore from Definition 2.4 we know that $H'_{j_1}(a_1)$ must also be protected at some label $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_2.\theta_1, m, H'_{j_1}(a_1)) \in W'_2.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_2.\theta_2, m, H'_{j_2}(a_2)) \in W'_2.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with m_1 and using Lemma 2.16 we get

$$(\theta'_1, m_1, H'_{j_1}(a_1)) \in \theta'_1(a_1)$$

Since from R2 we know that $(m_2+1, H'_2) \triangleright \theta'_2$ therefore from Definition 2.8

$$\text{we know that } (\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$$

Therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$\text{iv. } H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) = H'_2(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V:$$

$$\underline{i = 1}$$

This means that given some m we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

Like before we instantiate Equation 85 and Equation 86 with $m + 2 + t_1$ and $m + 2 + t_2$ respectively. This will give us

$$\exists \theta'_1. W'_2.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H'_{j_1}(a) \neq H'_1(a) \implies \exists \ell'. W'_2.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_2.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma))$$

and

$$\exists \theta'_2. W'_2.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_2) \in [\tau_2 \sigma]_V \wedge$$

$$(\forall a. H'_{j_2}(a) \neq H'_2(a) \implies \exists \ell'. W'_2.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_2.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma))$$

Since we have $(m+1, H'_1) \triangleright \theta'_1$ and $(m+1, H'_2) \triangleright \theta'_2$ therefore we get the desired from Definition 2.8

$$\underline{i = 2}$$

Symmetric to $i = 1$

$$(b) (W', n - n' - 1, v'_1, v'_2) \in [\tau_2 \sigma]_V^A:$$

Let $\tau_2 = \mathbf{A}^{\ell_i}$ Since $\tau_2 \sigma \searrow \ell \sigma$ and since $\ell \sigma \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From R1 and R2 we and Definition 2.4 we get the desired.

4. FG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp}$$

To prove: $(W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\perp \sigma]_E^A$

Say $e_1 = (e_1, e_2) (\gamma \downarrow_1)$ and $e_2 = (e_1, e_2) (\gamma \downarrow_2)$

From Definition of $[(\tau_1 \times \tau_2)^\perp \sigma]_E^A$ it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies$$

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^A$$

This means that given some H_1, H_2 and $n' < n$ s.t

$$(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp \sigma]_{\mathcal{V}}^A \quad (87)$$

$$\underline{\text{IH1}} (W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{p11}, H_{p12}. (n, H_{p11}, H_{p12}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11}) \wedge (H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12}) \implies$$

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{p11}, H'_{p12}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

Instantiating H_{p11} with H_1 and H_{p22} with H_2 in IH1 and since the (e_1, e_2) reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore we know that $\exists i < n' < n$ s.t $(H_{p11}, e_1 (\gamma \downarrow_1)) \downarrow_i (H'_{p11}, v'_{p11})$. Similarly since we know that (e_1, e_2) reduces to value with $\gamma \downarrow_2$ therefore we know that $(H_{p12}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p12}, v'_{p12})$. Hence we get

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{p11}, H'_{p12}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (88)$$

$$\underline{\text{IH2}} (W, n - i, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_2 \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{p21}, H_{p22}. (n - i, H_{p21}, H_{p22}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_{p21}, e_2 (\gamma \downarrow_1)) \downarrow_j (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_2 (\gamma \downarrow_2)) \downarrow (H'_{p22}, v'_{p22}) \implies$$

$$\exists W'_2 \sqsubseteq W'_1. (n - i - j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Instantiating H_{p21} with H'_{p11} and H_{p22} with H'_{p21} and in IH2. Since (e_1, e_2) reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps and e_1 has reduced with $i < n'$ steps. Therefore we know that $\exists j < n' - i < n - i$ s.t $(H_{p21}, e_2 (\gamma \downarrow_1)) \downarrow_j (H'_{p21}, v'_{p21})$. Similarly since we know that (e_1, e_2) reduces to value with $\gamma \downarrow_2$ therefore we know that $(H_{p22}, e_2 (\gamma \downarrow_2)) \downarrow (H'_{p22}, v'_{p22})$. Hence we get

since the (e_1, e_2) reduces to value with both $\gamma \downarrow_1$ and $\gamma \downarrow_2$ therefore we know that $(H_{p21}, e_2 (\gamma \downarrow_1)) \downarrow (H'_{p21}, v'_{p21}) \wedge (H_{p22}, e_1 (\gamma \downarrow_2)) \downarrow (H'_{p22}, v'_{p22})$. Hence we get

$$\exists W'_2 \sqsubseteq W'_1. (n - i - j, H'_{p21}, H'_{p22}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (89)$$

In order to prove Equation 87 we instantiate W' in Equation 87 with W'_2 we are required to show the following:

- $W \sqsubseteq W'_2$:
Since $W \sqsubseteq W'_1$ from Equation 88 and $W'_1 \sqsubseteq W'_2$ from Equation 89
Therefore, $W \sqsubseteq W'_2$ from Definition 2.3

- $(n - n', H'_1, H'_2) \overset{A}{\triangleright} W'$:

Here $n' = i + j + 1$

From evaluation rule of products we know that $H'_1 = H'_{p21}$ and $H'_2 = H'_{p22}$

From Equation 89 we know that $(n - i - j, H'_{p21}, H'_{p22}) \overset{A}{\triangleright} W'_2$

Therefore from Lemma 2.21 we get $(n - i - j - 1, H'_{p21}, H'_{p22}) \overset{A}{\triangleright} W'_2$

- $(W', n - i - j - 1, v'_1, v'_2) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^A$:

From evaluation rule of products we know that $v'_1 = (v'_{p11}, v'_{p21})$ and $v'_2 = (v'_{p12}, v'_{p22})$

We are required to show

$$- (W'_2, n - i - j - 1, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_V^A \wedge (W'_2, n - i - j - 1, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_V^A:$$

From Equation 88 and Equation 89 we know that

$$(W'_2, n - i - j, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_V^A \wedge (W'_2, n - i - j, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_V^A$$

Therefore from Lemma 2.17 we get

$$(W'_2, n - i - j - 1, v'_{p11}, v'_{p12}) \in [\tau_1 \sigma]_V^A \wedge (W'_2, n - i - j - 1, v'_{p21}, v'_{p22}) \in [\tau_2 \sigma]_V^A$$

5. FG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e_i) : \tau_1}$$

To prove: $(W, n, (\text{fst}(e_i)) (\gamma \downarrow_1), (\text{fst}(e_i)) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

Say $e_1 = (\text{fst}(e_i)) (\gamma \downarrow_1)$ and $e_2 = (\text{fst}(e_i)) (\gamma \downarrow_2)$

From Definition 2.5 it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \overset{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \overset{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A \quad (90)$$

IH1

$$(W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \overset{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \overset{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^A$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $\text{fst}(e_i)$ reduces to value reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore we know that $\exists i < n' < n$ s.t $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1})$. Similarly since we know that $\text{fst}(e_i)$ reduces to value with $\gamma \downarrow_2$ therefore we know that $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^A \quad (91)$$

We case analyze on $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^A$ from Equation 91

- Case $\ell \sigma \sqsubseteq \mathcal{A}$:

From Definition 2.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \times \tau_2) \sigma]_V^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\tau_1 \sigma \times \tau_2 \sigma)]_V^A$$

Let $v'_{i1} = (v_{i1}, v_{i2})$ and $v'_{i2} = (v_{j1}, v_{j2})$

Again from Definition 2.4 it means that

$$(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1 \sigma]_V^A \wedge (W'_1, n - i, v_{i2}, v_{j2}) \in [\tau_2 \sigma]_V^A \quad (F1)$$

In order to prove Equation 90 we choose W' as W'_1 and from the evaluation rule of fst we know that $H'_1 = H'_{i1}$ and $H'_2 = H'_{i2}$. Also, from reduction rules we know that $n' = i + 1$. And then we need to show:

- $W \sqsubseteq W'_1$:

Directly from Equation 91

- $(n - n', H'_1, H'_2) \triangleright^A W'_1$:

Since from Equation 91 we know that $(n - i, H'_1, H'_2) \triangleright^A W'_1$

Therefore from Lemma 2.21 we get $(n - i - 1, H'_1, H'_2) \triangleright^A W'_1$

- $(W'_1, n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A$:

From the evaluation rule we know that $v'_1 = v_{i1}$ and $v'_2 = v_{j1}$

From F1 we know that $(W'_1, n - i, v_{i1}, v_{j1}) \in [\tau_1 \sigma]_V^A$

Therefore from Lemma 2.17 we get $(W'_1, n - i - 1, v_{i1}, v_{j1}) \in [\tau_1 \sigma]_V^A$

- Case $\ell \sigma \not\sqsubseteq \mathcal{A}$:

In this case from Definition 2.6 we know that

$$(a) \forall m. (W'_1.\theta_1, m, v'_{i1}) \in [(\tau_1 \sigma \times \tau_2 \sigma)]_V \text{ and}$$

$$(b) \forall m. (W'_1.\theta_2, m, v'_{i2}) \in [(\tau_1 \sigma \times \tau_2 \sigma)]_V$$

where

$$v'_{i1} = (v_{i1}, v_{i2}) \text{ and } v'_{i2} = (v_{j1}, v_{j2})$$

In order to prove Equation 90 we choose W' as W'_1 and from the evaluation rule of fst we know that $H'_1 = H'_{i1}$ and $H'_2 = H'_{i2}$. And then we need to show:

- $W \sqsubseteq W'_1$:

Directly from Equation 91

- $(n - n', H'_1, H'_2) \triangleright^A W'_1$:

From Equation 91 we know that $(n - i, H'_1, H'_2) \triangleright^A W'_1$

Therefore from Lemma 2.21 we get

$$(n - i - 1, H'_1, H'_2) \triangleright^A W'_1$$

- $(W'_1, n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A$:
From the evaluation rule we know that $v'_1 = v_{i1}$ and $v'_2 = v_{j1}$
Let $\tau_1 = \mathbf{A}^{\ell_i}$ Since $\tau_1 \sigma \searrow \ell$ and since $\ell \sigma \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

Therefore from Definition 2.4 it suffices to prove that

$$\forall m_1. (W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V$$

and

$$\forall m_2. (W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V$$

This means given m_1 and it suffices to prove:

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V \quad (92)$$

Similarly given m_2 , it suffices to prove:

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V \quad (93)$$

Instantiating (a) with m_1

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V \wedge (W'_1.\theta_1, m_1, v_{i2}) \in [\tau_2 \sigma]_V \quad (94)$$

Instantiating (b) with m_2

$$(W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V \wedge (W'_1.\theta_2, m_2, v_{j2}) \in [\tau_2 \sigma]_V \quad (95)$$

From Equation 94 and Equation 95 we get

$$(W'_1.\theta_1, m_1, v_{i1}) \in [\tau_1 \sigma]_V \text{ and } (W'_1.\theta_2, m_2, v_{j1}) \in [\tau_1 \sigma]_V$$

6. FG-snd:

Symmetric case as FG-fst

7. FG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e_i) : (\tau_1 + \tau_2)^\perp}$$

To prove: $(W, n, (\text{inl}(e_i))(\gamma \downarrow_1), (\text{inl}(e_i))(\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\perp \sigma]_E^A$

Say $e_1 = (\text{inl}(e_i))(\gamma \downarrow_1)$ and $e_2 = (\text{inl}(e_i))(\gamma \downarrow_2)$

From Definition of $[(\tau_1 + \tau_2)^\perp \sigma]_E^A$ it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp \sigma]_V^A$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp \sigma]_V^A \quad (96)$$

IH1 $(W, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

This means from Definition 2.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A \end{aligned}$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $\text{inl}(e_i)$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore we know that $\exists i < n' < n$ s.t. $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$. Similarly since we know that $\text{inl}(e_i)$ reduces to value with $\gamma \downarrow_2$ therefore we know that $(H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A \quad (97)$$

Instantiating W' in Equation 96 with W'_1 . Also from reduction relation we know that $n' = i + 1$ we are required to show the following:

- $W \sqsubseteq W'_1$:

Directly from Equation 97

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$:

From Equation 97 we know that $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 2.21 we get

$$(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$$

- $(W'_1, n - n', v'_1, v'_2) \in [(\tau_1 + \tau_2)^\perp \sigma]_V^A$:

From evaluation rule of inl we know that $v'_1 = \text{inl}(v'_{i1})$ and $v'_2 = \text{inl}(v'_{i2})$

We are required to show

$$- (W'_1, n - n', v'_1, v'_2) \in [\tau_1 \sigma]_V^A:$$

From Equation 97 we know that $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A$

Therefore from Lemma 2.17 we get

$$(W'_1, n - i - 1, v'_{i1}, v'_{i2}) \in [\tau_1 \sigma]_V^A$$

8. FG-inr:

Symmetric case to FG-inl.

9. FG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{i1} : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_{i2} : \tau \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e_i, x.e_{i1}, y.e_{i2}) : \tau}$$

To prove: $(W, (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1), (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$

Say $e_1 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_1)$ and $e_2 = (\text{case}(e_i, x.e_{i1}, y.e_{i2})) (\gamma \downarrow_2)$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} & \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ & \exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A \quad (98)$$

$$\underline{\text{IH1}} \ (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^A$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$ reduces to value with both $\gamma \downarrow_1$ and $\gamma \downarrow_2$ therefore we know that $(H_{i1}, e_i (\gamma \downarrow_1)) \Downarrow (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$. Hence we get

$$\exists W'_1 \sqsupseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{s1}, v'_{s2}) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^A \quad (99)$$

IH2:

$$(W'_1, n - i, (e_{i1}) (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\}), (e_{i1}) (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \in [(\tau) \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{j1}, H_{j2}. (n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_1, e_{i1} (\gamma \downarrow_1 \cup \{x \mapsto v_{i1}\})) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2 \cup \{x \mapsto v_{i2}\})) \Downarrow (H'_{j2}, v'_{j2}) \implies$$

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^A$$

Instantiating H_{j1} with H'_1 and H_{j2} with H'_2 in IH2. Also instantiating W with W'_1 . Since the $(\text{case}(e_i, x.e_{i1}, y.e_{i2}))$ reduces to value in both runs therefore we know that $(H_1, e_{i1} (\gamma \downarrow_1)) \Downarrow (H'_{j1}, v'_{j1}) \wedge (H_2, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$. Hence we get

$$\exists W'_2 \sqsupseteq W'_1. (n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^A \quad (100)$$

IH3:

$$(W'_1, n - i, (e_{i2}) (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\}), (e_{i2}) (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \in [(\tau) \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{k1}, H_{k2}. (n - i, H_{k1}, H_{k2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall k < n - i. (H_1, e_{i2} (\gamma \downarrow_1 \cup \{y \mapsto v_{i1}\})) \Downarrow_k (H'_{k1}, v'_{k1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2 \cup \{y \mapsto v_{i2}\})) \Downarrow (H'_{k2}, v'_{k2}) \implies$$

$$\exists W'_3 \sqsupseteq W'_1. (n - i - k, H'_{k1}, H'_{k2}) \stackrel{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k1}, v'_{k2}) \in [(\tau) \sigma]_V^A$$

Instantiating H_{k1} with H'_1 and H_{k2} with H'_2 in IH2. Also instantiating W with W'_1 . Since the $(\text{case}(e_i, x.e_{i2}, y.e_{i2}))$ reduces to value in both runs therefore we know that $(H_1, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{k1}, v'_{k1}) \wedge (H_2, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{k2}, v'_{k2})$. Hence we get

$$\exists W'_3 \sqsupseteq W'_1.(n - i - k, H'_{k1}, H'_{k2}) \overset{A}{\triangleright} W'_3 \wedge (W'_3, n - i - k, v'_{k1}, v'_{k2}) \in [\tau] \sigma \uparrow_V^A \quad (101)$$

We case analyze $(W'_1, n - i, v'_1, v'_2) \in [(\tau_1 + \tau_2)^\ell \sigma] \uparrow_V^A$ from Equation 99

- Case $\ell \sigma \sqsubseteq \mathcal{A}$:

From Definition 2.4 2 further cases arise:

- $v'_1 = \text{inl}(v_{i1})$ and $v'_2 = \text{inl}(v_{i2})$:

In this case from Definition 2.4 we know that $(W, n - i, v_{i1}, v_{i2}) \in [\tau_1 \sigma] \uparrow_V^A$

In order to prove Equation 98 we choose W' as W'_2 from Equation 100 and from the first evaluation rule of case we know that $H'_1 = H'_{j1}$ and $H'_2 = H'_{j2}$. Also we know from the evaluation rule that $n' = i + j + 1$. And then we need to show:

- * $W \sqsubseteq W'_2$:

Since $W \sqsubseteq W'_1$ from Equation 99 and $W'_1 \sqsubseteq W'_2$ from Equation 100

Therefore, $W \sqsubseteq W'_2$ from Definition 2.3

- * $(n - n', H'_{j1}, H'_{j2}) \overset{A}{\triangleright} W'_2$:

From Equation 100 we know that $(n - i - j, H'_{j1}, H'_{j2}) \overset{A}{\triangleright} W'_2$

Therefore from Lemma 2.21 we get

$$(n - i - j - 1, H'_{j1}, H'_{j2}) \overset{A}{\triangleright} W'_2$$

- * $(W'_2, n - n', v'_1, v'_2) \in [\tau \sigma] \uparrow_V^A$:

From the evaluation rule we know that $v'_1 = v'_{j1}$ and $v'_2 = v'_{j2}$

From Equation 100 we know that $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [\tau \sigma] \uparrow_V^A$

Therefore from Lemma 2.17 we get

$$(W'_2, n - i - j - 1, v'_{j1}, v'_{j2}) \in [\tau \sigma] \uparrow_V^A$$

- $v'_1 = \text{inr}(v_{i1})$ and $v'_2 = \text{inr}(v_{i2})$:

In this case from Definition 2.4 we know that $(W, v_{i1}, v_{i2}) \in [\tau_2 \sigma] \uparrow_V^A$

In order to prove Equation 98 we choose W' as W'_3 from Equation 101 and from the second evaluation rule of case we know that $H'_1 = H'_{k1}$ and $H'_2 = H'_{k2}$. Also we know from the evaluation rule that $n' = i + k + 1$. And then we need to show:

- * $W \sqsubseteq W'_3$:

Since $W \sqsubseteq W'_1$ from Equation 99 and $W'_1 \sqsubseteq W'_3$ from Equation 101

Therefore, $W \sqsubseteq W'_3$ from Definition 2.3

- * $(n - n', H'_1, H'_2) \overset{A}{\triangleright} W'_3$:

From Equation 101 we know that $(n - i - k, H'_{k1}, H'_{k2}) \overset{A}{\triangleright} W'_3$

Therefore from Lemma 2.21 we get

$$(n - i - k - 1, H'_{k1}, H'_{k2}) \overset{A}{\triangleright} W'_3$$

- * $(W'_3, n - n', v'_1, v'_2) \in [\tau \sigma] \uparrow_V^A$:

From the evaluation rule we know that $v'_1 = v'_{k1}$ and $v'_2 = v'_{k2}$

From Equation 101 we know that $(W'_3, n - i - k, v'_{k1}, v'_{k2}) \in [\tau \sigma] \uparrow_V^A$

Therefore from Lemma 2.17 we get

$$(W'_3, n - i - k - 1, v'_{k1}, v'_{k2}) \in [\tau \sigma] \uparrow_V^A$$

- Case $\ell \sigma \not\sqsubseteq \mathcal{A}$:

The following cases arise:

- Reduction of e_1 happens via Case1 and Reduction of e_2 happens via Case1 :
Exactly the same reasoning as in the $v'_1 = \text{inl}(v_{i1})$ and $v'_2 = \text{inl}(v_{i2})$ subcase of the $\ell \sigma \not\sqsubseteq \mathcal{A}$ case before.
- Reduction of e_1 happens via Case2 and Reduction of e_2 happens via Case2 :
Exactly the same reasoning as in the $v'_1 = \text{inr}(v_{i1})$ and $v'_2 = \text{inr}(v_{i2})$ subcase of the $\ell \sigma \not\sqsubseteq \mathcal{A}$ case before.
- Reduction of e_1 happens via Case1 and Reduction of e_2 happens via Case2 :

From Equation 98 we know that we need to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$$

In this case since we know that $\ell \sigma \not\sqsubseteq \mathcal{A}$. Let $\tau \sigma = A^{\ell_i}$ and since $\tau \sigma \searrow \ell \sigma$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove $\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$

From Definition 2.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means given m_1 and m_2 it suffices to prove:

$$(\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V \quad (102)$$

Since we know that $(W, n, \gamma) \in [\Gamma]_V^A$ (given) therefore from Lemma 2.25 we know that $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

Therefore by instantiating it at $m_1 + 1 + j$ we know that

$$(W.\theta_1, m_1 + 1 + j, \gamma \downarrow_1) \in [\Gamma]_V \quad (103)$$

Next we apply Theorem 2.22 on $e_{i1} \gamma \downarrow_1$. Here j is the number of steps in which $e_{i1} \gamma \downarrow_1$ reduces. We use $\gamma \downarrow_1 \cup \{x \mapsto v'_{s1}\}$ as the unary substitution to get

$$(W.\theta_1, m_1 + 1 + j, e_{i1} \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \in [(\tau) \sigma]_E^{pc}$$

This means from Definition 2.7 we get

$$\begin{aligned} & \forall H_{c2}. (m_1 + 1 + j, H_{c1}) \triangleright W_1.\theta_1 \wedge \forall l_c < (m_1 + 1 + j). (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \downarrow_{k_c} \\ & (H'_{c2}, v'_c) \implies \\ & \exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1 + j - l_c, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1 + j - l_c, v'_c) \in [(\tau) \sigma]_V \wedge \\ & (\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W_1.\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

Since from Equation 99 we know that $(n - i, H'_1, H'_2) \triangleright W'_1$ therefore from Lemma 2.27 we get $\forall m. (m, H'_1) \triangleright W'_1.\theta_1$

Instantiating m with $m_1 + 1 + j$ we get $(m_1 + 1 + j, H'_1) \triangleright W'_1.\theta_1$

Instantiating H_{c2} with H'_1 from Equation 99 and l_c with j we get
 $\exists \theta'_1. W_1.\theta_1 \sqsubseteq \theta'_1 \wedge (m_1 + 1, H'_{c2}) \triangleright \theta'_1 \wedge (\theta'_1, m_1 + 1, v'_c) \in [(\tau) \sigma]_V \wedge$
 $(\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_1(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_1). \theta'_1(a) \searrow (pc \sqcup \ell) \sigma) \quad (\text{CC1})$

Similarly we apply Theorem 2.22 on $e_{i2} \gamma \downarrow_2$. Here j_2 is the number of steps in which $e_{i2} \gamma \downarrow_2$ reduces. We use $\gamma \downarrow_2 \cup \{y \mapsto v'_{s2}\}$ as the unary substitution to get
 $(W_1.\theta_2, m_2 + 1 + j_2, e_{i2} \gamma \downarrow_1 \cup \{y \mapsto v'_c\}) \in [(\tau) \sigma]_E^{pc}$

This means from Definition 2.7 we get

$\forall H_{c2}. (m_2 + 1 + j_2, H_{c1}) \triangleright W_1.\theta_2 \wedge \forall l_c < m_2 + 1 + j_2. (H_{c2}, (e_{i1}) \gamma \downarrow_1 \cup \{x \mapsto v'_c\}) \Downarrow_{k_c}$
 $(H'_{c2}, v'_c) \implies$
 $\exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1 + j_2 - l_c, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1 + j_2 - l_c, v'_c) \in [(\tau) \sigma]_V \wedge$
 $(\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma)$

Since from Equation 99 we know that $(n - i, H'_1, H'_2) \triangleright W'_1$ therefore from Lemma 2.27 we get $\forall m. (m, H'_2) \triangleright W'_1.\theta_2$

Instantiating m with $m_2 + 1 + j_2$ we get $(m_2 + 1 + j_2, H'_2) \triangleright W'_1.\theta_2$

Instantiating H_{c2} with H'_2 (from Equation 99) and l_c with j_2 to get
 $\exists \theta'_2. W_1.\theta_2 \sqsubseteq \theta'_2 \wedge (m_2 + 1, H'_{c2}) \triangleright \theta'_2 \wedge (\theta'_2, m_2 + 1, v'_c) \in [(\tau) \sigma]_V \wedge$
 $(\forall a. H_{c2}(a) \neq H'_{c2}(a) \implies \exists \ell'. W_1.\theta_2(a) = A^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(\theta_1). \theta'_2(a) \searrow (pc \sqcup \ell) \sigma) \quad (\text{CC2})$

We choose

$W_n.\theta_1 = \theta'_1$ (from CC1)
 $W_n.\theta_2 = \theta'_2$ (from CC2)
 $W_n.\hat{\beta} = W'_1.\hat{\beta}$ (from Equation 99)

In order to prove Equation 98 we choose W' as W_n

i. $(n - n', H'_1, H'_2) \triangleright W'$:

From Definition 2.9 it suffices to show that

– $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$:

From (CC1) we know that $(m_1 + 1, H'_1) \triangleright \theta'_1$, therefore from Definition 2.8 we get $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from (CC2) we know that $(m_2 + 1, H'_2) \triangleright \theta'_2$, therefore from Definition 2.8 we get $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

– $(W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2))$:

Since from Equation 99 we have $(n - i, H'_1, H'_2) \triangleright W'_1$ therefore from Definition 2.9 we get $(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$

From (CC1) and (CC2) we know that $W'_1.\theta_1 \sqsubseteq \theta'_1$ and $W'_1.\theta_2 \sqsubseteq \theta'_2$ therefore
 $(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

– $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge$
 $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'.\theta_1(a_1)]_V^A$:

4 cases arise for each a_1 and a_2

- A. $H'_{j_1}(a_1) = H'_1(a_1) \wedge H'_{j_2}(a_2) = H'_2(a_2)$:
 $\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2)}$:

We know from Equation 99 that $(n - i, H'_1, H'_2) \triangleright W'_1$

Therefore from Definition 2.9 we have

$$\forall(a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since $W'.\hat{\beta} = W'_1.\hat{\beta}$ by construction therefore

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From (CC1) and (CC2) we know that $W'_1.\theta_1 \sqsubseteq \theta'_1$ and $W'_1.\theta_2 \sqsubseteq \theta'_2$ respectively.

Therefore from Definition 2.2

$$\forall(a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}$$

From Equation 99 we know that $(n - i, H'_1, H'_2) \triangleright W'_1$

This means from Definition 2.9 that

$$\forall(a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) \wedge (W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Instantiating with a_1 and a_2 and since $W'_1 \sqsubseteq W'$ and $n - n' - 1 < n - i - 1$ (since $n' = i + t_1 + 1$ where t_1 is the number of steps taken by e_{i1} , i is the number of steps taken by $e_1 \gamma \downarrow_1$ to reduce) therefore from Lemma 2.17 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

- B. $H'_{j_1}(a_1) \neq H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$:

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2)}$$

Same as before

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}$$

From (CC1) and (CC2) we know that

$$(\forall a. H'_1(a) \neq H'_{c1}(a) \implies \exists \ell'. W'_1.\theta_1(a) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_2(a) \neq H'_{c2}(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = A^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell'$$

Since $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $(pc \sqcup \ell) \sigma \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Also from (CC1) and (CC2), $(m_1 + 1, H'_{c1}) \triangleright \theta'_1$ and $(m_2 + 1, H'_{c2}) \triangleright \theta'_2$.

Therefore from Definition 2.8 we have

$$(\theta'_1, m_1, H'_{c1}(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V$$

$$(\theta'_2, m_2, H'_{c2}(a_1)) \in \lfloor \theta'_2(a_2) \rfloor_V$$

Since m_1 and m_2 are arbitrary indices therefore from Definition 2.4 we get (here $H'_1 = H'_{c1}$ and $H'_2 = H'_{c2}$)

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

- C. $H'_{j_1}(a_1) = H'_1(a_1) \vee H'_{j_2}(a_2) \neq H'_2(a_2)$:

$$\underline{W'.\theta_1(a_1) = W'.\theta_2(a_2)}$$

Same as before

$$\underline{(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A}$$

From (CC2) we know that

$$(\forall a. H_2'(a) \neq H_{c2}'(a) \implies \exists \ell'. W_1'.\theta_2(a) = \mathbf{A}^{\ell'} \wedge ((pc \sqcup \ell) \sigma) \sqsubseteq \ell')$$

This means that a_2 was protected at $(pc \sqcup \ell) \sigma$ in the world before the modification. Since $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $(pc \sqcup \ell) \sigma \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 99 we know that $(n - i, H_1', H_2') \stackrel{\mathcal{A}}{\triangleright} W_1'$ that means from Definition 2.9 that $(W_1', n - i - 1, H_1'(a_1), H_2'(a_2)) \in \lceil W_1'.\theta_1(a_1) \rceil_V^{\mathcal{A}}$. Since $((pc \sqcup \ell) \sigma) \sqsubseteq \ell'$ therefore from Definition 2.4 we know that $H_1'(a_1)$ must also be protected at some label $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W_1'.\theta_1, m, H_1'(a_1)) \in W_1'.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W_1'.\theta_2, m, H_2'(a_2)) \in W_1'.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with m_1 and using Lemma 2.16 we get

$$(\theta_1', m_1, H_1'(a_1)) \in \theta_1'(a_1)$$

Since from (CC2) we know that $(m_2 + 1, H_{c2}') \triangleright \theta_2'$ therefore from Definition 2.8 we know that $(\theta_2', m_2, H_{c2}'(a_2)) \in \theta_2'(a_2)$

Therefore from Definition 2.4 we get

$$(W', n - n' - 1, H_{c1}'(a_1), H_{c2}'(a_2)) \in \lceil \theta_1'(a_1) \rceil_V^{\mathcal{A}}$$

$$\text{D. } H_{j1}'(a_1) \neq H_1'(a_1) \vee H_{j2}'(a_2) = H_2'(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H_i'(a_i)) \in \lfloor W'.\theta_i(a_i) \rfloor_V:$$

$$\underline{i = 1}$$

This means that given some m we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H_i'(a_i)) \in \lfloor W'.\theta_i(a_i) \rfloor_V$$

Like before we apply Theorem 2.22 on $e_{i1} \gamma_1$ and $e_{i2} \gamma_2$ but this time using $m + 1 + i$ and $m + 1 + j$ where i and j are the number of steps in which $e_{i1} \gamma_1$ and $e_{i2} \gamma_2$ reduces respectively. This will give us

$$\begin{aligned} & \exists \theta_1'. W_1.\theta_1 \sqsubseteq \theta_1' \wedge (m + 1, H_{c2}') \triangleright \theta_1' \wedge (\theta_1', m + 1, v'_c) \in \lfloor (\tau) \sigma \rfloor_V \wedge \\ & (\forall a. H_{c2}(a) \neq H_{c2}'(a) \implies \exists \ell'. W_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta_1') \setminus \text{dom}(\theta_1). \theta_1'(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

and

$$\begin{aligned} & \exists \theta_2'. W_1.\theta_2 \sqsubseteq \theta_2' \wedge (m + 1, H_{c2}') \triangleright \theta_2' \wedge (\theta_2', m + 1, v'_c) \in \lfloor (\tau) \sigma \rfloor_V \wedge \\ & (\forall a. H_{c2}(a) \neq H_{c2}'(a) \implies \exists \ell'. W_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (pc \sqcup \ell) \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta_2') \setminus \text{dom}(\theta_1). \theta_1'(a) \searrow (pc \sqcup \ell) \sigma) \end{aligned}$$

Since we have $(m + 1, H_{c1}') \triangleright \theta_1'$ and $(m + 1, H_{c2}') \triangleright \theta_2'$ therefore we get the desired from Definition 2.8

$$\underline{i = 2}$$

Symmetric to $i = 1$

$$\text{ii. } (W', n - n' - 1, v'_1, v'_2) \in \lceil \tau_2 \sigma \rceil_V^{\mathcal{A}}:$$

Let $\tau_2 = \mathbf{A}^{\ell_i}$ Since $\tau_2 \sigma \searrow \ell \sigma$ and since $\ell \sigma \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From CC1 and CC2 we and Definition 2.4 we get the desired.

- (d) Reduction of e_1 happens via Case2 and Reduction of e_2 happens via Case1 :
Symmetric case as before

10. FG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : \tau \quad \Sigma; \Psi \vdash \tau \searrow_{pc}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e_i : (\text{ref } \tau)^\perp}$$

To prove: $(W, (\text{new } (e_i)) (\gamma \downarrow_1), (\text{new } (e_i)) (\gamma \downarrow_2)) \in [(\text{ref } \tau)^\perp \sigma]_E^A$

Say $e_1 = (\text{new } (e_i)) (\gamma \downarrow_1)$ and $e_2 = (\text{new } (e_i)) (\gamma \downarrow_2)$

From Definition of $[(\text{ref } \tau)^\perp \sigma]_E^A$ it suffices to prove that

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A \end{aligned}$$

This means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A \quad (104)$$

$$\underline{\text{IH1}} (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$$

This means from Definition 2.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow \\ (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsubseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A \end{aligned}$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $\text{ref}(e_i)$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore $\exists i < n' < n$. s.t $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1})$. Similarly since $\text{ref}(e_i)$ reduces with $\gamma \downarrow_2$ therefore we know that $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsubseteq W. (n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A \quad (105)$$

From the evaluation rule of ref we know that $H'_1 = H'_{i1} \cup \{a_{n1} \mapsto v_{i1}\}$ and $H'_2 = H'_{i2} \cup \{a_{n2} \mapsto v_{i2}\}$

Inorder to prove Equation 104 we instantiate W' with W_n where W_n is

$$W_n.\theta_1 = W'_1.\theta_1 \cup \{a_{n1} \mapsto \tau\}$$

$$W_n.\theta_2 = W'_1.\theta_2 \cup \{a_{n2} \mapsto \tau\}$$

$$W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$$

Also we know that $n' = i + 1$

We are now required to prove

- $W \sqsubseteq W_n$:

From Equation 105 we know that $W \sqsubseteq W'_1$ and $W'_1 \sqsubseteq W_n$ by construction. Therefore from Definition 2.3, $W \sqsubseteq W_n$

- $(n - n', H'_1, H'_2) \hat{\triangleright}^A W_n$:

From Definition 2.9 it suffices to show that

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H'_2):$$

From Equation 105 and by construction of W_n

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_1)):$$

From Equation 105 and by construction of W_n

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, n - n', H'_1(a_1), H'_2(a_2)) \in [\![W_n.\theta_1(a_1)]\!]_V^A:$$

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

From Equation 105 and by construction of W_n

$$* \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\![W_n.\theta_1(a_1)]\!]_V^A:$$

From Equation 105 since we know that $(n - i, H'_{i1}, H'_{i2}) \hat{\triangleright}^A W'_1$ that means

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 1, H'_1(a_1), H'_2(a_2)) \in [\![W'_1.\theta_1(a_1)]\!]_V^A$$

Therefore from Lemma 2.17 we get $(n - i - 2 = n - n' - 1, \text{ since } n' = i + 1)$

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). (W'_1, n - i - 2, H'_1(a_1), H'_2(a_2)) \in [\![W'_1.\theta_1(a_1)]\!]_V^A$$

Since $W_n.\hat{\beta} = W'_1.\hat{\beta} \cup \{(a_{n1}, a_{n2})\}$ and from Equation 105 we know that $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\![\tau \sigma]\!]_V^A$

Therefore combining the two we get

$$\forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\![W_n.\theta_1(a_1)]\!]_V^A$$

$$- \forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in [\![W.\theta_i(a_i)]\!]_V:$$

From Equation 105 we have $(n - i, H'_{i1}, H'_{i2}) \hat{\triangleright}^A W'_1$ that means from Definition 2.9 we have

$$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W'_1.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in [\![W.\theta_i(a_i)]\!]_V$$

Also from Equation 105 we know that $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [\![\tau \sigma]\!]_V^A$

Therefore from Lemma 2.15 and Lemma 2.16 we get

$$\forall m. (W'_1.\theta_1, m, v'_{i1}) \in [\![\tau \sigma]\!]_V$$

and

$$\forall m. (W'_1.\theta_2, m, v'_{i2}) \in [\![\tau \sigma]\!]_V$$

Combining the two we get

$$\forall i \in \{1, 2\}. \forall a_i \in \text{dom}(W_n.\theta_i). \forall m. (W_n, m, H_i(a_i)) \in [\![W.\theta_i(a_i)]\!]_V$$

- $(W_n, n - n', v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A$:

Here $v'_1 = a_{n1}$ and $v'_2 = a_{n2}$

Since $(a_{n1}, a_{n2}) \in W_n$ and also $W_n.\theta_1(a_{n1}) = W_n.\theta_1(a_{n1}) = \tau$

Therefore from Definition 2.4 $(W_n, v'_1, v'_2) \in [(\text{ref } \tau)^\perp \sigma]_V^A$

11. FG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_i : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e_i : \tau'}$$

To prove: $(W, n, !(e_i)) (\gamma \downarrow_1), !(e_i)) (\gamma \downarrow_2) \in [(\tau') \sigma]_E^A$

Say $e_1 = !(e_i) (\gamma \downarrow_1)$ and $e_2 = !(e_i) (\gamma \downarrow_2)$

This means from Definition 2.5 we need to prove:

$$\begin{aligned} \forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !(e_i)(\gamma \downarrow_1)) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !(e_i)(\gamma \downarrow_2)) \downarrow \\ (H'_2, v'_2) \implies \\ \exists W' \sqsubseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau') \sigma]_{\mathcal{V}}^A \end{aligned}$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, !(e_i)(\gamma \downarrow_1)) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, !(e_i)(\gamma \downarrow_2)) \downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsubseteq W. (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau') \sigma]_{\mathcal{V}}^A \quad (106)$$

$$\underline{\text{IH1}} (W, n, (e_i) (\gamma \downarrow_1), (e_i) (\gamma \downarrow_2)) \in [(\text{ref } \tau)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\begin{aligned} \forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_1, v'_1) \wedge (H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_2, v'_2) \implies \\ \exists W'_1 \sqsubseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_{\mathcal{V}}^A \end{aligned}$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $!(e_i)$ reduces to value with both $\gamma \downarrow_1$ in $n' < n$ steps therefore $\exists i < n' < n$ s.t. $(H_{i1}, e_i (\gamma \downarrow_1)) \downarrow_i (H'_1, v'_1)$. Similarly since $!e_i$ reduces to value with $\gamma \downarrow_2$ therefore $(H_{i2}, e_i (\gamma \downarrow_2)) \downarrow (H'_2, v'_2)$. Hence we get

$$\exists W'_1 \sqsubseteq W. (n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_{\mathcal{V}}^A \quad (107)$$

We case analyze on $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau)^\ell \sigma]_{\mathcal{V}}^A$ from Equation 107

- Case $\ell \sigma \sqsubseteq \mathcal{A}$:

From Definition 2.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau) \sigma]_{\mathcal{V}}^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } (\tau \sigma))]_{\mathcal{V}}^A$$

Let $v'_{i1} = a_{i1}$ and $v'_{i2} = a_{i2}$

Again from Definition 2.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1.\hat{\beta} \wedge W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau \quad (\text{D1})$$

In order to prove Equation 106 we instantiate W' with W'_1 . Also we know that $n' = i + 1$

- $W'_1 \sqsubseteq W$:

From Equation 107

– $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$:

From Equation 107 we know that

$(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 2.21 we get

$(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

– $(W'_1, n - n', v'_1, v'_2) \in [(\tau') \sigma]_V^A$:

From the evaluation rule of deref we know that $v'_1 = H'_1(a_{i1})$ and $v'_2 = H'_2(a_{i2})$

Since from Equation 107 we know that $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$, therefore from Definition 2.9 we know that

$(W'_1, n - i - 1, H'_1(a_{i1}), H'_2(a_{i2})) \in [W'_1.\theta_1(a_{i1})]_V^A$

And from D1 we know that $W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau$

Therefore $(W'_1, v'_1, v'_2) \in [(\tau) \sigma]_V^A$

Since $\tau \sigma <: \tau' \sigma$ Therefore from Lemma 2.28, we get

$(W'_1, n - i - 1, v'_1, v'_2) \in [(\tau') \sigma]_V^A$

• Case $\ell \sigma \not\sqsubseteq \mathcal{A}$:

From the evaluation rule of deref we know that $v'_{i1} = a_1$ and $v'_{i2} = a_2$

In this case from Definition 2.4 we know that

$$\forall m_1. (W'_1.\theta_1, m_1, a_1) \in [(\text{ref } \tau) \sigma]_V \quad (108)$$

and

$$\forall m_2. (W'_1.\theta_2, m_2, a_2) \in [(\text{ref } \tau) \sigma]_V \quad (109)$$

In order to prove Equation 106 we choose W' as W'_1 . And then we need to show:

– $W \sqsubseteq W'_1$:

Directly from Equation 107

– $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$:

From Equation 107 we know that $(n - i, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

Therefore from Lemma 2.21 we get

$(n - i - 1, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1$

– $(W'_1, n - n', v'_1, v'_2) \in [\tau' \sigma]_V^A$:

Let $\tau' = A^{\ell_i}$ Since $\tau' \sigma \searrow_{\ell}$ and since $\ell \sigma \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

Therefore from Definition 2.4 it suffices to prove that

$\forall m_1. (W'_1.\theta_1, m_1, v'_1) \in [\tau' \sigma]_V$

and

$\forall m_2. (W'_1.\theta_2, m_2, v'_2) \in [\tau' \sigma]_V$

This means given m_1 and it suffices to prove:

$$(W'_1.\theta_1, m_1, v'_1) \in [\tau' \sigma]_V \quad (110)$$

Similarly given m_2 , it suffices to prove:

$$(W'_1.\theta_2, m_2, v'_2) \in [\tau' \sigma]_V \quad (111)$$

Since from Equation 107 we know that $(n-i, H'_1, H'_2) \triangleright W'_1$ therefore from Lemma 2.27 we get

$$\forall m_{h1}.(m_{h1}, H'_1) \triangleright W'_1.\theta_1 \quad (112)$$

$$\forall m_{h2}.(m_{h2}, H'_2) \triangleright W'_1.\theta_2 \quad (113)$$

Instantiating m_{h1} in Equation 112 with $m_1 + 1$ we get $(m_1, H'_1) \triangleright W'_1.\theta_1$

Therefore from Definition 2.8, we get

$$\forall a \in \text{dom}(W'_1.\theta_1).(W'_1.\theta_1, m_1, H'_1(a)) \in [W'_1.\theta_1(a)]_V$$

Instantiating a with a_1 we get $(W'_1.\theta_1, m_1, H'_1(a_1)) \in [W'_1.\theta_1(a)]_V$

Since $W'_1.\theta_1(a_{i1}) = \tau$ therefore we get

$$(W'_1.\theta_1, m_1, v'_1) \in [\tau \sigma]_V$$

and since $\tau \sigma <: \tau' \sigma$ therefore from Lemma 2.24 we get

$$(W'_1.\theta_1, m_1, v'_1) \in [\tau' \sigma]_V$$

Similarly we also get

$$(W'_1.\theta_2, m_2, v'_2) \in [\tau' \sigma]_V$$

Finally from Definition 2.4 we get

$$(W'_1, v'_1, v'_2) \in [(\tau') \sigma]_V^A$$

12. FG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{i1} : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{i2} : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_{i1} := e_{i2} : \text{unit}}$$

To prove: $(W, n, (e_{i1} := e_{i2}) (\gamma \downarrow_1), (e_{i1} := e_{i2}) (\gamma \downarrow_2)) \in [(\text{unit}) \sigma]_E^A$

Say $e_1 = (e_{i1} := e_{i2}) (\gamma \downarrow_1)$ and $e_2 = (e_{i1} := e_{i2}) (\gamma \downarrow_2)$

This means from Definition 2.5 we need to prove:

$$\forall H_1, H_2.(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n.(H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{unit}) \sigma]_V^A$$

This further means that given

$$\forall H_1, H_2.(n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n.(H_1, (e_{i1} := e_{i2})(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e_{i1} := e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\text{unit}) \sigma]_V^A \quad (114)$$

IH1 $(W, n, (e_{i1}) (\gamma \downarrow_1), (e_{i1}) (\gamma \downarrow_2)) \in [(\text{ref } \tau)^\ell \sigma]_E^A$

This means from Definition 2.5 we get

$\forall H_{i1}, H_{i2}.(n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_1) \wedge (H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_2) \implies$

$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_V^A$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH1 and since the $(e_{i1} := e_{i2})$ reduces to value with both $\gamma \downarrow_1$ in $n' < n$ steps therefore $\exists i < n' < n$ s.t $(H_{i1}, e_{i1} (\gamma \downarrow_1)) \Downarrow (H'_{i1}, v'_1)$. Similarly since $(e_{i1} := e_{i2})$ reduces to value with $\gamma \downarrow_2$ therefore we also have $(H_{i2}, e_{i1} (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_2)$. Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\text{ref } \tau)^\ell \sigma]_V^A \quad (115)$$

IH2 $(W, n - i, (e_{i2}) (\gamma \downarrow_1), (e_{i2}) (\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$

This means from Definition 2.5 we get

$\forall H_{j1}, H_{j2}.(n - i, H_{j1}, H_{j2}) \stackrel{A}{\triangleright} W'_1 \wedge \forall j < n - i. (H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow_j (H'_{j1}, v'_{j1}) \wedge (H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2}) \implies$

$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^A$

Instantiating H_{j1} with H'_{i1} and H_{j2} with H'_{i2} in IH2 and since the $(e_{i1} := e_{i2})$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps and e_1 reduces $\gamma \downarrow_1$ with $i < n'$ steps therefore $\exists j < (n' - i) < (n - i)$ s.t $(H_{j1}, e_{i2} (\gamma \downarrow_1)) \Downarrow (H'_{j1}, v'_{j1})$. Similarly we also have $(H_{j2}, e_{i2} (\gamma \downarrow_2)) \Downarrow (H'_{j2}, v'_{j2})$. Hence we get

$$\exists W'_2 \sqsupseteq W'_1.(n - i - j, H'_{j1}, H'_{j2}) \stackrel{A}{\triangleright} W'_2 \wedge (W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^A \quad (116)$$

We case analyze on $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau)^\ell \sigma]_V^A$ from Equation 115

- Case $\ell \sigma \sqsubseteq \mathcal{A}$:

From Definition 2.4 we know that this would mean that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } \tau) \sigma]_V^A$$

This means

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\text{ref } (\tau \sigma))]_V^A$$

Let $v'_{i1} = a_{i1}$ and $v'_{i2} = a_{i2}$

Again from Definition 2.4 it means that

$$(a_{i1}, a_{i2}) \in W'_1.\hat{\beta} \wedge W'_1.\theta_1(a_{i1}) = W'_1.\theta_2(a_{i2}) = \tau \sigma \quad (\text{A1})$$

In order to prove Equation 114 we instantiate W' with W'_2

- $W'_2 \sqsupseteq W$:

Since $W'_1 \sqsupseteq W$ from Equation 115 and $W'_2 \sqsupseteq W'_1$ from Equation 116

Therefore from Definition 2.3 we get $W'_2 \sqsupseteq W$

- $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$:

From the evaluation rule assign we know that

$$H'_1 = H'_{j1}[a_{i1} \mapsto v'_{j1}] \text{ and } H'_2 = H'_{j2}[a_{i2} \mapsto v'_{j2}]$$

Inorder to prove $(n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W'_2$ we need to show:

- * $dom(W'_2.\theta_1) \subseteq dom(H'_1) \wedge dom(W'_2.\theta_2) \subseteq dom(H'_2)$:
Directly from Equation 116
- * $W'_2.\hat{\beta} \subseteq (dom(W'_2.\theta_1) \times dom(W'_2.\theta_1))$:
Directly from Equation 116
- * $\forall(a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge$
 $(W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$:
- (a) $\forall(a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2)$:
 $\forall(a_1, a_2) \in (W'_2.\hat{\beta})$.
 - i. When $a_1 = a_{i1}$ and $a_2 = a_{i2}$:
From A1 we know that $W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$
and since $W'_1 \sqsubseteq W'_2$ therefore from Lemma 2.16 we get $W'_2.\theta_1(a_1) =$
 $W'_2.\theta_2(a_2) = \tau$
 - ii. When $a_1 = a_{i1}$ and $a_2 \neq a_{i2}$: This case cannot arise
 - iii. When $a_1 \neq a_{i1}$ and $a_2 = a_{i2}$: This case cannot arise
 - iv. When $a_1 \neq a_{i1}$ and $a_2 \neq a_{i2}$: From Equation 116 and Lemma 2.17
- (b) $\forall(a_1, a_2) \in (W'_2.\hat{\beta}). (W'_2, n - n', H'_1(a_1), H'_2(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$:
 $\forall(a_1, a_2) \in (W'_2.\hat{\beta})$.
 - i. When $a_1 = a_{i1}$ and $a_2 = a_{i2}$:
Since $H'_1(a_{i1}) = v'_{j1}$ and $H'_1(a_{i2}) = v'_{j2}$
From A1 we know that $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$
And since from Equation 116 we know that $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in$
 $\lceil (\tau) \sigma \rceil_V^A$
Therefore from Lemma 2.17 we get
 $(W'_2, n - j - i - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'_2.\theta_1(a_1) \rceil_V^A$
 - ii. When $a_1 = a_{i1}$ and $a_2 \neq a_{i2}$: This case cannot arise
 - iii. When $a_1 \neq a_{i1}$ and $a_2 = a_{i2}$: This case cannot arise
 - iv. When $a_1 \neq a_{i1}$ and $a_2 \neq a_{i2}$: From Equation 116 and from Lemma 2.17
- * $\forall i \in \{1, 2\}. \forall m. \forall a_i \in dom(W'_2.\theta_i). (W'_2.\theta_i, m, H'_i(a_i)) \in \lfloor W'_2.\theta_i(a_i) \rfloor_V$:
When $i = 1$
Given some m
 $\forall a_1 \in dom(W'_2.\theta_1)$.
 - when $a_1 = a_{i1}$:
From Equation 116 we know that $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in \lceil (\tau) \sigma \rceil_V^A$ thus
from Lemma 2.15 we know that
 $\forall m_1. (W'_2.\theta_1, m_1, H'_1(a_1)) \in \lfloor W'_2.\theta_1(a_1) \rfloor_V$

Instantiating with m we get
 $(W'_2.\theta_1, m, H'_1(a_1)) \in \lfloor W'_2.\theta_1(a_1) \rfloor_V$
 - Otherwise:
From Equation 116 and Lemma 2.27
- When $i = 2$
Similar reasoning as with $i = 1$
- $(W'_1, n - n', val'_1, v'_2) \in \lceil (unit) \sigma \rceil_V^A$:
From evaluation rule assign we know that $v'_1 = v'_2 = ()$
Directly from Definition 2.4

- Case $\ell \sigma \not\sqsubseteq \mathcal{A}$:

From Definition 2.4 we know that this would mean that

$$\forall m_1. (W'_1.\theta_1, m_1, a_{i1}) \in [(\text{ref } \tau) \sigma]_V \quad (117)$$

$$\forall m_2. (W'_1.\theta_2, m_2, a_{i2}) \in [(\text{ref } \tau) \sigma]_V \quad (118)$$

In order to prove Equation 114 we instantiate W' with W'_2 and then we need to show that:

- $W'_2 \sqsupseteq W$:

Since $W'_1 \sqsupseteq W$ from Equation 115 and $W'_2 \sqsupseteq W'_1$ from Equation 116

Therefore from Definition 2.3 we get $W'_2 \sqsupseteq W$

- $(n - n', H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W'_2$:

From the evaluation rule assign we know that

$$H'_1 = H'_{j1}[a_{i1} \mapsto v'_{j1}] \text{ and } H'_2 = H'_{j2}[a_{i2} \mapsto v'_{j2}]$$

In order to prove $(n - n', H'_1, H'_2) \stackrel{\mathcal{A}}{\triangleright} W'_2$ we need to show:

$$* \text{ dom}(W'_2.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'_2.\theta_2) \subseteq \text{dom}(H'_2):$$

Directly from Equation 116

$$* W'_2.\hat{\beta} \subseteq (\text{dom}(W'_2.\theta_1) \times \text{dom}(W'_2.\theta_2)):$$

Directly from Equation 116

$$* \forall (a_1, a_2) \in (W'_2.\hat{\beta}). W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) \wedge (W'_2, n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\![W'_2.\theta_1(a_1)]\!]_V^{\mathcal{A}}:$$

- (a) When $(a_{i1}, a_{i2}) \in W'_2.\hat{\beta}$:

$$\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$$

- i. When $a_1 = a_{i1}$ and $a_2 = a_{i2}$:

Instantiating Equation 117 and Equation 118 with $n - n' - 1$ we get

$$W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) = \tau$$

and since $W'_1 \sqsubseteq W'_2$ therefore from Definition 2.3 we get $W'_2.\theta_1(a_1) = W'_2.\theta_2(a_2) = \tau$

From Equation 116 we know that $(W'_2, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_V^{\mathcal{A}}$

Therefore $(W'_2, H_1(a_{i1})', H_2(a_{i2})') \in [(\tau) \sigma]_V^{\mathcal{A}}$

- ii. When $a_1 = a_{i1}$ and $a_2 \neq a_{i2}$: This case cannot arise

- iii. When $a_1 \neq a_{i1}$ and $a_2 = a_{i2}$: This case cannot arise

- iv. When $a_1 \neq a_{i1}$ and $a_2 \neq a_{i2}$: From Equation 116

- (b) When $(a_{i1}, a_{i2}) \notin W'_2.\hat{\beta}$:

$$\forall (a_1, a_2) \in (W'_2.\hat{\beta}).$$

- i. When $a_1 = a_{i1}$ and $a_2 = a_{i2}$: This case cannot arise

- ii. When $a_1 = a_{i1}$ and $a_2 \neq a_{i2}$:

From Equation 116 we know that $(n - i - j, H'_{j1}, H'_{j2}) \stackrel{\mathcal{A}}{\triangleright} W'_2$ and since

$(a_{i1}, a_2) \in W'_2.\hat{\beta}$ therefore from Definition 2.9 we know that

$$(W'_2.\theta_1(a_{i1}) = W'_2.\theta_2(a_2) \wedge (W'_2, n - i - j - 1, H'_{j1}(a_{i1}), H'_{j2}(a_2)) \in [W'_2.\theta_1(a_{i1})]_V^{\mathcal{A}}) \quad (119)$$

Instantiating Equation 117 and Equation 118 with $n - i - j - 1$ we get $W'_1.\theta_1(a_{i1}) = \tau \sigma$ therefore from monotonicity we also have $W'_2.\theta_1(a_{i1}) = \tau \sigma$.

As a result from Equation 119 we get $W'_2.\theta_2(a_2) = \tau \sigma$

Also since from Equation 119 $(W'_2, n - i - j - 1, H'_{j1}(a_{i1}), H'_{j2}(a_2)) \in [\tau \sigma]_{\mathcal{V}}^A$ and $\tau \sigma \searrow \ell, \ell \sigma \not\sqsubseteq \mathcal{A}$ therefore from Lemma 2.15 we know that

$$\forall m.(W'_2.\theta_1, m, H'_{j1}(a_{i1})) \in \lfloor \tau \sigma \rfloor_V \quad (120)$$

$$\forall m.(W'_2.\theta_2, m, H'_{j2}(a_2)) \in \lfloor \tau \sigma \rfloor_V \quad (121)$$

Instantiating m with $n - i - j - 1$ in Equation 120 and Equation 121 to get

$$(W'_2.\theta_1, n - i - j - 1, H'_{j1}(a_{i1})) \in \lfloor \tau \sigma \rfloor_V$$

and

$$(W'_2.\theta_2, n - i - j - 1, H'_{j2}(a_2)) \in \lfloor \tau \sigma \rfloor_V$$

Since $H'_1(a_{i1}) = v'_{j1}$ and $H'_2(a_2) = H'_{j2}(a_2)$

Again from Equation 116 we know that $(W'_2, n - i - j, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_{\mathcal{V}}^A$. This means from Lemma 2.15 and instantiating it with $n - i - j - 1$ we get

$$(W'_2.\theta_1, n - i - j - 1, v'_{j1}) \in \lfloor (\tau) \sigma \rfloor_V \quad (122)$$

Therefore from Equation 121 and Equation 122 we have

$$(W'_2, n - i - j - 1, H'_1(a_{i1}), H'_2(a_2)) \in [\tau \sigma]_{\mathcal{V}}^A$$

iii. When $a_1 \neq a_{i1}$ and $a_2 = a_{i2}$:

Symmetric case as (ii)

iv. When $a_1 \neq a_{i1}$ and $a_2 \neq a_{i2}$:

From Equation 116 and Definition 2.9

* $\forall i \in \{1, 2\}.\forall m.\forall a_i \in \text{dom}(W'_2.\theta_i).(W'_2.\theta_i, m, H'_i(a_i)) \in \lfloor W'_2.\theta_i(a_i) \rfloor_V$:

When $i = 1$

Given some m

$\forall a_1 \in \text{dom}(W'_2.\theta_1)$.

· when $a_1 = a_{i1}$:

From Equation 116 we know that $(W'_2, v'_{j1}, v'_{j2}) \in [(\tau) \sigma]_{\mathcal{V}}^A$ thus from Lemma 2.15 we know that

$$(W'_2.\theta_1, H'_1(a_1)) \in \lfloor W'_2.\theta_1(a_1) \rfloor_V$$

· Otherwise:

From Equation 116 and Lemma 2.27

When $i = 2$

Similar reasoning as with $i = 1$

– $(W'_1, n - n', v'_1, v'_2) \in [(\text{unit}) \sigma]_{\mathcal{V}}^A$:

From evaluation rule assign we know that $v'_1 = v'_2 = ()$

Directly from Definition 2.4

13. FG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e_i : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_i : (\forall \alpha. (\ell_e, \tau))^\perp}$$

To prove: $(W, n, \Lambda e_i (\gamma \downarrow_1), \Lambda e_i (\gamma \downarrow_2)) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_E^A$

Say $e_1 = \Lambda e_i (\gamma \downarrow_1)$ and $e_2 = \Lambda e_i (\gamma \downarrow_2)$

From Definition of $[(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_E^A$ it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_V^A$$

This means that given $\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2)$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_V^A \quad (123)$$

IH1 $(W, n, (e_i (\gamma \downarrow_1), (e_i (\gamma \downarrow_2))) \in [\tau \sigma]_E^A$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e (\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e (\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A$$

We know from the evaluation rules that $H'_1 = H_1$, $H'_2 = H_2$, $v'_1 = e_1 = \Lambda e_i (\gamma \downarrow_1)$ and $v'_2 = e_2 = \Lambda e_i (\gamma \downarrow_2)$. We choose $W' = W$ and we know that $n' = 0$ we need to show the following:

- $W \sqsubseteq W$: From Definition 2.3
- $(n, H_1, H_2) \stackrel{A}{\triangleright} W$: Given
- $(W, n, v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\perp \sigma]_V^A$
Here $v'_1 = \Lambda e_i (\gamma \downarrow_1)$ and $v'_2 = \Lambda e_i (\gamma \downarrow_2)$

From Definition 2.4 it suffices to prove

$$\begin{aligned} & \forall W' \sqsupseteq W. \forall \ell' \in \mathcal{L}. \forall j < n. \\ & ((W', j, e_i(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau[\ell'/\alpha]]_E^A) \\ & \wedge \forall \theta_l \sqsupseteq W. \theta_l, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma) \\ & \wedge \forall \theta_l \sqsupseteq W. \theta_l, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma) \end{aligned}$$

This means given some $W' \sqsupseteq W$, $\ell' \in \mathcal{L}$ and $j < n$ we need to show that

$$\begin{aligned} & - \forall W' \sqsupseteq W. \forall \ell' \in \mathcal{L}. \forall j < n. \\ & ((W', j, e_i(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau[\ell'/\alpha]]_E^A): \end{aligned}$$

This means that given some $W' \sqsupseteq W$, $\ell' \in \mathcal{L}$, $j < n$ we need to prove

$$((W', j, e_i(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau[\ell'/\alpha]]_E^A)$$

From Definition 2.5 it suffices to show that

$$\forall H_{s1}, H_{s2}. (j, H_{s1}, H_{s2}) \triangleright^A W \wedge \forall m < j. (H_{s1}, e (\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau[\ell'/\alpha] \sigma]_V^A$$

This means for some H_{s1} and H_{s2} and some $m < j$ we are given $(j, H_{s1}, H_{s2}) \triangleright^A W \wedge m < j. (H_{s1}, e (\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2})$

And we need to show that

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau[\ell'/\alpha] \sigma]_V^A$$

We instantiate IH1 with H_{s1} , H_{s2} , m and $\sigma \cup \{\alpha \mapsto \ell'\}$ to obtain

$$\exists W'_1 \sqsupseteq W. (n - m, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - m, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A \cup \{\alpha \mapsto \ell'\}$$

Since $j < n$ therefore from Lemma 2.21 and Lemma 2.17 we get

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau[\ell'/\alpha] \sigma]_V^A$$

$$- \forall \theta_l \sqsupseteq W. \theta_l, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma):$$

From Lemma 2.25 we know that $(W'.\theta_1, \gamma \downarrow_1) \in [\Gamma]_V$. Therefore, we can apply Theorem 2.22 with $\sigma \cup \{\alpha \mapsto \ell''\}$

$$\forall k. (W'.\theta_1, k, e \gamma \downarrow_1) \in [\tau (\sigma \cup \{\alpha \mapsto \ell''\})]_E^{\ell_e (\sigma \cup \{\alpha \mapsto \ell''\})}$$

From Lemma 2.16 we get

$$\forall \theta_l \sqsupseteq W'. \theta_1. \forall k. (\theta_l, k, e \gamma \downarrow_1) \in [\tau (\sigma \cup \{\alpha \mapsto \ell''\})]_E^{\ell_e (\sigma \cup \{\alpha \mapsto \ell''\})}$$

$$- \forall \theta_l \sqsupseteq W. \theta_2, k, \ell'' \in \mathcal{L}. ((\theta_l, k, e_i[\ell''/\alpha]) \in [\tau]_E^{\ell_e} \sigma):$$

Similar reasoning as in the previous case

14. FG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \quad \ell'' \in \text{FV}(\Sigma) \quad \Sigma; \Psi \vdash_{pc} \ell \sqsubseteq \ell_e[\ell''/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell''/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell''/\alpha]}$$

To prove: $(W, n, (e[]) (\gamma \downarrow_1), (e[]) (\gamma \downarrow_2)) \in [(\tau[\ell''/\alpha]) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e[]) (\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e[]) (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

This further means that given

$$\forall H_1, H_2. (n, H_1, H_2) \triangleright^A W \wedge \forall n' < n. (H_1, (e[]) (\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e[]) (\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W. (n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A \quad (124)$$

$$\underline{\text{IH}} (W, n, (e) (\gamma \downarrow_1), (e) (\gamma \downarrow_2)) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}.(n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n. (H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W.(n - i, H'_1, H'_2) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_1, v'_2) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_V^A$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH and since the $(e[])$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore $\exists i < n' < n$ s.t $(H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$. Similarly $(e[])$ also reduces to value with $\gamma \downarrow_2$ therefore we also have $(H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_V^A \quad (125)$$

We case analyze on $(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\forall \alpha. (\ell_e, \tau))^\ell \sigma]_V^A$ from Equation 125

- Case $\ell \sigma \sqsubseteq \mathcal{A}$:

In this case from Definition 2.4 we know that

$$(W'_1, n - i, v'_{i1}, v'_{i2}) \in [(\forall \alpha. (\ell_e, \tau)) \sigma]_V^A$$

Here $v'_{i1} = \Lambda e_{i1}$ and $v'_{i2} = \Lambda e_{i2}$

This further means that we have

$$\begin{aligned} \forall W'' \sqsupseteq W'_1. \forall \ell' \in \mathcal{L}. \forall j < n - i. ((W'', j, e_{i1}, e_{i2}) \in [\tau[\ell'/\alpha]]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W'_1. \theta_l, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_{i1}) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]} \sigma) \\ \wedge \forall \theta_l \sqsupseteq W'_1. \theta_l, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_{i2}) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]} \sigma) \end{aligned} \quad (E1)$$

Instantiating the first conjunct of (E1) with W'_1 , ℓ'' and $n - i - 1$ we get

$$((W'_1, n - i - 1, e_{i1}, e_{i2}) \in [\tau[\ell''/\alpha]]_E^A)$$

Therefore from Definition 2.5 we get

$$\forall H_1, H_2. (n - i - 1, H_1, H_2) \triangleright^A W'_1 \wedge \forall k < (n - i - 1). (H_1, (e_{i1})(\gamma \downarrow_1)) \Downarrow_k (H'_1, v'_1) \wedge (H_2, (e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W''' \sqsupseteq W'_1. ((n - i - 1) - k, H'_1, H'_2) \triangleright^A W'_1 \wedge (W'_1, (n - i - 1) - k, v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

Instantiating H_1 and H_2 with H'_{i1} and H'_{i2} and since $e[]$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps and e with $\gamma \downarrow_1$ reduces in $i < n' < n$ steps. Therefore $\exists k < (n' - i - 1)$ steps in which e_{i1} reduces. Also since $e[]$ reduces to value with $\gamma \downarrow_2$ therefore e_{i2} must also reduce. As a result we get

$$\exists W''' \sqsupseteq W'_1. ((n - i - 1) - k, H'_1, H'_2) \triangleright^A W'_1 \wedge (W'_1, (n - i - 1) - k, v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

Since $n' = i + k + 1$ therefore we are done

- Case $\ell \sigma \not\sqsubseteq \mathcal{A}$:

From Equation 124 we know that we need to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

In this case since we know that $\ell \sigma \not\sqsubseteq \mathcal{A}$. Let $\tau[\ell''/\alpha] \sigma = \mathbf{A}^{\ell_i}$ and since $\tau[\ell''/\alpha] \sigma \searrow \ell \sigma$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove $\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$

From Definition 2.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau[\ell''/\alpha]) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau[\ell''/\alpha]) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V)$$

This means given m_1 and m_2 it suffices to prove:

$$(\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau[\ell''/\alpha]) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V \quad (126)$$

In this case from Definition 2.6 we know that

$$\forall m. (W'_1.\theta_1, m, \Lambda e_{h1}) \in [\forall \alpha. (\ell_e, \tau) \sigma]_V \quad (127)$$

$$\forall m. (W'_1.\theta_2, m, \Lambda e_{h2}) \in [\forall \alpha. (\ell_e, \tau) \sigma]_V \quad (128)$$

Applying Definition 2.6 on Equation 127 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h1}) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \text{ where } \theta = W'_1.\theta_1$$

We instantiate m with $m_1 + 2 + t_1$ where t_1 is the number of steps in which e_{h1} reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 2 + t_1). \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h1}) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \quad (\text{FB-FE1})$$

Instantiating θ' with $W'_1.\theta_1$, j_1 with $m_1 + t_1 + 1$ and ℓ' with ℓ''

$$\text{Therefore we get } (W'_1.\theta_1, m_1 + t_1 + 1, e_{h1}) \in [\tau[\ell''/\alpha] \sigma]_E^{\ell_e \sigma}$$

From Definition 2.7, we get

$$\begin{aligned} \forall H. (m_1 + t_1 + 1, H) \triangleright W'_1.\theta_1 \wedge \forall k_c < (m_1 + t_1 + 1). (H, e_{h1}) \downarrow_{k_c} (H'_1, v'_1) \implies \\ \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + t_1 + 1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + t_1 + 1 - k_c), v'_1) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

Since from Equation 125 we have

$$(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$$

Therefore from Lemma 2.27 we get

$$\forall m. (m, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating m with $m_1 + 1 + t_1$ we get

$$(m_1 + 1 + t_1, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating H with H'_{j1} from Equation 125 and k_c with t_1 , we get

$$\begin{aligned} \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \quad (\text{CF1}) \end{aligned}$$

Similarly applying Definition 2.6 to Equation 128 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h_2}[v/x]) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \text{ where } \theta = W'_1.\theta_2$$

We instantiate m with $m_2 + 1 + t_2$ where t_2 is the number of steps in which e_{h_2} reduces

$$\forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h_2}) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \quad (\text{FB-FE2})$$

Instantiating θ' with $W'_1.\theta_2$, j_1 with $m_2 + t_2 + 1$ and ℓ' with ℓ''

$$\text{Therefore we get } (W'_1.\theta_2, m_2 + t_2 + 1, e_{h_2}) \in [\tau[\ell''/\alpha]]_E^{\ell_e[\ell''/\alpha]} \sigma$$

From Definition 2.7, we get

$$\begin{aligned} \forall H. (m_2 + t_2 + 1, H) \triangleright W'_1.\theta_2 \wedge \forall k_c < (m_2 + t_2 + 1). (H, e_{h_2}) \Downarrow_{k_c} (H'_2, v'_1) \implies \\ \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + t_2 + 1 - k_c), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + t_2 + 1 - k_c), v'_1) \in [\tau[\ell''/\alpha]]_V \sigma \wedge \\ (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

Since from Equation 125 we have

$$(n - i, H'_{i_1}, H'_{i_2}) \stackrel{A}{\triangleright} W'_1$$

Therefore from Lemma 2.27 we get

$$\forall m. (m, H'_{i_2}) \triangleright W'_1.\theta_2$$

Instantiating m with $m_2 + 1 + t_2$ we get

$$(m_2 + 1 + t_2, H'_{i_2}) \triangleright W'_1.\theta_2$$

Instantiating H with H'_{j_2} from Equation 125 and k_c with t_2 , we get

$$\begin{aligned} \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in [\tau[\ell''/\alpha]]_V \sigma \wedge \\ (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \quad (\text{CF2}) \end{aligned}$$

In order to prove Equation 124 we choose W' to be $(\theta'_1, \theta'_2, W'_1.\beta)$. Now we need to show two things:

$$(a) (n - n', H'_1, H'_2) \triangleright W':$$

From Definition 2.9 it suffices to show that

$$- \text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2):$$

From CF1 we know that $(m_1 + 1, H'_1) \triangleright \theta'_1$, therefore from Definition 2.8 we get $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from CF2 we know that $(m_2 + 1, H'_2) \triangleright \theta'_2$, therefore from Definition 2.8 we get $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

$$- (W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2)):$$

Since $(n - i, H'_{j_1}, H'_{j_2}) \triangleright W'_1$ therefore from Definition 2.9 we know that

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$$

From CF1 and CF2 we know that $W'_1.\theta_1 \sqsubseteq \theta'_1$ and $W'_1.\theta_2 \sqsubseteq \theta'_2$ therefore

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$$

$$- \forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge \\ (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [W'.\theta_1(a_1)]_V^A:$$

4 cases arise for each a_1 and a_2

$$i. H'_{i_1}(a_1) = H'_1(a_1) \wedge H'_{i_2}(a_2) = H'_2(a_2):$$

* $W'.\theta_1(a_1) = W'.\theta_2(a_2)$:

We know from Equation 125 that $(n - i, H'_{i1}, H'_{i2}) \triangleright W'_1$

Therefore from Definition 2.9 we have

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since $W'.\hat{\beta} = W'_1.\hat{\beta}$ by construction therefore

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From CF1 and CF2 we know that $W'_1.\theta_1 \sqsubseteq \theta'_1$ and $W'_1.\theta_2 \sqsubseteq \theta'_2$ respectively.

Therefore from Definition 2.2

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

* $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

From Equation 125 we know that $(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$

This means from Definition 2.9 that

$$\forall (a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) \wedge (W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Instantiating with a_1 and a_2 and since $W'_1 \sqsubseteq W'$ and $n - n' - 1 < n - i - 1$ (since $i < n'$) therefore from Lemma 2.17 we get

$$(W', n - n' - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

ii. $H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2)$:

* $W'.\theta_1(a_1) = W'.\theta_2(a_2)$:

Same as in the previous case

* $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

From CF1 and CF2 we know that

$$(\forall a. H'_{j1}(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_{j2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell'$$

Since $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$ (given) and $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_e[\ell''/\alpha] \sigma \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Also from CF1 and CF2, $(m_1 + 1, H'_1) \triangleright \theta'_1$ and $(m_2 + 1, H'_2) \triangleright \theta'_2$. Therefore from Definition 2.8 we have

$$(\theta'_1, m_1, H'_1(a_1)) \in \lceil \theta'_1(a_1) \rceil_V$$

$$(\theta'_2, m_2, H'_2(a_1)) \in \lceil \theta'_2(a_2) \rceil_V$$

Since m_1 and m_2 are arbitrary indices therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

iii. $H'_{i1}(a_1) = H'_1(a_1) \vee H'_{i2}(a_2) \neq H'_2(a_2)$:

* $W'.\theta_1(a_1) = W'.\theta_2(a_2)$:

Same as in the previous case

* $(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

From CF2 we know that

$$(\forall a. H'_{i2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell')$$

This means that a_2 was protected at $\ell_e[\ell''/\alpha]$ σ in the world before the modification. Since $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e[\ell''/\alpha] \sigma$ (given) and $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_e[\ell''/\alpha] \sigma \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 125 we know that $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$ that means from Definition 2.9 that $(W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in [W'_1.\theta_1(a_1)]_V^A$. Since $(\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell'$ therefore from Definition 2.4 we know that $H'_{i1}(a_1)$ must also have a label $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_1.\theta_1, m, H'_{i1}(a_1)) \in W'_1.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_1.\theta_2, m, H'_{i2}(a_2)) \in W'_1.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with m_1 and using Lemma 2.16 we get

$$(\theta'_1, m_1, H'_{i1}(a_1)) \in \theta'_1(a_1)$$

Since from CF2 we know that $(m_2 + 1, H'_2) \triangleright \theta'_2$ therefore from Definition 2.8 we know that $(\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$\text{iv. } H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) = H'_2(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

$i = 1$

This means that given some m we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in [W.\theta_i(a_i)]_V$$

Like before we apply Theorem 2.22 on e_{h1} and e_{h2} but this time $m + 2 + t_1$ and $m + 2 + t_2$ where t_1 and t_2 are the number of steps in which e_{h1} and e_{h2} reduces respectively. This will give us

$$\begin{aligned} & \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

and

$$\begin{aligned} & \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_2) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in [\tau[\ell''/\alpha] \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e[\ell''/\alpha] \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e[\ell''/\alpha] \sigma)) \end{aligned}$$

Since we have $(m + 1, H'_1) \triangleright \theta'_1$ and $(m + 1, H'_2) \triangleright \theta'_2$ therefore we get the desired from Definition 2.8

$i = 2$

Symmetric to $i = 1$

$$(b) (W', n - n' - 1, v'_1, v'_2) \in [\tau[\ell''/\alpha] \sigma]_V^A:$$

Let $\tau[\ell''/\alpha] = \mathbf{A}^{\ell_i}$ Since $\tau[\ell''/\alpha] \sigma \searrow \ell \sigma$ and since $\ell \sigma \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From CF1 and CF2 we and Definition 2.4 we get the desired.

15. FG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp}$$

To prove: $(W, n, \nu e (\gamma \downarrow_1), \nu e (\gamma \downarrow_2)) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_E^A$

Say $e_1 = \nu e (\gamma \downarrow_1)$ and $e_2 = \nu e (\gamma \downarrow_2)$

From Definition of $[(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_E^A$ it suffices to prove that

$$\forall H_1, H_2. (n, H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2) \implies \\ \exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_V^A$$

This means that given $\forall H_1, H_2. (n', H_1, H_2) \stackrel{A}{\triangleright} W \wedge \forall n' < n. (H_1, e_1) \downarrow_{n'} (H'_1, v'_1) \wedge (H_2, e_2) \downarrow (H'_2, v'_2)$

We are required to prove:

$$\exists W'. W \sqsubseteq W' \wedge (n - n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n - n', v'_1, v'_2) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_V^A \quad (129)$$

IH1 $(W, n, (e) (\gamma \downarrow_1), (e) (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}. (n, H_{i1}, H_{i2}) \stackrel{A}{\triangleright} W \wedge \forall i < n. (H_{i1}, e (\gamma \downarrow_1)) \downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e (\gamma \downarrow_2)) \downarrow (H'_{i2}, v'_{i2}) \implies \\ \exists W'_1 \sqsupseteq W. (n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [\tau \sigma]_V^A$$

We know from the evaluation rules that $H'_1 = H_1$, $H'_2 = H_2$, $v'_1 = e_1 = \nu e (\gamma \downarrow_1)$ and $v'_2 = e_2 = \nu e (\gamma \downarrow_2)$. We choose $W' = W$ and we know that $n' = 0$. We need to show the following:

- $W \sqsubseteq W$: From Definition 2.3

- $(n, H_1, H_2) \stackrel{A}{\triangleright} W$: Given

- $(W, n, v'_1, v'_2) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \sigma]_V^A$

Here $v'_1 = \nu e (\gamma \downarrow_1)$ and $v'_2 = \nu e (\gamma \downarrow_2)$

From Definition 2.4 it suffices to prove

$$\forall W' \sqsupseteq W. \forall j < n. \mathcal{L} \models c \sigma \implies (W', j, e \gamma \downarrow_1, e \gamma \downarrow_2) \in [\tau \sigma]_E^A \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c \implies (\theta_l, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma} \wedge$$

$$\forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c \implies (\theta_l, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$$

We need to prove:

$$- \forall W' \sqsupseteq W. \forall j < n. \mathcal{L} \models c \sigma \implies (W', j, e \gamma \downarrow_1, e \gamma \downarrow_2) \in [\tau \sigma]_E^A:$$

This means given some $W' \sqsupseteq W$, $j < n$ and given that $\mathcal{L} \models c \sigma$ we need to show that

$$(W', j, e \gamma \downarrow_1, e \gamma \downarrow_2) \in [\tau \sigma]_E^A$$

From Definition 2.5 it suffices to show that

$$\forall H_{s1}, H_{s2}. (j, H_{s1}, H_{s2}) \stackrel{A}{\triangleright} W \wedge \forall m < j. (H_{s1}, e (\gamma \downarrow_1)) \downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e (\gamma \downarrow_2)) \downarrow (H'_{s2}, v'_{s2}) \implies$$

$$\exists W'_1 \sqsupseteq W. (j - m, H'_{s1}, H'_{s2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_V^A$$

This means for some $H_{s1}, H_{s2}, m < j$ s.t

$$(H_{s1}, H_{s2}) \triangleright^A W \wedge (H_{s1}, e(\gamma \downarrow_1)) \Downarrow_m (H'_{s1}, v'_{s1}) \wedge (H_{s2}, e(\gamma \downarrow_2)) \Downarrow (H'_{s2}, v'_{s2})$$

And we need to show that

$$\exists W'_1 \sqsupseteq W.(j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

We instantiate IH1 with H_{s1}, H_{s2} and m to obtain

$$\exists W'_1 \sqsupseteq W.(n - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, n - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

Since $j < n$ therefore from Lemma 2.21 and Lemma 2.17 we get

$$\exists W'_1 \sqsupseteq W.(j - m, H'_{s1}, H'_{s2}) \triangleright^A W'_1 \wedge (W'_1, j - m, v'_{s1}, v'_{s2}) \in [\tau \sigma]_{\mathcal{V}}^A$$

$$- \forall \theta_l \sqsupseteq W.\theta_1, j, \mathcal{L} \models c \implies (\theta_l, j, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}:$$

This means given $\theta_l \sqsupseteq W.\theta_1, j, \mathcal{L} \models c$

We need to prove: $(\theta_l, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$

From Lemma 2.25 we know that $\forall m_1. (W'.\theta_1, m_1, \gamma \downarrow_1) \in [\Gamma]_{\mathcal{V}}$. Therefore by instantiating m_1 at j we can apply Theorem 2.22 to get

$$(\theta_l, j, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}$$

$$- \forall \theta_l \sqsupseteq W.\theta_2, j, \mathcal{L} \models c \implies (\theta_l, j, e \gamma \downarrow_1) \in [\tau \sigma]_E^{\ell_e \sigma}:$$

Symmetric reasoning as in the previous case

16. FG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau}$$

To prove: $(W, n, (e \bullet)(\gamma \downarrow_1), (e \bullet)(\gamma \downarrow_2)) \in [(\tau) \sigma]_E^A$

This means from Definition 2.5 we need to prove:

$$\forall H_1, H_2.(n, H_1, H_2) \triangleright^A W \wedge \forall n' < n.(H_1, (e \bullet)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e \bullet)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) \implies$$

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_{\mathcal{V}}^A$$

This further means that given

$$\forall H_1, H_2.(n, H_1, H_2) \triangleright^A W \wedge \forall n' < n.(H_1, (e \bullet)(\gamma \downarrow_1)) \Downarrow_{n'} (H'_1, v'_1) \wedge (H_2, (e \bullet)(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2)$$

It suffices to prove

$$\exists W' \sqsupseteq W.(n - n', H'_1, H'_2) \triangleright^A W' \wedge (W', n - n', v'_1, v'_2) \in [(\tau) \sigma]_{\mathcal{V}}^A \quad (130)$$

$$\underline{\text{IH}} (W, n, (e)(\gamma \downarrow_1), (e)(\gamma \downarrow_2)) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_E^A$$

This means from Definition 2.5 we get

$$\forall H_{i1}, H_{i2}.(n, H_{i1}, H_{i2}) \triangleright^A W \wedge \forall i < n.(H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1}) \wedge (H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2}) \implies$$

$$\exists W'_1 \sqsupseteq W.(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1 \wedge (W'_1, n - i, v'_{i1}, v'_{i2}) \in [(c \xrightarrow{\ell_e} \tau)^\ell \sigma]_{\mathcal{V}}^A$$

Instantiating H_{i1} with H_1 and H_{i2} with H_2 in IH and since the $(e \bullet)$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps therefore $\exists i < n' < n$ s.t $(H_{i1}, e(\gamma \downarrow_1)) \Downarrow_i (H'_{i1}, v'_{i1})$. Similarly since

(e•) reduces to value with $\gamma \downarrow_2$ therefore also have $(H_{i2}, e(\gamma \downarrow_2)) \Downarrow (H'_{i2}, v'_{i2})$. Hence we get

$$\exists W'_1 \sqsupseteq W.(n-i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, n-i, v'_{i1}, v'_{i2}) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\ell \sigma]_V^A \quad (131)$$

We case analyze on $(W'_1, n-i, v'_{i1}, v'_{i2}) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\ell \sigma]_V^A$ from Equation 131

- Case $\ell \sigma \sqsubseteq \mathcal{A}$:

In this case from Definition 2.4 we know that

$$(W'_1, n-i, v'_{i1}, v'_{i2}) \in [(c \stackrel{\ell_e}{\Rightarrow} \tau)^\ell \sigma]_V^A$$

Here $v'_{i1} = \nu e_{i1}$ and $v'_{i2} = \nu e_{i2}$

This further means that we have

$$\begin{aligned} \forall W' \sqsupseteq W. \forall j < n-i. \mathcal{L} \models c \sigma &\implies ((W', j, e_{i1}, e_{i2}) \in [\tau \sigma]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c &\implies ((\theta_l, j, e_{i1}) \in [\tau \sigma]_E^{\ell_e} \sigma) \\ \wedge \forall \theta_l \sqsupseteq W. \theta_l, j. \mathcal{L} \models c &\implies ((\theta_l, j, e_{i2}) \in [\tau \sigma]_E^{\ell_e} \sigma) \end{aligned} \quad (CE1)$$

Instantiating the first conjunct of (CE1) with W'_1, ℓ'' and $n-i-1$ we get

$$((W'_1, n-i-1, e_{i1}, e_{i2}) \in [\tau \sigma]_E^A)$$

Therefore from Definition 2.5 we get

$$\begin{aligned} \forall H_1, H_2. (n-i-1, H_1, H_2) \stackrel{A}{\triangleright} W'_1 \wedge \forall k < (n-i-1). (H_1, (e_{i1})(\gamma \downarrow_1)) \Downarrow_k (H'_1, v'_1) \wedge \\ (H_2, (e_{i2})(\gamma \downarrow_2)) \Downarrow (H'_2, v'_2) &\implies \\ \exists W''' \sqsupseteq W'_1. ((n-i-1) - k, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, (n-i-1) - k, v'_1, v'_2) &\in [(\tau) \sigma]_V^A \end{aligned}$$

Instantiating H_1 and H_2 with H'_{i1} and H'_{i2} and since $e[]$ reduces to value with $\gamma \downarrow_1$ in $n' < n$ steps and e with $\gamma \downarrow_1$ reduces in $i < n' < n$ steps. Therefore $\exists k < (n' - i - 1)$ steps in which e_{i1} reduces. Also since $e[]$ reduces to value with $\gamma \downarrow_2$ therefore e_{i2} must also reduce. As a result we get

$$\exists W''' \sqsupseteq W'_1. ((n-i-1) - k, H'_1, H'_2) \stackrel{A}{\triangleright} W'_1 \wedge (W'_1, (n-i-1) - k, v'_1, v'_2) \in [(\tau[\ell''/\alpha]) \sigma]_V^A$$

Since $n' = i + k + 1$ therefore we are done

- Case $\ell \sigma \not\sqsubseteq \mathcal{A}$:

From Equation 130 we know that we need to prove

$$\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n-n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$$

In this case since we know that $\ell \sigma \not\sqsubseteq \mathcal{A}$. Let $\tau \sigma = \mathbf{A}^{\ell_i}$ and since $\tau \sigma \searrow \ell \sigma$ therefore $\ell_i \not\sqsubseteq \mathcal{A}$

This means in order to prove $\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (W', n-n', v'_1, v'_2) \in [(\tau) \sigma]_V^A$

From Definition 2.4 it will suffice to prove

$$\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \stackrel{A}{\triangleright} W' \wedge (\forall m_1. (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (\forall m_2. (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means it suffices to prove

$$(\forall m_1, m_2. \exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge ((W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V)$$

This means given m_1 and m_2 it suffices to prove:

$$(\exists W' \sqsupseteq W.(n-n', H'_1, H'_2) \triangleright^A W' \wedge (W'.\theta_1, m_1, v'_1) \in [(\tau) \sigma]_V) \wedge (W'.\theta_1, m_2, v'_2) \in [(\tau) \sigma]_V \quad (132)$$

In this case from Definition 2.6 we know that

$$\forall m. (W'_1.\theta_1, m, \nu e_{h1}) \in [(c \xrightarrow{\ell_e} \tau) \sigma]_V \quad (133)$$

$$\forall m. (W'_1.\theta_2, m, \nu e_{h2}) \in [(c \xrightarrow{\ell_e} \tau) \sigma]_V \quad (134)$$

Applying Definition 2.6 to Equation 133 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \mathcal{L} \models c \sigma \implies (\theta', j_1, e_{h1}) \in [\tau \sigma]_E^{\ell_e \sigma} \text{ where } \theta = W'_1.\theta_1$$

We instantiate m with $m_1 + 2 + t_1$ where t_1 is the number of steps in which e_{h1} reduces

$$\forall \theta'. W'_1.\theta_1 \sqsubseteq \theta' \wedge \forall j_1 < (m_1 + 2 + t_1). \mathcal{L} \models c \sigma \implies (\theta', j_1, e_{h1}) \in [\tau[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]} \quad (\text{FB-CE1})$$

Instantiating θ' with $W'_1.\theta_1$, j_1 with $m_1 + t_1 + 1$ and since we know that $\mathcal{L} \models c \sigma$. Therefore we get

$$(W'_1.\theta_1, m_1 + t_1 + 1, e_{h1}) \in [\tau \sigma]_E^{\ell_e \sigma}$$

From Definition 2.7, we get

$$\begin{aligned} \forall H. (m_1 + t_1 + 1, H) \triangleright W'_1.\theta_1 \wedge \forall k_c < (m_1 + t_1 + 1). (H, e_{h1}) \Downarrow_{k_c} (H'_1, v'_1) \implies \\ \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + t_1 + 1 - k_c), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + t_1 + 1 - k_c), v'_1) \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since from Equation 131 we have

$$(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$$

Therefore from Lemma 2.27 we get

$$\forall m. (m, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating m with $m_1 + 1 + t_1$ we get

$$(m_1 + 1 + t_1, H'_{i1}) \triangleright W'_1.\theta_1$$

Instantiating H with H'_{i1} from Equation 131 and k_c with t_1 , we get

$$\begin{aligned} \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_{i1}) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'_{i1}(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \quad (\text{CCE1}) \end{aligned}$$

Similarly applying Definition 2.6 to Equation 134 we get

$$\forall m. \forall \theta'. \theta \sqsubseteq \theta' \wedge \forall j_1 < m. \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h2}) \in [\tau \sigma]_E^{\ell_e[\ell'/\alpha]} \text{ where } \theta = W'_1.\theta_2$$

We instantiate m with $m_2 + 2 + t_2$ where t_2 is the number of steps in which e_{h_2} reduces $\forall \theta'. W'_1.\theta_2 \sqsubseteq \theta' \wedge \forall j_1 < (m_2 + 2 + t_2). \forall \ell' \in \mathcal{L}. (\theta', j_1, e_{h_2}) \in \lfloor \tau \rfloor_E^{\ell_e[\ell'/\alpha]}$ (FB-CE2)

Instantiating θ' with $W'_1.\theta_2$, j_1 with $m_2 + t_2 + 1$ and ℓ' with ℓ''

Therefore we get $(W'_1.\theta_2, m_2 + t_2 + 1, e_{h_2}) \in \lfloor \tau \sigma \rfloor_E^{\ell_e \sigma}$

From Definition 2.7, we get

$$\begin{aligned} \forall H.(m_2 + t_2, H) \triangleright W'_1.\theta_2 \wedge \forall k_c < (m_2 + t_2 + 1).(H, e_{h_2}) \Downarrow_{k_c} (H'_1, v'_1) \implies \\ \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + t_2 + 1 - k_c), H'_1) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + t_2 + 1 - k_c), v'_1) \in \lfloor \tau \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since from Equation 131 we have

$$(n - i, H'_{i1}, H'_{i2}) \stackrel{A}{\triangleright} W'_1$$

Therefore from Lemma 2.27 we get

$$\forall m. (m, H'_{i2}) \triangleright W'_1.\theta_2$$

Instantiating m with $m_2 + 1 + t_2$ we get

$$(m_2 + 1 + t_2, H'_{i2}) \triangleright W'_1.\theta_2$$

Instantiating H with H'_{i2} from Equation 125 and k_c with t_2 , we get

$$\begin{aligned} \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_1) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in \lfloor \tau \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \quad (\text{CCE2}) \end{aligned}$$

In order to prove Equation 130 we choose W' to be $(\theta'_1, \theta'_2, W'_1.\beta)$. Now we need to show two things:

(a) $(n - n', H'_1, H'_2) \triangleright W'$:

From Definition 2.9 it suffices to show that

– $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$:

From CCE1 we know that $(m_1 + 1, H'_1) \triangleright \theta'_1$, therefore from Definition 2.8 we get $\text{dom}(W'.\theta_1) \subseteq \text{dom}(H'_1)$

Similarly, from CCE2 we know that $(m_2 + 1, H'_2) \triangleright \theta'_2$, therefore from Definition 2.8 we get $\text{dom}(W'.\theta_2) \subseteq \text{dom}(H'_2)$

– $(W'.\hat{\beta}) \subseteq (\text{dom}(W'.\theta_1) \times \text{dom}(W'.\theta_2))$:

Since $(n - i, H'_{j1}, H'_{j2}) \triangleright W'_1$ therefore from Definition 2.9 we know that

$$(W'_1.\hat{\beta}) \subseteq (\text{dom}(W'_1.\theta_1) \times \text{dom}(W'_1.\theta_2))$$

From CCE1 and CCE2 we know that $W'_1.\theta_1 \sqsubseteq \theta'_1$ and $W'_1.\theta_2 \sqsubseteq \theta'_2$ therefore $(W'_1.\hat{\beta}) \subseteq (\text{dom}(\theta'_1) \times \text{dom}(\theta'_2))$

– $\forall (a_1, a_2) \in (W'.\hat{\beta}). W'.\theta_1(a_1) = W'.\theta_2(a_2) \wedge (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$:

4 cases arise for each a_1 and a_2

i. $H'_{i1}(a_1) = H'_1(a_1) \wedge H'_{i2}(a_2) = H'_2(a_2)$:

* $W'.\theta_1(a_1) = W'.\theta_2(a_2)$

We know from Equation 125 that $(n - i, H'_{i1}, H'_{i2}) \triangleright W'_1$

Therefore from Definition 2.9 we have

$$\forall (a_1, a_2) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

Since $W'.\hat{\beta} = W'_1.\hat{\beta}$ by construction therefore

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2)$$

From CCE1 and CCE2 we know that $W'_1.\theta_1 \sqsubseteq \theta'_1$ and $W'_1.\theta_2 \sqsubseteq \theta'_2$ respectively.

Therefore from Definition 2.2

$$\forall (a_1, a_2) \in (W'.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

From Equation 131 we know that $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$

This means from Definition 2.9 that

$$\forall (a_{i1}, a_{i2}) \in (W'_1.\hat{\beta}). W'_1.\theta_1(a_1) = W'_1.\theta_2(a_2) \wedge (W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$$

Instantiating with a_1 and a_2 and since $W'_1 \sqsubseteq W'$ and $n - n' - 1 < n - i - 1$

(since $i < n'$) therefore from Lemma 2.17 we get

$$(W', n - n' - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A$$

$$\text{ii. } H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) \neq H'_2(a_2):$$

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2)$$

Same as in the previous case

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

From CCE1 and CCE2 we know that

$$(\forall a. H'_{j1}(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_{j2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W'_1.\theta_1(a_1) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W'_1.\theta_2(a_2) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell'$$

Since $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$ (given) and $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_e \sigma \not\sqsubseteq \mathcal{A}$. And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Also from CCE1 and CCE2, $(m_1 + 1, H'_1) \triangleright \theta'_1$ and $(m_2 + 1, H'_2) \triangleright \theta'_2$.

Therefore from Definition 2.8 we have

$$(\theta'_1, m_1, H'_1(a_1)) \in \lceil \theta'_1(a_1) \rceil_V$$

$$(\theta'_2, m_2, H'_2(a_1)) \in \lceil \theta'_2(a_2) \rceil_V$$

Since m_1 and m_2 are arbitrary indices therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

$$\text{iii. } H'_{i1}(a_1) = H'_1(a_1) \vee H'_{i2}(a_2) \neq H'_2(a_2):$$

$$* W'.\theta_1(a_1) = W'.\theta_2(a_2)$$

Same as in the previous case

$$* (W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W'.\theta_1(a_1) \rceil_V^A:$$

From CCE2 we know that

$$(\forall a. H'_{i2}(a) \neq H'_2(a) \implies \exists \ell'. W'_1.\theta_2(a) = A^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell')$$

This means that a_2 was protected at $\ell_e \sigma$ in the world before the modification. Since $pc \sigma \sqcup \ell \sigma \sqsubseteq \ell_e \sigma$ (given) and $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $\ell_e \sigma \not\sqsubseteq \mathcal{A}$.

And thus, $\ell' \not\sqsubseteq \mathcal{A}$

Since from Equation 131 we know that $(n - i, H'_{i1}, H'_{i2}) \triangleright^A W'_1$ that means from Definition 2.9 that $(W'_1, n - i - 1, H'_{i1}(a_1), H'_{i2}(a_2)) \in \lceil W'_1.\theta_1(a_1) \rceil_V^A$. Since $(\ell_e \sigma) \sqsubseteq \ell'$ therefore from Definition 2.4 we know that $H'_{i1}(a_1)$ must have a label $\not\sqsubseteq \mathcal{A}$

Therefore

$$\forall m. (W'_1.\theta_1, m, H'_{i1}(a_1)) \in W'_1.\theta_1(a_1) \quad (\text{F})$$

and

$$\forall m. (W'_1.\theta_2, m, H'_{i2}(a_2)) \in W'_1.\theta_2(a_1) \quad (\text{S})$$

Instantiating the (F) with m_1 and using Lemma 2.16 we get

$$(\theta'_1, m_1, H'_{i1}(a_1)) \in \theta'_1(a_1)$$

Since from CCE2 we know that $(m_2 + 1, H'_2) \triangleright \theta'_2$ therefore from Definition 2.8 we know that $(\theta'_2, m_2, H'_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 2.4 we get

$$(W', n - n' - 1, H'_1(a_1), H'_2(a_2)) \in \lceil \theta'_1(a_1) \rceil_V^A$$

$$\text{iv. } H'_{j1}(a_1) \neq H'_1(a_1) \vee H'_{j2}(a_2) = H'_2(a_2):$$

Symmetric case as above

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$$

$i = 1$

This means that given some m we need to prove

$$\forall a_i \in \text{dom}(W'.\theta_i). (W'.\theta_i, m, H'_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$$

Like before we apply Theorem 2.22 on e_{h1} and e_{h2} but this time $m + 2 + t_1$ and $m + 2 + t_2$ where t_1 and t_2 are the number of steps in which e_{h1} and e_{h2} reduces respectively. This will give us

$$\begin{aligned} & \exists \theta'_1. W'_1.\theta_1 \sqsubseteq \theta'_1 \wedge ((m_1 + 1), H'_1) \triangleright \theta'_1 \wedge (\theta'_1, (m_1 + 1), v'_1) \in \lceil \tau \sigma \rceil_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_1(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(W'_1.\theta_1). \theta'_1(a) \searrow (\ell_e \sigma)) \end{aligned}$$

and

$$\begin{aligned} & \exists \theta'_2. W'_1.\theta_2 \sqsubseteq \theta'_2 \wedge ((m_2 + 1), H'_1) \triangleright \theta'_2 \wedge (\theta'_2, (m_2 + 1), v'_1) \in \lceil \tau \sigma \rceil_V \wedge \\ & (\forall a. H(a) \neq H'_1(a) \implies \exists \ell'. W'_1.\theta_2(a) = \mathbf{A}^{\ell'} \wedge (\ell_e \sigma) \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta'_2) \setminus \text{dom}(W'_1.\theta_2). \theta'_2(a) \searrow (\ell_e \sigma)) \end{aligned}$$

Since we have $(m + 1, H'_1) \triangleright \theta'_1$ and $(m + 1, H'_2) \triangleright \theta'_2$ therefore we get the desired from Definition 2.8

$i = 2$

Symmetric to $i = 1$

$$(b) (W', n - n' - 1, v'_1, v'_2) \in \lceil \tau \sigma \rceil_V^A:$$

Let $\tau = \mathbf{A}^{\ell_i}$ Since $\tau \sigma \searrow \ell \sigma$ and since $\ell \sigma \not\sqsubseteq \mathcal{A}$ therefore $\ell_i \sigma \not\sqsubseteq \mathcal{A}$

From CCE1 and CCE2 we and Definition 2.4 we get the desired.

□

Lemma 2.27 (FG: Binary heap well formedness implies unary heap well formedness). $\forall H_1, H_2, W. (n, H_1, H_2) \triangleright W \implies \forall i \in \{1, 2\}. \forall m. (m, H_i) \triangleright W.\theta_i$

Proof. Directly from Definition 2.9

□

Lemma 2.28 (FG: Subtyping binary). *The following holds:*

$\forall \Sigma, \Psi, \sigma.$

1. $\forall A, A'.$

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (A \sigma) \rrbracket_V^A \subseteq \llbracket (A' \sigma) \rrbracket_V^A$$

2. $\forall \tau, \tau'.$

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_V^A \subseteq \llbracket (\tau' \sigma) \rrbracket_V^A$$

$$(b) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies \llbracket (\tau \sigma) \rrbracket_E^A \subseteq \llbracket (\tau' \sigma) \rrbracket_E^A$$

Proof. Proof by simultaneous induction on $A <: A'$ and $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of A in the last step:

1. FGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

To prove: $\llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V^A \subseteq \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V^A$

IH1: $\llbracket (\tau'_1 \sigma) \rrbracket_V^A \subseteq \llbracket (\tau_1 \sigma) \rrbracket_V^A$

IH2: $\llbracket (\tau_2 \sigma) \rrbracket_E^A \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_E^A$

It suffices to prove:

$$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V^A$$

This means that given: $(W, n, \lambda x.e_1, \lambda x.e_2) \in \llbracket ((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma) \rrbracket_V^A$

And it suffices to prove: $(W, n, \lambda x.e_1, \lambda x.e_2) \in \llbracket ((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma) \rrbracket_V^A$

From Definition 2.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in \llbracket \tau_1 \sigma \rrbracket_V^A \implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j, v_c. ((\theta_l, j, v_c) \in \llbracket \tau_1 \sigma \rrbracket_V \implies (\theta_l, j, e_1[v_1/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^{\ell_e} \sigma) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j, v_c. ((\theta_l, j, v_c) \in \llbracket \tau_1 \sigma \rrbracket_V \implies (\theta_l, j, e_2[v_c/x]) \in \llbracket \tau_2 \sigma \rrbracket_E^{\ell_e} \sigma) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 2.4 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in \llbracket \tau'_1 \sigma \rrbracket_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ \llbracket \tau'_2 \sigma \rrbracket_E^A) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in \llbracket \tau'_1 \sigma \rrbracket_V \implies (\theta'_l, k, e_1[v'_c/x]) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\ell'_e} \sigma) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in \llbracket \tau'_1 \sigma \rrbracket_V \implies (\theta'_l, k, e_2[v'_c/x]) \in \llbracket \tau'_2 \sigma \rrbracket_E^{\ell'_e} \sigma) \end{aligned}$$

This means given some $W'' \sqsupseteq W, k < n$ and v'_1, v'_2 we need to prove:

(a) $\forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A)$:

Given: $W'' \sqsupseteq W, k < n$ and v'_1, v'_2 . We are also given $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$

To prove: $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

Instantiating the first conjunct of Sub-A1 with W'', k, v'_1 and v'_2 we get

$$((W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A) \quad (135)$$

Since $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$ therefore from IH1 we know that $(W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A$

Thus from Equation 135 we get $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A$

Finally using IH2 we get $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

(b) $\forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma})$:

Given: $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$. We are also given $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$

To prove: $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$

Since we are given $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$ and since $\tau'_1 \sigma <: \tau_1 \sigma$ therefore from Lemma 2.24 we get

$$(\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \quad (136)$$

Instantiating the second conjunct of Sub-A1 with θ'_l, k, v'_1 and v'_2 we get

$$((\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}) \quad (137)$$

Therefore from Equation 136 and 137 we get $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2 \sigma]_E^{\ell_e \sigma}$

Since $\tau_2 \sigma <: \tau'_2 \sigma$ and $\ell'_e \sigma \sqsubseteq \ell_e \sigma$ therefore from Lemma 2.24 and 2.23 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma}$$

(c) $\forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E^{\ell'_e \sigma})$:

Similar reasoning as in the previous case

2. FGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove: $\lceil ((\tau_1 \times \tau_2) \sigma) \rceil_V^A \subseteq \lceil ((\tau'_1 \times \tau'_2) \sigma) \rceil_V^A$

IH1: $\lceil (\tau_1 \sigma) \rceil_V^A \subseteq \lceil (\tau'_1 \sigma) \rceil_V^A$

IH2: $\lceil (\tau_2 \sigma) \rceil_V^A \subseteq \lceil (\tau'_2 \sigma) \rceil_V^A$

It suffices to prove: $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau_1 \times \tau_2) \sigma) \rceil_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau'_1 \times \tau'_2) \sigma) \rceil_V^A$

This means that given: $(W, n, (v_1, v_2), (v'_1, v'_2)) \in \lceil ((\tau_1 \times \tau_2) \sigma) \rceil_V^A$

Therefore from Definition 2.4 we are given:

$$(W, n, v_1, v'_1) \in [\tau_1 \sigma]_{\mathcal{V}}^A \wedge (W, n, v_2, v'_2) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (138)$$

And it suffices to prove: $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_{\mathcal{V}}^A$

Again from Definition 2.4, it suffices to prove:

$$(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_{\mathcal{V}}^A \wedge (W, n, v_2, v'_2) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

Since from Equation 138 we know that $(W, n, v_1, v'_1) \in [\tau_1 \sigma]_{\mathcal{V}}^A$ therefore from IH1 we have $(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$

Similarly since $(W, n, v_2, v'_2) \in [\tau_2 \sigma]_{\mathcal{V}}^A$ from Equation 138 therefore from IH2 we have $(W, n, v_2, v'_2) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$

3. FGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove: $[((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A \subseteq [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

IH1: $[(\tau_1 \sigma)]_{\mathcal{V}}^A \subseteq [(\tau'_1 \sigma)]_{\mathcal{V}}^A$

IH2: $[(\tau_2 \sigma)]_{\mathcal{V}}^A \subseteq [(\tau'_2 \sigma)]_{\mathcal{V}}^A$

It suffices to prove: $\forall (W, n, v_{s1}, v_{s2}) \in [((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A. (W, n, v_{s1}, v_{s2}) \in [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

This means that given: $(W, n, v_{s1}, v_{s2}) \in [((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A$

And it suffices to prove: $(W, n, v_{s1}, v_{s2}) \in [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

2 cases arise

(a) $v_{s1} = \text{inl } v_{i1}$ and $v_{s1} = \text{inl } v_{i2}$:

From Definition 2.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (139)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$$

From Equation 139 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$$

(b) $v_s = \text{inr } v_{i1}$ and $v_{s2} = \text{inr } v_{i2}$:

From Definition 2.4 we are given:

$$(W, n, v_{i1}, v_{i2}) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (140)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

From Equation 140 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

4. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

To prove: $\lceil ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rceil_V^A \subseteq \lceil (\forall \alpha. (\ell'_e, \tau_2)) \sigma \rceil_V^A$

IH1: $\lceil (\tau_1 \sigma) \rceil_V^A \subseteq \lceil (\tau_2 \sigma) \rceil_V^A$

IH2: $\lceil (\tau_1 \sigma) \rceil_E^A \subseteq \lceil (\tau_2 \sigma) \rceil_E^A$

It suffices to prove: $\forall (W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rceil_V^A$.

$(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rceil_V^A$

This means that given: $(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell_e, \tau_1)) \sigma) \rceil_V^A$

Therefore from Definition 2.4 we are given:

$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau_1[\ell'/\alpha] \sigma]_E^A) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau_1[\ell'/\alpha]]_E^{\ell_e[\ell'/\alpha]}) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau_1[\ell''/\alpha]]_E^{\ell_e[\ell'/\alpha]})$ (Sub-F1)

And it suffices to prove: $(W, n, \Lambda e_1, \Lambda e_2) \in \lceil ((\forall \alpha. (\ell'_e, \tau_2)) \sigma) \rceil_V^A$

Again from Definition 2.4, it suffices to prove:

$\forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A) \wedge$

$\forall \theta'_l \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]}) \wedge$

$\forall \theta'_l \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]})$

This means we are required to show:

(a) $\forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$:

By instantiating the first conjunct of Sub-F1 with W'', n'' and ℓ'' we know that the following holds

$$((W'', n'', e_1, e_2) \in [\tau_1[\ell''/\alpha] \sigma]_E^A)$$

Therefore from IH1 instantiated at $\sigma \cup \{\alpha \mapsto \ell''\}$

$$((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$$

(b) $\forall \theta'_l \sqsupseteq W.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]})$:

By instantiating the second conjunct of Sub-F1 with θ'_l and ℓ'' we know that the following holds

$$((\theta'_l, k, e_1) \in [\tau_1[\ell''/\alpha] \sigma]_E^{\ell_e[\ell''/\alpha] \sigma})$$

Since $\tau_1 \sigma <: \tau_2 \sigma$ and $\ell'_e \sigma \sqsubseteq \ell_e \sigma$ therefore from Lemma 2.24 and Lemma 2.23 we know that

$$((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha] \sigma]_E^{\ell'_e[\ell''/\alpha] \sigma})$$

(c) $\forall \theta'_l \sqsupseteq W.\theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E^{\ell'_e[\ell''/\alpha]})$:

Similar reasoning as in the previous case

5. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \Longrightarrow c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell_e} \tau_1 <: c_2 \xrightarrow{\ell'_e} \tau_2} \text{FGsub-constraint}$$

To prove: $\lceil ((c_1 \xrightarrow{\ell_e} \tau_1) \sigma) \rceil_V^A \subseteq \lceil ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rceil_V^A$

IH: $\lceil (\tau_1 \sigma) \rceil_E^A \subseteq \lceil (\tau_2 \sigma) \rceil_E^A$

It suffices to prove: $\forall (W, n, \nu e_1, \nu e_2) \in \lceil ((c_1 \xrightarrow{\ell_e} \tau_1) \sigma) \rceil_V^A. (W, n, \nu e_1, \nu e_2) \in \lceil ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rceil_V^A$

This means that given: $(W, n, \nu e_1, \nu e_2) \in \lceil ((c_1 \xrightarrow{\ell_e} \tau_1) \sigma) \rceil_V^A$

Therefore from Definition 2.4 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c_1 \sigma &\Longrightarrow (W', n', e_1, e_2) \in \lceil \tau_1 \sigma \rceil_E^A \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_1 &\Longrightarrow (\theta_l, k, e_1) \in \lceil \tau_1 \sigma \rceil_E^{\ell'_e \sigma} \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, k. \mathcal{L} \models c_1 &\Longrightarrow (\theta_l, k, e_2) \in \lceil \tau_1 \sigma \rceil_E^{\ell_e \sigma} \quad (\text{Sub-C1}) \end{aligned}$$

And it suffices to prove: $(W, n, \nu e_1, \nu e_2) \in \lceil ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rceil_V^A$

Again from Definition 2.4, it suffices to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma &\Longrightarrow (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c_2 &\Longrightarrow (\theta'_l, j, e_1) \in \lceil \tau_2 \sigma \rceil_E^{\ell'_e \sigma} \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 &\Longrightarrow (\theta'_l, j, e_2) \in \lceil \tau_2 \sigma \rceil_E^{\ell_e \sigma} \end{aligned}$$

This means that we are required to show the following:

(a) $\forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \Longrightarrow (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A$:

We are given $W'' \sqsupseteq W, n'' < n$ also we know that $\mathcal{L} \models c_2 \sigma$ and $c_2 \sigma \Longrightarrow c_1 \sigma$ therefore we also know that $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the first conjunct of Sub-C1 with W'' and n'' we know that the following holds

$$(W'', n'', e_1, e_2) \in \lceil \tau_1 \sigma \rceil_E^A$$

Therefore from IH we get $(W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A$

(b) $\forall \theta'_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_2 \Longrightarrow (\theta'_l, k, e_1) \in \lceil \tau_2 \sigma \rceil_E^{\ell'_e \sigma}$:

We are given some $\theta'_l \sqsupseteq W.\theta_1, k$, also we know that $\mathcal{L} \models c_2 \sigma$ and $c_2 \sigma \Longrightarrow c_1 \sigma$ therefore we also know that $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the second conjunct of Sub-C1 with θ'_l we know that the following holds

$$(\theta'_l, k, e_1) \in \lceil \tau_1 \sigma \rceil_E^{\ell_e \sigma}$$

Since $\tau_1 \sigma <: \tau_2 \sigma$ and $\ell'_e \sigma \sqsubseteq \ell_e \sigma$ therefore from Lemma 2.23 and Lemma 2.24 we get

$$(\theta'_l, k, e_1) \in \lceil \tau_2 \sigma \rceil_E^{\ell'_e \sigma}$$

(c) $\forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \Longrightarrow (\theta'_l, j, e_2) \in \lceil \tau_2 \sigma \rceil_E^{\ell_e \sigma}$:

Similar reasoning as in the previous case

6. FGsub-ref:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove: $\lceil ((\text{ref } \tau) \sigma) \rceil_V^A \subseteq \lceil ((\text{ref } \tau) \sigma) \rceil_V^A$

Directly from Definition 2.4

7. FGsub-base:

Given:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

To prove: $\lceil ((\mathbf{b}) \sigma) \rceil_V^A \subseteq \lceil ((\mathbf{b}) \sigma) \rceil_V^A$

Directly from Definition 2.4

8. FGsub-unit:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

To prove: $\lceil ((\text{unit}) \sigma) \rceil_V^A \subseteq \lceil ((\text{unit}) \sigma) \rceil_V^A$

Directly from Definition 2.4

Proof of statement 2(a)

Given:

$$\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash A <: A'}{\Sigma; \Psi \vdash A^\ell <: A^{\ell'}} \text{FGsub-label}$$

To prove: $\lceil ((A^\ell) \sigma) \rceil_V^A \subseteq \lceil ((A^{\ell'}) \sigma) \rceil_V^A$

2 cases arise

1. $\ell \sigma \sqsubseteq \ell' \sigma$:

From Definition 2.4 it suffices to prove: $\lceil ((A) \sigma) \rceil_V^A \subseteq \lceil ((A') \sigma) \rceil_V^A$

This we get directly from IH (Statement (1))

2. $\ell \sigma \not\sqsubseteq \ell' \sigma$:

We need to prove that

$$\forall (W, n, v_1, v_2) \in \lceil A \sigma \rceil_V^A. (W, n, v_1, v_2) \in \lceil A' \sigma \rceil_V^A$$

From Definition 2.4 it suffices to prove:

$$\forall i \in \{1, 2\}. \forall m. (W(n). \theta_i, m, v_i) \in \lceil A \sigma \rceil_V. (W(n). \theta_i, m, v_i) \in \lceil A \rceil_V \in \lceil A' \sigma \rceil_V$$

Since $A \sigma <: A' \sigma$ therefore from Lemma 2.24 we get the desired

Proof of statement 2(b)

Given: $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma$

To prove: $\lceil (\tau \sigma) \rceil_E^A \subseteq \lceil (\tau' \sigma) \rceil_E^A$

This means we need to prove that

$\forall (W, n, e_1, e_2) \in \lceil (\tau \sigma) \rceil_E^A. (W, n, e_1, e_2) \in \lceil (\tau' \sigma) \rceil_E^A$

This means given $\forall (W, n, e_1, e_2) \in \lceil (\tau \sigma) \rceil_E^A$

It suffices to prove that $(W, n, e_1, e_2) \in \lceil (\tau' \sigma) \rceil_E^A$

From Definition 2.5 we know we are given:

$$\forall H_1, H_2, j < n. (n, H_1, H_2) \triangleright^A W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \\ \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \triangleright^A W' \wedge (W', n - j, v'_1, v'_2) \in \lceil \tau \sigma \rceil_V^A \quad (\text{Sub-exp1})$$

And we need prove that

$$\forall H_{21}, H_{22}, k < n. (n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22}) \implies \\ \exists W'' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in \lceil \tau \sigma \rceil_V^A$$

This means that we are given some H_{21}, H_{22} and $k < n$ such that $(n, H_{21}, H_{22}) \triangleright^A W \wedge (H_{21}, e_1) \Downarrow_k (H'_{21}, v'_{21}) \wedge (H_{22}, e_2) \Downarrow (H'_{22}, v'_{22})$

It suffices to prove:

$$\exists W'' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W'' \wedge (W'', n - k, v'_{21}, v'_{22}) \in \lceil \tau \sigma \rceil_V^A \quad (141)$$

Instantiating (Sub-exp1) with H_{21}, H_{22} and k we get

$$\exists W' \sqsupseteq W. (n - k, H'_{21}, H'_{22}) \triangleright^A W' \wedge (W', n - k, v'_{21}, v'_{22}) \in \lceil \tau \sigma \rceil_V^A \quad (142)$$

We choose W'' in Equation 141 as W' from Equation 142 and we are done

□

Theorem 2.29 (FG: NI). *Say* $\text{bool} = (\text{unit} + \text{unit})$

$\forall v_1, v_2, e, \tau, n_1.$

$\emptyset; \emptyset; \emptyset \vdash_{\perp} v_1 : \text{bool}^{\top} \wedge \emptyset; \emptyset; \emptyset \vdash_{\perp} v_2 : \text{bool}^{\top}$

$\emptyset; \emptyset; x : \text{bool}^{\top} \vdash_{\perp} e : \text{bool}^{\perp} \wedge$

$(\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow (-, v'_2) \implies$

$v'_1 = v'_2$

Proof. Given some

$\emptyset; \emptyset; \emptyset \vdash_{\perp} v_1 : \text{bool}^{\top} \wedge \emptyset; \emptyset; \emptyset \vdash_{\perp} v_2 : \text{bool}^{\top}$

$\emptyset; \emptyset; x : \text{bool}^{\top} \vdash_{\perp} e : \text{bool}^{\perp} \wedge$

$(\emptyset, e[v_1/x]) \Downarrow_{n_1} (-, v'_1) \wedge (\emptyset, e[v_2/x]) \Downarrow (-, v'_2)$

We need to prove

$v'_1 = v'_2$

From Theorem 2.26 we have

$\forall n. (\emptyset, n, v_1, v_2) \in \lceil \text{bool}^{\top} \rceil_E^{\perp}$

Therefore from Theorem 2.26 and from Definition 2.14 we have

$\forall n. (\emptyset, n, e[v_1/x], e[v_1/x]) \in \lceil \text{bool}^{\perp} \rceil_E^{\perp}$

Therefore from Definition 2.5 we know that

$$\forall n. (\forall H_1, H_2, j < n. (n, H_1, H_2) \overset{A}{\triangleright} W \wedge (H_1, e_1) \Downarrow_j (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow (H'_2, v'_2) \implies \exists W' \sqsupseteq W. (n - j, H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', n - j, v'_1, v'_2) \in \lceil (\text{unit} + \text{unit})^\perp \rceil_V^A)$$

Instantiating with $n_1 + 1$ and then with $\emptyset, \emptyset, n_1$ we get

$$\exists W' \sqsupseteq W. (1, H'_1, H'_2) \overset{A}{\triangleright} W' \wedge (W', 1, v'_1, v'_2) \in \lceil (\text{unit} + \text{unit})^\perp \rceil_V^A$$

Since we have $(W', 1, v'_1, v'_2) \in \lceil (\text{unit} + \text{unit})^\perp \rceil_V^A$ therefore from Definition 2.4 we get $v'_1 = v'_2$
 \square

2.2 Coarse-grained IFC enforcement (CG)

2.2.1 CG type system

Syntax, types, constraints:

Expressions	$e ::= x \mid \lambda x.e \mid e e \mid (e, e) \mid \text{fst}(e) \mid \text{snd}(e) \mid \text{inl}(e) \mid \text{inr}(e) \mid \text{case}(e, x.e, y.e) \mid \text{new } e \mid !e \mid e := e \mid () \mid \Lambda e \mid e [] \mid \nu e \mid e \bullet \mid \text{Lb}(e) \mid \text{unlabel}(e) \mid \text{toLabeled}(e) \mid \text{ret}(e) \mid \text{bind}(e, x.e)$
Labels	$\ell ::= l \mid \alpha \mid \ell \sqcup \ell \mid \ell \sqcap \ell$
Types	$\tau ::= \mathbf{b} \mid \tau \rightarrow \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{ref } \ell \tau \mid \text{unit} \mid \forall \alpha. \tau \mid c \Rightarrow \tau \mid \text{Labeled } \ell \tau \mid \mathbb{C} \ell_i \ell_o \tau$
Constraints	$c ::= \ell \sqsubseteq \ell \mid (c, c)$

Type system: $\boxed{\Sigma; \Psi; \Gamma \vdash e : \tau}$

(All rules of the simply typed lambda-calculus pertaining to the types $\mathbf{b}, \tau \rightarrow \tau, \tau \times \tau, \tau + \tau, \text{unit}$ are included.)

$\frac{\Sigma; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e) : \text{Labeled } \ell \tau} \text{CG-label}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e) : \mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau} \text{CG-unlabel}$
$\frac{\Sigma; \Psi; \Gamma \vdash e : \mathbb{C} \ell_i \ell_o \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)} \text{CG-toLabeled}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e) : \mathbb{C} \ell_i \ell_i \tau} \text{CG-ret}$
$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \mathbb{C} \ell_i \ell \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell \ell_o \tau'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell_i \ell_o \tau'} \text{CG-bind}$	
$\frac{\Sigma; \Psi; \Gamma \vdash e : \tau' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash e : \tau} \text{CG-sub}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e : \mathbb{C} \ell \ell (\text{ref } \ell' \tau)} \text{CG-ref}$
$\frac{\Sigma; \Psi; \Gamma \vdash e : \text{ref } \ell \tau}{\Sigma; \Psi; \Gamma \vdash !e : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)} \text{CG-deref}$	
$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \ell \text{ unit}} \text{CG-assign}$	
$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e : \tau}{\Sigma; \Gamma \vdash \Lambda e : \forall \alpha. \tau} \text{CG-FI}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e [] : \tau[\ell/\alpha]} \text{CG-FE}$
$\frac{\Sigma; \Psi, c; \Gamma \vdash e : \tau}{\Sigma; \Gamma \vdash \nu e : c \Rightarrow \tau} \text{CG-CI}$	$\frac{\Sigma; \Psi; \Gamma \vdash e : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e \bullet : \tau} \text{CG-CE}$

Figure 12: Type system for CG

2.2.2 CG semantics

Judgement: $e \Downarrow_i v$ and $(H, e) \Downarrow_i^f (H', v)$

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \tau <: \tau} \text{CGsub-ref} \qquad \frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2} \text{CGsub-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{CGsub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{CGsub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'} \text{CGsub-labeled} \\
\\
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'} \text{CGsub-monad} \\
\\
\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2} \text{CGsub-forall} \qquad \frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \implies \tau_1 <: c_2 \implies \tau_2} \text{CGsub-constraint}
\end{array}$$

Figure 13: CG subtyping

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \mathbf{b} \text{ } WF} \text{CG-wff-base} \qquad \frac{}{\Sigma; \Psi \vdash \text{unit } WF} \text{CG-wff-unit} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 \xrightarrow{\ell_e} \tau_2) \text{ } WF} \text{CG-wff-arrow} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 \times \tau_2) \text{ } WF} \text{CG-wff-times} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 \text{ } WF \quad \Sigma; \Psi \vdash \tau_2 \text{ } WF}{\Sigma; \Psi \vdash (\tau_1 + \tau_2) \text{ } WF} \text{CG-wff-sum} \qquad \frac{\text{FV}(\ell) = \emptyset \quad \text{FV}(\tau) = \emptyset}{\Sigma; \Psi \vdash (\text{ref } \ell \tau) \text{ } WF} \text{CG-wff-ref} \\
\\
\frac{\Sigma, \alpha; \Psi \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (\forall \alpha. (\tau)) \text{ } WF} \text{CG-wff-forall} \qquad \frac{\Sigma; \Psi, c \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (c \xrightarrow{\ell_c} \tau) \text{ } WF} \text{CG-wff-constraint} \\
\\
\frac{\Sigma; \Psi \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (\text{Labeled } \ell \tau) \text{ } WF} \text{CG-wff-labeled} \qquad \frac{\Sigma; \Psi \vdash \tau \text{ } WF}{\Sigma; \Psi \vdash (\mathbb{C} \ell_i \ell_o \tau) \text{ } WF} \text{CG-wff-monad}
\end{array}$$

Figure 14: Well-formedness relation for CG

$$\begin{array}{c}
\frac{e_1 \Downarrow_i \lambda x. e_i \quad e_2 \Downarrow_j v_2 \quad e_i[v_2/x] \Downarrow_k v_3}{e_1 e_2 \Downarrow_{i+j+k+1} v_3} \text{cg-app} \qquad \frac{e_1 \Downarrow_i v_1 \quad e_2 \Downarrow_j v_2}{(e_1, e_2) \Downarrow_{i+j+1} (v_1, v_2)} \text{cg-prod} \\
\\
\frac{e \Downarrow_i (v_1, v_2)}{\text{fst}(e) \Downarrow_{i+1} v_1} \text{cg-fst} \qquad \frac{e \Downarrow_i (v_1, v_2)}{\text{snd}(e) \Downarrow_{i+1} v_2} \text{cg-snd} \qquad \frac{e \Downarrow_i v}{\text{inl}(e) \Downarrow_{i+1} \text{inl}(v)} \text{cg-inl} \\
\\
\frac{e \Downarrow_i v}{\text{inr}(e) \Downarrow_{i+1} \text{inr}(v)} \text{cg-inr} \qquad \frac{e \Downarrow_i \text{inl } v \quad e_1[v/x] \Downarrow_j v_1}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_1} \text{cg-case1} \\
\\
\frac{e \Downarrow_i \text{inr } v \quad e_2[v/x] \Downarrow_j v_2}{\text{case}(e, x.e_1, y.e_2) \Downarrow_{i+j+1} v_2} \text{cg-case2} \qquad \frac{e \Downarrow_i v}{\text{Lb}(e) \Downarrow_{i+1} \text{Lb}(v)} \text{cg-Lb} \\
\\
\frac{e \Downarrow_i \Lambda e_i \quad e_i \Downarrow_j v}{e[] \Downarrow_{i+j+1} v} \text{cg-FE} \qquad \frac{e \Downarrow_i \nu e_i \quad e_i \Downarrow_j v}{e\bullet \Downarrow_{i+j+1} v} \text{cg-CE} \qquad \frac{e \Downarrow_i v}{(H, \text{ret}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-ret} \\
\\
\frac{e_1 \Downarrow_i v_1 \quad (H, v_1) \Downarrow_j^f (H', v'_1) \quad e_2[v'_1/x] \Downarrow_k v_2 \quad (H', v_2) \Downarrow_l^f (H'', v'_2)}{(H, \text{bind}(e_1, x.e_2)) \Downarrow_{i+j+k+l+1}^f (H'', v'_2)} \text{cg-bind} \\
\\
\frac{e \Downarrow_i \text{Lb}(v)}{(H, \text{unlabel}(e)) \Downarrow_{i+1}^f (H, v)} \text{cg-unlabel} \qquad \frac{e \Downarrow_i v \quad (H, v) \Downarrow_j^f (H', v')}{(H, \text{toLabeled}(e)) \Downarrow_{i+j+1}^f (H', \text{Lb}(v'))} \text{cg-toLabeled} \\
\\
\frac{e \Downarrow_i \text{Lb}v \quad a \notin \text{dom}(H)}{(H, \text{new } (e)) \Downarrow_{i+1}^f (H[a \mapsto \text{Lb}v], a)} \text{cg-ref} \qquad \frac{e \Downarrow_i a}{(H, !e) \Downarrow_{i+1}^f (H, H(a))} \text{cg-deref} \\
\\
\frac{e_1 \Downarrow_i a \quad e_2 \Downarrow_j \text{Lb}v}{(H, e_1 := e_2) \Downarrow_{i+j+1}^f (H[a \mapsto \text{Lb}v], ())} \text{cg-assign} \\
\\
\frac{e \in \{x, \lambda y. -, \Lambda -, \nu -, \text{ret} -, \text{bind}(-, -, -), \text{unlabel}(-), \text{toLabeled}(-), \text{new }(-), !-, - := -\}}{e \Downarrow_0 e} \text{cg-val}
\end{array}$$

2.2.3 Logical relation for CG

$W : ((\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \leftrightarrow \text{Loc}))$

Definition 2.30 (CG: θ_2 extends θ_1). $\theta_1 \sqsubseteq \theta_2 \triangleq$

$$\forall a \in \theta_1. \theta_1(a) = \tau \implies \theta_2(a) = \tau$$

Definition 2.31 (CG: W_2 extends W_1). $W_1 \sqsubseteq W_2 \triangleq$

1. $\forall i \in \{1, 2\}. W_1.\theta_i \sqsubseteq W_2.\theta_i$
2. $\forall p \in (W_1.\hat{\beta}). p \in (W_2.\hat{\beta})$

Definition 2.32 (CG: Value Equivalence).

$$\text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau) \triangleq \begin{cases} (W, n, v_1, v_2) \in [\tau]_V^{\mathcal{A}} & \ell \sqsubseteq \mathcal{A} \\ \forall j. (W.\theta_1, j, v_1) \in [\tau]_V \wedge (W.\theta_2, j, v_2) \in [\tau]_V & \ell \not\sqsubseteq \mathcal{A} \end{cases}$$

Definition 2.33 (CG: Binary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^A &\triangleq \{(W, n, v_1, v_2) \mid v_1 = v_2 \wedge \{v_1, v_2\} \in \llbracket \mathbf{b} \rrbracket\} \\
[\mathbf{unit}]_V^A &\triangleq \{(W, n, (), ()) \mid () \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^A &\triangleq \{(W, n, (v_1, v_2), (v'_1, v'_2)) \mid (W, n, v_1, v'_1) \in [\tau_1]_V^A \wedge (W, n, v_2, v'_2) \in [\tau_2]_V^A\} \\
[\tau_1 + \tau_2]_V^A &\triangleq \{(W, n, \text{inl } v, \text{inl } v') \mid (W, n, v, v') \in [\tau_1]_V^A\} \cup \\
&\quad \{(W, n, \text{inr } v, \text{inr } v') \mid (W, n, v, v') \in [\tau_2]_V^A\} \\
[\tau_1 \rightarrow \tau_2]_V^A &\triangleq \{(W, n, \lambda x. e_1, \lambda x. e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, v_1, v_2. \\
&\quad ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, v_c, j. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, v_c, j. \\
&\quad ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)\} \\
[\forall \alpha. \tau]_V^A &\triangleq \{(W, n, \Lambda e_1, \Lambda e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n, \ell' \in \mathcal{L}. \\
&\quad ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha]]_E \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha]]_E\} \\
[c \Rightarrow \tau]_V^A &\triangleq \{(W, n, \nu e_1, \nu e_2) \mid \\
&\quad \forall W' \sqsupseteq W, j < n. \\
&\quad \mathcal{L} \models c \implies (W', j, e_1, e_2) \in [\tau]_E^A \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E \wedge \\
&\quad \forall \theta_l \sqsupseteq W. \theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E\} \\
[\text{ref } \ell \tau]_V^A &\triangleq \{(W, n, a_1, a_2) \mid \\
&\quad (a_1, a_2) \in W. \hat{\beta} \wedge W. \theta_1(a_1) = W. \theta_2(a_2) = \text{Labeled } \ell \tau\} \\
[\text{Labeled } \ell \tau]_V^A &\triangleq \{(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \mid \text{ValEq}(\mathcal{A}, W, \ell, n, v_1, v_2, \tau)\} \\
[\mathbb{C} \ell_1 \ell_2 \tau]_V^A &\triangleq \{(W, n, v_1, v_2) \mid \\
&\quad (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \\
&\quad \forall v'_1, v'_2, j. (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\
&\quad \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\
&\quad \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\
&\quad \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1))\}
\end{aligned}$$

Definition 2.34 (CG: Binary expression relation).

$$[\tau]_E^A \triangleq \{(W, n, e_1, e_2) \mid \forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow_i v_2 \implies (W, n - i, v_1, v_2) \in [\tau]_V^A\}$$

Definition 2.35 (CG: Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V &\triangleq \{(\theta, m, v) \mid v \in \llbracket \mathbf{b} \rrbracket\} \\
\llbracket \mathbf{unit} \rrbracket_V &\triangleq \{(\theta, m, v \mid v \in \llbracket \mathbf{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V &\triangleq \{(\theta, m, (v_1, v_2)) \mid (\theta, m, v_1) \in \llbracket \tau_1 \rrbracket_V \wedge (\theta, m, v_2) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \mathbf{inl} \ v) \mid (\theta, m, v) \in \llbracket \tau_1 \rrbracket_V\} \cup \{(\theta, m, \mathbf{inr} \ v) \mid (\theta, m, v) \in \llbracket \tau_2 \rrbracket_V\} \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket_V &\triangleq \{(\theta, m, \lambda x. e) \mid \forall \theta' \supseteq \theta, v, j < m. (\theta', j, v) \in \llbracket \tau_1 \rrbracket_V \implies (\theta', j, e[v/x]) \in \llbracket \tau_2 \rrbracket_E\} \\
\llbracket \forall \alpha. \tau \rrbracket_V &\triangleq \{(\theta, m, \Lambda e) \mid \forall \theta'. \theta \sqsubseteq \theta', j < m. \forall \ell' \in \mathcal{L}. (\theta', j, e) \in \llbracket \tau[\ell'/\alpha] \rrbracket_E\} \\
\llbracket c \Rightarrow \tau \rrbracket_V &\triangleq \{(\theta, m, \nu e) \mid \mathcal{L} \models c \implies \forall \theta'. \theta \sqsubseteq \theta', j < m. (\theta', j, e) \in \llbracket \tau \rrbracket_E\} \\
\llbracket \mathbf{ref} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, a) \mid \theta(a) = \mathbf{Labeled} \ \ell \ \tau\} \\
\llbracket \mathbf{Labeled} \ \ell \ \tau \rrbracket_V &\triangleq \{(\theta, m, \mathbf{Lb}(v)) \mid (\theta, m, v) \in \llbracket \tau \rrbracket_V\} \\
\llbracket \mathbb{C} \ \ell_1 \ \ell_2 \ \tau \rrbracket_V &\triangleq \{(\theta, m, e) \mid \\
&\quad \forall k \leq m, \theta_e \supseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k \implies \\
&\quad \exists \theta' \supseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge \\
&\quad (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \mathbf{Labeled} \ \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
&\quad (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)\}
\end{aligned}$$

Definition 2.36 (CG: Unary expression relation).

$$\llbracket \tau \rrbracket_E \triangleq \{(\theta, n, e) \mid \forall i < n. e \Downarrow_i v \implies (\theta, n - i, v) \in \llbracket \tau \rrbracket_V\}$$

Definition 2.37 (CG: Unary heap well formedness).

$$(n, H) \triangleright \theta \triangleq \text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in \llbracket \theta(a) \rrbracket_V$$

Definition 2.38 (CG: Binary heap well formedness).

$$\begin{aligned}
(n, H_1, H_2) \overset{A}{\triangleright} W &\triangleq \text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\
&\quad (W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\
&\quad \forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\
&\quad (W, n - 1, H_1(a_1), H_2(a_2)) \in \llbracket W.\theta_1(a_1) \rrbracket_V^A \wedge \\
&\quad \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \llbracket W.\theta_i(a_i) \rrbracket_V
\end{aligned}$$

Definition 2.39 (CG: Label substitution). $\sigma : \text{Lvar} \mapsto \text{Label}$

Definition 2.40 (CG: Value substitution to value pairs). $\gamma : \text{Var} \mapsto (\text{Val}, \text{Val})$

Definition 2.41 (CG: Value substitution to values). $\delta : \text{Var} \mapsto \text{Val}$

Definition 2.42 (CG: Unary interpretation of Γ).

$$\llbracket \Gamma \rrbracket_V \triangleq \{(\theta, n, \delta) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta) \wedge \forall x \in \text{dom}(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V\}$$

Definition 2.43 (CG: Binary interpretation of Γ).

$$\llbracket \Gamma \rrbracket_V^A \triangleq \{(W, n, \gamma) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A\}$$

2.2.4 Soundness proof for CG

Lemma 2.44 (CG: Binary value relation subsumes unary value relation). $\forall W, v_1, v_2, \mathcal{A}, n, \tau.$

$$(W, n, v_1, v_2) \in \llbracket \tau \rrbracket_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in \llbracket \tau \rrbracket_V$$

Proof. Proof by induction on τ

1. Case b:

From Definition 2.35

2. Case $\tau_1 \times \tau_2$:

Given: $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$ (P01)

and

$\forall m. (W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$ (P02)

From Definition 2.33 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$ (P1)

IH1a: $\forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V$ and

IH1b: $\forall m_1. (W.\theta_2, m_1, v_{j1}) \in [\tau_1]_V$

IH2a: $\forall m_2. (W.\theta_1, m_2, v_{i2}) \in [\tau_2]_V$ and

IH2b: $\forall m_2. (W.\theta_2, m_2, v_{j2}) \in [\tau_2]_V$

From (P01) we know that given some m we need to prove

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly from (P02) we know that given some m we need to prove

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

We instantiate IH1a and IH2a with the given m from (P01) to get

$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$ and $(W.\theta_1, m, v_{i2}) \in [\tau_2]_V$

Then from Definition 2.35, we get

$(W.\theta_1, m, (v_{i1}, v_{i2})) \in [\tau_1 \times \tau_2]_V$

Similarly we instantiate IH1b and IH2b with the given m from (P02) to get

$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$ and $(W.\theta_2, m, v_{j2}) \in [\tau_2]_V$

Then from Definition 2.35, we get

$(W.\theta_2, m, (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V$

3. Case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v_1 = \text{inl}(v_{i1})$ and $v_2 = \text{inl}(v_{j1})$

Given: $(W, n, \text{inl}(v_{i1}), \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V^A$

To prove:

$\forall m. (W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$ (S01)

and

$\forall m. (W.\theta_2, m, \text{inl}(v_{i2})) \in [\tau_1 + \tau_2]_V$ (S02)

From Definition 2.33 we know that we are given

$$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \quad (\text{S0})$$

$$\text{IH1: } \forall m_1. (W.\theta_1, m_1, v_{i1}) \in [\tau_1]_V \text{ and}$$

$$\text{IH2: } \forall m_2. (W.\theta_2, m_2, v_{j1}) \in [\tau_1]_V$$

From (S01) we know that given some m and we are required to prove:

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

Also from (S02) we know that given some m and we are required to prove:

$$(W.\theta_2, m, \text{inl}(v_{i2})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH1 with m from (S01) to get

$$(W.\theta_1, m, v_{i1}) \in [\tau_1]_V$$

Therefore from Definition 2.35, we get

$$(W.\theta_1, m, \text{inl}(v_{i1})) \in [\tau_1 + \tau_2]_V$$

We instantiate IH2 with m from (S02) to get

$$(W.\theta_2, m, v_{j1}) \in [\tau_1]_V$$

Therefore from Definition 2.35, we get

$$(W.\theta_2, m, \text{inl}(v_{j1})) \in [\tau_1 + \tau_2]_V$$

$$(b) \ v_1 = \text{inr}(v_{i2}) \text{ and } v_2 = \text{inr}(v_{j2})$$

Symmetric reasoning as in the (a) case above

4. Case $\tau_1 \rightarrow \tau_2$:

$$\text{Given: } (W, n, \lambda x.e_1, \lambda x.e_2) \in [\tau_1 \rightarrow \tau_2]_V^A$$

This means from Definition 2.33 we know that

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, v_c. ((\theta_l, i, v_c) \in [\tau_1]_V \implies (\theta_l, i, e_1[v_c/x]) \in [\tau_2]_E) \\ \wedge \forall \theta_l \sqsupseteq W.\theta_2, k, v_c. ((\theta_l, k, v_c) \in [\tau_1]_V \implies (\theta_l, k, e_2[v_c/x]) \in [\tau_2]_E) \quad (\text{L0}) \end{aligned}$$

To prove:

$$(a) \ \forall m. (W.\theta_1, m, \lambda x.e_1) \in [\tau_1 \rightarrow \tau_2]_V:$$

This means from Definition 2.35 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta' \wedge \forall j < m. \forall v. (\theta', j, v) \in [\tau_1]_V \implies (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

This further means that we have some θ', j and v s.t

$$W.\theta_1 \sqsubseteq \theta' \wedge j < m \wedge (\theta', j, v) \in [\tau_1]_V$$

$$\text{And we need to prove: } (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

Instantiating θ_l, i and v_c in the second conjunct of L0 with θ', j and v respectively and since we know that $W.\theta_1 \sqsubseteq \theta'$ and $(\theta', j, v) \in [\tau_1]_V$

$$\text{Therefore we get } (\theta', j, e_1[v/x]) \in [\tau_2]_E$$

$$(b) \ \forall m. (W.\theta_2, m, \lambda x.e_2) \in [\tau_1 \rightarrow \tau_2]_V:$$

Similar reasoning with e_2

5. Case $\forall\alpha.\tau$:

Given: $(W, n, \Lambda e_1, \Lambda e_2) \in [\forall\alpha.\tau]_V^A$

This means from Definition 2.33 we know that

$$\begin{aligned} & \forall W_b \sqsupseteq W, n_b < n, \ell' \in \mathcal{L}. ((W_b, n_b, e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \\ & \wedge \forall \theta_l \sqsupseteq W.\theta_1, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_1) \in [\tau[\ell''/\alpha]]_E) \\ & \wedge \forall \theta_l \sqsupseteq W.\theta_2, i, \ell'' \in \mathcal{L}. ((\theta_l, i, e_2) \in [\tau[\ell''/\alpha]]_E) \end{aligned} \quad (\text{F0})$$

To prove:

(a) $\forall m. (W.\theta_1, m, \Lambda e_1) \in [\forall\alpha.\tau]_V$:

This means from Definition 2.35 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \forall \ell_u \in \mathcal{L}. (\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$$

This further means that we are given some θ', m' and ℓ_u s.t $W.\theta_1 \sqsubseteq \theta', m' < m$ and $\ell_u \in \mathcal{L}$

And we need to prove: $(\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$

Instantiating θ_l, i and ℓ'' in the second conjunct of F0 with θ', m' and ℓ_u respectively and since we know that $W.\theta_1 \sqsubseteq \theta'$ and $\ell_u \in \mathcal{L}$

Therefore we get $(\theta', m', e_1) \in [\tau[\ell_u/\alpha]]_E$

(b) $\forall m. (W.\theta_2, m, \Lambda e_2) \in [\forall\alpha.\tau]_V$:

Symmetric reasoning for e_2

6. Case $c \Rightarrow \tau$:

Given: $(W, n, \nu e_1, \nu e_2) \in [c \Rightarrow \tau]_V^A$

This means from Definition 2.33 we know that

$$\begin{aligned} & \forall W_b \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W_b, n', e_1, e_2) \in [\tau]_E^A \\ & \wedge \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau]_E \\ & \wedge \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau]_E \end{aligned} \quad (\text{C0})$$

To prove:

(a) $\forall m. (W.\theta_1, m, \nu e_1) \in [c \Rightarrow \tau]_V$:

This means from Definition 2.35 we need to prove:

$$\forall \theta'. W.\theta_1 \sqsubseteq \theta'. \forall m' < m. \mathcal{L} \models c \implies (\theta', m', e_1) \in [\tau]_E$$

This further means that we are given some θ' and m' s.t $W.\theta_1 \sqsubseteq \theta', m' < m$ and $\mathcal{L} \models c$

And we need to prove: $(\theta', m', e_1) \in [\tau]_E$

Instantiating θ_l, j in the second conjunct of C0 with θ', m' respectively and since we know that $W.\theta_1 \sqsubseteq \theta'$ and $\mathcal{L} \models c$

Therefore we get $(\theta', m', e_1) \in [\tau]_E$

(b) $\forall m. (W.\theta_2, m, \nu e_2) \in [c \Rightarrow \tau]_V$:

Symmetric reasoning for e_2

7. Case ref $\ell \tau$:

From Definition 2.33 and 2.35

8. Case Labeled $\ell \tau$:

Given $(W, n, \text{Lb}v_1, \text{Lb}v_2) \in [\text{Labeled } \ell \tau]_V^A$

2 cases arise:

(a) $\ell \sqsubseteq \mathcal{A}$:

From Definition 2.32 we know that

$$(W, n, v_1, v_2) \in [\tau]_V^A$$

Therefore from IH we get $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$ and $\forall m. (W.\theta_2, m, v_2) \in [\tau]_V$

(b) $\ell \not\sqsubseteq \mathcal{A}$:

Directly from Definition 2.32

9. Case $\mathbb{C} \ell_1 \ell_2 \tau$:

Given: $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V^A$

This means from Definition 2.33 we know that

$$\begin{aligned} & (\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \quad (\text{CG0}) \end{aligned}$$

To prove: $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\mathbb{C} \ell_1 \ell_2 \tau]_V$

This means from Definition 2.35 we need to prove

$$\begin{aligned} & \forall l \in \{1, 2\}. \forall m. (\forall k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1)) \end{aligned}$$

Case $l = 1$

And given some m and $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove that

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \end{aligned}$$

Instantiating (CG0) with $l = 1$ and the given $k \leq m, \theta_e \sqsupseteq W.\theta_l, H, j$ we get the desired.

Case $l = 2$

Symmetric reasoning as in the previous case above

□

Lemma 2.45 (CG: Monotonicity Unary). *The following holds:*

$\forall \theta, \theta', v, m, m', \tau.$

$(\theta, m, v) \in \lfloor \tau \rfloor_V \wedge m' < m \wedge \theta \sqsubseteq \theta' \implies (\theta', m', v) \in \lfloor \tau \rfloor_V$

Proof. Proof by induction on τ

1. case b:

Directly from Definition 2.35

2. case $\tau_1 \times \tau_2$:

Given: $(\theta, m, (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

To prove: $(\theta', m', (v_1, v_2)) \in \lfloor \tau_1 \times \tau_2 \rfloor_V$

This means from Definition 2.35 we know that

$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V \wedge (\theta, m, v_2) \in \lfloor \tau_2 \rfloor_V$

IH1 : $(\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$

IH2 : $(\theta', m', v_2) \in \lfloor \tau_2 \rfloor_V$

We get the desired from IH1, IH2 and Definition 2.35

3. case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v = \text{inl}(v_1)$:

Given: $(\theta, m, (\text{inl } v_1)) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

To prove: $(\theta', m', \text{inl } v_1) \in \lfloor \tau_1 + \tau_2 \rfloor_V$

This means from Definition 2.35 we know that

$(\theta, m, v_1) \in \lfloor \tau_1 \rfloor_V$

IH : $(\theta', m', v_1) \in \lfloor \tau_1 \rfloor_V$

Therefore from IH and Definition 2.35 we get the desired

(b) $v = \text{inr}(v_2)$

Symmetric case

4. case $\tau_1 \rightarrow \tau_2$:

Given: $(\theta, m, (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

To prove: $(\theta', m', (\lambda x. e_1)) \in \lfloor \tau_1 \rightarrow \tau_2 \rfloor_V$

This means from Definition 2.35 we know that

$$\forall \theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall v. (\theta'', j, v) \in \lfloor \tau_1 \rfloor_V \implies (\theta'', j, e_1[v/x]) \in \lfloor \tau_2 \rfloor_E \quad (143)$$

Similarly from Definition 2.35 we know that we are required to prove

$$\forall \theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall v_1. (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V \implies (\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$$

This means that given some θ''', k and v_1 such that $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge (\theta''', k, v_1) \in \lfloor \tau_1 \rfloor_V$

And we are required to prove $(\theta''', k, e_1[v_1/x]) \in \lfloor \tau_2 \rfloor_E$

Instantiating Equation 143 with θ''', k and v_1 and since we know that $\theta' \sqsubseteq \theta'''$ and $\theta \sqsubseteq \theta'$ therefore we have $\theta \sqsubseteq \theta'''$. Also, we know that $k < m' < m$ and $(\theta''', k, v_1) \in [\tau_1]_V$

Therefore we get $(\theta''', k, e_1[v_1/x]) \in [\tau_2]_E$

5. case ref $\ell \tau$:

From Definition 2.35 and Definition 2.30

6. case $\forall\alpha.\tau$:

Given: $(\theta, m, (\Lambda e_1)) \in [\forall\alpha.\tau]_V$

To prove: $(\theta', m', (\Lambda e_1)) \in [\forall\alpha.\tau]_V$

This means from Definition 2.35 we know that

$$\forall\theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \forall \ell_i \in \mathcal{L}. (\theta'', j, e_1) \in [\tau[\ell_i/\alpha]]_E \quad (144)$$

Similarly from Definition 2.35 we know that we are required to prove

$$\forall\theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \forall \ell_j \in \mathcal{L}. (\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$$

This means that given some θ''', k and ℓ_j such that $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$

Instantiating Equation 144 with θ''', k and ℓ_j and since we know that $\theta' \sqsubseteq \theta'''$ and $\theta \sqsubseteq \theta'$ therefore we have $\theta \sqsubseteq \theta'''$. Also, we know that $k < m' < m$ and $\ell_j \in \mathcal{L}$

Therefore we get $(\theta''', k, e_1) \in [\tau[\ell_j/\alpha]]_E$

7. case $c \Rightarrow \tau$:

Given: $(\theta, m, (\nu e_1)) \in [c \Rightarrow \tau]_V$

To prove: $(\theta', m', (\nu e_1)) \in [c \Rightarrow \tau]_V$

This means from Definition 2.35 we know that

$$\forall\theta''. \theta \sqsubseteq \theta'' \wedge \forall j < m. \mathcal{L} \models c \implies (\theta'', j, e_1) \in [\tau]_E \quad (145)$$

Similarly from Definition 2.35 we know that we are required to prove

$$\forall\theta'''. \theta' \sqsubseteq \theta''' \wedge \forall k < m'. \mathcal{L} \models c \implies (\theta''', k, e_1) \in [\tau]_E$$

This means that given some θ''', k and ℓ_j such that $\theta' \sqsubseteq \theta''' \wedge k < m' \wedge \ell_j \in \mathcal{L}$

And we are required to prove $(\theta''', k, e_1) \in [\tau]_E$

Instantiating Equation 145 with θ''', k and since we know that $\theta' \sqsubseteq \theta'''$ and $\theta \sqsubseteq \theta'$ therefore we have $\theta \sqsubseteq \theta'''$. Also, we know that $k < m' < m$ and $\mathcal{L} \models c$

Therefore we get $(\theta''', k, e_1) \in [\tau]_E$

8. case Labeled $\ell \tau$:

Given: $(\theta, m, (\text{Lb } v)) \in \llbracket \text{Labeled } \ell \tau \rrbracket_V$

To prove: $(\theta', m', (\text{Lb } v)) \in \llbracket \text{Labeled } \ell \tau \rrbracket_V$

This means from Definition 2.35 we know that $(\theta, m, v) \in \llbracket \tau \rrbracket_V$

IH: $(\theta', m', v) \in \llbracket \tau \rrbracket_V$

Therefore from IH and Definition 2.35 we get the desired

9. case $\mathbb{C} \ell_1 \ell_2 \tau$:

Given: $(\theta, m, e) \in \llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket_V$

To prove: $(\theta', m', e) \in \llbracket \mathbb{C} \ell_1 \ell_2 \tau \rrbracket_V$

This means from Definition 2.35 we know that

$$\begin{aligned} \forall k \leq m, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) & \quad (\text{LB0}) \end{aligned}$$

Similarly from Definition 2.35 we are required to prove

$$\begin{aligned} \forall k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1. (k_1, H_1) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1 &\implies \\ \exists \theta' \sqsupseteq \theta_e. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

This means we are given

$$k_1 \leq m', \theta_{e1} \sqsupseteq \theta', H_1, j_1 \text{ s.t. } (k_1, H) \triangleright \theta_{e1} \wedge (H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1) \wedge j_1 < k_1$$

And we are required to prove:

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

Instantiating (LB0), k with k_1 , θ_e with θ_{e1} , H with H_1 and j with j_1 . We know that $k_1 < m' < m$, $\theta \sqsubseteq \theta' \sqsubseteq \theta_{e1}$, $(k_1, H_1) \triangleright \theta_{e1}$, $(H_1, v_1) \Downarrow_{j_1}^f (H'_1, v'_1)$ and $i_1 + j_1 < k_1$. Therefore we get

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k_1 - j_1, H') \triangleright \theta' \wedge (\theta'_1, k_1 - j_1, v') \in \llbracket \tau \rrbracket_V \wedge & \\ (\forall a. H_1(a) \neq H'_1(a) \implies \exists \ell'. \theta_{e1}(a) = \text{Labeled } \ell' \tau \wedge \ell_1 \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta'_1) \setminus \text{dom}(\theta_{e1}). \theta'_1(a) \searrow \ell_1) & \end{aligned}$$

□

Lemma 2.46 (CG: Monotonicity binary). *The following holds:*

$$\forall W, W', v_1, v_2, \mathcal{A}, n, n', \tau.$$

$$(W, n, v_1, v_2) \in \llbracket \tau \rrbracket_V^{\mathcal{A}} \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', v_1, v_2) \in \llbracket \tau \rrbracket_V^{\mathcal{A}}$$

Proof. Proof by induction on τ

1. Case **b**, unit:

From Definition 2.33

2. Case $\tau_1 \times \tau_2$:

Given: $(W, n, (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

To prove: $(W', n', (v_{i1}, v_{i2}), (v_{j1}, v_{j2})) \in [\tau_1 \times \tau_2]_V^A$

From Definition 2.33 we know that we are given

$(W, n, v_{i1}, v_{j1}) \in [\tau_1]_V^A \wedge (W, n, v_{i2}, v_{j2}) \in [\tau_2]_V^A$

IH1 : $(W', n', v_{i1}, v_{j1}) \in [\tau_1]_V^A$

IH2 : $(W', n', v_{i2}, v_{j2}) \in [\tau_2]_V^A$

From IH1, IH2 and Definition 2.33 we get the desired.

3. Case $\tau_1 + \tau_2$:

2 cases arise:

(a) $v_1 = \text{inl } v_{i1}$ and $v_2 = \text{inl } v_{i2}$:

Given: $(W, n, (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

To prove: $(W', n', (\text{inl } v_{i1}, \text{inl } v_{i2})) \in [\tau_1 + \tau_2]_V^A$

From Definition 2.33 we know that we are given

$(W, n, v_{i1}, v_{i2}) \in [\tau_1]_V^A$

IH : $(W', n', v_{i1}, v_{i2}) \in [\tau_1]_V^A$

Therefore from Definition 2.33 we get

$(W', n', \text{inl } v_{i1}, \text{inl } v_{i2}) \in [\tau_1 + \tau_2]_V^A$

(b) $v_1 = \text{inr}(v_{i1})$ and $v_2 = \text{inr}(v_{i2})$:

Symmetric case

4. Case $\tau_1 \rightarrow \tau_2$:

Given: $(W, n, (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_V^A$

To prove: $(\theta', n', (\lambda x.e_1), (\lambda x.e_2)) \in [\tau_1 \rightarrow \tau_2]_V^A$

This means from Definition 2.33 we know that the following holds

$\forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1]_V^A \implies (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2]_E^A)$
(BM-A0)

$\forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2]_E)$ (BM-A1)

$\forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2]_E)$ (BM-A2)

Similarly from Definition 2.33 we know that we are required to prove

(a) $\forall W'' \sqsupseteq W', k < n', v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau_1]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A)$:

This means that we are given some $W'' \sqsupseteq W', k < n'$ and v'_1, v'_2 s.t

$(W'', k, v'_1, v'_2) \in [\tau_1]_V^A$

And we are required to prove: $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

Instantiating BM-A0 with W'', k and v'_1, v'_2 we get

$(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2]_E^A$

(b) $\forall \theta'_l \sqsupseteq W'.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E)$:

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_1, k$ and v'_c s.t

$$(\theta'_l, k, v'_c) \in [\tau_1]_V$$

And we are required to prove: $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with θ'_l, k and v'_c we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2]_E$$

(c) $\forall \theta'_l \sqsupseteq W'.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau_1]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E)$:

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_2, k$ and v'_c s.t

$$(\theta'_l, k, v'_c) \in [\tau_1]_V$$

And we are required to prove: $(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E$

Instantiating BM-A1 with θ'_l, k and v'_c we get

$$(\theta'_l, k, e_2[v'_c/x]) \in [\tau_2]_E$$

5. Case ref $\ell \tau$:

From Definition 2.33 and Definition 2.31

6. Case $\forall \alpha. \tau$:

Given: $(W, n, (\Lambda e_1), (\Lambda e_2)) \in [\forall \alpha. \tau]_V^A$

To prove: $(\theta', n', (\Lambda e_1), (\Lambda e_1)) \in [\forall \alpha. \tau]_V^A$

This means from Definition 2.33 we know that the following holds

$$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau[\ell'/\alpha]]_E^A) \quad (\text{BM-F0})$$

$$\forall \theta_l \sqsupseteq W.\theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau[\ell'/\alpha]]_E) \quad (\text{BM-F1})$$

$$\forall \theta_l \sqsupseteq W.\theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau[\ell'/\alpha]]_E) \quad (\text{BM-F2})$$

Similarly from Definition 2.33 we know that we are required to prove

(a) $\forall W'' \sqsupseteq W', n'' < n', \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$:

This means that we are given some $W'' \sqsupseteq W', n'' < n'$ and $\ell'' \in \mathcal{L}$

And we are required to prove: $((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$

Instantiating BM-F0 with W'', n'' and ℓ'' . And since $W'' \sqsupseteq W'$ and $W' \sqsupseteq W$ therefore $W'' \sqsupseteq W$. Also since $n'' < n'$ and $n' < n$ therefore $n'' < n$. And finally since $\ell'' \in \mathcal{L}$ therefore we get

$$((W'', n'', e_1, e_2) \in [\tau[\ell''/\alpha]]_E^A)$$

(b) $\forall \theta'_l \sqsupseteq W'.\theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$:

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_1, k$ and $\ell'' \in \mathcal{L}$

And we are required to prove: $((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$

Instantiating BM-F1 with θ'_l, k and ℓ'' . And since $\theta'_l \sqsupseteq W'.\theta_1$ and $W' \sqsupseteq W$ therefore $\theta'_l \sqsupseteq W.\theta_1$. And since $\ell'' \in \mathcal{L}$ therefore we get

$$((\theta'_l, k, e_1) \in [\tau[\ell''/\alpha]]_E)$$

(c) $\forall \theta_l \sqsupseteq W.\theta_2, j, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$:

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_2, k$ and $\ell'' \in \mathcal{L}$

And we are required to prove: $((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$

Instantiating BM-F1 with θ'_l, k and ℓ'' . And since $\theta'_l \sqsupseteq W'.\theta_2$ and $W' \sqsupseteq W$ therefore $\theta'_2 \sqsupseteq W.\theta_2$. And since $\ell'' \in \mathcal{L}$ therefore we get

$((\theta'_l, k, e_2) \in \lfloor \tau[\ell''/\alpha] \rfloor_E)$

7. Case $c \Rightarrow \tau$:

Given: $(W, n, (\nu e_1), (\nu e_2)) \in \lceil c \Rightarrow \tau \rceil_V^A$

To prove: $(\theta', n', (\nu e_1), (\nu e_2)) \in \lceil c \Rightarrow \tau \rceil_V^A$

This means from Definition 2.33 we know that the following holds

$\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c \implies (W', n', e_1, e_2) \in \lceil \tau \rceil_E^A$ (BM-C0)

$\forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in \lfloor \tau \rfloor_E$ (BM-C1)

$\forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in \lfloor \tau \rfloor_E$ (BM-C2)

Similarly from Definition 2.33 we know that we are required to prove

(a) $\forall W'' \sqsupseteq W', n'' < n. \mathcal{L} \models c \implies (W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$:

This means that we are given some $W'' \sqsupseteq W', n'' < n'$ and $\mathcal{L} \models c$

And we are required to prove: $(W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$

Instantiating BM-C0 with W'', n'' . And since $W'' \sqsupseteq W'$ and $W' \sqsupseteq W$ therefore $W'' \sqsupseteq W$. And since $\mathcal{L} \models c$ therefore we get

$(W'', n'', e_1, e_2) \in \lceil \tau \rceil_E^A$

(b) $\forall \theta'_l \sqsupseteq W'.\theta_1, k. \mathcal{L} \models c \implies (\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$:

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_1, k$ and $\mathcal{L} \models c$

And we are required to prove: $(\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$

Instantiating BM-F1 with θ'_l, k . And since $\theta'_l \sqsupseteq W'.\theta_1$ and $W' \sqsupseteq W$ therefore $\theta'_1 \sqsupseteq W.\theta_1$. And since $\mathcal{L} \models c$ therefore we get

$(\theta'_l, k, e_1) \in \lfloor \tau \rfloor_E$

(c) $\forall \theta'_l \sqsupseteq W'.\theta_2, k. \mathcal{L} \models c \implies (\theta_l, k, e_2) \in \lfloor \tau \rfloor_E$:

This means that we are given some $\theta'_l \sqsupseteq W'.\theta_2, k$ and $\mathcal{L} \models c$

And we are required to prove: $(\theta'_l, k, e_2) \in \lfloor \tau \rfloor_E$

Instantiating BM-F1 with θ'_l, k . And since $\theta'_l \sqsupseteq W'.\theta_2$ and $W' \sqsupseteq W$ therefore $\theta'_2 \sqsupseteq W.\theta_2$. And since $\mathcal{L} \models c$ therefore we get

$(\theta'_l, k, e_2) \in \lfloor \tau \rfloor_E$

8. Case Labeled $\ell \tau$:

Given: $(W, n, (\text{Lb } v_1), (\text{Lb } v_2)) \in \lceil \text{Labeled } \ell \tau \rceil_V^A$

To prove: $(W', n', (\text{Lb } v_1), (\text{Lb } v_2)) \in \lceil \text{Labeled } \ell \tau \rceil_V^A$

From Definition 2.33 2 cases arise:

(a) $\ell \sqsubseteq \mathcal{A}$:

In this case we know that $(W, n, v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Therefore from IH we know that $(W', n', v_1, v_2) \in [\tau]_{\mathcal{V}}^{\mathcal{A}}$

Hence from Definition 2.33 we get $(W', n', (\text{Lb}v_1), (\text{Lb}v_2)) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\mathcal{A}}$

(b) $\ell \not\sqsubseteq \mathcal{A}$:

In this case we know that $\forall m. (W.\theta_1, m, v_1) \in [\tau]_V$ and $(W.\theta_2, m, v_2) \in [\tau]_V$

Since $W.\theta_1 \sqsubseteq W'.\theta_1$ (from Definition 2.31). Therefore from Lemma 2.45 we know that $\forall m' < m. (W'.\theta_1, m', v_1) \in [\tau]_V$

Similarly since $W.\theta_2 \sqsubseteq W'.\theta_2$ (from Definition 2.31). Therefore from Lemma 2.45 we know that

$\forall m' < m. (W'.\theta_2, m', v_2) \in [\tau]_V$

Finally from Definition 2.33 we get $(W', n', (\text{Lb}v_1), (\text{Lb}v_2)) \in [\text{Labeled } \ell \tau]_{\mathcal{V}}^{\mathcal{A}}$

9. Case $\mathbb{C} \ell_1 \ell_2 \tau$:

Given: $(W, n, v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_{\mathcal{V}}^{\mathcal{A}}$

To prove: $(W', n', v_1, v_2) \in [\mathbb{C} \ell_1 \ell_2 \tau]_{\mathcal{V}}^{\mathcal{A}}$

From Definition 2.33 we are given that

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \Big) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \quad (\text{BM-M0}) \end{aligned}$$

Similarly from Definition 2.33 it suffices to prove that

$$\begin{aligned} & \text{(a) } \left(\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \right. \\ & \forall v'_1, v'_2, j.(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau) \right): \\ & \text{This means that given some } k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j \text{ s.t} \\ & (k, H_1, H_2) \triangleright W_e \wedge (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \end{aligned}$$

It suffices to prove that

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, \tau)$$

Instantiating the first conjunct of (BM-M0) with the given $k, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$ and since we know that $n' \leq n$ and $W \sqsubseteq W'$ we get the desired

$$\begin{aligned} & \text{(b) } \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right): \end{aligned}$$

Similar reasoning as in the previous case but using Lemma 2.45

□

Lemma 2.47 (CG: Unary monotonicity for Γ). $\forall \theta, \theta', \delta, \Gamma, n, n'$.
 $(\theta, n, \delta) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge \theta \sqsubseteq \theta' \implies (\theta', n', \delta) \in \llbracket \Gamma \rrbracket_V$

Proof. Given: $(\theta, n, \delta) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge \theta \sqsubseteq \theta'$
 To prove: $(\theta', n', \delta) \in \llbracket \Gamma \rrbracket_V$

From Definition 2.42 it is given that
 $dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

And again from Definition 2.42 we are required to prove that
 $dom(\Gamma) \subseteq dom(\delta) \wedge \forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

- $dom(\Gamma) \subseteq dom(\delta)$:
 Given
- $\forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$:
 Since we know that $\forall x \in dom(\Gamma). (\theta, n, \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$ (given)
 Therefore from Lemma 2.45 we get
 $\forall x \in dom(\Gamma). (\theta', n', \delta(x)) \in \llbracket \Gamma(x) \rrbracket_V$

□

Lemma 2.48 (CG: Binary monotonicity for Γ). $\forall W, W', \delta, \Gamma, n, n'$.
 $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge W \sqsubseteq W' \implies (W', n', \gamma) \in \llbracket \Gamma \rrbracket_V$

Proof. Given: $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V \wedge n' < n \wedge W \sqsubseteq W'$
 To prove: $(W', n', \gamma) \in \llbracket \Gamma \rrbracket_V$

From Definition 2.43 it is given that
 $dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

And again from Definition 2.42 we are required to prove that
 $dom(\Gamma) \subseteq dom(\gamma) \wedge \forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

- $dom(\Gamma) \subseteq dom(\gamma)$:
 Given
- $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$:
 Since we know that $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$ (given)
 Therefore from Lemma 2.46 we get
 $\forall x \in dom(\Gamma). (W', n', \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in \llbracket \Gamma(x) \rrbracket_V^A$

□

Lemma 2.49 (CG: Unary monotonicity for H). $\forall \theta, H, n, n'$.
 $(n, H) \triangleright \theta \wedge n' < n \implies (n', H) \triangleright \theta$

Proof. Given: $(n, H) \triangleright \theta \wedge n' < n$

To prove: $(n', H) \triangleright \theta$

From Definition 2.37 it is given that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$$

And again from Definition 2.42 we are required to prove that

$$\text{dom}(\theta) \subseteq \text{dom}(H) \wedge \forall a \in \text{dom}(\theta). (\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

- $\text{dom}(\theta) \subseteq \text{dom}(H)$:

Given

- $\forall a \in \text{dom}(\theta). (\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$:

Since we know that $\forall a \in \text{dom}(\theta). (\theta, n - 1, H(a)) \in \lfloor \theta(a) \rfloor_V$ (given)

Therefore from Lemma 2.45 we get

$$\forall a \in \text{dom}(\theta). (\theta, n' - 1, H(a)) \in \lfloor \theta'(a) \rfloor_V$$

□

Lemma 2.50 (CG: Binary monotonicity for heaps). $\forall W, H_1, H_2, n, n'$.

$$(n, H_1, H_2) \triangleright W \wedge n' < n \implies (n', H_1, H_2) \triangleright W$$

Proof. Given: $(n, H_1, H_2) \triangleright W \wedge n' < n \wedge W \sqsubseteq W'$

To prove: $(n', H_1, H_2) \triangleright W$

From Definition 2.38 it is given that

$$\begin{aligned} \text{dom}(W.\theta_1) &\subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2) \wedge \\ (W.\hat{\beta}) &\subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2)) \wedge \\ \forall (a_1, a_2) &\in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2)) \wedge \\ (W, n - 1, H_1(a_1), H_2(a_2)) &\in \lceil W.\theta_1(a_1) \rceil_V^A \wedge \\ \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). &(W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V \end{aligned}$$

And again from Definition 2.38 we are required to prove:

- $\text{dom}(W.\theta_1) \subseteq \text{dom}(H_1) \wedge \text{dom}(W.\theta_2) \subseteq \text{dom}(H_2)$:

Given

- $(W.\hat{\beta}) \subseteq (\text{dom}(W.\theta_1) \times \text{dom}(W.\theta_2))$:

Given

- $\forall (a_1, a_2) \in (W.\hat{\beta}). (W.\theta_1(a_1) = W.\theta_2(a_2))$ and $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$:

$$\forall (a_1, a_2) \in (W.\hat{\beta}).$$

– $(W.\theta_1(a_1) = W.\theta_2(a_2))$: Given

– $(W, n' - 1, H_1(a_1), H_2(a_2)) \in \lceil W.\theta_1(a_1) \rceil_V^A$:

Given and from Lemma 2.46

- $\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W.\theta_i). (W.\theta_i, m, H_i(a_i)) \in \lfloor W.\theta_i(a_i) \rfloor_V$:

Given

□

Theorem 2.51 (CG: Fundamental theorem unary). $\forall \Sigma, \Psi, \Gamma, \theta, \mathcal{L}, e, \tau, \sigma, \delta, n.$

$$\begin{aligned} & \Sigma; \Psi; \Gamma \vdash e : \tau \wedge \\ & \mathcal{L} \models \Psi \sigma \wedge \\ & (\theta, n, \delta) \in [\Gamma \sigma]_V \implies \\ & (\theta, n, e \delta) \in [\tau \sigma]_E \end{aligned}$$

Proof. Proof by induction on CG typing derivation

1. CG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, x \delta) \in [\tau \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. x \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau]_V$$

This means that given some $i < n$ s.t $x \delta \Downarrow_i v$

(from *cg - val* we know that $v = x \delta$ and $i = 0$)

It suffices to prove $(\theta, n, x \delta) \in [\tau \sigma]_V$ (FU-V0)

Since $(\theta, n, \delta) \in [\Gamma' \sigma]_V$ where $\Gamma' = \Gamma \cup \{x : \tau\}$. Therefore from Definition 2.42 we know that $(\theta, n, \delta(x)) \in [\Gamma'(x) \sigma]_V$

So we are done.

2. CG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e' : \tau_2}{\Sigma; \Psi; \Gamma \vdash \lambda x. e' : (\tau_1 \rightarrow \tau_2)}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, \lambda x. e_i \delta) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \lambda x. e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

This means that given some $i < n$ s.t $\lambda x. e' \delta \Downarrow_i v$

(from *cg - val* we know that $v = \lambda x. e' \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \lambda x. e' \delta) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V \quad (\text{FU-L0})$$

From Definition 2.35 it further suffices to prove

$$\forall \theta'' \sqsupseteq \theta, v', j < n. (\theta'', j, v') \in [\tau_1 \sigma]_V \implies (\theta'', j, (e' \delta)[v'/x]) \in [\tau_2 \sigma]_E$$

This means given some θ'', v', j s.t $\theta'' \sqsupseteq \theta, j < n$ and $(\theta'', j, v') \in [\tau_1 \sigma]_V$ (FU-L1)

We are required to prove

$$(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2 \sigma]_E$$

Since $(\theta, n, \delta) \in [\Gamma \sigma]_V$ therefore from Lemma 2.47 we know that $(\theta, j, \delta) \in [\Gamma \sigma]_V$ where $j < n$ (from FU-L1)

IH:

$$\forall \theta_h, v_x. (\theta_h, j, e' \delta \cup \{x \mapsto v_x\}) \in [\tau_2 \sigma]_E, \text{ s.t. } (\theta_i, j, v_x) \in [\tau_1 \sigma]_V$$

Instantiating IH with θ'' and v' from (FU-L1) we get $(\theta'', j, (e' \delta)[v'/x]) \in [\tau_2 \sigma]_E$

3. CG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_1}{\Sigma; \Psi; \Gamma \vdash e_1 e_2 : \tau_2}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, (e_1 e_2) \delta) \in [\tau_2 \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. (e_1 e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau_2 \sigma]_V$$

This means that given some $i < n$ s.t. $(e_1 e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau_2 \sigma]_V \quad (\text{FU-P0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

Since we know that $(e_1 e_2) \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t. $e_1 \delta \Downarrow_j v_1$. This means we have

$$(\theta, n - j, v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V$$

From *cg - app* we know that $v_1 = \lambda x. e'$. Therefore we have

$$(\theta, n - j, \lambda x. e') \in [(\tau_1 \rightarrow \tau_2) \sigma]_V \quad (\text{FU-P1})$$

This means from Definition 2.35 we have

$$\forall \theta'' \sqsupseteq \theta \wedge I < (n - j), v. (\theta'', I, v) \in [\tau_1 \sigma]_V \implies (\theta'', I, e'[v/x]) \in [\tau_2 \sigma]_E \quad (146)$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in [\tau_1 \sigma]_V$$

Since we know that $(e_1 e_2) \delta \Downarrow_i v$ therefore $\exists k < i - j$ (since $i < n$ therefore $i - j < n - j$) s.t. $e_2 \delta \Downarrow_k v_2$. This means we have

$$(\theta, n - j - k, v_2) \in [\tau_1 \sigma]_V \quad (\text{FU-P2})$$

Instantiating Equation 146 with $\theta, (n - j - k), v_2$ and since we know that $(\theta, n - j - k, v_2) \in [\tau_1 \sigma]_V$ therefore we get

$$(\theta, n - j - k, e'[v_2/x]) \in [\tau_2 \sigma]_E$$

This means from Definition 2.36 we have

$$\forall J < n - j - k. e'[v_2/x] \Downarrow_J v_f \implies (\theta, n - j - k - J, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$$

Since we know that $(e_1 \ e_2) \delta \Downarrow_i v$ therefore we know that $\exists J < i < n$ s.t $i = j + k + J$ (since $j + k + J < n$ therefore $J < n - j - k$) and $e'[v_2/x] \Downarrow_J v_f$

Therefore we have $(\theta, n - j - k - J, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$

Since we know that $i = j + k + J$ and $v = v_J$ therefore we get $(\theta, n - i, v_J) \in \lfloor \tau_2 \sigma \rfloor_E$ (so FU-P0 is proved)

4. CG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove: $(\theta, n, (e_1, e_2) \delta) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. (e_1, e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V$$

This means that given some $i < n$ s.t $(e_1, e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V \quad (\text{FU-PA0})$$

IH1:

$$\forall j < n. e_1 \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$$

Since we know that $(e_1, e_2) \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t $e_1 \delta \Downarrow_j v_1$. This means we have

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V \quad (\text{FU-PA1})$$

IH2:

$$\forall k < (n - j). e_2 \delta \Downarrow_k v_2 \implies (\theta, n - j - k, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$$

Since we know that $(e_1 \ e_2) \delta \Downarrow_i v$ therefore $\exists k < i - j$ (since $i < n$ therefore $i - j < n - j$) s.t $e_2 \delta \Downarrow_k v_2$. This means we have

$$(\theta, n - j - k, v_2) \in \lfloor \tau_2 \sigma \rfloor_V \quad (\text{FU-PA2})$$

In order to prove (FU-PA0) from *cg - prod* we know that $i = j + k + 1$ and $v = (v_1, v_2)$ therefore from Definition 2.35 it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in \lfloor \tau_1 \sigma \rfloor_V \text{ and } (\theta, n - j - k - 1, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$$

We get this from (FU-PA1) and Lemma 2.45 and from (FU-PA2) and Lemma 2.45

5. CG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e') : \tau_1}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove: $(\theta, n, \text{fst}(e') \delta) \in \lfloor \tau_1 \sigma \rfloor_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{fst}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau_1 \sigma \rfloor_V$

This means that given some $i < n$ s.t $\text{fst}(e') \delta \Downarrow_i v$

It suffices to prove

$(\theta, n - i, v) \in \lfloor \tau_1 \sigma \rfloor_V$ (FU-F0)

IH1:

$\forall j < n. e' \delta \Downarrow_j (v_1, v_2) \implies (\theta, n - j, (v_1, v_2)) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V$

Since we know that $\text{fst}(e') \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t $e' \delta \Downarrow_j (v_1, v_2)$. This means we have

$(\theta, n - j, (v_1, v_2)) \in \lfloor (\tau_1 \times \tau_2) \sigma \rfloor_V$

From Definition 2.35 we know the following holds

$(\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$ and $(\theta, n - j, v_2) \in \lfloor \tau_2 \sigma \rfloor_V$ (FU-F1)

From *cg-fst* we know that $v = v_1$ and $i = j + 1$. Therefore from (FU-F0), we are required to prove

$(\theta, n - j - 1, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$

We get this from (FU-F1) and Lemma 2.45

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau_1}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove: $(\theta, n, \text{inl}(e') \delta) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{inl}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$

This means that given some $i < n$ s.t $\text{inl}(e') \delta \Downarrow_i v$

It suffices to prove

$(\theta, n - i, v) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$ (FU-LE0)

IH1:

$\forall j < n. e' \delta \Downarrow_j v_1 \implies (\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$

Since we know that $\text{inl}(e') \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t $e' \delta \Downarrow_j v_1$. This means we have

$$(\theta, n - j, v_1) \in \lfloor \tau_1 \sigma \rfloor_V \quad (\text{FU-LE1})$$

From *cg - inl* we know that $v = v_1$ and $i = j + 1$. Therefore from (FU-LE0) we are required to prove

$$(\theta, n - j - 1, v_1) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$$

From Definition 2.35 it suffices to prove

$$(\theta, n - j - 1, v_1) \in \lfloor \tau_1 \sigma \rfloor_V$$

We get this from (FU-LE1) and Lemma 2.45

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Sigma; \Psi; \Gamma \vdash \text{case}(e_c, x.e_1, y.e_2) : \tau}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in \lfloor \Gamma \sigma \rfloor_V$

To prove: $(\theta, n, (\text{case } e_c, x.e_1, y.e_2) \delta) \in \lfloor \tau \sigma \rfloor_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. (\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V$$

This means that given some $i < n$ s.t. $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V \quad (\text{FU-C0})$$

IH1:

$$\forall j < n. e_c \delta \Downarrow_j v_c \implies (\theta, n - j, v_1) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V$$

Since we know that $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t. $e_c \delta \Downarrow_j v_c$. This means we have

$$(\theta, n - j, v_c) \in \lfloor (\tau_1 + \tau_2) \sigma \rfloor_V \quad (\text{FU-C1})$$

2 cases arise:

(a) $v_c = \text{inl}(v_l)$:

IH2:

$$\forall k < (n - j). e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1 \implies (\theta, n - j - k, v_1) \in \lfloor \tau \sigma \rfloor_V$$

Since we know that $(\text{case } e_c, x.e_1, y.e_2) \delta \Downarrow_i v$ therefore $\exists k < i - j$ (since $i < n$ therefore $i - j < n - j$) s.t. $e_1 \delta \cup \{x \mapsto v_l\} \Downarrow_k v_1$. This means we have

$$(\theta, n - j - k, v_1) \in \lfloor \tau \sigma \rfloor_V \quad (\text{FU-C2})$$

From *cg - case1* we know that $i = j + k + 1$ and $v = v_1$. Therefore from (FU-C0) it suffices to prove

$$(\theta, n - j - k - 1, v_1) \in \lfloor \tau \sigma \rfloor_V$$

We get this from (FU-C2) and Lemma 2.45

(b) $v_c = \text{inr}(v_r)$:

Symmetric reasoning as in the previous case

10. CG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \Lambda e' : \forall \alpha. \tau}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, \Lambda e' \delta) \in [(\forall \alpha. (\ell_e, \tau)) \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \Lambda e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\forall \alpha. \tau) \sigma]_V$$

This means that given some $i < n$ s.t $\lambda x. e' \delta \Downarrow_i v$

(from *cg - val* we know that $v = \Lambda e' \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \Lambda e' \delta) \in [(\forall \alpha. \tau) \sigma]_V \quad (\text{FU-FI0})$$

From Definition 2.35 it further suffices to prove

$$\forall \theta'. \theta \sqsubseteq \theta', j < n. \forall \ell' \in \mathcal{L}. (\theta', j, e' \delta) \in [\tau[\ell'/\alpha]]_E$$

This means given some $\theta', j, \ell' \in \mathcal{L}$ s.t $\theta' \sqsupseteq \theta, j < n$ (FU-FI1)

We are required to prove

$$(\theta', j, (e' \delta)) \in [\tau[\ell'/\alpha] \sigma]_E \quad (\text{FU-FI2})$$

Since $(\theta, n, \delta) \in [\Gamma \sigma]_V$ therefore from Lemma 2.47 we know that $(\theta, j, \delta) \in [\Gamma \sigma]_V$ where $j < n$ (from FU-L1)

$$\underline{\text{IH}}: (\theta', j, e' \delta) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_E$$

(FU-FI2) is obtained directly from IH

11. CG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \nu e' : c \Rightarrow \tau}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, \nu e' \delta) \in [(c \Rightarrow \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \nu e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(c \Rightarrow \tau) \sigma]_V$$

This means that given some $i < n$ s.t $\nu e' \delta \Downarrow_i v$

(from *cg - val* we know that $v = \nu e' \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \nu e' \delta) \in [(c \Rightarrow \tau) \sigma]_V \quad (\text{FU-CI0})$$

From Definition 2.35 it further suffices to prove

$$\mathcal{L} \models c \implies \forall \theta'. \theta \sqsubseteq \theta', j < n. (\theta', j, e' \delta) \in [\tau]_E$$

This means given $\mathcal{L} \models c$ and some θ', j s.t $\theta' \sqsupseteq \theta, j < n$ (FU-CI1)

We are required to prove

$$(\theta', j, (e' \delta)) \in [\tau \sigma]_E \quad (\text{FU-CI2})$$

Since $(\theta, n, \delta) \in [\Gamma \sigma]_V$ therefore from Lemma 2.47 we know that $(\theta, j, \delta) \in [\Gamma \sigma]_V$ where $j < n$ (from FU-L1). Also we know that $\mathcal{L} \models c \sigma$ therefore $\mathcal{L} \models (\Sigma \cup \{c\}) \sigma$

$$\underline{\text{IH}}: (\theta', j, e' \delta) \in [\tau \sigma]_E$$

(FU-CI2) is obtained directly from IH

12. CG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e' [] : \tau[\ell/\alpha]}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, e' [] \delta) \in [\tau[\ell/\alpha] \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. e' [] \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau[\ell/\alpha] \sigma]_V$$

This means that given some $i < n$ s.t $e' [] \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau[\ell/\alpha] \sigma]_V \quad (\text{FU-FE0})$$

$$\underline{\text{IH}}: (\theta, n, e' \delta) \in [\forall \alpha. \tau]_E$$

From Definition 2.36 we know that

$$\forall h_1 < n. e' \delta \Downarrow_{h_1} \Lambda e_{h_1} \implies (\theta, n - h_1, \Lambda e_{h_1}) \in [(\forall \alpha. \tau) \sigma]_V$$

Since $e' [] \delta$ reduces therefore we know that $\exists h_1 < i < n$ such that $e' \delta \Downarrow_{h_1} \Lambda e_i$

Therefore we know that $(\theta, n - h_1, \Lambda e_{h_1}) \in [(\forall \alpha. \tau) \sigma]_V$

From Definition 2.35 we know that

$$\forall \theta'' \sqsupseteq \theta, x < (n - h_1), \ell_h \in \mathcal{L}. (\theta'', x, e_{h_1}) \in [(\tau[\ell_h/\alpha]) \sigma]_E$$

Instantiating θ'' with θ, x with $n - h_1 - 1$ and ℓ_h with ℓ . So, we get

$$(\theta, n - h_1 - 1, e_{h_1}) \in [(\tau[\ell/\alpha]) \sigma]_E$$

From Definition 2.36 we know that the following holds

$$\forall h_2 < n - h_1 - 1. e_{h_1} \delta \Downarrow_{h_2} v \implies (\theta, n - h_1 - 1 - h_2, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

Since $e' [] \delta$ reduces in i steps therefore from cg-FE we know that $(i = h_1 + h_2 + 1)$ and since we know that $i < n$ therefore we have $h_2 < n - h_1 - 1$ such that $e_{h_1} \delta \Downarrow_{h_2} v$. Therefore we get

$$(\theta, n - h_1 - 1 - h_2, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

Since $i = h_1 + h_2 + 1$ therefore we get

$$(\theta, n - i, v) \in [(\tau[\ell/\alpha]) \sigma]_V$$

13. CG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e' \bullet : \tau}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, e' \bullet \delta) \in [\tau \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. e' \bullet \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\tau \sigma]_V$$

This means that given some $i < n$ s.t. $e' \bullet \delta \Downarrow_i v$

It suffices to prove

$$(\theta, n - i, v) \in [\tau \sigma]_V \quad (\text{FU-CE0})$$

$$\underline{\text{IH}}: (\theta, n, e' \delta) \in [c \Rightarrow \tau \sigma]_E$$

From Definition 2.36 we know that

$$\forall h_1 < n. e' \delta \Downarrow_{h_1} \nu e_{h_1} \implies (\theta, n - h_1, \nu e_{h_1}) \in [c \Rightarrow \tau \sigma]_V$$

Since $e' \bullet \delta$ reduces therefore we know that $\exists h_1 < i < n$ such that $e' \delta \Downarrow_{h_1} \nu e_{h_1}$

Therefore we know that $(\theta, n - h_1, \nu e_{h_1}) \in [c \Rightarrow \tau \sigma]_V$

From Definition 2.35 we know that

$$\mathcal{L} \models c \sigma \implies \forall \theta'' \sqsupseteq \theta, x < (n - h_1). (\theta'', x, e_{h_1}) \in [\tau \sigma]_E$$

Since we know that $\mathcal{L} \models c \sigma$ and then we instantiate θ'' with θ , x with $n - h_1 - 1$. So, we get

$$(\theta, n - h_1 - 1, e_{h_1}) \in [\tau \sigma]_E$$

From Definition 2.36 we know that the following holds

$$\forall h_2 < n - h_1 - 1. e_{h_1} \delta \Downarrow_{h_2} v \implies (\theta, n - h_1 - 1 - h_2, v) \in [\tau \sigma]_V$$

Since $e' \bullet \delta$ reduces in i steps therefore from cg-CE we know that $(i = h_1 + h_2 + 1)$ and since we know that $i < n$ therefore we have $h_2 < n - h_1 - 1$ such that $e_{h_1} \delta \Downarrow_{h_2} v$. Therefore we get

$$(\theta, n - h_1 - 1 - h_2, v) \in [\tau \sigma]_V$$

Since we know that $i = h_1 + h_2 + 1$ therefore we get

$$(\theta, n - i, v) \in [\tau \sigma]_V$$

14. CG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } (e') : \mathbb{C} \ell \ell (\text{ref } \ell' \tau)}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, \text{new } (e') \delta) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \text{new } (e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \lfloor \mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma \rfloor_V$$

This means that given some $i < n$ s.t $\text{new } (e') \delta \Downarrow_i v$

(from $cg - val$ we know that $v = \text{new } (e') \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \text{new } (e') \delta) \in \lfloor \mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma \rfloor_V$$

From Definition 2.35 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor (\text{ref } \ell' \tau) \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, \text{new } (e') \delta) \Downarrow_j^f (H', v') \wedge j < k$.
Also from $cg - ref$ we know that $v' = a$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, a) \in \lfloor (\text{ref } \ell' \tau) \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_E$$

From Definition 2.36 this means we have

$$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, n - l, v_h) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_V$$

Since we know that $(H, \text{new } (e') \delta) \Downarrow_j^f (H', a)$ therefore from $cg - ref$ we know that

$$\exists l < j < k \text{ s.t } e' \delta \Downarrow_l v_h$$

Therefore we have

$$(\theta_e, n - l, v_h) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_V \quad (\text{FU-R2})$$

In order to prove (FU-R0) we choose θ' as $\theta_n = \theta_e \cup \{a \mapsto \text{Labeled } \ell' \tau\}$

Now we need to prove:

(a) $(k - j, H') \triangleright \theta_n$:

From Definition 2.37 it suffices to prove that

$$\text{dom}(\theta_n) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in \lfloor \theta_n(a) \rfloor_V$$

• $\text{dom}(\theta_n) \subseteq \text{dom}(H')$:

We know that $\text{dom}(H') = \text{dom}(H) \cup \{a\}$

We know that $\text{dom}(\theta_n) = \text{dom}(\theta_e) \cup \{a\}$

And $(k, H) \triangleright \theta_e$ therefore from Definition 2.37 we know that $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

• $\forall a \in \text{dom}(\theta_n). (\theta_n, (k - j) - 1, H'(a)) \in \lfloor \theta_n(a) \rfloor_V$:

Since from (FU-R2) we know that $(\theta_h, n - l, v_h) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_V$

Since $\theta_h \sqsubseteq \theta_n$ and $k - j - 1 < n - l$ (since $k < n$ and $l < j$) therefore from

Lemma 2.45 we know that $(\theta_n, k - j - 1, v_h) \in \lfloor (\text{Labeled } \ell' \tau) \sigma \rfloor_V$

- (b) $(\theta_n, k - j - 1, a) \in [(\text{ref } \ell' \tau)]_V$:
 From Definition 2.35 it suffices to prove that $\theta_n(a) = \text{Labeled } \ell' \tau$
 We get this by construction of θ_n
- (c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$:
 Holds vacuously
- (d) $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow \ell)$:
 From *CG-ref* we know that $\ell \sqsubseteq \ell'$

15. CG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{ref } \ell \tau}{\Sigma; \Psi; \Gamma \vdash !e' : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, (!e') \delta) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. !e' \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V$

(From *cg-val* we know that $v = !e' \delta$ and $i = 0$)

This means that given some $i < n$ s.t. $!e' \delta \Downarrow_i !e' \delta$

It suffices to prove

$(\theta, n, !e' \delta) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V$

From Definition 2.35 it suffices to prove

$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k \implies$
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sqsubseteq \ell'') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell')$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t. $(k, H) \triangleright \theta_e \wedge (H, (!e' \delta)) \Downarrow_j^f (H', v') \wedge j < k$.

It suffices to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{Labeled } \ell \tau)]_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell''. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sqsubseteq \ell'') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell') \quad (\text{FU-D0})$

IH:

$(\theta_e, k, e' \delta) \in [(\text{ref } \ell \tau) \sigma]_E$

From Definition 2.36 this means we have

$\forall l < k. e' \delta \Downarrow_l v_h \implies (\theta_e, k - l, v_h) \in [(\text{ref } \ell \tau) \sigma]_V$

Since we know that $(H, !e') \Downarrow_j^f (H', a)$ therefore from *cg-deref* we know that

$\exists l < j < k$ s.t. $e' \delta \Downarrow_l v_h, v_h = a$

Therefore we have

$(\theta_e, k - l, a) \in [(\text{ref } \ell \tau) \sigma]_V \quad (\text{FU-D1})$

In order to prove (FU-D0) we choose θ' as θ_e

Now we need to prove:

(a) $(k - j, H') \triangleright \theta_e$:

From Definition 2.37 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e).(\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

- $\text{dom}(\theta_e) \subseteq \text{dom}(H')$:

And $(k, H) \triangleright \theta_e$ therefore from Definition 2.37 we know that $\text{dom}(\theta_e) \subseteq \text{dom}(H)$
And since $H' = H$ (from *cg - deref*) so we are done

- $\forall a \in \text{dom}(\theta_e).(\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$:

Since we know that $(k, H) \triangleright \theta_e$ therefore from Definition 2.37 we know that

$$\forall a \in \text{dom}(\theta_e).(\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since $H' = H$ and from Lemma 2.45 we get

$$\forall a \in \text{dom}(\theta_e).(\theta_e, (k - j) - 1, H'(a)) \in [\theta_e(a)]_V$$

(b) $(\theta_e, k - j, v') \in [(\text{Labeled } \ell \tau)]_V$:

From *cg - deref* we know that $H = H'$ and $v' = H(a)$

From (FU-D1) and Definition 2.35 we know that $\theta_e(a) = \text{Labeled } \ell \tau$

Since we know that $(k, H) \triangleright \theta_e$ therefore from Definition 2.37 we know that

$$\forall a \in \text{dom}(\theta_e).(\theta_e, k - 1, H(a)) \in [\theta_e(a)]_V$$

Since from *cg - deref* we know that $j \geq 1$. Therefore from Lemma 2.45 we get

$$(\theta_e, k - j, H(a)) \in [(\text{Labeled } \ell \tau)]_V$$

(c) $(\forall a. H(a) \neq H'(a)) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell'$:

Holds vacuously

(d) $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \ell)$:

Holds vacuously

16. CG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \ell \text{ unit}}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, (e_1 := e_2) \delta) \in [(\mathbb{C} \ell \ell \text{ unit}) \sigma]_E^{pc}$

This means that from Definition 2.7 we need to prove

$$\forall i < n. (e_1 := e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell \ell \text{ unit}) \sigma]_V$$

This means that given some $i < n$ s.t $(e_1 := e_2) \delta \Downarrow_i v$.

It suffices to prove

$$(\theta, n - i, ()) \in [(\mathbb{C} \ell \ell \text{ unit}) \sigma]_V$$

From Definition 2.35 it suffices to prove

$$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies$$

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\text{ref } \ell' \tau)]_V \wedge$$

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, (e_1 := e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$. Also from *cg - assign* we know that $v' = ()$

It suffices to prove

$$\begin{aligned}
& \exists \theta' \sqsupseteq \theta_e.(k-j, H') \triangleright \theta' \wedge (\theta', k-j, ()) \in [\text{unit}]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-A0})
\end{aligned}$$

IH1:

$$\forall l < k. e_1 \delta \Downarrow_l v_1 \implies (\theta, k-l, a) \in [(\text{ref } \ell' \tau) \sigma]_V$$

Since we know that $(e_1 := e_2) \delta \Downarrow_j^f v$ therefore $\exists l < j < k$ s.t $e_1 \delta \Downarrow_l a$. This means we have

$$(\theta, k-l, a) \in [(\text{ref } \ell' \tau) \sigma]_V \quad (\text{FU-A1})$$

IH2:

$$\forall m < (k-l). e_2 \delta \Downarrow_m v_2 \implies (\theta, k-l-m, v_2) \in [\text{Labeled } \ell' \tau \sigma]_V$$

Since we know that $(e_1 := e_2) \delta \Downarrow_j^f v$ therefore $\exists m < j-l$ (since $j < k$ therefore $j-l < k-l$) s.t $e_2 \delta \Downarrow_m v_2$. This means we have

$$(\theta, k-l-m, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V \quad (\text{FU-A2})$$

In order to prove (FU-A0) we choose θ' as θ_e

Now we need to prove:

(a) $(k-j, H') \triangleright \theta_e$:

From Definition 2.37 it suffices to prove that

$$\text{dom}(\theta_e) \subseteq \text{dom}(H') \wedge \forall a \in \text{dom}(\theta_e). (\theta_e, (k-j) - 1, H'(a)) \in [\theta_e(a)]_V$$

• $\text{dom}(\theta_e) \subseteq \text{dom}(H')$:

We know that $\text{dom}(H') = \text{dom}(H)$

And $(k, H) \triangleright \theta_e$ therefore from Definition 2.37 we know that $\text{dom}(\theta_e) \subseteq \text{dom}(H)$

So we are done

• $\forall a \in \text{dom}(\theta_e). (\theta_e, (k-j) - 1, H'(a)) \in [\theta_e(a)]_V$:

$\forall a \in \text{dom}(\theta_e)$.

i. $H(a) = H'(a)$:

Since $(k, H) \triangleright \theta_e$ therefore from Definition 2.37 we know that

$$(\theta_e, k-1, H(a)) \in [\theta_e(a)]_V$$

Therefore from Lemma 2.45 we get

$$(\theta_e, k-1-j, H(a)) \in [\theta_e(a)]_V$$

ii. $H(a) \neq H'(a)$:

From *cg-assign* we know that $H'(a) = v_2$

From (FU-A1) we know that $\theta_e(a) = \text{Labeled } \ell' \tau$

Also we know that $j = l + m + 1$

Since from (FU-A2) we know that

$$(\theta, k-l-m, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V$$

Therefore we get

$$(\theta, k-j+1, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V$$

Therefore from Lemma 2.45 we get

$$(\theta, k-j-1, v_2) \in [(\text{Labeled } \ell' \tau) \sigma]_V$$

(b) $(\theta_e, k-j-1, ()) \in [\text{unit}]_V$:

From Definition 2.35

(c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell \sqsubseteq \ell')$:

From *CG – assign* we know that $\ell \sqsubseteq \ell'$

(d) $(\forall a \in \text{dom}(\theta_e) \setminus \text{dom}(\theta_e). \theta_e(a) \searrow \ell)$:

Holds vacuously

17. CG-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, \text{Lb}(e') \delta) \in [\text{Labeled } \ell \tau \sigma]_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{Lb}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\text{Labeled } \ell \tau \sigma]_V$

This means we are given some $i < n$ s.t $\text{Lb}(e') \delta \Downarrow_i v$ and we are required to prove

$(\theta, n - i, v) \in [\text{Labeled } \ell \tau \sigma]_V$

Let $v = \text{Lb}(v_i)$. This means from Definition 2.35 we are required to prove

$(\theta, n - i, v_i) \in [\tau \sigma]_V$

IH: $(\theta, n, e' \delta) \in [\tau \sigma]_E$

This means from Definition 2.36 we have

$\forall j < n. e' \delta \Downarrow_j v_i \implies (\theta, n - j, v_i) \in [\tau \sigma]_V$

Since we know that $\text{Lb}(e') \delta \Downarrow_i v$ therefore $\exists j < i < n$ s.t $e' \delta \Downarrow_j v_i$

Therefore we have $(\theta, n - j, v_i) \in [\tau \sigma]_V$

From *cg – label* we know that $i = j + 1$ therefore from Lemma 2.45 we have

$(\theta, n - i, v_i) \in [\tau \sigma]_V$

18. CG-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e') : \mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, \text{unlabel}(e') \delta) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{unlabel}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V$

This means that given some $i < n$ s.t $\text{unlabel}(e') \delta \Downarrow_i v$

(from *cg – val* we know that $v = \text{unlabel}(e') \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \text{unlabel}(e') \delta) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V$$

From Definition 2.35 it suffices to prove

$$\begin{aligned} & \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t. $(k, H) \triangleright \theta_e \wedge (H, \text{unlabel}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$. Also from *cg - unlabel* we know that $H' = H$

It suffices to prove

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-U0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [(\text{Labeled } \ell \tau) \sigma]_E$$

This means that from Definition 2.36 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in [(\text{Labeled } \ell \tau) \sigma]_V$$

Since we know that $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$ therefore from *cg - unlabel* we know that

$$\exists h_1 < j < k \text{ s.t. } e' \delta \Downarrow_{h_1} \text{Lb } v'$$

This means we have

$$(\theta_e, k - h_1, \text{Lb } v') \in [(\text{Labeled } \ell \tau) \sigma]_V$$

This means from Definition 2.35 we have

$$(\theta_e, k - h_1, v') \in [\tau \sigma]_V \quad (\text{FU-U1})$$

In order to prove (FU-U0) we choose θ' as θ_e . And we are required to prove:

(a) $(k - j, H) \triangleright \theta_e$:

Since have $(k, H) \triangleright \theta_e$ therefore from Lemma 2.49 we get $(k - j, H) \triangleright \theta_e$

(b) $(\theta', k - j, v') \in [\tau \sigma]_V$:

Since from (FU-U1) we know that $(\theta_e, k - h_1, v') \in [\tau \sigma]_V$

And since $j = h_1 + 1$, therefore from Lemma 2.45 we get $(\theta_e, k - j, v') \in [\tau \sigma]_V$

(c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$:

Holds vacuously

(d) $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$:

Holds vacuously

19. CG-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e') : \mathbb{C} \ell_i \ell_i \tau}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \text{ret}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V$$

This means we are given some $i < n$ s.t $\text{ret}(e') \delta \Downarrow_i v$ and we are required to prove

$$(\theta, n - i, v) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V$$

(from *cg - val* we know that $v = \text{ret}(e') \delta$ and $i = 0$)

It suffices to prove

$$(\theta, n, \text{ret}(e') \delta) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V$$

From Definition 2.35 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, \text{ret}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$.
Also from *cg - ret* we know that $H' = H$

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H) \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-R0}) \end{aligned}$$

IH:

$$(\theta_e, k, e' \delta) \in [\tau \sigma]_E$$

This means that from Definition 2.36 we need to prove

$$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_h \implies (\theta_e, k - h_1, v_h) \in [\tau \sigma]_V$$

Since we know that $(H, \text{unlabel}(e')) \Downarrow_j^f (H, v')$ therefore from *cg - ret* we know that

$$\exists h_1 < j < k \text{ s.t } e' \delta \Downarrow_{h_1} v'$$

This means we have

$$(\theta_e, k - h_1, v') \in [\tau \sigma]_V \quad (\text{FU-R1})$$

In order to prove (FU-U0) we choose θ' as θ_e . And we are required to prove:

(a) $(k - j, H) \triangleright \theta_e$:

Since have $(k, H) \triangleright \theta_e$ therefore from Lemma 2.49 we get $(k - j, H) \triangleright \theta_e$

(b) $(\theta', k - j, v') \in [\tau \sigma]_V$:

Since from (FU-R1) we know that $(\theta_e, k - h_1, v') \in [\tau \sigma]_V$

And since $j = h_1 + 1$, therefore from Lemma 2.45 we get $(\theta_e, k - j, v') \in [\tau \sigma]_V$

(c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$:

Holds vacuously

(d) $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$:

Holds vacuously

20. CG-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \mathbb{C} \ell_i \ell \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell \ell_o \tau'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell_i \ell_o \tau'}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in [\Gamma \sigma]_V$

To prove: $(\theta, n, \text{bind}(e_1, x.e_2) \delta) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_E$

This means that from Definition 2.36 we need to prove

$$\forall i < n. \text{bind}(e_1, x.e_2) \delta \Downarrow_i v \implies (\theta, n - i, v) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_V$$

This means we are given some $i < n$ s.t $\text{bind}(e_1, x.e_2) \delta \Downarrow_i v$ and we are required to prove

$$(\theta, n - i, v) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_V$$

(from *cg - val* we know that $v = \text{bind}(e_1, x.e_2) \delta$ and $i = 0$)

Therefore we need to prove

$$(\theta, n, v) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_V$$

From Definition 2.35 it suffices to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \end{aligned}$$

This means we are given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, \text{bind}(e_1, x.e_2) \delta) \Downarrow_j^f (H', v') \wedge j < k$.

It suffices to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [\tau \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-B0}) \end{aligned}$$

IH1:

$$(\theta_e, k, e_1 \delta) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_E$$

This means that from Definition 2.36 we need to prove

$$\forall h_1 < k. e_1 \delta \Downarrow_{h_1} v_1 \implies (\theta_e, k - h_1, v_1) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_V$$

Since we know that $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$ therefore from *cg - bind* we know that

$$\exists h_1 < j < k \text{ s.t } e_1 \delta \Downarrow_{h_1} v_1$$

This means we have

$$(\theta_e, k - h_1, v_1) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_V$$

From Definition 2.35 we know that

$$\begin{aligned}
& \forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H, J.(k_{h_1}, H) \triangleright \theta'_e \wedge (H, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies \\
& \exists \theta'' \sqsupseteq \theta'_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell)
\end{aligned}$$

Instantiating k_{h_1} with $k - h_1$, θ'_e with θ_e . Since we know that $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$ therefore $\exists J < j - h_1 < k - h_1$ s.t $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$. And since we already know that $(k, H) \triangleright \theta_e$ therefore from Lemma 2.49 we get $(k - h_1, H) \triangleright \theta_e$

This means we have

$$\begin{aligned}
& \exists \theta'' \sqsupseteq \theta_e.(k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v') \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta_e). \theta''(a) \searrow \ell) \quad (\text{FU-B1})
\end{aligned}$$

IH2:

$$(\theta'', k - h_1 - J, e_2 \delta \cup \{x \mapsto v'\}) \in [(\mathbb{C} \ell_i \ell \tau') \sigma]_E$$

This means that from Definition 2.36 we need to prove

$$\forall h_2 < k - h_1 - J. e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v'' \implies (\theta'', k - h_1 - J - h_2, v'') \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_V$$

Since we know that $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H, v_1)$ therefore from $cg - bind$ we know that $\exists h_2 < j - h_1 - J < k - h_1 - J$ s.t $e_2 \delta \cup \{x \mapsto v'\} \Downarrow_{h_2} v''$

This means we have

$$(\theta'', k - h_1 - J - h_2, v'') \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_V$$

From Definition 2.35 we know that

$$\begin{aligned}
& \forall k_{h_2} \leq (k - h_1 - J - h_2), \theta'_e \sqsupseteq \theta'', H, J'.(k_{h_2}, H) \triangleright \theta'_e \wedge (H, v'') \Downarrow_{J'}^f (H'', v'_{h_2}) \wedge J' < k_{h_2} \implies \\
& \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v'') \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell)
\end{aligned}$$

Since we know that $(H, \text{bind}(e_1, x.e_2)) \Downarrow_j^f (H_1, v_1)$ therefore $\exists v_{h_2}, i$ s.t $(v'' \Downarrow_i v_{h_2})$. From $cg - val$ we know that $v_{h_2} = v''$ and $i = 0$. Instantiating k_{h_2} with $k - h_1 - J - h_2$, θ'_e with θ'' , H with H' (from FU-B1) and $\exists J' < j - h_1 - J - h_2 < k - h_1 - J - h_2$ s.t $(H', v_{h_2}) \Downarrow_{J'}^f (H'', v'_{h_2})$. And since we already know that $(k - h_1, H') \triangleright \theta''$ therefore from Lemma 2.49 we get $(k - h_1 - J - h_2, H') \triangleright \theta''$

This means we have

$$\begin{aligned}
& \exists \theta''' \sqsupseteq \theta'_e.(k_{h_2} - J', H'') \triangleright \theta''' \wedge (\theta''', k_{h_2} - J', v'') \in [\tau \sigma]_V \wedge \\
& (\forall a. H(a) \neq H''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}(\theta''') \setminus \text{dom}(\theta'_e). \theta'''(a) \searrow \ell) \quad (\text{FU-B2})
\end{aligned}$$

We get (FU-B0) by choosing θ' as θ'' (from FU-B2)

21. CG-toLabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \mathbb{C} \ell_i \ell_o \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)}$$

Also given is $\mathcal{L} \models \Psi \sigma \wedge$ and $(\theta, n, \delta) \in \llbracket \Gamma \sigma \rrbracket_V$

To prove: $(\theta, n, \text{toLabeled}(e') \delta) \in \llbracket (\mathbb{C} \ell_i \ell_i \text{Labeled} \ell_o \tau) \sigma \rrbracket_E$

This means that from Definition 2.36 we need to prove

$\forall i < n. \text{toLabeled}(e') \delta \Downarrow_i v \implies (\theta, n - i, v) \in \llbracket (\mathbb{C} \ell_i \ell_i \text{Labeled} \ell_o \tau) \sigma \rrbracket_V$

This means that given some $i < n$ s.t $\text{toLabeled}(e') \delta \Downarrow_i v$

(from *cg - val* we know that $v = \text{toLabeled}(e') \delta$ and $i = 0$)

It suffices to prove

$(\theta, n, \text{toLabeled}(e') \delta) \in \llbracket (\mathbb{C} \ell_i \ell_i \text{Labeled} \ell_o \tau) \sigma \rrbracket_V$

From Definition 2.35 it suffices to prove

$\forall k \leq n, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k \implies$
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket (\text{Labeled} \ell_o \tau) \sigma \rrbracket_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled} \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell)$

And given some $k \leq n, \theta_e \sqsupseteq \theta, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, \text{toLabeled}(e') \delta) \Downarrow_j^f (H', v') \wedge j < k$.
 Also from *cg - tolabeled* we know that $H' = H$

It suffices to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \llbracket (\text{Labeled} \ell_o \tau) \sigma \rrbracket_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled} \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \quad (\text{FU-TL0})$

IIH:

$(\theta_e, k, e' \delta) \in \llbracket (\mathbb{C} \ell_i \ell_o \tau) \sigma \rrbracket_E$

This means that from Definition 2.36 we need to prove

$\forall h_1 < k. e' \delta \Downarrow_{h_1} v_1 \implies (\theta, k - h_1, v_1) \in \llbracket (\mathbb{C} \ell_i \ell_o \tau) \sigma \rrbracket_V$

Since $H, \text{toLabeled}(e') \Downarrow_j^f H', v'$ therefore from *cg - tolabeled* we know that $\exists h_1 < j < k$
 s.t $e' \delta \Downarrow_{h_1} v_1$

Therefore we get $(\theta, k - h_1, v_1) \in \llbracket (\mathbb{C} \ell_i \ell_o \tau) \sigma \rrbracket_V$

From Definition 2.35 we know that

$\forall k_{h_1} \leq (k - h_1), \theta'_e \sqsupseteq \theta_e, H_h, J. (k_{h_1}, H_h) \triangleright \theta'_e \wedge (H_h, v_1) \Downarrow_J^f (H', v'_{h_1}) \wedge J < k_{h_1} \implies$
 $\exists \theta'' \sqsupseteq \theta'_e. (k_{h_1} - J, H') \triangleright \theta'' \wedge (\theta'', k_{h_1} - J, v_1) \in \llbracket \tau \sigma \rrbracket_V \wedge$
 $(\forall a. H_h(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled} \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell)$

Instantiating k_{h_1} with $k - h_1$, H_h with H , θ'_e with θ_e . Since we know that $(H, \text{toLabeled}(e')) \Downarrow_j^f$
 (H', v_1) therefore $\exists J < j - h_1 < k - h_1$ s.t $(H, v_1) \Downarrow_J^f (H', v'_{h_1})$. And since we already
 know that $(k, H) \triangleright \theta_e$ therefore from Lemma 2.49 we get $(k - h_1, H) \triangleright \theta_e$

This means we have

$\exists \theta'' \sqsupseteq \theta'_e. (k - h_1 - J, H') \triangleright \theta'' \wedge (\theta'', k - h_1 - J, v_1) \in \llbracket \tau \sigma \rrbracket_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled} \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta'') \setminus \text{dom}(\theta'_e). \theta''(a) \searrow \ell) \quad (\text{FU-TL1})$

In order to prove (FU-TL0) we choose θ' as θ'' . Now we need to prove the following

- (a) $(k - j, H') \triangleright \theta''$:
 Since $(k - h_1 - J, H') \triangleright \theta''$ and $j = h_1 + J + 1$ therefore from Lemma 2.49 we get
 $(k - j, H') \triangleright \theta''$
- (b) $(\theta'', k - j - 1, v') \in [(\text{Labeled } \ell_o \tau \sigma)]_V$:
 From *cg - tolabeled* we know that $v' = \text{toLabeled}(v_1)$
 From Definition 2.33 it suffices to prove that $(\theta'', k - j - 1, v_1) \in [\tau \sigma]_V$
 We get this from (FU-TL1) and Lemma 2.45
- (c) $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell')$:
 Directly from (FU-TL1)
- (d) $(\forall a \in \text{dom}(\theta_n) \setminus \text{dom}(\theta_e). \theta_n(a) \searrow \ell)$:
 Directly from (FU-TL1)

□

Lemma 2.52 (CG: Subtyping unary). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \tau, \tau'$.

1. $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$
2. $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E \subseteq [(\tau' \sigma)]_E$

Proof. Proof of Statement (1)

Proof by induction on $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove: $[((\tau_1 \rightarrow \tau_2) \sigma)]_V \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

IH1: $[(\tau'_1 \sigma)]_V \subseteq [(\tau_1 \sigma)]_V$ (Statement (1))

$[(\tau_2 \sigma)]_E \subseteq [(\tau'_2 \sigma)]_E$ (Sub-A0, From Statement (2))

It suffices to prove: $\forall (\theta, n, \lambda x. e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V. (\theta, n, \lambda x. e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

This means that given some θ, n and $\lambda x. e_i$ s.t $(\theta, n, \lambda x. e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V$

Therefore from Definition 2.35 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall v. (\theta_1, i, v) \in [\tau_1 \sigma]_V \implies (\theta_1, i, e_i[v/x]) \in [\tau_2 \sigma]_E \quad (147)$$

And it suffices to prove: $(\theta, n, \lambda x. e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V$

Again from Definition 2.35, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall v. (\theta_2, j, v) \in [\tau'_1 \sigma]_V \implies (\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$$

This means that given some $\theta_2, j < n, v$ s.t $\theta \sqsubseteq \theta_2$ and $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$

And we are required to prove: $(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$

Since $(\theta_2, j, v) \in [\tau'_1 \sigma]_V$ therefore from IH1 we know that $(\theta_2, j, v) \in [\tau_1 \sigma]_V$

As a result from Equation 147 we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau_2 \sigma]_E$$

From (Sub-A0), we know that

$$(\theta_2, j, e_i[v/x]) \in [\tau'_2 \sigma]_E$$

2. CGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove: $[\tau_1 \times \tau_2 \sigma]_V \subseteq [(\tau'_1 \times \tau'_2) \sigma]_V$

IH1: $[\tau_1 \sigma]_V \subseteq [\tau'_1 \sigma]_V$ (Statement (1))

IH2: $[\tau_2 \sigma]_V \subseteq [\tau'_2 \sigma]_V$ (Statement (1))

It suffices to prove: $\forall(\theta, n, (v_1, v_2)) \in [(\tau_1 \times \tau_2) \sigma]_V. (\theta, n, (v_1, v_2)) \in [(\tau'_1 \times \tau'_2) \sigma]_V$

This means that given some θ, n and (v_1, v_2) $(\theta, (v_1, v_2)) \in [(\tau_1 \times \tau_2) \sigma]_V$

Therefore from Definition 2.35 we are given:

$$(\theta, n, v_1) \in [\tau_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau_2 \sigma]_V \tag{148}$$

And it suffices to prove: $(\theta, (v_1, v_2)) \in [(\tau'_1 \times \tau'_2) \sigma]_V$

Again from Definition 2.35, it suffices to prove:

$$(\theta, n, v_1) \in [\tau'_1 \sigma]_V \wedge (\theta, n, v_2) \in [\tau'_2 \sigma]_V$$

Since from Equation 148 we know that $(\theta, n, v_1) \in [\tau_1 \sigma]_V$ therefore from IH1 we have $(\theta, n, v_1) \in [\tau'_1 \sigma]_V$

Similarly since $(\theta, n, v_2) \in [\tau_2 \sigma]_V$ from Equation 148 therefore from IH2 we have $(\theta, n, v_2) \in [\tau'_2 \sigma]_V$

3. CGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove: $[(\tau_1 + \tau_2) \sigma]_V \subseteq [(\tau'_1 + \tau'_2) \sigma]_V$

IH1: $[\tau_1 \sigma]_V \subseteq [\tau'_1 \sigma]_V$ (Statement (1))

IH2: $[\tau_2 \sigma]_V \subseteq [\tau'_2 \sigma]_V$ (Statement (1))

It suffices to prove: $\forall(\theta, n, v_s) \in [(\tau_1 + \tau_2) \sigma]_V. (\theta, v_s) \in [(\tau'_1 + \tau'_2) \sigma]_V$

This means that given: $(\theta, n, v_s) \in [((\tau_1 + \tau_2) \sigma)]_V$

And it suffices to prove: $(\theta, n, v_s) \in [((\tau'_1 + \tau'_2) \sigma)]_V$

2 cases arise

(a) $v_s = \text{inl } v_i$:

From Definition 2.35 we are given:

$$(\theta, n, v_i) \in [\tau_1 \sigma]_V \quad (149)$$

And we are required to prove that:

$$(\theta, n, v_i) \in [\tau'_1 \sigma]_V$$

From Equation 149 and IH1 we know that

$$(\theta, n, v_i) \in [\tau'_1 \sigma]_V$$

(b) $v_s = \text{inr } v_i$:

From Definition 2.35 we are given:

$$(\theta, n, v_i) \in [\tau_2 \sigma]_V \quad (150)$$

And we are required to prove that:

$$(\theta, n, v_i) \in [\tau'_2 \sigma]_V$$

From Equation 150 and IH2 we know that

$$(\theta, n, v_i) \in [\tau'_2 \sigma]_V$$

4. CGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove: $[((\forall \alpha. \tau_1) \sigma)]_V \subseteq [(\forall \alpha. \tau_2) \sigma]_V$

It suffices to prove: $\forall (\theta, n, \Lambda e_i) \in [((\forall \alpha. \tau_1) \sigma)]_V. (\theta, n, \Lambda e_i) \in [((\forall \alpha. \tau_2) \sigma)]_V$

This means that given: $(\theta, n, \Lambda e_i) \in [((\forall \alpha. \tau_1) \sigma)]_V$

Therefore from Definition 2.35 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \forall \ell' \in \mathcal{L} \implies (\theta_1, i, e_i) \in [\tau_1 (\sigma \cup [\alpha \mapsto \ell'])]_E \quad (151)$$

And it suffices to prove: $(\theta, n, \Lambda e_i) \in [((\forall \alpha. \tau_2) \sigma)]_V$

Again from Definition 2.35, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \forall \ell' \in \mathcal{L} \implies (\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E$$

This means that given some $\theta_2, j < n, \ell' \in \mathcal{L}$ s.t $\theta \sqsubseteq \theta_2$

And we are required to prove: $(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E$

Since we are given $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \ell' \in \mathcal{L}$ therefore from Equation 151 we have

$$(\theta_2, j, e_i) \in [\tau_1 (\sigma \cup [\alpha \mapsto \ell'])]_E$$

$$[(\tau_1 (\sigma \cup [\alpha \mapsto \ell']))]_E \subseteq [(\tau_2 (\sigma \cup [\alpha \mapsto \ell']))]_E \text{ (Sub-F0, Statement (2))}$$

From (Sub-F0), we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma \cup [\alpha \mapsto \ell'])]_E$$

5. CGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2}$$

To prove: $[\!(c_1 \Rightarrow \tau_1) \sigma\!]_V \subseteq [\!(c_2 \Rightarrow \tau_2) \sigma\!]_V$

It suffices to prove: $\forall (\theta, n, \nu e_i) \in [\!(c_1 \Rightarrow \tau_1) \sigma\!]_V. (\theta, n, \nu e_i) \in [\!(c_2 \Rightarrow \tau_2) \sigma\!]_V$

This means that given: $(\theta, n, \nu e_i) \in [\!(c_1 \Rightarrow \tau_1) \sigma\!]_V$

Therefore from Definition 2.35 we are given:

$$\exists \theta_1. \theta \sqsubseteq \theta_1 \wedge \forall i < n. \mathcal{L} \models c_1 \sigma \implies (\theta_1, i, e_i) \in [\tau_1 (\sigma)]_E \quad (152)$$

And it suffices to prove: $(\theta, n, \nu e_i) \in [\!(c_2 \Rightarrow \tau_2) \sigma\!]_V$

Again from Definition 2.35, it suffices to prove:

$$\exists \theta_2. \theta \sqsubseteq \theta_2 \wedge \forall j < n. \mathcal{L} \models c_2 \sigma \implies (\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E$$

This means that given some θ_2, j s.t $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$

And we are required to prove: $(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E$

Since we are given $\theta \sqsubseteq \theta_2 \wedge j < n \wedge \mathcal{L} \models c_2 \sigma$ and $\mathcal{L} \models c_2 \sigma \implies c_1 \sigma$ therefore from Equation 152 we have

$$(\theta_2, j, e_i) \in [\tau_1 (\sigma)]_E$$

$$[(\tau_1 \sigma)]_E \subseteq [(\tau_2 \sigma)]_E \text{ (Sub-C0, Statement (2))}$$

From (Sub-C0), we know that

$$(\theta_2, j, e_i) \in [\tau_2 (\sigma)]_E$$

6. CGsub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove: $[\!(\text{Labeled } \ell \tau) \sigma\!]_V \subseteq [\!(\text{Labeled } \ell' \tau') \sigma\!]_V$

IH: $[(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$ (Statement (1))

It suffices to prove:

$$\forall(\theta, n, \text{Lb}(v_i)) \in [((\text{Labeled } \ell \tau) \sigma)]_V. (\theta, n, \text{Lb}(v_i)) \in [((\text{Labeled } \ell' \tau') \sigma)]_V$$

This means that given some θ, n and $\text{Lb}(v_i)$ s.t $(\theta, n, \text{Lb}(v_i)) \in [((\text{Labeled } \ell \tau) \sigma)]_V$

Therefore from Definition 2.35 we are given:

$$(\theta, n, v_i) \in [(\tau \sigma)]_V \quad (\text{SL})$$

And we are required to prove that

$$(\theta, n, \text{Lb}(v_i)) \in [((\text{Labeled } \ell' \tau') \sigma)]_V$$

From Definition 2.35 it suffices to prove

$$(\theta, n, v_i) \in [(\tau' \sigma)]_V$$

We get this directly from (SL) and IH

7. CGsub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove: $[((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V \subseteq [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V$

IH: $[(\tau \sigma)]_V \subseteq [(\tau' \sigma)]_V$ (Statement (1))

It suffices to prove:

$$\forall(\theta, n, e) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V. (\theta, n, e) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V$$

This means that given some θ, n and e s.t $(\theta, n, e) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V$

Therefore from Definition 2.35 we are given:

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\tau \sigma)]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) & \quad (\text{SC0}) \end{aligned}$$

And we are required to prove

$$(\theta, n, e) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V$$

So again from Definition 2.35 we need to prove

$$\begin{aligned} \forall k \leq n, \theta_e \sqsupseteq \theta, H, j.(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v') \wedge j < k &\implies \\ \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\tau' \sigma)]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma) & \end{aligned}$$

This means we are given some $k \leq n, \theta_e \sqsupseteq \theta, H, j < k$ s.t $(k, H) \triangleright \theta_e \wedge (H, e) \Downarrow_j^f (H', v')$
(SC1)

And we need to prove

$$\begin{aligned} \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in [(\tau' \sigma)]_V \wedge & \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge & \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma) & \end{aligned}$$

We instantiate (SC0) with k, θ_e, H, j from (SC1) and we get

$$\begin{aligned} & \exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau \sigma \rfloor_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \end{aligned}$$

Since $\tau \sigma <: \tau' \sigma$ therefore from IH we get

$$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v') \in \lfloor \tau' \sigma \rfloor_V$$

And since $\ell'_i \sqsubseteq \ell_i$ therefore we also have

$$\begin{aligned} & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma) \end{aligned}$$

8. CGsub-base:

Trivial

Proof of Statement(2)

It suffice to prove that

$$\forall (\theta, n, e) \in \lfloor (\tau \sigma) \rfloor_E. (\theta, n, e) \in \lfloor (\tau' \sigma) \rfloor_E$$

This means that we are given $(\theta, n, e) \in \lfloor (\tau \sigma) \rfloor_E$

From Definition 2.36 it means we have

$$\forall i < n. e \downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V \quad (\text{Sub-E0})$$

And we need to prove

$$(\theta, n, e) \in \lfloor (\tau' \sigma) \rfloor_E$$

From Definition 2.36 we need to prove

$$\forall i < n. e \downarrow_i v \implies (\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$$

This further means that given some $i < n$ s.t $e \downarrow_i v$, it suffices to prove that

$$(\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$$

Instantiating (Sub-E0) with the given i we get $(\theta, n - i, v) \in \lfloor \tau \sigma \rfloor_V$

Finally from Statement(1) we get $(\theta, n - i, v) \in \lfloor \tau' \sigma \rfloor_V$

□

Lemma 2.53 (CG: Binary interpretation of Γ implies Unary interpretation of Γ). $\forall W, \gamma, \Gamma, n.$

$$(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$$

Proof. Given: $(W, n, \gamma) \in [\Gamma]_V^A$

To prove: $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

From Definition 2.43 we know that we are given:

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma) \wedge \forall x \in \text{dom}(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$$

And we are required to prove:

$$\forall i \in \{1, 2\}. \forall m.$$

$$\text{dom}(\Gamma) \subseteq \text{dom}(\gamma \downarrow_i) \wedge \forall x \in \text{dom}(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$$

Case $i = 1$

Given some m we need to show:

- $dom(\Gamma) \subseteq dom(\gamma \downarrow_i)$:
 $dom(\gamma) = dom(\gamma \downarrow_i)$
Therefore, $dom(\Gamma) \subseteq (dom(\gamma) = dom(\gamma \downarrow_i))$ (Given)
- $\forall x \in dom(\Gamma). (W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$:
We are given: $\forall x \in dom(\Gamma). (W, n, \pi_1(\gamma(x)), \pi_2(\gamma(x))) \in [\Gamma(x)]_V^A$
Therefore from Lemma 2.44 we know that
 $\forall m'. (W.\theta_i, m', \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$
Instantiating m' with m we get
 $(W.\theta_i, m, \gamma \downarrow_i(x)) \in [\Gamma(x)]_V$

Case $i = 2$

Symmetric reasoning as in the $i = 1$ case above

□

Theorem 2.54 (CG: Fundamental theorem binary). $\forall \Sigma, \Psi, \Gamma, pc, W, \mathcal{A}, \mathcal{L}, e, \tau, \sigma, \gamma, n.$
 $\Sigma; \Psi; \Gamma \vdash e : \tau \wedge \mathcal{L} \models \Psi \sigma \wedge (W, n, \gamma) \in [\Gamma]_V^A \implies$
 $(W, n, e(\gamma \downarrow_1), e(\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

Proof. Proof by induction on the typing derivation

1. CG-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau} \text{CG-var}$$

To prove: $(W, n, x(\gamma \downarrow_1), x(\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

Say $e_1 = x(\gamma \downarrow_1)$ and $e_2 = x(\gamma \downarrow_2)$

From Definition 2.34 it suffices to prove that

$$\forall i < n. e_1 \downarrow_i v'_1 \wedge e_2 \downarrow v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

This means given some $i < n$ s.t $e_1 \downarrow_i v'_1 \wedge e_2 \downarrow v'_2$

We are required to prove: $(W, n - i, v'_1, v'_2) \in [\tau]_V^A$

From *cg-val* we know that $x(\gamma \downarrow_1) \downarrow x(\gamma \downarrow_1)$ and $x(\gamma \downarrow_2) \downarrow x(\gamma \downarrow_2)$

This means $v'_1 = x(\gamma \downarrow_1)$ and $v'_2 = x(\gamma \downarrow_2)$

Since $(W, n, \gamma) \in [\tau]_V^A$. Therefore from Definition 2.43 we know that

$$(W, n, v'_1, v'_2) \in [\tau]_V^A$$

From Lemma 2.46 we get

$$(W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

2. CG-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_i : \tau_2}{\Sigma; \Psi; \Gamma \vdash \lambda x. e_i : (\tau_1 \rightarrow \tau_2)}$$

To prove: $(W, n, \lambda x. e (\gamma \downarrow_1), \lambda x. e (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^A$

Say $e_1 = \lambda x. e (\gamma \downarrow_1)$ and $e_2 = \lambda x. e (\gamma \downarrow_2)$

From Definition of $[(\tau_1 \rightarrow \tau_2) \sigma]_E^A$ it suffices to prove that

$$\forall i < n. e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2 \implies (W, n - i, v'_1, v'_2) \in [\tau]_V^A$$

This means given some $i < n$ s.t $e_1 \Downarrow_i v'_1 \wedge e_2 \Downarrow v'_2$

From *cg - val* we know that $v'_1 = (\lambda x. e_i) \gamma \downarrow_1$ and $v'_2 = (\lambda x. e_i) \gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\lambda x. e_i) \gamma \downarrow_1, (\lambda x. e_i) \gamma \downarrow_2) \in [\tau]_V^A$$

From Definition 2.33 it suffices to prove

$\forall W' \sqsupseteq W, j < n, v_1, v_2.$

$$((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E^A) \wedge$$

$\forall \theta_l \sqsupseteq W. \theta_1, v_c, j.$

$$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E) \wedge$$

$\forall \theta_l \sqsupseteq W. \theta_2, v_c, j.$

$$((\theta_l, j, v_c) \in [\tau_1]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E) \quad (\text{FB-L0})$$

IH:

$$\forall W, n. (W, n, e_i (\gamma \downarrow_1 \cup \{x \mapsto v_1\}), e_i (\gamma \downarrow_2 \cup \{x \mapsto v_2\})) \in [\tau_2 \sigma]_E^A$$

s.t

$$(W, n, (\gamma \cup \{x \mapsto (v_1, v_2)\})) \in [\Gamma]_V^A$$

In order to prove (FB-L0) we need to prove the following:

(a) $\forall W' \sqsupseteq W, j < n, v_1, v_2.$

$$((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E^A):$$

This means given some $W' \sqsupseteq W, j < n, v_1, v_2$ s.t. $(W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A$

We need to prove $(W', j, e_1[v_1/x] \gamma \downarrow_1, e_2[v_2/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E^A$

We get this by instantiating IH with W' and j

(b) $\forall \theta_l \sqsupseteq W. \theta_1, v_c, j.$

$$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E):$$

This means given some $\theta_l \sqsupseteq W. \theta_1, v_c, j$ s.t $(\theta_l, j, v_c) \in [\tau_1 \sigma]_V$

We need to prove: $(\theta_l, j, e_1[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_E$

It is given to us that

$$(W, n, \gamma) \in [\Gamma]_V^A$$

Therefore from Lemma 2.53 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$

Instantiating m with j we get

$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$

From Lemma 2.48 we know that

$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$

Since we know that $(\theta_l, j, v_c) \in [\tau_1 \sigma]_V$

Therefore we also have

$(\theta_l, j, \gamma \downarrow_1 \cup \{x \mapsto v_c\}) \in [\Gamma \cup \{x \mapsto \tau_1 \sigma\}]_V$

Therefore, we can apply Theorem 2.51 to obtain

$(\theta_l, j, e[v_c/x] \gamma \downarrow_1) \in [\tau_2 \sigma]_V$

(c) $\forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x] \gamma \downarrow_2) \in [\tau_2 \sigma]_E):$

Similar reasoning as in the previous case

3. CG-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : (\tau_1 \rightarrow \tau_2) \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_1}{\Sigma; \Psi; \Gamma \vdash e_1 e_2 : \tau_2}$$

To prove: $(W, n, (e_1 e_2) (\gamma \downarrow_1), (e_1 e_2) (\gamma \downarrow_2)) \in [(\tau_2) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$\forall i < n. (e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$

This further means that given some $i < n$ s.t $(e_1 e_2) \gamma \downarrow_i v_{f1} \wedge e_2 \downarrow v_{f2}$

It suffices to prove:

$(W, n - i, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$

IH1: $(W, n, (e_1) (\gamma \downarrow_1), (e_1) (\gamma \downarrow_2)) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^A$

This means from Definition 2.34 we know that

$\forall j < n. e_1 \gamma \downarrow_1 \downarrow_j v_{h1} \wedge e_1 \gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - j, v_{h1}, v_{h2}) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^A$

Since we know that $(e_1 e_2) \gamma \downarrow_1 \downarrow_i v_{f1}$. Therefore $\exists j < i < n$ s.t $e_1 \gamma \downarrow_1 \downarrow_j v_{h1}$. Similarly since $(e_1 e_2) \gamma \downarrow_2 \downarrow v_{f2}$ therefore $e_1 \gamma \downarrow_2 \downarrow v_{h2}$

This means we have $(W, n - j, v_{h1}, v_{h2}) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^A$

From *cg - app* we know that $val_{h1} = \lambda x. e_{h1}$ and $val_{h2} = \lambda x. e_{h2}$

From Definition 2.33 this further means

$\forall W' \sqsupseteq W, J < (n - j), v_1, v_2.$

$((W', J, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies (W', J, e_{h1}[v_1/x], e_{h2}[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_1, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_c/x]) \in [\tau_2 \sigma]_E) \wedge$

$\forall \theta_l \sqsupseteq W.\theta_2, v_c, j.$

$((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2 \sigma]_E) \quad \text{(FB-A1)}$

IH2: $(W, n - j, (e_2) (\gamma \downarrow_1), (e_2) (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

This means from Definition 2.34 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \Downarrow_j v_{h1'} \wedge e_2 \gamma \downarrow_2 \Downarrow v_{h2'} \implies (W, n - j - k, v_{h1'}, v_{h2'}) \in [\tau_1 \sigma]_V^A$$

Since we know that $(e_1 e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$. Therefore $\exists k < i - j < n - j$ s.t $e_2 \gamma \downarrow_1 \Downarrow_k v_{h1'}$. Similarly since $(e_1 e_2) \gamma \downarrow_2 \Downarrow v_{f2}$ therefore $e_2 \gamma \downarrow_2 \Downarrow v_{h2'}$

$$\text{This means we have } (W, n - j - k, v_{h1'}, v_{h2'}) \in [\tau_1 \sigma]_V^A \quad (\text{FB-A2})$$

Instantiating the first conjunct of (FB-A1) as follows W' with W , J with $n - j - k$, v_1 and v_2 with v'_{h1} and v'_{h2} respectively, we obtain

$$(W, n - j - k, e_{h1}[v'_{h1}/x], e_{h2}[v'_{h2}/x]) \in [\tau_2 \sigma]_E^A$$

From Definition 2.34

$$\forall l < n - j - k. (e_{h1}[v'_{h1}/x]) \gamma \downarrow_l \Downarrow v_{f1} \wedge e_{h2}[v'_{h2}/x] \Downarrow v_{f2} \implies (W, n - j - k - l, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

Since we know that $(e_1 e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$. Therefore $\exists l < i - j - k < n - j - k$ s.t $e_{h1}[v'_{h1}/x] \Downarrow_l v_{f1}$. Similarly since $(e_1 e_2) \gamma \downarrow_2 \Downarrow v_{f2}$ therefore $e_{h2}[v'_{h2}/x] \Downarrow v_{f2}$

$$\text{Therefore we have } (W, n - j - k - l, v_{f1}, v_{f2}) \in [\tau_2 \sigma]_V^A$$

Since $i = j + k + l$ threfore we are done

4. CG-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$

To prove: $(W, n, (e_1, e_2) (\gamma \downarrow_1), (e_1, e_2) (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. (e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f1}, v_{f2}) \wedge (e_1, e_2) \gamma \downarrow_2 \Downarrow (v'_{f1}, v'_{f2}) \implies (W, n - i, (v_{f1}, v_{f1}), (v'_{f1}, v'_{f2})) \in [(\tau_1 \times \tau_2) \sigma]_V^A$$

This means that given some $i < n$ s.t $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f1}, v_{f2}) \wedge (e_1, e_2) \gamma \downarrow_2 \Downarrow (v'_{f1}, v'_{f2})$

We are required to prove

$$(W, n - i, (v_{f1}, v_{f1}), (v'_{f1}, v'_{f2})) \in [(\tau_1 \times \tau_2) \sigma]_V^A \quad (\text{FB-P0})$$

IH1: $(W, n, e_1 (\gamma \downarrow_1), e_1 (\gamma \downarrow_2)) \in [\tau_1 \sigma]_E^A$

This means from Definition 2.34 we know that

$$\forall j < n. e_1 \gamma \downarrow_1 \Downarrow_j v_{f1} \wedge e_1 \gamma \downarrow_2 \Downarrow v'_{f1} \implies (W, n - j, (v_{f1}, v'_{f1})) \in [\tau_1 \sigma]_V^A$$

Since we know that $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f1}, v_{f2})$. Therefore $\exists j < i < n$ s.t $e_1 \gamma \downarrow_1 \Downarrow_j v_{f1}$. Similarly since $(e_1 e_2) \gamma \downarrow_2 \Downarrow v_{f2}$ therefore $e_1 \gamma \downarrow_2 \Downarrow v'_{f1}$

This means we have

$$(W, n - j, (v_{f1}, v'_{f1})) \in [\tau_1 \sigma]_V^A \quad (\text{FB-P1})$$

IH2: $(W, n - j, e_2 (\gamma \downarrow_1), e_2 (\gamma \downarrow_2)) \in [\tau_2 \sigma]_E^A$

This means from Definition 2.34 we know that

$$\forall k < n - j. e_2 \gamma \downarrow_1 \Downarrow_i v_{f_2} \wedge e_2 \gamma \downarrow_2 \Downarrow v'_{f_2} \implies (W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Since we know that $(e_1, e_2) \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2})$. Therefore $\exists k < i - j < n - j$ s.t $e_2 \gamma \downarrow_1 \Downarrow_j v_{f_2}$. Similarly since $(e_1, e_2) \gamma \downarrow_2 \Downarrow v_{f_2}$ therefore $e_2 \gamma \downarrow_2 \Downarrow v'_{f_2}$

This means we have

$$(W, n - j - k, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (\text{FB-P2})$$

In order to prove (FB-P0) from Definition 2.33 it suffices to prove that

$$(W, n - i, (v_{f_1}, v'_{f_1})) \in [\tau_1 \sigma]_{\mathcal{V}}^A \text{ and } (W, n - i, (v_{f_2}, v'_{f_2})) \in [\tau_2 \sigma]_{\mathcal{V}}^A$$

Since $i = j + k + 1$ therefore from (FB-P1) and (FB-P2) and from Lemma 2.46 we get

$$(W, n - i, (v_{f_1}, v_{f_1}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

5. CG-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : (\tau_1 \times \tau_2)}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e') : \tau_1}$$

To prove: $(W, n, \text{fst}(e') (\gamma \downarrow_1), \text{fst}(e') (\gamma \downarrow_2)) \in [(\tau_1) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. \text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1} \implies (W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

This means that given some $i < n$ s.t $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1} \wedge \text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$

We are required to prove

$$(W, n - i, v_{f_1}, v_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (\text{FB-F0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall j < n. e' \gamma \downarrow_1 \Downarrow_i (v_{f_1}, v_{f_2}) \wedge e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, v'_{f_2}) \implies (W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

Since we know that $\text{fst}(e') \gamma \downarrow_1 \Downarrow_i v_{f_1}$. Therefore $\exists j < i < n$ s.t $e' \gamma \downarrow_1 \Downarrow_j (v_{f_1}, -)$. Similarly since $\text{fst}(e') \gamma \downarrow_2 \Downarrow v'_{f_1}$ therefore $e' \gamma \downarrow_2 \Downarrow (v'_{f_1}, -)$

This means we have

$$(W, n - j, (v_{f_1}, v_{f_2}), (v'_{f_1}, v'_{f_2})) \in [(\tau_1 \times \tau_2) \sigma]_{\mathcal{V}}^A$$

From Definition 2.33 we know that

$$(W, n - j, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

Since from $cg - fst$ $i = j + 1$ therefore from Lemma 2.46 we get

$$(W, n - i, v_{f_1}, v'_{f_1}) \in [\tau_1 \sigma]_{\mathcal{V}}^A$$

6. CG-snd:

Symmetric reasoning as in the CG-fst case above

7. CG-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau_1}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e') : (\tau_1 + \tau_2)}$$

To prove: $(W, n, \text{inl}(e') (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. \text{inl}(e') \gamma \downarrow_1 \downarrow_i \text{inl}(v_{f1}) \wedge \text{inl}(e') \gamma \downarrow_2 \downarrow \text{inl}(v'_{f1}) \implies \\ (W, n - i, \text{inl}(v_{f1}), \text{inl}(v'_{f1})) \in [(\tau_1 + \tau_2) \sigma]_V^A$$

This means that given some $i < n$ s.t $\text{inl}(e') \gamma \downarrow_1 \downarrow_i \text{inl}(v_{f1}) \wedge \text{fst}(e') \gamma \downarrow_2 \downarrow \text{inl}(v'_{f1})$

We are required to prove

$$(W, n - i, \text{inl}(v_{f1}), \text{inl}(v'_{f1})) \in [(\tau_1 + \tau_2) \sigma]_V^A \quad (\text{FB-IL0})$$

IH:

$$(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\tau_1 \times \tau_2) \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall j < n. e' \gamma \downarrow_1 \downarrow_i v_{f1} \wedge e' \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - j, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

Since we know that $\text{inl}(e') \gamma \downarrow_1 \downarrow_i \text{inl}(v_{f1})$. Therefore $\exists j < i < n$ s.t $e' \gamma \downarrow_1 \downarrow_j v_{f1}$. Similarly since $\text{fst}(e') \gamma \downarrow_2 \downarrow \text{inl}(v'_{f1})$ therefore $e' \gamma \downarrow_2 \downarrow v'_{f1}$

This means we have

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A \quad (\text{FB-IL1})$$

In order to prove (FB-IL0) from Definition 2.33 it suffices to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau_1 \sigma]_V^A$$

From $cg - \text{inl}$ since $i = j + 1$ therefore from (FB-IL1) and Lemma 2.46 we get (FB-IL0)

8. CG-inr:

Symmetric reasoning as in the CG-inl case above

9. CG-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_c : (\tau_1 + \tau_2) \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Sigma; \Psi; \Gamma \vdash \text{case}(e_c, x.e_1, y.e_2) : \tau}$$

To prove: $(W, n, \text{case}(e_c, x.e_1, y.e_2) (\gamma \downarrow_1), \text{inl}(e') (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \downarrow v_{f2} \implies \\ (W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A$$

This means that given some $i < n$ s.t $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \downarrow v_{f2}$

We are required to prove

$$(W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A \quad (\text{FB-C0})$$

IH1:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall j < n. e_c \gamma \downarrow_1 \Downarrow_i v_{h1} \wedge e_c \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\ (W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2) \sigma]_V^A$$

Since we know that $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$. Therefore $\exists j < i < n$ s.t. $e_c \gamma \downarrow_1 \Downarrow_j v_{h1}$.

Similarly since $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \Downarrow v'_{h1}$ therefore $e_c \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W, n - j, v_{h1}, v'_{h1}) \in [(\tau_1 + \tau_2) \sigma]_V^A \quad (\text{FB-C1})$$

2 cases arise

- (a) $v_{h1} = \text{inl}(v_1)$ and $v'_{h1} = \text{inl}(v'_1)$:

IH2:

$$(W, n, e_c (\gamma \downarrow_1), e_c (\gamma \downarrow_2)) \in [(\tau_1 + \tau_2) \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall k < n - j. e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_i v_{h2} \wedge e_1 \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2} \implies \\ (W, n - j - k, v_{h2}, v'_{h2}) \in [\tau \sigma]_V^A$$

Since we know that $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_1 \Downarrow_i v_{f1}$. Therefore $\exists k < i - j < n - j$ s.t. $e_1 \gamma \downarrow_1 \cup \{x \mapsto v_1\} \Downarrow_j v_{h2}$. Similarly since $\text{case}(e_c, x.e_1, y.e_2) \gamma \downarrow_2 \cup \{x \mapsto v'_1\} \Downarrow v'_{h2}$ therefore $e_1 \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$$(W, n - j - k, v_{h2}, v'_{h2}) \in [\tau \sigma]_V^A$$

From *cg - case1* we know that $i = j + k + 1$ therefore from Lemma 2.46 we get (FB-C0)

- (b) $v_{h1} = \text{inr}(v_1)$ and $v'_{h1} = \text{inr}(v'_1)$:

Symmetric case

10. CG-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \Lambda e' : \forall \alpha. \tau}$$

To prove: $(W, n, \Lambda e' (\gamma \downarrow_1), \Lambda e' (\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$

From Definition 2.34 it suffices to prove that

$$\forall i < n. (\Lambda e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (\Lambda e') \gamma \downarrow_2 \Downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [(\forall \alpha. \tau) \sigma]_V^A$$

This means given some $i < n$ s.t. $(\Lambda e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge (\Lambda e') \gamma \downarrow_2 \Downarrow v_{f2}$

From *cg - val* we know that $v_{f1} = (\Lambda e') \gamma \downarrow_1$ and $v_{f2} = (\Lambda e') \gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\Lambda e')\gamma \downarrow_1, (\Lambda e')\gamma \downarrow_2) \in [(\forall \alpha. \tau) \sigma]_V^A$$

Let $e_1 = (\Lambda e')\gamma \downarrow_1$ and $e_2 = (\Lambda e')\gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} & \forall W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}. ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A) \wedge \\ & \forall \theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E \wedge \\ & \forall \theta_l \sqsupseteq W.\theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha] \sigma]_E \quad (\text{FB-FI0}) \end{aligned}$$

$$\underline{\text{IH}}: \forall W, n. (W, n, e'(\gamma \downarrow_1), e'(\gamma \downarrow_2)) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_E^A$$

In order to prove (FB-FI0) we need to prove the following

$$(a) \quad \forall W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}. ((W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A):$$

This means given $W' \sqsupseteq W, j < (n - i), \ell' \in \mathcal{L}$ and we are required to prove

$$(W', j, e_1, e_2) \in [\tau[\ell'/\alpha] \sigma]_E^A$$

Instantiating IH with W' and j we get the desired

$$(b) \quad \forall \theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E:$$

This means given $\theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, j$ and we are required to prove

$$(\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E$$

Since from Lemma 2.53

$$(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$$

Therefore we get

$$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$$

And from Lemma 2.46 we also get

$$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$$

Therefore we can apply Theorem 2.51 to get

$$(\theta_l, j, e_1) \in [\tau[\ell''/\alpha] \sigma]_E$$

$$(c) \quad \forall \theta_l \sqsupseteq W.\theta_2, \ell'' \in \mathcal{L}, j. (\theta_l, j, e_2) \in [\tau[\ell''/\alpha] \sigma]_E:$$

Symmetric reasoning as before

11. CG-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \forall \alpha. \tau \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e' [] : \tau[\ell/\alpha]}$$

$$\text{To prove: } (W, n, e'[](\gamma \downarrow_1), e'[](\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$$

From Definition 2.34 it suffices to prove that

$$\forall i < n. (e'[]) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e'[]) \gamma \downarrow_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A$$

This means given some $i < n$ s.t $(e'[]) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e'[]) \gamma \downarrow_2 \downarrow v_{f2}$

We are required to prove:

$$(W, n - i, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A \quad (\text{FB-FE0})$$

$$\underline{\text{IH}}: (W, n, e'(\gamma \downarrow_1), e'(\gamma \downarrow_2)) \in [(\forall \alpha. \tau) \sigma]_E^A$$

From Definition 2.34 it suffices to prove that

$$\forall i < n. (e')\gamma \downarrow_1 \downarrow_i v_{h1} \wedge (e')\gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - i, v_{h1}, v_{h2}) \in [(\forall \alpha. \tau) \sigma]_V^A$$

Since we know that $(e')\gamma \downarrow_1 \downarrow_i v_{f1}$. Therefore $\exists j < i < n$ s.t $e' \gamma \downarrow_1 \downarrow_j v_{h1}$. Similarly since $(e')\gamma \downarrow_2 \downarrow v_{f2}$ therefore $e' \gamma \downarrow_2 \downarrow v_{h2}$

This means we have $(W, n - j, v_{h1}, v_{h2}) \in [(\forall \alpha. \tau) \sigma]_V^A$

From $cg - FE$ we know that $v_{h1} = \Lambda e_{h1}$ and $v_{h2} = \Lambda e_{h2}$

From Definition 2.33 this further means

$$\begin{aligned} \forall W' \sqsupseteq W, k < (n - j), \ell' \in \mathcal{L}. ((W', k, e_{h1}, e_{h2}) \in [\tau[\ell'/\alpha] \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, \ell'' \in \mathcal{L}, k. (\theta_l, k, e_{h1}) \in [\tau[\ell''/\alpha] \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, \ell'' \in \mathcal{L}, k. (\theta_l, k, e_{h2}) \in [\tau[\ell''/\alpha] \sigma]_E \quad (\text{FB-FE1}) \end{aligned}$$

Instantiating the first conjunct of (FB-FE1) with $W, n - j - 1$ and ℓ we get

$$(W, n - j - 1, e_{h1}, e_{h2}) \in [\tau[\ell/\alpha] \sigma]_E^A$$

This means from Definition 2.34 we know that

$$\forall l < n - j - 1. (e_{h1}) \downarrow_l v_{f1} \wedge e_{h2} \downarrow v_{f2} \implies (W, n - j - 1 - l, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A$$

Since we know that $(e')\gamma \downarrow_1 \downarrow_i v_{f1}$ therefore from $cg - FE$ we know that $(i = j + l + 1)$ and since we know that $i < n$ therefore we have $l < n - j - 1$ s.t $e_{h1} \gamma \downarrow_1 \downarrow_j v_{f1}$. Similarly since $(e')\gamma \downarrow_2 \downarrow v_{f2}$ therefore $e_{h2} \gamma \downarrow_2 \downarrow v_{f2}$

Therefore we get

$$(W, n - j - 1 - l, v_{f1}, v_{f2}) \in [(\tau[\ell/\alpha]) \sigma]_V^A \quad (\text{FB-FE2})$$

Since we know that $i = j + l + 1$ therefore from (FB-FE2) we get (FB-FE0)

12. CG-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e' : \tau}{\Sigma; \Gamma \vdash \nu e' : c \Rightarrow \tau}$$

To prove: $(W, n, \nu e' (\gamma \downarrow_1), \nu e' (\gamma \downarrow_2)) \in [(c \Rightarrow \tau) \sigma]_E^A$

From Definition 2.34 it suffices to prove that

$$\forall i < n. (\nu e')\gamma \downarrow_1 \downarrow_i v_{f1} \wedge (\nu e')\gamma \downarrow_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [(c \Rightarrow \tau) \sigma]_V^A$$

This means given some $i < n$ s.t $(\nu e')\gamma \downarrow_1 \downarrow_i v_{f1} \wedge (\nu e')\gamma \downarrow_2 \downarrow v_{f2}$

From $cg - val$ we know that $v_{f1} = (\nu e')\gamma \downarrow_1$ and $v_{f2} = (\nu e')\gamma \downarrow_2$

We are required to prove:

$$(W, n - i, (\nu e')\gamma \downarrow_1, (\nu e')\gamma \downarrow_2) \in [(c \Rightarrow \tau) \sigma]_V^A$$

Let $e_1 = (\nu e')\gamma \downarrow_1$ and $e_2 = (\nu e')\gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} \forall W' \sqsupseteq W, j < n. \mathcal{L} \models c \implies (W', j, e_1, e_2) \in [\tau \sigma]_E^A \wedge \\ \forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \implies (\theta_l, j, e_1) \in [\tau \sigma]_E \wedge \\ \forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau \sigma]_E \quad (\text{FB-CI0}) \end{aligned}$$

IH: $\forall W, n. (W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

In order to prove (FB-CI0) we need to prove the following

(a) $\forall W' \sqsupseteq W, j < n. \mathcal{L} \models c \sigma \implies (W', j, e_1, e_2) \in [\tau \sigma]_E^A$:

This means given $W' \sqsupseteq W, j < n, \mathcal{L} \models c \sigma$ and we are required to prove

$(W', j, e_1, e_2) \in [\tau \sigma]_E^A$

Instantiating IH with W' and j we get the desired

(b) $\forall \theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \sigma \implies (\theta_l, j, e_1) \in [\tau \sigma]_E$:

This means given $\theta_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c \sigma$ and we are required to prove

$(\theta_l, j, e_1) \in [\tau \sigma]_E$

Since from Lemma 2.53 $(W, n, \gamma) \in [\Gamma]_V^A \implies \forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, \gamma \downarrow_i) \in [\Gamma]_V$

Therefore we get

$(W.\theta_1, j, \gamma \downarrow_1) \in [\Gamma]_V$

And from Lemma 2.46 we also get

$(\theta_l, j, \gamma \downarrow_1) \in [\Gamma]_V$

Therefore we can apply Theorem 2.51 to get

$(\theta_l, j, e_1) \in [\tau \sigma]_E$

(c) $\forall \theta_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c \implies (\theta_l, j, e_2) \in [\tau \sigma]_E$:

Symmetric reasoning as before

13. CG-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : c \Rightarrow \tau \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e' \bullet : \tau}$$

To prove: $(W, n, e' \bullet (\gamma \downarrow_1), e' \bullet (\gamma \downarrow_2)) \in [\tau \sigma]_E^A$

From Definition 2.34 it suffices to prove that

$\forall i < n. (e' \bullet) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e' \bullet) \gamma \downarrow_2 \downarrow v_{f2} \implies (W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A$

This means given some $i < n$ s.t. $(e' \bullet) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e' \bullet) \gamma \downarrow_2 \downarrow v_{f2}$

We are required to prove:

$(W, n - i, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A$ (FB-CE0)

IH: $(W, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(c \Rightarrow \tau) \sigma]_E^A$

From Definition 2.34 it suffices to prove that

$\forall i < n. e' \gamma \downarrow_1 \downarrow_i v_{h1} \wedge e' \gamma \downarrow_2 \downarrow v_{h2} \implies (W, n - i, v_{h1}, v_{h2}) \in [(c \Rightarrow \tau) \sigma]_V^A$

Since we know that $(e' \bullet) \gamma \downarrow_1 \downarrow_i v_{f1}$. Therefore $\exists j < i < n$ s.t. $e' \gamma \downarrow_1 \downarrow_j v_{h1}$. Similarly since $(e' \bullet) \gamma \downarrow_2 \downarrow v_{f2}$ therefore $e' \gamma \downarrow_2 \downarrow v_{h2}$

This means we have $(W, n - j, v_{h1}, v_{h2}) \in [(c \Rightarrow \tau) \sigma]_V^A$

From $cg - CE$ we know that $v_{h1} = \nu e_{h1}$ and $v_{h2} = \nu e_{h2}$

From Definition 2.33 this further means

$\forall W' \sqsupseteq W, k < n - j. \mathcal{L} \models c \sigma \implies (W', k, e_1, e_2) \in [\tau \sigma]_E^A \wedge$

$\forall \theta_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c \sigma \implies (\theta_l, k, e_1) \in [\tau \sigma]_E \wedge$

$\forall \theta_l \sqsupseteq W.\theta_2, k. \mathcal{L} \models c \sigma \implies (\theta_l, k, e_2) \in [\tau \sigma]_E$ (FB-CE1)

Instantiating the first conjunct of (FB-CE1) with W , $n - j - 1$ and since we know that $\mathcal{L} \models c \sigma$ therefore we get

$$(W, n - j - 1, e_{h1}, e_{h2}) \in [\tau \sigma]_E^A$$

This means from Definition 2.34 we know that

$$\forall l < n - j - 1. (e_{h1}) \Downarrow_l v_{f1} \wedge e_{h2} \Downarrow v_{f2} \implies (W, n - j - 1 - l, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A$$

Since we know that $(e' \bullet) \gamma \Downarrow_1 \Downarrow_i v_{f1}$ therefore from cg-CE we know that $(i = j + l + 1)$ and since we know that $i < n$ therefore we have $l < n - j - 1$ s.t $e_{h1} \gamma \Downarrow_1 \Downarrow_l v_{f1}$. Similarly since $(e' \bullet) \gamma \Downarrow_2 \Downarrow v_{f2}$ therefore $e_{h2} \gamma \Downarrow_2 \Downarrow v_{f2}$

Therefore we get

$$(W, n - j - 1 - l, v_{f1}, v_{f2}) \in [\tau \sigma]_V^A \quad (\text{FB-CE2})$$

Since we know that $i = j + l + 1$ therefore from (FB-CE2) we get (FB-CE0)

14. CG-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(e') : \text{Labeled } \ell \tau}$$

To prove: $(W, n, \text{Lb}(e') (\gamma \Downarrow_1), \text{Lb}(e') (\gamma \Downarrow_2)) \in [\text{Labeled } \ell \tau \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\forall i < n. \text{Lb}(e') \gamma \Downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \Downarrow_2 \Downarrow \text{Lb}(v'_{f1}) \implies (W, n - i, \text{Lb}(v_{f1}), \text{Lb}(v'_{f1})) \in [\text{Labeled } \ell \tau \sigma]_V^A$$

This means that given some $i < n$ s.t $\text{Lb}(e') \gamma \Downarrow_1 \Downarrow_i \text{Lb}(v_{f1}) \wedge \text{Lb}(e') \gamma \Downarrow_2 \Downarrow \text{Lb}(v'_{f1})$

We are required to prove

$$(W, n - i, v_{f1}, v'_{f1}) \in [\text{Labeled } \ell \tau \sigma]_V^A \quad (\text{FB-LB0})$$

IH:

$$(W, n, e' (\gamma \Downarrow_1), e' (\gamma \Downarrow_2)) \in [\tau \sigma]_E^A$$

This means from Definition 2.34 we have:

$$\forall j < n. e' \gamma \Downarrow_1 \Downarrow_i v_{f1} \wedge e' \gamma \Downarrow_2 \Downarrow v'_{f1} \implies (W, n - j, v_{f1}, v'_{f1}) \in [\tau \sigma]_V^A$$

Since we know that $\text{Lb}(e') \gamma \Downarrow_1 \Downarrow_i \text{Lb}(v_{f1})$. Therefore $\exists j < i < n$ s.t $e' \gamma \Downarrow_1 \Downarrow_j v_{f1}$. Similarly since $\text{Lb}(e') \gamma \Downarrow_2 \Downarrow \text{Lb}(v'_{f1})$ therefore $e' \gamma \Downarrow_2 \Downarrow v'_{f1}$

This means we have

$$(W, n - j, v_{f1}, v'_{f1}) \in [\tau \sigma]_V^A \quad (\text{FB-LB1})$$

In order to prove (FB-LB0) from Definition 2.33 it suffices to prove that

$$(W, n - i, v_{f1}, v'_{f1}) \in [\tau \sigma]_V^A$$

From *cg-label* we know that $i = j + 1$. Therefore we get the desired from (FB-LB1) and Lemma 2.46

15. CG-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell \tau}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e') : \mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau}$$

To prove: $(W, n, \text{unlabel}(e') (\gamma \downarrow_1), \text{unlabel}(e') (\gamma \downarrow_2)) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. \text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V^A & \end{aligned}$$

This means that given some $i < n$ s.t $\text{unlabel}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{unlabel}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From *cg - val* we know that $v_{f1} = \text{unlabel}(e') \gamma \downarrow_1$ and $v'_{f1} = \text{unlabel}(e') \gamma \downarrow_2$. Also $i = 0$

We are required to prove

$$(W, n, \text{unlabel}(e') \gamma \downarrow_1, \text{unlabel}(e') \gamma \downarrow_2) \in [(\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau) \sigma]_V^A$$

This means from Definition 2.33 we need to prove

Let $e_1 = \text{unlabel}(e') \gamma \downarrow_1$ and $e_2 = \text{unlabel}(e') \gamma \downarrow_2$

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, (\ell_i \sqcup \ell) \sigma, v'_1, v'_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau']_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right) \end{aligned}$$

We need to show

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, (\ell_i \sqcup \ell) \sigma, v'_1, v'_2, \tau \sigma): \end{aligned}$$

Also given is some $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j$ s.t $(k, H_1, H_2) \triangleright W_e$ and $(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k$

And we are required to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, (\ell_i \sqcup \ell) \sigma, v'_1, v'_2, \tau \sigma) \quad (\text{FB-U0})$$

$$\underline{\text{IH}}: (W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{Labeled } \ell \tau) \sigma]_E^A$$

This means from Definition 2.34 we are given

$$\begin{aligned} \forall I < k. e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1}) &\implies \\ (W_e, k - I, \text{Lb}(v_{h1}), \text{Lb}(v'_{h1})) \in [(\text{Labeled } \ell \tau) \sigma]_V^A & \end{aligned}$$

Since we know that

$$\begin{aligned} (H_1, \text{unlabel}(e') \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{unlabel}(e') \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2) \wedge j < k &\text{ therefore} \\ \exists I < j < k \text{ s.t } e' \gamma \downarrow_1 \Downarrow_I \text{Lb}(v_{h1}) \wedge e' \gamma \downarrow_2 \Downarrow \text{Lb}(v'_{h1}) & \end{aligned}$$

Therefore we have

$$(W_e, k - I, \text{Lb}(v_{h_1}), \text{Lb}(v'_{h_1})) \in \lceil (\text{Labeled } \ell \tau) \sigma \rceil_V^A$$

This means from Definition 2.33 we have

$$\text{ValEq}(\mathcal{A}, W_e, k - I, \ell \sigma, v_{h_1}, v'_{h_1}, \tau \sigma) \quad (\text{FB-U1})$$

In order to prove (FB-U0) we choose W' as W_e and from *cg - unlabel* we know that $H'_1 = H_1$ and $H'_2 = H_2$. And we already know that $(k, H_1, H_2) \triangleright W_e$. Therefore from Lemma 2.50 we get $(k - j, H_1, H_2) \triangleright W_e$

From *cg - unlabel* we know that v'_1, v'_2 in (FB-U0) is v_{h_1}, v'_{h_1} respectively. And since from (FB-U1) we know that $\text{ValEq}(\mathcal{A}, W_e, k - I, \ell \sigma, v_{h_1}, v'_{h_1}, \tau \sigma)$. Therefore from Lemma 2.55 we get

$$\text{ValEq}(\mathcal{A}, W_e, k - j, (\ell_i \sqcup \ell) \sigma, v_{h_1}, v'_{h_1}, \tau \sigma)$$

$$(b) \quad \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \sigma \rfloor_V \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right):$$

Case $l = 1$

$$\text{Given some } k, \theta_e \sqsupseteq W.\theta_l, H, j \text{ s.t. } (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k$$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \sigma \rfloor_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$$

Since $(W, n, \gamma) \in \lceil \Gamma \rceil_V^A$ therefore from Lemma 2.53 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in \lfloor \Gamma \rfloor_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in \lfloor \Gamma \rfloor_V$$

$$\text{Instantiating } m \text{ with } k \text{ we get } (W.\theta_1, k, \gamma \downarrow_1) \in \lfloor \Gamma \rfloor_V$$

Now we can apply Theorem 2.51 to get

$$(W.\theta_1, k, (\text{unlabel } e')\gamma \downarrow_1) \in \lfloor (\mathbb{C} \ell_i \ell_i \sqcup \ell \tau) \sigma \rfloor_E$$

This means from Definition 2.36 we get

$$\forall c < k. (\text{unlabel } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in \lfloor (\mathbb{C} \ell_i \ell_i \sqcup \ell \tau) \sigma \rfloor_V$$

This further means that given some $c < k$ s.t $(\text{unlabel } e')\gamma \downarrow_1 \Downarrow_c v$. From *cg - val* we know that $c = 0$ and $v = (\text{unlabel } e')\gamma \downarrow_1$

$$\text{And we have } (W.\theta_1, k, (\text{unlabel } e')\gamma \downarrow_1) \in \lfloor (\mathbb{C} \ell_i \ell_i \sqcup \ell \tau) \sigma \rfloor_V$$

From Definition 2.35 we have

$$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{unlabel } e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in \lfloor \tau \rfloor_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_1)$$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

16. CG-tolabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \mathbb{C} \ell_i \ell_o \tau}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e') : \mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)}$$

To prove: $(W, n, \text{toLabeled}(e') (\gamma \downarrow_1), \text{toLabeled}(e') (\gamma \downarrow_2)) \in [\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. \text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau) \sigma]_V^A & \end{aligned}$$

This means that given some $i < n$ s.t $\text{toLabeled}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{toLabeled}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$
From *cg - val* we know that $v_{f1} = \text{toLabeled}(e') \gamma \downarrow_1$, $v_{f2} = \text{toLabeled}(e') \gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, \text{toLabeled}(e') \gamma \downarrow_1, \text{toLabeled}(e') \gamma \downarrow_2) \in [\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau) \sigma]_V^A$$

Let $v_1 = \text{toLabeled}(e') \gamma \downarrow_1$ and $v_2 = \text{toLabeled}(e') \gamma \downarrow_2$

This means from Definition 2.33 we are required to prove

$$\begin{aligned} &(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, (\text{Labeled } \ell_o \tau) \sigma)) \wedge \\ &\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge \\ &(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \\ &(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)) \end{aligned}$$

We need to prove:

$$\begin{aligned} \text{(a)} \quad &\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_2, v'_1, v'_2, (\text{Labeled } \ell_o \tau) \sigma): \end{aligned}$$

This means that we are given some $k \leq n$, $W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$ s.t

$$(k, H_1, H_2) \triangleright W_e \text{ and } (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$$

And we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, (\text{Labeled } \ell_o \tau) \sigma) \quad (\text{FB-TL0})$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\mathbb{C} \ell_i \ell_o \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \Downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, n - J, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_i \ell_o \tau \sigma]_V^A$$

Since we know that $(H_1, \text{toLabeled}(e') \gamma \downarrow_1) \Downarrow_j (H'_1, v'_1)$ and $(H_2, \text{toLabeled}(e') \gamma \downarrow_1) \Downarrow_j (H'_2, v'_2)$. Therefore from *cg - val* we know that $\exists J < j < k \leq n$ s.t $e' \gamma \downarrow_1 \Downarrow_J v_{h1}$ and similarly we also know that $e' \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - J, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_i \ell_o \tau \sigma]_V^A$$

From Definition 2.33 we know that

$$\begin{aligned} & (\forall k_1 \leq (k - J), W_e'' \sqsupseteq W_e. \forall H_1'', H_2''. (k_1, H_1'', H_2'') \triangleright W_e'' \wedge \forall v_1'', v_2'', m. \\ & (H_1'', v_{h1}) \Downarrow_m^f (H_1', v_1'') \wedge (H_2'', v'_{h1}) \Downarrow^f (H_2', v_2'') \wedge m < k_1 \implies \\ & \exists W' \sqsupseteq W_e''. (k_1 - m, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k_1 - m, \ell_o, v_1'', v_2'', \tau \sigma)) \wedge \\ & \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v_l') \in [\tau \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)) \quad (\text{FB-TL1}) \end{aligned}$$

We instantiate W_e'' with W_e , H_1'' with H_1 , H_2'' with H_2 and k_1 with k in (FB-TL1).

Since we know that $(H_1, \text{toLabeled}(e')\gamma \downarrow_1) \Downarrow_j^f (H_1', v_1') \wedge (H_2, \text{toLabeled}(e')\gamma \downarrow_2) \Downarrow^f (H_2', v_2')$, therefore $\exists m < j < k \leq n$ s.t $(H_1, v_{h1}) \Downarrow_m^f (H_1', v_1') \wedge (H_2, v'_{h1}) \Downarrow^f (H_2', v_2')$

This means we have

$$\begin{aligned} & \exists W' \sqsupseteq W_e. (k - m, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - m, \ell_o, v_1'', v_2'', \tau \sigma) \\ & (\text{FB-TL2}) \end{aligned}$$

In order to prove (FB-TL0) we choose W' as W' from (FB-TL2). Since from *cg - tolabeled* we know that $v_1' = \text{Lb}_{\ell_o}(v_1'')$, $v_2' = \text{Lb}_{\ell_o}(v_2'')$ and $j = m + 1$, therefore from Lemma 2.50 we get $(k - j, H_1', H_2') \triangleright W'$.

Since we have by assumption that $\ell_i \sqsubseteq \ell_o$ therefore the following cases arise

i. $\ell_i \sqsubseteq \ell_o \sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 it suffices to prove that

$$(W', k - j, v_1', v_2') \in [(\text{Labeled } \ell_o \tau) \sigma]_V^A$$

Since $v_1' = \text{Lb}_{\ell_o}(v_1'')$ and $v_2' = \text{Lb}_{\ell_o}(v_2'')$. Therefore from Definition 2.33 it suffices to prove that

$$\text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v_1'', v_2'', \tau \sigma)$$

We get this from (FB-TL2) and Lemma 2.55

ii. $(\ell_i \sqsubseteq \ell_o) \not\sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 it suffices to prove that

$$\forall m. (W', m, v_1') \in [(\text{Labeled } \ell_o \tau) \sigma]_V \text{ and } \forall m. (W', m, v_2') \in [(\text{Labeled } \ell_o \tau) \sigma]_V$$

Since $\ell_o \not\sqsubseteq \mathcal{A}$ therefore we get this from (FB-TL2), Definition 2.32 and Definition 2.35

iii. $(\ell_i \sqsubseteq \mathcal{A} \sqsubseteq \ell_o)$:

In this case from Definition 2.32 it suffices to prove that

$$(W', k - j, v_1', v_2') \in [(\text{Labeled } \ell_o \tau) \sigma]_V^A$$

Since $v_1' = \text{Lb}_{\ell_o}(v_1'')$ and $v_2' = \text{Lb}_{\ell_o}(v_2'')$. Therefore from Definition 2.33 it suffices to prove that

$$\forall m. (W', m, v_1'') \in [\tau \sigma]_V \text{ and } \forall m. (W', m, v_2'') \in [\tau \sigma]_V$$

We obtain this directly from (FB-TL2) and Definition 2.32

$$\begin{aligned} & (\text{b}) \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \\ & \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v_l') \in [(\text{Labeled } \ell_o \tau) \sigma]_V \wedge \\ & (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sqsubseteq \ell') \wedge \\ & (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i)) \end{aligned}$$

Case $l = 1$

Given some $k, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\text{Labeled } \ell_o \tau \sigma]_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$

Since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 2.53 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$ and $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_E$

This means from Definition 2.36 we get

$\forall c < k. (\text{toLabeled } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_V$

Instantiating c with 0 and from $cg - val$ we know $v = (\text{toLabeled } e')\gamma \downarrow_1$

And we have $(W.\theta_1, k, (\text{toLabeled } e')\gamma \downarrow_1) \in [(\mathbb{C} \ell_i \ell_i \text{ Labeled } \ell_o \tau) \sigma]_V$

From Definition 2.35 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, (\text{toLabeled } e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies$

$\exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [\text{Labeled } \ell_o \tau \sigma]_V \wedge$
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_i \sigma)$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

17. CG-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \tau}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e') : \mathbb{C} \ell_i \ell_i \tau}$$

To prove: $(W, n, \text{ret}(e') (\gamma \downarrow_1), \text{ret}(e') (\gamma \downarrow_2)) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$\forall i < n. \text{ret}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \Downarrow v'_{f1} \implies$
 $(W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^A$

This means that given some $i < n$ s.t $\text{ret}(e') \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{ret}(e') \gamma \downarrow_2 \Downarrow v'_{f1}$

From $cg - val$ we know that $v_{f1} = \text{ret}(e')\gamma \downarrow_1$, $v_{f2} = \text{ret}(e')\gamma \downarrow_2$ and $i = 0$

We are required to prove

$(W, n, \text{ret}(e')\gamma \downarrow_1, \text{ret}(e')\gamma \downarrow_2) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^A$

Let $v_1 = \text{ret}(e')\gamma \downarrow_1$ and $v_2 = \text{ret}(e')\gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, \tau) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall v, i. (e_l \Downarrow_i v_l) \implies \right. \\ & \forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \rfloor_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right) \end{aligned}$$

It suffices to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, \tau): \end{aligned}$$

We are given is some $k \leq n$, $W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$ s.t $(k, H_1, H_2) \triangleright W_e$ and $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

From $cg - \text{ret}$ we know that $H'_1 = H_1$ and $H'_2 = H_2$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H_1, H_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_i, v'_1, v'_2, \tau) \quad (\text{FB-R0})$$

$$\underline{\text{IH}}: (W_e, n, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in \lceil \tau \sigma \rceil_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall J < k. e' \gamma \downarrow_1 \Downarrow_J v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies (W_e, k - J, v_{h1}, v'_{h1}) \in \lceil \tau \sigma \rceil_V^A$$

Since we know that $(H_1, \text{ret}(e')\gamma \downarrow_1) \Downarrow_j^f (H_1, v'_1) \wedge (H_2, \text{ret}(e')\gamma \downarrow_2) \Downarrow_j^f (H_2, v'_2)$, therefore $\exists J < j < k$ s.t $e' \gamma \downarrow_1 \Downarrow_J v_{h1}$ and similarly $e' \gamma \downarrow_2 \Downarrow v'_{h1}$.

$$\text{Therefore we have } (W_e, k - J, v_{h1}, v'_{h1}) \in \lceil \tau \sigma \rceil_V^A \quad (\text{FB-R1})$$

In order to prove (FB-R0) we choose W' as W_e and from $cg - \text{ret}$ we know that $v'_1 = v_{h1}$ and $v'_2 = v'_{h1}$. We need to prove the following:

i. $(k - j, H_1, H_2) \triangleright W_e$:

Since we have $(k, H_1, H_2) \triangleright W_e$ therefore from Lemma 2.50 we get

$$(k - j, H_1, H_2) \triangleright W_e$$

ii. $\text{ValEq}(\mathcal{A}, W_e, k - j, \ell_i, v'_1, v'_2, \tau)$:

2 cases arise:

A. $\ell_i \sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 it suffices to prove

$$(W_e, k - j, v'_1, v'_2) \in \lceil \tau \sigma \rceil_V^A$$

Since $j = J + 1$ therefore we get this from (FB-R1) and Lemma 2.46

B. $\ell_i \not\sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 it suffices to prove that

$$\forall m. (W_e, m, v'_1) \in \lfloor \tau \sigma \rfloor_V \text{ and } \forall m. (W_e, m, v'_2) \in \lfloor \tau \sigma \rfloor_V$$

We get this From (FB-R1) and Lemma 2.44

(b) $\forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W.\theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \sigma \rfloor_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma):$

Case $l = 1$

Given some $k, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t. $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor \tau \sigma \rfloor_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$

Since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 2.53 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$ and $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in \lfloor (\mathbb{C} \ell_i \ell_i \tau) \sigma \rfloor_E$

This means from Definition 2.36 we get

$\forall c < k. (\text{ret } e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in \lfloor (\mathbb{C} \ell_i \ell_i \tau) \sigma \rfloor_V$

Instantiating c with 0 and from $cg - val$ we know that $v = (\text{ret } e')\gamma \downarrow_1$

And we have $(W.\theta_1, k, (\text{ret } e')\gamma \downarrow_1) \in \lfloor (\mathbb{C} \ell_i \ell_i \tau) \sigma \rfloor_V$

From Definition 2.35 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies$
 $\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in \lfloor \tau \sigma \rfloor_V \wedge$
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_i \sigma)$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

18. CG-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_l : \mathbb{C} \ell_i \ell \tau \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_b : \mathbb{C} \ell \ell_o \tau'}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_l, x.e_b) : \mathbb{C} \ell_i \ell_o \tau'}$$

To prove: $(W, n, \text{bind}(e_l, x.e_b) (\gamma \downarrow_1), \text{bind}(e_l, x.e_b) (\gamma \downarrow_2)) \in \lfloor \mathbb{C} \ell_i \ell_o \tau' \sigma \rfloor_E^A$

This means from Definition 2.34 we need to prove:

$\forall i < n. \text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1} \implies$
 $(W, n - i, v_{f1}, v'_{f1}) \in \lfloor \mathbb{C} \ell_i \ell_o \tau' \sigma \rfloor_V^A$

This means that given some $i < n$ s.t. $\text{bind}(e_l, x.e_b) \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge \text{bind}(e_l, x.e_b) \gamma \downarrow_2 \Downarrow v'_{f1}$

From $cg - val$ we know that $v_{f1} = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$, $v_{f2} = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, \text{bind}(e_l, x.e_b)\gamma \downarrow_1, \text{bind}(e_l, x.e_b)\gamma \downarrow_2) \in [\mathbb{C} \ell_i \ell_o \tau' \sigma]_V^A$$

Let $v_1 = \text{bind}(e_l, x.e_b)\gamma \downarrow_1$ and $v_2 = \text{bind}(e_l, x.e_b)\gamma \downarrow_2$

This means from Definition 2.33 we need to prove

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \sigma \wedge \ell_i \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i) \right) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma): \end{aligned}$$

This means we are given some $k \leq n, W_e \sqsupseteq W, H_1, H_2$ s.t. $(k, H_1, H_2) \triangleright W_e$

Also given some $v'_1, v'_2, j < k$ s.t. $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau \sigma) \quad (\text{FB-B0})$$

IH1:

$$(W_e, k, e_l (\gamma \downarrow_1), e_l (\gamma \downarrow_2)) \in [\mathbb{C} \ell_i \ell \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} & \forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\ & (W_e, k - f, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_i \ell \tau \sigma]_V^A \end{aligned}$$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$ therefore $\exists f < j < k$ s.t. $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\mathbb{C} \ell_i \ell \tau \sigma]_V^A$$

This means from Definition 2.33 we have

$$\begin{aligned} & \left(\forall K \leq (k - f), W_e \sqsupseteq W_e. \forall H''_1, H''_2. (K, H''_1, H''_2) \triangleright W_e \wedge \forall v''_1, v''_2, J. \right. \\ & (H''_1, v_{h1}) \Downarrow_J^f (H''_1, v''_1) \wedge (H''_2, v'_{h1}) \Downarrow^f (H''_2, v''_2) \wedge J < K \implies \\ & \left. \exists W'' \sqsupseteq W_e. (K - J, H''_1, H''_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', K - J, \ell \sigma, v''_1, v''_2, \tau \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \sigma \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right) \end{aligned}$$

Instantiating K with $(k - f)$, W'_e with W_e , H'_1 with H_1 and H'_2 with H_2 in the first conjunct of the above equation. Since we know that $(k, H_1, H_2) \triangleright W_e$ therefore from Lemma 2.50 we also have $(k - f, H_1, H_2) \triangleright W_e$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$ therefore $\exists J < j - f < k - f$ s.t $(H_1, v_{h1}) \Downarrow_J^f (H'_1, v'_{h1}) \wedge (H_2, v'_{h1}) \Downarrow^f (H'_2, v'_2)$

This means we have

$$\exists W'' \sqsupseteq W'_e.(k - f - J, H'_1, H'_2) \triangleright W'' \wedge \text{ValEq}(\mathcal{A}, W'', k - f - J, \ell \sigma, v'_1, v'_2, \tau \sigma) \quad (\text{FB-B1})$$

From Definition 2.32 two cases arise:

i. $\ell \sigma \sqsubseteq \mathcal{A}$:

In this case we know that $(W'', k - f - J, v'_1, v'_2) \in [\tau \sigma]_V^A$

IH2:

$$(W'', k - f - J, e_b(\gamma \downarrow_1 \cup \{x \mapsto v'_1\}), e_b(\gamma \downarrow_2 \cup \{x \mapsto v'_2\})) \in [\mathbb{C} \ell \ell_o \tau' \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall s < k - f - J. e_b(\gamma \downarrow_1 \cup \{x \mapsto v'_1\}) \Downarrow_s v_{h2} \wedge e_b(\gamma \downarrow_2 \cup \{x \mapsto v'_2\}) \Downarrow v'_{h2} \implies (W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\mathbb{C} \ell \ell_o \tau' \sigma]_V^A$$

Since we know that $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$ therefore $\exists s < j - f - J < k - f - J$ s.t $e_b(\gamma \downarrow_1 \cup \{x \mapsto v'_1\}) \Downarrow_s v_{h2} \wedge e_b(\gamma \downarrow_2 \cup \{x \mapsto v'_2\}) \Downarrow v'_{h2}$

This means we have

$$(W'', k - f - J - s, v_{h2}, v'_{h2}) \in [\mathbb{C} \ell \ell_o \tau' \sigma]_V^A$$

This means from Definition 2.33 we know that

$$\left(\forall K_s \leq (k - f - J - s), W_s \sqsupseteq W''. \forall H_1, H_2. (K_s, H_1, H_2) \triangleright W_s \wedge \forall v'_{s1}, v'_{s2}, J_s. \right.$$

$$(H_1, v_{h2}) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H_2, v'_{h2}) \Downarrow^f (H'_{s2}, v'_{s2}) \wedge J_s < K_s \implies$$

$$\left. \exists W'_s \sqsupseteq W_s. (K_s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, K_s - J_s, \ell_i, v'_1, v'_2, \tau' \sigma) \right) \wedge$$

$$\forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$$

$$\left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge \right.$$

$$\left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \sigma \wedge \ell_1 \sqsubseteq \ell') \wedge \right.$$

$$\left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1) \right)$$

Instantiating K_s with $(k - f - J - s)$, W_s with W'' , H_1 with H'_1 and H_2 with H_2 . Since we know that $(k - f - J, H'_1, H'_2) \triangleright W''$ therefore from Lemma 2.50 we also have $(k - f - J - s, H'_1, H'_2) \triangleright W''$

Since we know that $(H_1, \text{bind}(e_l, x.e_b) \gamma \downarrow_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, \text{bind}(e_l, x.e_b) \gamma \downarrow_2) \Downarrow^f (H'_2, v'_2)$ therefore $\exists J_s < j - f - J - s < k - f - J - s$ s.t $(H'_1, v'_{s1}) \Downarrow_{J_s}^f (H'_{s1}, v'_{s1}) \wedge (H'_2, v'_{s2}) \Downarrow^f (H'_{s2}, v'_{s2})$

This means we have

$$\exists W'_s \sqsupseteq W_s. (k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s \wedge \text{ValEq}(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma) \quad (\text{FB-B2})$$

In order to prove (FB-B0) we choose W' as W'_s . From $cg - bind$ we know that $H'_1 = H'_{s1}$, $H'_2 = H'_{s2}$, $v'_1 = v'_{s1}$, $v'_2 = v'_{s2}$ and $j = f + J + s + J_s + 1$. And we need to prove:

- A. $(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$:
 Since from (FB-B2) we know that $(k - f - J - s - J_s, H'_{s1}, H'_{s2}) \triangleright W'_s$ therefore from Lemma 2.50 we get
 $(k - j, H'_{s1}, H'_{s2}) \triangleright W'_s$
- B. $ValEq(\mathcal{A}, W'_s, k - j, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma)$:
 Since from (FB-B2) we know that $ValEq(\mathcal{A}, W'_s, k - f - J - s - J_s, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma)$ therefore from Lemma 2.55 we get
 $ValEq(\mathcal{A}, W'_s, k - j, \ell_o, v'_{s1}, v'_{s2}, \tau' \sigma)$

ii. $\ell \sigma \not\sqsubseteq \mathcal{A}$:

From (FB-B0) we know that we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell_o, v'_1, v'_2, \tau' \sigma)$$

Since $\ell_i \sigma \sqsubseteq \ell \sigma \sqsubseteq \ell_o \sigma$ (by assumption) and $\ell \sigma \not\sqsubseteq \mathcal{A}$ therefore we have $\ell_o \sigma \not\sqsubseteq \mathcal{A}$

This means that from Definition 2.32 it suffices to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \forall m_{u1}. (W'.\theta_1, m_{u1}, v'_1) \in \llbracket \tau' \sigma \rrbracket_V \wedge \forall m_{u2}. (W'.\theta_2, m_{u2}, v'_2) \in \llbracket \tau' \sigma \rrbracket_V$$

This means given some m_{u1}, m_{u2} and we need to prove

$$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge (W'.\theta_1, m_{u1}, v'_1) \in \llbracket \tau' \sigma \rrbracket_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in \llbracket \tau' \sigma \rrbracket_V \quad (\text{FB-B01})$$

In this case we know that

$$\forall m. (W''.\theta_1, m, v''_1) \in \llbracket \tau \sigma \rrbracket_V \text{ and } \forall m. (W''.\theta_2, m, v''_2) \in \llbracket \tau \sigma \rrbracket_V \quad (\text{FB-B3})$$

Since $\text{bind}(e_l, x.e_b)\gamma \downarrow_1 \Downarrow_j v'_1$ therefore $\exists J_1 < j - f - J < k - f - J$ s.t $(e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \Downarrow_{J_1} v'_1$. Similarly, $\exists J'_1 < j - f - J - J_1 < k - f - J - J_1$ s.t $(H'_1, v'_1) \Downarrow_{J'_1}^f -$

Instantiating m with $m_{u1} + 1 + J_1 + J'_1$ in the first conjunct of (FB-B3)

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, v''_1) \in \llbracket \tau \sigma \rrbracket_V$$

Since $(W, n, \gamma) \in \llbracket \Gamma \rrbracket_V^A$ therefore from Lemma 2.53 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V$$

Instantiating m with $m_{u1} + 1 + J_1 + J'_1$ we get $(W.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V$

From Lemma 2.47 we know that

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, \gamma \downarrow_1) \in \llbracket \Gamma \rrbracket_V \quad (\text{FB-B4})$$

Now we can apply Theorem 2.51 to get

$$(W''.\theta_1, m_{u1} + 1 + J_1 + J'_1, (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\}) \in \llbracket (\mathbb{C} \ell \ell_o \tau') \sigma \rrbracket_E$$

This means from Definition 2.36 we get

$$\forall c_1 < m_{u1} + 1 + J_1 + J'_1. (e_b)\gamma \downarrow_1 \cup \{x \mapsto v''_1\} \Downarrow_{c_1} v_{o1} \implies (W''.\theta_1, m_{u1} + 1 + J_1 + J'_1 - c_1, v_{o1}) \in \llbracket (\mathbb{C} \ell \ell_o \tau') \sigma \rrbracket_V \quad (\text{FB-B5})$$

Instantiating c_1 with J_1 in (FB-B5)

$$\text{Therefore we have } (W''.\theta_1, m_{u1} + 1 + J'_1, v_{o1}) \in \llbracket (\mathbb{C} \ell \ell_o \tau') \sigma \rrbracket_V$$

From Definition 2.35 we have

$$\forall K \leq (m_{u1} + 1 + J'_1), \theta'_e \sqsupseteq W''.\theta_1, H_1, J_2. (K, H_1) \triangleright \theta'_e \wedge (H_1, v_{o1}) \Downarrow_{J_2}^f (H''_1, v'_1) \wedge J_2 < K \implies$$

$$\exists \theta'_1 \sqsupseteq \theta'_e. (K - J_2, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, K - J_2, v'_1) \in \llbracket \tau' \sigma \rrbracket_V \wedge$$

$$\begin{aligned} (\forall a. H_1(a) \neq H_1''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1)/\text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_i \sigma) \end{aligned}$$

Instantiating K with $m_{u1} + 1 + J'_1$, θ'_e with $W''.\theta_1$, H_1 with H'_1 (from FB-B1) and J_2 with J'_1 we get

$$\begin{aligned} \exists \theta'_1 \sqsupseteq W''.\theta_1.(m_{u1} + 1, H_1'') \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in [\tau' \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H_1''(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1)/\text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_i \sigma) \quad (\text{FB-B6}) \end{aligned}$$

Since we know that $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \downarrow v'_2$. Say this reduction happens in t steps. Therefore $\exists t_1 < t < k \leq n$ s.t $(e_l)\gamma \downarrow_2 \cup \{x \mapsto v'_2\} \downarrow_{t_1} v_{t_2}$ and similarly $\exists t_2 < t - t_1 < k - t_1$ s.t $(H, v_{t_2})\gamma \downarrow_2 \downarrow_{t_2}^f (H_2'', v''_2)$

Again since $\text{bind}(e_l, x.e_b)\gamma \downarrow_2 \downarrow_t v'_2$ therefore $\exists J_2 < t - t_1 - t_2 < k - t_1 - t_2$ s.t $(e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{J_2} v'_2$. Similarly $\exists J'_2 < t - t_1 - t_2 - J_2 < k - t_1 - t_2 - J_2$ s.t $(H'_2, v'_2) \downarrow_{J'_2}^f -$

Instantiating the second conjunct of (FB-B3) with $m_{u2} + 1 + J_2 + J'_2$ we get $(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, v''_2) \in [\tau \sigma]_V$

Again since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 2.53 we know that $\forall m. (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating m with $m_{u2} + 1 + J_2 + J'_2$ we get $(W.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V$

From Lemma 2.47 we know that

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, \gamma \downarrow_2) \in [\Gamma]_V \quad (\text{FB-B7})$$

Now we can apply Theorem 2.51 to get

$$(W''.\theta_2, m_{u2} + 1 + J_2 + J'_2, (e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_E$$

This means from Definition 2.36 we get

$$\forall c_2 < (m_{u2} + 1 + J_2 + J'_2). (e_b)\gamma \downarrow_2 \cup \{x \mapsto v''_2\} \downarrow_{c_2} v_{o2} \implies (W''.\theta_2, m_{u2} + 1 + J_2 - c_2, v_{o2}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_V \quad (\text{FB-B8})$$

Instantiating c_2 with J_2 in (FB-B8) we get

$$(W''.\theta_2, m_{u2} + 1 + J'_2, v_{o2}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_V$$

From Definition 2.35 we have

$$\forall K \leq (m_{u2} + 1 + J'_2), \theta'_e \sqsupseteq W''.\theta_2, H_2, J_3. (K, H_2) \triangleright \theta'_e \wedge (H_2, v_{o2}) \downarrow_{J_3}^f (H_2'', v'_2) \wedge J_3 < K \implies$$

$$\begin{aligned} \exists \theta'_2 \sqsupseteq \theta'_e. (K - J_3, H_2'') \triangleright \theta'_2 \wedge (\theta'_2, K - J_3, v'_2) \in [\tau' \sigma]_V \wedge \\ (\forall a. H_2(a) \neq H_2''(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2)/\text{dom}(\theta'_e). \theta'_2(a) \searrow \ell \sigma) \end{aligned}$$

Instantiating K with $m_{u2} + 1 + J'_2$, θ'_e with $W''.\theta_2$, H_2 with H'_2 (from FB-B1) and J_3 with J'_2 , we get

$$\begin{aligned} \exists \theta'_2 \sqsupseteq W''.\theta_2.(m_{u2} + 1, H_2'') \triangleright \theta'_2 \wedge (\theta'_2, m_{u2} + 1, v'_2) \in [\tau' \sigma]_V \wedge \\ (\forall a. H_2(a) \neq H_2''(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge \ell \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_2)/\text{dom}(\theta'_e). \theta'_2(a) \searrow \ell \sigma) \quad (\text{FB-B9}) \end{aligned}$$

In order to prove (FB-B01) we chose W' as W_n where W_n is defined as follows: $W_n.\theta_1 = \theta'_1$ (From (FB-B6))

$$W_n.\theta_2 = \theta'_2 \text{ (From (FB-B9))}$$

$$W_n.\hat{\beta} = W''.\hat{\beta} \text{ (From (FB-B1))}$$

It suffices to prove

- $(k - j, H_1'', H_2'') \triangleright W_n$:

From Definition 2.38 we need to prove the following

$$- \text{dom}(W_n.\theta_1) \subseteq \text{dom}(H_1'') \wedge \text{dom}(W_n.\theta_2) \subseteq \text{dom}(H_2'')$$

From (FB-B6) we know that $(m_{u1}+1, H_1'') \triangleright \theta'_1$ therefore from Definition 2.37 we know that $\text{dom}(W_n.\theta_1) \subseteq \text{dom}(H_1'')$

Similarly from (FB-B9) we know that $(m_{u2} + 1, H_2'') \triangleright \theta'_2$ therefore from Definition 2.37 we know that $\text{dom}(W_n.\theta_2) \subseteq \text{dom}(H_2'')$

$$- (W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2)):$$

Since from (FB-B1) we know that $(k - f - J, H_1', H_2') \triangleright W''$ therefore from Definition 2.38 we know that $(W''.\hat{\beta}) \subseteq (\text{dom}(W''.\theta_1) \times \text{dom}(W''.\theta_2))$

Since from (FB-B6) and (FB-B9) we know that $W''.\theta_1 \sqsubseteq W_n.\theta_1$ and $W''.\theta_2 \sqsubseteq W_n.\theta_2$

Therefore we get

$$(W_n.\hat{\beta}) \subseteq (\text{dom}(W_n.\theta_1) \times \text{dom}(W_n.\theta_2))$$

$$- \forall (a_1, a_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a_1) = W_n.\theta_2(a_2) \wedge (W_n, k-j-1, H_1''(a_1), H_2''(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A):$$

4 cases arise for each $(a_1, a_2) \in W_n.\hat{\beta}$

$$A. H_1'(a_1) = H_1''(a_1) \wedge H_2'(a_2) = H_2''(a_2):$$

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2):$$

We know from that $(k - f - J, H_1', H_2') \triangleright W''$

Therefore from Definition 2.38 we have

$$\forall (a'_1, a'_2) \in (W''.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$$

Since $W_n.\hat{\beta} = W''.\hat{\beta}$ by construction therefore

$$\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). W''.\theta_1(a'_1) = W''.\theta_2(a'_2)$$

From (FB-B6) and (FB-B9) we know that $W''.\theta_1 \sqsubseteq \theta'_1$ and $W''.\theta_2 \sqsubseteq \theta'_2$ respectively.

Therefore from Definition 2.30

$$\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). \theta'_1(a_1) = \theta'_2(a_2)$$

To prove:

$$(W_n, k - j - 1, H_1''(a_1), H_2''(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A:$$

From (FB-B1) we know that $(k - f - J, H_1', H_2') \triangleright W''$

This means from Definition 2.38 we know that

$$\forall (a_{i1}, a_{i2}) \in (W''.\hat{\beta}). W''.\theta_1(a_{i1}) = W''.\theta_2(a_{i2}) \wedge (W'', k - f - J - 1, H_1'(a_{i1}), H_2'(a_{i2})) \in \lceil W''.\theta_1(a_{i1}) \rceil_V^A$$

Instantiating with a_1 and a_2 and since $W'' \sqsubseteq W_n$ and $k - j - 1 < k - f - J - 1$ (since $j = f + J + J_1 + 1$ therefore from Lemma 2.46 we get

$$(W_n, k - j - 1, H'_1(a_1), H'_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

B. $H'_1(a_1) \neq H''_1(a_1) \wedge H'_2(a_2) \neq H''_2(a_2)$:

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2)$$

Same reasoning as in the previous case

To prove:

$$(W_n, k - j - 1, H''_1(a_1), H''_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

From (FB-B6) and (FB-B9) we know that

$$(\forall a. H'_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell')$$

$$(\forall a. H'_2(a) \neq H''_2(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W''.\theta_1(a_1) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell' \text{ and}$$

$$\exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell'$$

Since $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $\ell' \not\sqsubseteq \mathcal{A}$.

Also from (FB-B6) and (FB-B9), $(m_{u1}+1, H''_1) \triangleright \theta'_1$ and $(m_{u2}+1, H''_2) \triangleright \theta'_2$.

Therefore from Definition 2.37 we have

$$(\theta'_1, m_{u1}, H''_1(a_1)) \in [\theta'_1(a_1)]_V \text{ and}$$

$$(\theta'_2, m_{u2}, H''_2(a_1)) \in [\theta'_2(a_2)]_V$$

Since m_{u1} and m_{u2} are arbitrary indices therefore from Definition 2.33 we get

$$(W_n, k - j - 1, H''_1(a_1), H''_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

C. $H'_1(a_1) = H''_1(a_1) \wedge H'_2(a_2) \neq H''_2(a_2)$:

To prove:

$$W_n.\theta_1(a_1) = W_n.\theta_2(a_2)$$

Same reasoning as in the previous case

To prove:

$$(W_n, k - j - 1, H''_1(a_1), H''_2(a_2)) \in [W_n.\theta_1(a_1)]_V^A$$

From (FB-B9) we know that

$$(\forall a. H'_2(a) \neq H''_2(a) \implies \exists \ell'. W''.\theta_2(a) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell')$$

This means we have

$$\exists \ell'. W''.\theta_2(a_2) = \text{Labeled } \ell' \tau'' \wedge (\ell \sigma) \sqsubseteq \ell'$$

Since $\ell \sigma \not\sqsubseteq \mathcal{A}$. Therefore, $\ell' \not\sqsubseteq \mathcal{A}$.

Since from (FB-B1) we know that $(k - f - J, H'_1, H'_2) \triangleright^A W''$ that means from Definition 2.38 that $(W'', k - f - J - 1, H'_1(a_1), H'_2(a_2)) \in [W''.\theta_1(a_1)]_V^A$. Since $W''.\theta_1(a_1) = W''.\theta_2(a_2) = \text{Labeled } \ell' \tau''$ and since $\ell' \not\sqsubseteq \mathcal{A}$ therefore from Definition 2.33 and Definition 2.32 we know that

Therefore

$$\forall m. (W''.\theta_1, m, H'_1(a_1)) \in W''.\theta_1(a_1) \quad (\text{F})$$

Instantiating the (F) with m_{u1} and using Lemma 2.45 we get

$$(\theta'_1, m_{u1}, H'_1(a_1)) \in \theta'_1(a_1)$$

Since from (FB-B9) we know that $(m_{u2} + 1, H''_2) \triangleright \theta'_2$ therefore from Definition 2.37 we know that $(\theta'_2, m_{u2}, H''_2(a_2)) \in \theta'_2(a_2)$

Therefore from Definition 2.33 we get

$$(W', k - j - 1, H'_1(a_1), H''_2(a_2)) \in [\theta'_1(a_1)]_V^A$$

$$D. H'_1(a_1) \neq H''_1(a_1) \wedge H'_2(a_2) = H''_2(a_2):$$

Symmetric reasoning as in the previous case

$$- \forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V:$$

Case $i = 1$

Given some m we need to prove

$$\forall a_i \in \text{dom}(W_n.\theta_i). (W_n.\theta_i, m, H''_i(a_i)) \in \lfloor W_n.\theta_i(a_i) \rfloor_V$$

This further means that given some $a_1 \in \text{dom}(W_n.\theta_1)$ we need to show

$$(W_n.\theta_1, m, H''_1(a_1)) \in \lfloor W_n.\theta_1(a_1) \rfloor_V$$

Since $W_n.\theta_1 = \theta'_1$, it suffices to prove

$$(\theta'_1, m, H''_1(a_1)) \in \lfloor \theta'_1(a_1) \rfloor_V$$

Like before we apply Theorem 2.51 on $e_b \gamma \downarrow_1 \cup \{x \mapsto v'_1\}$ but this time at $m + 1 + J_1 + J'_1$ to get

$$\begin{aligned} \exists \theta'_1 \sqsupseteq W''.\theta_1.(m + 1, H''_1) \triangleright \theta'_1 \wedge (\theta'_1, m_{u1} + 1, v'_1) \in \lfloor \tau' \sigma \rfloor_V \wedge \\ (\forall a. H_1(a) \neq H''_1(a) \implies \exists \ell'. W''.\theta_1(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \\ (\forall a \in \text{dom}(\theta'_1) / \text{dom}(\theta'_e). \theta'_1(a) \searrow \ell_i \sigma) \end{aligned}$$

Since we have $(m + 1, H''_1) \triangleright \theta'_1$ therefore from Definition 2.37 we get the desired.

Case $i = 2$

Similar reasoning as in the $i = 1$ case

$$\bullet (W'.\theta_1, m_{u1}, v'_1) \in \lfloor \tau' \sigma \rfloor_V \wedge (W'.\theta_2, m_{u2}, v'_2) \in \lfloor \tau' \sigma \rfloor_V:$$

We get this from (FB-B6), (FB-B9) and Lemma 2.45 we get the desired

19. CG-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } (e') : \mathbb{C} \ell \ell (\text{ref } \ell' \tau)}$$

To prove: $(W, n, \text{new } (e') (\gamma \downarrow_1), \text{new } (e') (\gamma \downarrow_2)) \in \llbracket (\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma \rrbracket_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. \text{new } (e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{new } (e') \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in \llbracket (\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma \rrbracket_V^A \end{aligned}$$

This means that given some $i < n$ s.t $\text{new } (e') \gamma \downarrow_1 \downarrow_i v_{f1} \wedge \text{new } (e') \gamma \downarrow_2 \downarrow v'_{f1}$

From *cg-val* we know that $v_{f1} = \text{new } (e') \gamma \downarrow_1$, $v_{f2} = \text{new } (e') \gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, \text{new } (e') \gamma \downarrow_1, \text{new } (e') \gamma \downarrow_2) \in \llbracket (\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma \rrbracket_V^A$$

Let $v_1 = \text{new } (e')\gamma \downarrow_1$ and $v_2 = \text{new } (e')\gamma \downarrow_2$

From Definition 2.33 we are required to prove

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau)]_V \sigma \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove the following:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ & (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma): \end{aligned}$$

This means we are given some $k \leq n, W_e \sqsupseteq W, H_1, H_2$ s.t $(k, H_1, H_2) \triangleright W_e$

Also we are given some $v'_1, v'_2, j < k$ s.t $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma) \quad \text{(FB-R0)}$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} & \forall f < k. e' \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1} \implies \\ & (W_e, k - f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau \sigma]_V^A \end{aligned}$$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$ therefore $\exists f < j < k$ s.t $e' \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e' \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{Labeled } \ell' \tau \sigma]_V^A \quad \text{(FB-R1)}$$

In order to prove (FB-R0) we choose W' as W_n where

$$W_n. \theta_1 = W_e. \theta_1 \cup \{a_1 \mapsto (\text{Labeled } \ell' \tau) \sigma\}$$

$$W_n. \theta_2 = W_e. \theta_2 \cup \{a_2 \mapsto (\text{Labeled } \ell' \tau) \sigma\}$$

$$W_n. \hat{\beta} = W_e. \hat{\beta} \cup \{a_1, a_2\}$$

Now we need to prove:

$$\text{i. } (k - j, H'_1, H'_2) \triangleright W_n:$$

From Definition 2.38 it suffices to prove:

$$\text{dom}(W_n. \theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_n. \theta_2) \subseteq \text{dom}(H'_2) \wedge$$

$$(W_n. \hat{\beta}) \subseteq (\text{dom}(W_n. \theta_1) \times \text{dom}(W_n. \theta_2)) \wedge$$

$$\forall (a_1, a_2) \in (W_n. \hat{\beta}). (W_n. \theta_1(a_1) = W_n. \theta_2(a_2) \wedge$$

$$(W_n, (k - j) - 1, H'_1(a_1), H'_2(a_2)) \in [\text{Labeled } \ell' \tau \sigma]_V^A) \wedge$$

$$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_n. \theta_i). (W_n. \theta_i, m, H_i(a_i)) \in [W_n. \theta_i(a_i)]_V$$

This means we need to prove

- $dom(W_n.\theta_1) \subseteq dom(H'_1) \wedge dom(W_n.\theta_2) \subseteq dom(H'_2) \wedge (W_n.\hat{\beta}) \subseteq (dom(W_n.\theta_1) \times dom(W_n.\theta_2))$:

We know that $dom(W_n.\theta_1) = dom(W_e.\theta_1) \cup \{a_1\}$ and $dom(W_n.\theta_2) = dom(W_e.\theta_2) \cup \{a_2\}$

Also $dom(H'_1) = dom(H_1) \cup \{a_1\}$ and $dom(H'_2) = dom(H_2) \cup \{a_2\}$

Therefore from $(k, H_1, H_2) \triangleright W_e$ and from construction of W_n we get the desired.

- $\forall (a'_1, a'_2) \in (W_n.\hat{\beta}). (W_n.\theta_1(a'_1) = W_n.\theta_2(a'_2) \wedge (W_n, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \lceil W_n.\theta_1(a'_1) \rceil_V^A)$:

$\forall (a'_1, a'_2) \in (W_n.\hat{\beta})$.

- A. When $a'_1 = a_1$ and $a'_2 = a_2$:

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$$

Since from (FB-R1) we know that $(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \sigma \rceil_V^A$

And since from $cg - ref$ we know that $H'_1(a_1) = v_{h1}$, $H'_2(a_2) = v'_{h1}$ and $j = f + 1$ therefore from Lemma 2.46 we get

$$(W_n, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \lceil W_n.\theta_1(a_1) \rceil_V^A$$

- B. When $a'_1 = a_1$ and $a'_2 \neq a_2$: This case cannot arise

- C. When $a'_1 \neq a_1$ and $a'_2 = a_2$: This case cannot arise

- D. When $a'_1 \neq a_1$ and $a'_2 \neq a_2$:

Since $(k, H_1, H_2) \triangleright W_e$ therefore the desired is obtained directly from Definition 2.38

- $\forall i \in \{1, 2\}. \forall m. \forall a'_i \in dom(W_n.\theta_i). (W_n.\theta_i, m, H_i(a'_i)) \in \lfloor W_n.\theta_i(a'_i) \rfloor_V$:

When $i = 1$

Given some m

$\forall a'_1 \in dom(W_n.\theta_1)$.

- when $a'_1 = a_1$:

From construction

$$(W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$$

And from (FB-R1) we know that $(W_e, k - f, v_{h1}, v'_{h1}) \in \lceil \text{Labeled } \ell' \tau \sigma \rceil_V^A$

Therefore from Lemma 2.44 get the desired

- Otherwise:

Since $(k, H_1, H_2) \triangleright W_e$ therefore the desired is obtained directly from Definition 2.38

When $i = 2$

Similar reasoning as with $i = 1$

- ii. $ValEq(\mathcal{A}, W_n, k - j, \ell, v'_1, v'_2, (\text{ref } \ell' \tau) \sigma)$:

From $cg - ref$ we know that $v'_1 = a_1$ and $v'_2 = a_2$

2 cases arise:

- A. $\ell \sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 it suffices to prove that

$$(W_n, k - j, a_1, a_2) \in (\text{ref } \ell' \tau) \sigma$$

From Definition 2.33 it suffices to prove

$$(a_1, a_2) \in W_n.\hat{\beta} \wedge W_n.\theta_1(a_1) = W_n.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$$

This holds from construction of W_n

B. $\ell \not\sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 it suffices to prove that

$\forall m. (W_n.\theta_1, m, a_1) \in (\text{ref } \ell' \tau) \sigma$ and $(W_n.\theta_2, m, a_2) \in (\text{ref } \ell' \tau) \sigma$

From Definition 2.35 this means for any given m we need to prove that

$W_n.\theta_1(a_1) \in (\text{Labeled } \ell' \tau) \sigma$ and $W_n.\theta_2(a_2) \in (\text{Labeled } \ell' \tau) \sigma$

This holds from construction of W_n

- (b) $\forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right.$
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau) \sigma]_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_1):$

Case $l = 1$

Given some $k, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t. $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{ref } \ell' \tau) \sigma]_V \wedge$
 $(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$

Since $(W, n, \gamma) \in [\Gamma]_V^{\mathcal{A}}$ therefore from Lemma 2.53 we know that

$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$ and $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$(W.\theta_1, k, (\text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma]_E$

This means from Definition 2.36 we get

$\forall c < k. \text{ref } (e')\gamma \downarrow_1 \Downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma]_V$

This further means that given some $c < k$ s.t. $\text{ref } (e')\gamma \downarrow_1 \Downarrow_c v$. From $cg - val$ we know that $c = 0$ and $v = \text{ref } (e')\gamma \downarrow_1$

And we have $(W.\theta_1, k, \text{ref } (e')\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{ref } \ell' \tau)) \sigma]_V$

From Definition 2.35 we have

$\forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J. (K, H_1) \triangleright \theta'_e \wedge (H_1, \text{ref } (e')\gamma \downarrow_1) \Downarrow_J^f (H', v') \wedge J < K \implies$
 $\exists \theta' \sqsupseteq \theta'_e. (K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{ref } \ell' \tau) \sigma]_V \wedge$
 $(\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell_i \sigma)$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

20. CG-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e' : \text{ref } \ell \tau}{\Sigma; \Psi; \Gamma \vdash !e' : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)}$$

To prove: $(W, n, !e' (\gamma \downarrow_1), !e' (\gamma \downarrow_2)) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. !e' \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge !e' \gamma \downarrow_2 \Downarrow v'_{f1} &\implies \\ (W, n - i, v_{f1}, v'_{f1}) &\in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V^A \end{aligned}$$

This means that given some $i < n$ s.t $!e' \gamma \downarrow_1 \Downarrow_i v_{f1} \wedge !e' \gamma \downarrow_2 \Downarrow v'_{f1}$

From $cg - val$ we know that $v_{f1} = !e' \gamma \downarrow_1$, $v_{f2} = !e' \gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, !e' \gamma \downarrow_1, !e' \gamma \downarrow_2) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V^A$$

Let $v_1 = !e' \gamma \downarrow_1$ and $v_2 = !e' \gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} &(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma)) \wedge \\ &\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ &\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \lfloor (\text{Labeled } \ell \tau) \sigma \rfloor_V \wedge \\ &(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ &(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell' \sigma)) \end{aligned}$$

This means we need to prove:

$$\begin{aligned} \text{(a)} \quad &\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \\ &(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma): \end{aligned}$$

This means we are given is some $k \leq n$, $W_e \sqsupseteq W, H_1, H_2$ s.t $(k, H_1, H_2) \triangleright W_e$

Also given some $v'_1, v'_2, j < k$ s.t $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$

And we are required to prove:

$$\begin{aligned} &\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge ValEq(\mathcal{A}, W', k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma) \\ &(\text{FB-D0}) \end{aligned}$$

IH:

$$(W_e, k, e' (\gamma \downarrow_1), e' (\gamma \downarrow_2)) \in [(\text{ref } \ell \tau) \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} &\implies \\ (W_e, k - f, v_{h1}, v'_{h1}) &\in [(\text{ref } \ell \tau) \sigma]_V^A \end{aligned}$$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow^f (H'_2, v'_2)$ therefore $\exists f < j < k$ s.t $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [(\text{ref } \ell \tau) \sigma]_V^A \quad (\text{FB-D1})$$

In order to prove (FB-D0) we choose W' as W_e . Also from $cg - deref$ we know that $H'_1 = H_1$ and $H'_2 = H_2$. Also we know that $v_{h1} = a_1$ and $v'_{h1} = a_2$.

- $(k - j, H_1, H_2) \triangleright W_e$:

Since we know that $(k, H_1, H_2) \triangleright W_e$ therefore from Lemma 2.50 we get
 $(k - j, H_1, H_2) \triangleright W_e$

- $ValEq(\mathcal{A}, W_e, k - j, \ell' \sigma, v'_1, v'_2, (\text{Labeled } \ell \tau) \sigma)$:

From *cg - ref* we know that $v'_1 = H_1(a_1)$ and $v'_2 = H_2(a_2)$

2 cases arise:

- $\ell' \sigma \sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 it suffices to prove that

$$(W_e, k - j, v'_1, v'_2) \in (\text{Labeled } \ell \tau) \sigma$$

Since from (FB-D1) we know that $(W_e, k - f, a_1, a_2) \in [\text{ref } \ell \tau \sigma]_V^{\mathcal{A}}$

Therefore from Definition 2.33 we know that $(a_1, a_2) \in W_e.\hat{\beta} \wedge W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = \text{Labeled } \ell \tau \sigma$

And since we know that $(k, H_1, H_2) \triangleright W_e$ therefore from Definition we know that $(W_e, k, H_1(a_1), H_2(a_2)) \in [\text{Labeled } \ell \tau \sigma]_V^{\mathcal{A}}$

From Lemma 2.46 we get $(W_e, k - j, H_1(a_1), H_2(a_2)) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\mathcal{A}}$

- $\ell' \not\sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 it suffices to prove that

$$\forall m. (W_e.\theta_1, m, H_1(a_1)) \in (\text{Labeled } \ell \tau) \sigma \text{ and } (W_e.\theta_2, m, H_2(a_2)) \in (\text{Labeled } \ell \tau) \sigma$$

(FB-B2)

Since from (FB-D1) we know that $(W_e, k - f, a_1, a_2) \in [\text{ref } \ell \tau \sigma]_V^{\mathcal{A}}$

Therefore from Definition 2.33 we know that $(a_1, a_2) \in W_e.\hat{\beta} \wedge W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = \text{Labeled } \ell \tau \sigma$

And since we know that $(k, H_1, H_2) \triangleright W_e$ therefore from Definition we know that $(W_e, k, H_1(a_1), H_2(a_2)) \in [\text{Labeled } \ell \tau \sigma]_V^{\mathcal{A}}$

Finally from Lemma 2.44 we get (FB-B2)

$$(b) \forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \\ \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell' \sigma):$$

Case $l = 1$

Given some $k, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$

We need to prove

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge \\ (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell' \sigma)$$

Since $(W, n, \gamma) \in [\Gamma]_V^{\mathcal{A}}$ therefore from Lemma 2.53 we know that

$$\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V \text{ and } (W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$$(W.\theta_1, k, (!e' \gamma \downarrow_1) \in [(\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)) \sigma]_E$$

This means from Definition 2.36 we get

$$\forall c < k. !e'\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)) \sigma]_V$$

Instantiating c with 0 and from $cg - val$ we know that $v = !e'\gamma \downarrow_1$

$$\text{And we have } (W.\theta_1, k, !e'\gamma \downarrow_1) \in [(\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau)) \sigma]_V$$

From Definition 2.35 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell' \sigma) \end{aligned}$$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

21. CG-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_l : \text{ref } \ell' \tau \quad \Sigma; \Psi; \Gamma \vdash e_r : \text{Labeled } \ell' \tau \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_l := e_r : \mathbb{C} \ell \ell \text{ unit}}$$

To prove: $(W, n, (e_l := e_r) (\gamma \downarrow_1), (e_l := e_r) (\gamma \downarrow_2)) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_E^A$

This means from Definition 2.34 we need to prove:

$$\begin{aligned} \forall i < n. (e_l := e_r) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \downarrow v'_{f1} \implies \\ (W, n - i, v_{f1}, v'_{f1}) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_V^A \end{aligned}$$

This means that given some $i < n$ s.t $(e_l := e_r) \gamma \downarrow_1 \downarrow_i v_{f1} \wedge (e_l := e_r) \gamma \downarrow_2 \downarrow v'_{f1}$

From $cg - val$ we know that $v_{f1} = (e_l := e_r) \gamma \downarrow_1$, $v_{f2} = (e_l := e_r) \gamma \downarrow_2$ and $i = 0$

We are required to prove

$$(W, n, (e_l := e_r) \gamma \downarrow_1, (e_l := e_r) \gamma \downarrow_2) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_V^A$$

Let $e_1 = (e_l := e_r) \gamma \downarrow_1$ and $e_2 = (e_l := e_r) \gamma \downarrow_2$

From Definition 2.33 it suffices to prove

$$\begin{aligned} \left(\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2. \right. \\ (H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies \\ \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \text{unit}) \right) \wedge \\ \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\text{unit}]_V \wedge \right. \\ \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sqsubseteq \ell') \wedge \right. \\ \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell) \right) \end{aligned}$$

This means we need to prove:

- (a) $\forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2.$
 $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies$
 $\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \text{unit}):$

This means we are given some $k \leq n, W_e \sqsupseteq W, H_1, H_2$ s.t $(k, H_1, H_2) \triangleright W_e$

And finally given some $v'_1, v'_2, j < k$ s.t $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$

And we are required to prove:

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell, v'_1, v'_2, \text{unit})$$

(FB-A0)

IH1:

$$(W_e, k, e_l (\gamma \downarrow_1), e_l (\gamma \downarrow_2)) \in [\text{ref } \ell' \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall f < k. e_l \gamma \downarrow_1 \Downarrow_f v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1} \implies$$

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau \sigma]_V^A$$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$ therefore $\exists f < j < k$ s.t $e_l \gamma \downarrow_f \Downarrow_j v_{h1} \wedge e_l \gamma \downarrow_2 \Downarrow v'_{h1}$

This means we have

$$(W_e, k - f, v_{h1}, v'_{h1}) \in [\text{ref } \ell' \tau \sigma]_V^A \quad (\text{FB-A1})$$

IH2:

$$(W_e, k - f, e_r (\gamma \downarrow_1), e_r (\gamma \downarrow_2)) \in [\text{Labeled } \ell' \tau \sigma]_E^A$$

This means from Definition 2.34 we need to prove:

$$\forall s < k - f. e_r \gamma \downarrow_1 \Downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \Downarrow v'_{h2} \implies$$

$$(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau \sigma]_V^A$$

Since we know that $(H_1, v_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, v_2) \Downarrow_j^f (H'_2, v'_2)$ therefore $\exists s < j - f < k - f$ s.t $e_r \gamma \downarrow_1 \Downarrow_s v_{h2} \wedge e_r \gamma \downarrow_2 \Downarrow v'_{h2}$

This means we have

$$(W_e, k - f - s, v_{h2}, v'_{h2}) \in [\text{Labeled } \ell' \tau \sigma]_V^A \quad (\text{FB-A2})$$

In order to prove (FB-A0) we choose W' as W_e . Also from *cg - assign* we know that $H'_1 = H_1[v_{h1} \mapsto v_{h2}]$ and $H'_2 = H_2[v'_{h1} \mapsto v'_{h2}]$, and $j = f + s + 1$

We need to prove the following:

- i. $(k - j, H'_1, H'_2) \triangleright W_e$:

Say $v_{h1} = a_1$ and $v'_{h1} = a_2$

From Definition 2.38 it suffices to prove:

$$\text{dom}(W_e.\theta_1) \subseteq \text{dom}(H'_1) \wedge \text{dom}(W_e.\theta_2) \subseteq \text{dom}(H'_2) \wedge$$

$$(W_e.\hat{\beta}) \subseteq (\text{dom}(W_e.\theta_1) \times \text{dom}(W_e.\theta_2)) \wedge$$

$$\forall (a_1, a_2) \in (W_e.\hat{\beta}). (W_e.\theta_1(a_1) = W_e.\theta_2(a_2) \wedge$$

$$(W_e, (k - j) - 1, H'_1(a_1), H'_2(a_2)) \in [W_e.\theta_1(a_1)]_V^A) \wedge$$

$$\forall i \in \{1, 2\}. \forall m. \forall a_i \in \text{dom}(W_e.\theta_i). (W_e.\theta_i, m, H_i(a_i)) \in [W_e.\theta_i(a_i)]_V$$

This means we need to prove

- $dom(W_e.\theta_1) \subseteq dom(H'_1) \wedge dom(W_e.\theta_2) \subseteq dom(H'_2) \wedge (W_e.\hat{\beta}) \subseteq (dom(W_e.\theta_1) \times dom(W_e.\theta_2))$:
Since $dom(H_1) = dom(H'_1)$ and $dom(H_2) = dom(H'_2)$, and also we know that $(k, H_1, H_2) \triangleright W_e$. Therefore we obtain the desired directly from Definition 2.38
 - $\forall (a'_1, a'_2) \in (W_e.\hat{\beta}). (W_e.\theta_1(a'_1) = W_e.\theta_2(a'_2) \wedge (W_e, k - j - 1, H'_1(a'_1), H'_2(a'_2)) \in \llbracket W_e.\theta_1(a'_1) \rrbracket_V^A)$:
 $\forall (a'_1, a'_2) \in (W_e.\hat{\beta})$.
 - A. When $a'_1 = a_1$ and $a'_2 = a_2$:
From (FB-A1) and from Definition 2.33 we get
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$
Since from (FB-A2) we know that $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \llbracket \text{Labeled } \ell' \tau \sigma \rrbracket_V^A$
And since from *cg - assign* we know that $H'_1(a_1) = v_{h2}, H'_2(a_2) = v'_{h2}$ and $j = f + s + 1$ threfore from Lemma 2.46 we get
 $(W_e, k - j - 1, H'_1(a_1), H'_2(a_2)) \in \llbracket W_e.\theta_1(a_1) \rrbracket_V^A$
 - B. When $a'_1 = a_1$ and $a'_2 \neq a_2$: This case cannot arise
 - C. When $a'_1 \neq a_1$ and $a'_2 = a_2$: This case cannot arise
 - D. When $a'_1 \neq a_1$ and $a'_2 \neq a_2$:
Since $(k, H_1, H_2) \triangleright W_e$ therefore the desired is obtained directly from Definition 2.38
 - $\forall i \in \{1, 2\}. \forall m. \forall a'_i \in dom(W_e.\theta_i). (W_e.\theta_i, m, H_i(a'_i)) \in \llbracket W_e.\theta_i(a'_i) \rrbracket_V$:
When $i = 1$
Given some m
 $\forall a'_1 \in dom(W_e.\theta_1)$.
 - when $a'_1 = a_1$:
From (FB-A1) and from Definition 2.33 we get
 $(W_e.\theta_1(a_1) = W_e.\theta_2(a_2) = (\text{Labeled } \ell' \tau) \sigma$
Since from (FB-A2) we know that $(W_e, k - f - s, v_{h2}, v'_{h2}) \in \llbracket \text{Labeled } \ell' \tau \sigma \rrbracket_V^A$
Therefore from Lemma 2.44 get the desired
 - Otherwise:
Since $(k, H_1, H_2) \triangleright W_e$ therefore the desired is obtained directly from Definition 2.38
 - When $i = 2$
Similar reasoning as with $i = 1$
- ii. $ValEq(\mathcal{A}, W_e, k - j, \ell, (), (), \text{unit})$:
Holds directly from Definition 2.32 and Definition 2.33
- (b) $\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket \text{unit} \rrbracket_V \wedge (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell \sigma \sqsubseteq \ell') \wedge (\forall a \in dom(\theta') \setminus dom(\theta_e). \theta'(a) \searrow \ell \sigma)$
- Case $l = 1$
Given some $k, \theta_e \sqsupseteq W.\theta_l, H, j$ s.t $(k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k$
- We need to prove
 $\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in \llbracket (\text{unit}) \sigma \rrbracket_V \wedge (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell \sigma \sqsubseteq \ell'') \wedge (\forall a \in dom(\theta') \setminus dom(\theta_e). \theta'(a) \searrow \ell \sigma)$

Since $(W, n, \gamma) \in [\Gamma]_V^A$ therefore from Lemma 2.53 we know that
 $\forall m. (W.\theta_1, m, \gamma \downarrow_1) \in [\Gamma]_V$ and $(W.\theta_2, m, \gamma \downarrow_2) \in [\Gamma]_V$

Instantiating m with k we get $(W.\theta_1, k, \gamma \downarrow_1) \in [\Gamma]_V$

Now we can apply Theorem 2.51 to get

$$(W.\theta_1, k, ((e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{unit})) \sigma]_E$$

This means from Definition 2.36 we get

$$\forall c < k. (e_l := e_r)\gamma \downarrow_1 \downarrow_c v \implies (W.\theta_1, k - c, v) \in [(\mathbb{C} \ell \ell (\text{unit})) \sigma]_V$$

Instantiating c with 0 and from $cg - val$ we know that $v = (e_l := e_r)\gamma \downarrow_1$

And we have $(W.\theta_1, k, (e_l := e_r)\gamma \downarrow_1) \in [(\mathbb{C} \ell \ell (\text{unit})) \sigma]_V$

From Definition 2.35 we have

$$\begin{aligned} \forall K \leq k, \theta'_e \sqsupseteq W.\theta_1, H_1, J.(K, H_1) \triangleright \theta'_e \wedge (H_1, v) \Downarrow_J^f (H', v') \wedge J < K \implies \\ \exists \theta' \sqsupseteq \theta'_e.(K - J, H') \triangleright \theta' \wedge (\theta', K - J, v') \in [(\text{Labeled } \ell \tau) \sigma]_V \wedge \\ (\forall a. H_1(a) \neq H'(a) \implies \exists \ell'. \theta'_e(a) = \text{Labeled } \ell'' \tau'' \wedge \ell' \sigma \sqsubseteq \ell'') \wedge \\ (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta'_e). \theta'(a) \searrow \ell' \sigma) \end{aligned}$$

Instantiating K with k , θ'_e with θ_e , H_1 with H and J with j we get the desired

Case $l = 2$

Symmetric reasoning as in the $l = 1$ case above

□

Lemma 2.55 (CG: Equivalence of values). $\forall \mathcal{A}, W, W', \ell, \ell', v_1, v_2, \tau, i, j.$

$$\begin{aligned} \text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau) \wedge j < i \wedge \ell \sqsubseteq \ell' \wedge W \sqsubseteq W' \implies \\ \text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau) \end{aligned}$$

Proof. Given that $\text{ValEq}(\mathcal{A}, W, \ell, i, v_1, v_2, \tau)$. From Definition 2.32 two cases arise

1. $\ell \sqsubseteq \mathcal{A}$:

In this case we know that $(W, i, v_1, v_2) \in [\tau]_V^A$

2 cases arise

(a) $\ell' \sqsubseteq \mathcal{A}$:

Since $(W, i, v_1, v_2) \in [\tau]_V^A$ therefore from Lemma 2.46 we know that $(W', j, v_1, v_2) \in [\tau]_V^A$

And thus from Definition 2.32 we know that $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

(b) $\ell' \not\sqsubseteq \mathcal{A}$:

Since $(W, i, v_1, v_2) \in [\tau]_V^A$ therefore from Lemma 2.44 we know that $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 2.45 we know that $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 2.32 we know that $\text{ValEq}(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

2. $\ell \not\sqsubseteq \mathcal{A}$:

Given is $\ell \sqsubseteq \ell' \not\sqsubseteq \mathcal{A}$

In this case we know that $\forall i \in \{1, 2\}. \forall m. (W.\theta_i, m, v_i) \in [\tau]_V$

And from Lemma 2.45 we know that $\forall i \in \{1, 2\}. \forall m. (W'.\theta_i, m, v_i) \in [\tau]_V$

Hence from Definition 2.32 we know that $ValEq(\mathcal{A}, W', \ell', j, v_1, v_2, \tau)$

□

Lemma 2.56 (CG: Subtyping binary). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \tau, \tau'$.

1. $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^A \subseteq [(\tau' \sigma)]_V^A$
2. $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^A \subseteq [(\tau' \sigma)]_E^A$

Proof. Proof of statement (1)

Proof by induction on the $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove: $[((\tau_1 \rightarrow \tau_2) \sigma)]_V^A \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^A$

IH1: $[(\tau'_1 \sigma)]_V^A \subseteq [(\tau_1 \sigma)]_V^A$ (Statement 1)

$[(\tau_2 \sigma)]_E^A \subseteq [(\tau'_2 \sigma)]_E^A$ (Sub-A0 From Statement 2)

It suffices to prove:

$\forall (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^A. (W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^A$

This means that given: $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^A$

And it suffices to prove: $(W, n, \lambda x.e_1, \lambda x.e_2) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^A$

From Definition 2.33 we are given:

$$\begin{aligned} \forall W' \sqsupseteq W, j < n, v_1, v_2. ((W', j, v_1, v_2) \in [\tau_1 \sigma]_V^A \implies \\ (W', j, e_1[v_1/x], e_2[v_2/x]) \in [\tau_2 \sigma]_E^A) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_1[v_1/x]) \in [\tau_2 \sigma]_E) \wedge \\ \forall \theta_l \sqsupseteq W.\theta_l, j, v_c. ((\theta_l, j, v_c) \in [\tau_1 \sigma]_V \implies (\theta_l, j, e_2[v_c/x]) \in [\tau_2 \sigma]_E) \end{aligned} \quad (\text{Sub-A1})$$

Again from Definition 2.33 we are required to prove:

$$\begin{aligned} \forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in \\ [\tau'_2 \sigma]_E^A) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_l, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E) \wedge \\ \forall \theta'_l \sqsupseteq W.\theta_l, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E) \end{aligned}$$

This means need to prove:

(a) $\forall W'' \sqsupseteq W, k < n, v'_1, v'_2. ((W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A)$:

Given: $W'' \sqsupseteq W, k < n$ and v'_1, v'_2 . We are also given $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$

To prove: $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

Instantiating the first conjunct of Sub-A1 with W'', k, v'_1 and v'_2 we get

$$((W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A \implies (W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A) \quad (153)$$

Since $(W'', k, v'_1, v'_2) \in [\tau'_1 \sigma]_V^A$ therefore from IH1 we know that $(W'', k, v'_1, v'_2) \in [\tau_1 \sigma]_V^A$

Thus from Equation 153 we get $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau_2 \sigma]_E^A$

Finally using (Sub-A0) we get $(W'', k, e_1[v'_1/x], e_2[v'_2/x]) \in [\tau'_2 \sigma]_E^A$

(b) $\forall \theta'_l \sqsupseteq W.\theta_1, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E)$:

Given: $\theta'_l \sqsupseteq W.\theta_1, k, v'_c$. We are also given $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$

To prove: $(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E$

Since we are given $(\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V$ and since $\tau'_1 \sigma <: \tau_1 \sigma$ therefore from Lemma 2.52 we get

$$(\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \quad (154)$$

Instantiating the second conjunct of Sub-A1 with θ'_l, k, v'_1 and v'_2 we get

$$((\theta'_l, k, v'_c) \in [\tau_1 \sigma]_V \implies (\theta'_l, e_1[v'_c/x]) \in [\tau_2 \sigma]_E) \quad (155)$$

Therefore from Equation 154 and 155 we get $(\theta'_l, k, e_1[v'_c/x]) \in [\tau_2 \sigma]_E$

Since $\tau_2 \sigma <: \tau'_2 \sigma$ therefore from Lemma 2.52 we get

$$(\theta'_l, k, e_1[v'_c/x]) \in [\tau'_2 \sigma]_E$$

(c) $\forall \theta'_l \sqsupseteq W.\theta_2, k, v'_c. ((\theta'_l, k, v'_c) \in [\tau'_1 \sigma]_V \implies (\theta'_l, k, e_2[v'_c/x]) \in [\tau'_2 \sigma]_E)$:

Similar reasoning as in the previous case

2. CGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove: $[((\tau_1 \times \tau_2) \sigma)]_V^A \subseteq [((\tau'_1 \times \tau'_2) \sigma)]_V^A$

IH1: $[(\tau_1 \sigma)]_V^A \subseteq [(\tau'_1 \sigma)]_V^A$ (Statement (1))

IH2: $[(\tau_2 \sigma)]_V^A \subseteq [(\tau'_2 \sigma)]_V^A$ (Statement (1))

It suffices to prove: $\forall (W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^A. (W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_V^A$

This means that given: $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau_1 \times \tau_2) \sigma)]_V^A$

Therefore from Definition 2.33 we are given:

$$(W, n, v_1, v'_1) \in [\tau_1 \sigma]_{\mathcal{V}}^A \wedge (W, n, v_2, v'_2) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (156)$$

And it suffices to prove: $(W, n, (v_1, v_2), (v'_1, v'_2)) \in [((\tau'_1 \times \tau'_2) \sigma)]_{\mathcal{V}}^A$

Again from Definition 2.33, it suffices to prove:

$$(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_{\mathcal{V}}^A \wedge (W, n, v_2, v'_2) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

Since from Equation 156 we know that $(W, n, v_1, v'_1) \in [\tau_1 \sigma]_{\mathcal{V}}^A$ therefore from IH1 we have $(W, n, v_1, v'_1) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$

Similarly since $(W, n, v_2, v'_2) \in [\tau_2 \sigma]_{\mathcal{V}}^A$ from Equation 156 therefore from IH2 we have $(W, n, v_2, v'_2) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$

3. CGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove: $[((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A \subseteq [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

IH1: $[(\tau_1 \sigma)]_{\mathcal{V}}^A \subseteq [(\tau'_1 \sigma)]_{\mathcal{V}}^A$ (Statement (1))

IH2: $[(\tau_2 \sigma)]_{\mathcal{V}}^A \subseteq [(\tau'_2 \sigma)]_{\mathcal{V}}^A$ (Statement (1))

It suffices to prove: $\forall (W, n, v_{s1}, v_{s2}) \in [((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A. (W, n, v_{s1}, v_{s2}) \in [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

This means that given: $(W, n, v_{s1}, v_{s2}) \in [((\tau_1 + \tau_2) \sigma)]_{\mathcal{V}}^A$

And it suffices to prove: $(W, n, v_{s1}, v_{s2}) \in [((\tau'_1 + \tau'_2) \sigma)]_{\mathcal{V}}^A$

2 cases arise

(a) $v_{s1} = \text{inl } v_{i1}$ and $v_{s1} = \text{inl } v_{i2}$:

From Definition 2.33 we are given:

$$(W, n, v_{i1}, v_{i2}) \in [\tau_1 \sigma]_{\mathcal{V}}^A \quad (157)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$$

From Equation 157 and IH1 we know that

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_1 \sigma]_{\mathcal{V}}^A$$

(b) $v_s = \text{inr } v_{i1}$ and $v_{s2} = \text{inr } v_{i2}$:

From Definition 2.33 we are given:

$$(W, n, v_{i1}, v_{i2}) \in [\tau_2 \sigma]_{\mathcal{V}}^A \quad (158)$$

And we are required to prove that:

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

From Equation 158 and IH2 we know that

$$(W, n, v_{i1}, v_{i2}) \in [\tau'_2 \sigma]_{\mathcal{V}}^A$$

4. CGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove: $[(\forall \alpha. \tau_1) \sigma]_{\mathcal{V}}^A \subseteq [(\forall \alpha. \tau_2) \sigma]_{\mathcal{V}}^A$

$\forall \sigma. [(\tau_1 \sigma)]_E^A \subseteq [(\tau_2 \sigma)]_E^A$ (Sub-F2, From Statement (2))

It suffices to prove: $\forall (W, n, \Lambda e_1, \Lambda e_2) \in [(\forall \alpha. \tau_1) \sigma]_{\mathcal{V}}^A$.

$(W, n, \Lambda e_1, \Lambda e_2) \in [(\forall \alpha. \tau_2) \sigma]_{\mathcal{V}}^A$

This means that given: $(W, n, \Lambda e_1, \Lambda e_2) \in [(\forall \alpha. (\tau_1)) \sigma]_{\mathcal{V}}^A$

Therefore from Definition 2.33 we are given:

$\forall W' \sqsupseteq W, n' < n, \ell' \in \mathcal{L}. ((W', n', e_1, e_2) \in [\tau_1[\ell'/\alpha] \sigma]_E^A) \wedge$
 $\forall \theta_l \sqsupseteq W. \theta_1, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_1) \in [\tau_1[\ell'/\alpha]]_E) \wedge$
 $\forall \theta_l \sqsupseteq W. \theta_2, j, \ell' \in \mathcal{L}. ((\theta_l, j, e_2) \in [\tau_1[\ell''/\alpha]]_E)$ (Sub-F1)

And it suffices to prove: $(W, n, \Lambda e_1, \Lambda e_2) \in [(\forall \alpha. \tau_2) \sigma]_{\mathcal{V}}^A$

Again from Definition 2.33, it suffices to prove:

$\forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A) \wedge$
 $\forall \theta'_l \sqsupseteq W. \theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E) \wedge$
 $\forall \theta'_l \sqsupseteq W. \theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E)$

This means we are required to show:

(a) $\forall W'' \sqsupseteq W, n'' < n, \ell'' \in \mathcal{L}. ((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$:

By instantiating the first conjunct of Sub-F1 with W'', n'' and ℓ'' we know that the following holds

$$((W'', n'', e_1, e_2) \in [\tau_1[\ell''/\alpha] \sigma]_E^A)$$

Therefore from Sub-F2 instantiated at $\sigma \cup \{\alpha \mapsto \ell''\}$

$$((W'', n'', e_1, e_2) \in [\tau_2[\ell''/\alpha] \sigma]_E^A)$$

(b) $\forall \theta'_l \sqsupseteq W. \theta_1, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha]]_E)$:

By instantiating the second conjunct of Sub-F1 with θ'_l and ℓ'' we know that the following holds

$$((\theta'_l, k, e_1) \in [\tau_1[\ell''/\alpha] \sigma]_E)$$

Since $\tau_1 \sigma \cup \{\alpha \mapsto \ell''\} <: \tau_2 \sigma \cup \{\alpha \mapsto \ell''\}$ therefore from Lemma 2.52 we know that

$$((\theta'_l, k, e_1) \in [\tau_2[\ell''/\alpha] \sigma]_E)$$

(c) $\forall \theta'_l \sqsupseteq W. \theta_2, k, \ell'' \in \mathcal{L}. ((\theta'_l, k, e_2) \in [\tau_2[\ell''/\alpha]]_E)$:

Similar reasoning as in the previous case

5. CGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \implies \tau_1 <: c_2 \implies \tau_2}$$

To prove: $\lceil ((c_1 \Rightarrow \tau_1) \sigma) \rceil_V^A \subseteq \lceil ((c_2 \Rightarrow \tau_2) \sigma) \rceil_V^A$

$\lceil (\tau_1 \sigma) \rceil_E^A \subseteq \lceil (\tau_2 \sigma) \rceil_E^A$ (Sub-C0, From Statement (2))

It suffices to prove: $\forall (W, n, \nu e_1, \nu e_2) \in \lceil ((c_1 \Rightarrow \tau_1) \sigma) \rceil_V^A. (W, n, \nu e_1, \nu e_2) \in \lceil ((c_2 \Rightarrow \tau_2) \sigma) \rceil_V^A$

This means that given: $(W, n, \nu e_1, \nu e_2) \in \lceil ((c_1 \Rightarrow \tau_1) \sigma) \rceil_V^A$

Therefore from Definition 2.33 we are given:

$\forall W' \sqsupseteq W, n' < n. \mathcal{L} \models c_1 \sigma \implies (W', n', e_1, e_2) \in \lceil \tau_1 \sigma \rceil_E^A \wedge$

$\forall \theta_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_1 \implies (\theta_l, k, e_1) \in \lfloor \tau_1 \sigma \rfloor_E \wedge$

$\forall \theta_l \sqsupseteq W.\theta_2, k. \mathcal{L} \models c_1 \implies (\theta_l, k, e_2) \in \lfloor \tau_1 \sigma \rfloor_E$ (Sub-C1)

And it suffices to prove: $(W, n, \nu e_1, \nu e_2) \in \lceil ((c_2 \Rightarrow \tau_2) \sigma) \rceil_V^A$

Again from Definition 2.33, it suffices to prove:

$\forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \implies (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A \wedge$

$\forall \theta'_l \sqsupseteq W.\theta_1, j. \mathcal{L} \models c_2 \implies (\theta'_l, j, e_1) \in \lfloor \tau_2 \sigma \rfloor_E \wedge$

$\forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \implies (\theta'_l, j, e_2) \in \lfloor \tau_2 \sigma \rfloor_E$

This means that we are required to show the following:

(a) $\forall W'' \sqsupseteq W, n'' < n. \mathcal{L} \models c_2 \sigma \implies (W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A$:

We are given $W'' \sqsupseteq W, n'' < n$ also we know that $\mathcal{L} \models c_2 \sigma$ and $c_2 \sigma \implies c_1 \sigma$ therefore we also know that $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the first conjunct of Sub-C1 with W'' and n'' we know that the following holds

$(W'', n'', e_1, e_2) \in \lceil \tau_1 \sigma \rceil_E^A$

Therefore from (Sub-C0) we get $(W'', n'', e_1, e_2) \in \lceil \tau_2 \sigma \rceil_E^A$

(b) $\forall \theta'_l \sqsupseteq W.\theta_1, k. \mathcal{L} \models c_2 \implies (\theta'_l, k, e_1) \in \lfloor \tau_2 \sigma \rfloor_E$:

We are given some $\theta'_l \sqsupseteq W.\theta_1, k$, also we know that $\mathcal{L} \models c_2 \sigma$ and $c_2 \sigma \implies c_1 \sigma$ therefore we also know that $\mathcal{L} \models c_1 \sigma$

Hence by instantiating the second conjunct of Sub-C1 with θ'_l we know that the following holds

$(\theta'_l, k, e_1) \in \lfloor \tau_1 \sigma \rfloor_E$

Since $\tau_1 \sigma <: \tau_2 \sigma$ therefore from Lemma 2.52 we get

$(\theta'_l, k, e_1) \in \lfloor \tau_2 \sigma \rfloor_E$

(c) $\forall \theta'_l \sqsupseteq W.\theta_2, j. \mathcal{L} \models c_2 \implies (\theta'_l, j, e_2) \in \lfloor \tau_2 \sigma \rfloor_E$:

Similar reasoning as in the previous case

6. CGsub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove: $\lceil ((\text{Labeled } \ell \tau) \sigma) \rceil_V^A \subseteq \lceil ((\text{Labeled } \ell' \tau') \sigma) \rceil_V^A$

IH: $\lceil (\tau \sigma) \rceil_V^A \subseteq \lceil (\tau' \sigma) \rceil_V^A$

It suffices to prove: $\forall (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell \tau) \sigma)]_V^A. (W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell' \tau') \sigma)]_V^A$

This means we are given $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell \tau) \sigma)]_V^A$

From Definition 2.33 it means we have $\text{ValEq}(\mathcal{A}, W, \ell \sigma, n, v_1, v_2, \tau \sigma)$ (Sub-L0)

and it suffices to prove $(W, n, \text{Lb}(v_1), \text{Lb}(v_2)) \in [((\text{Labeled } \ell' \tau') \sigma)]_V^A$

Again from Definition 2.33 it means we need to prove that

$\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau' \sigma)$

Since we have (Sub-L0) and $\ell \sigma \sqsubseteq \ell' \sigma$ therefore from Lemma 2.55 we have

$\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau \sigma)$

2 cases arise:

(a) $\ell' \sigma \sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 we know that $(W, n, v_1, v_2) \in [\tau \sigma]_V^A$

From IH we also know that $(W, n, v_1, v_2) \in [\tau' \sigma]_V^A$

And from Definition 2.33 we get $\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau' \sigma)$

(b) $\ell' \sigma \not\sqsubseteq \mathcal{A}$:

In this case from Definition 2.32 we know that $\forall j. (W.\theta_1, j, v_1) \in [\tau \sigma]_V$ and $(W.\theta_2, j, v_2) \in [\tau \sigma]_V$

Since $\tau \sigma <: \tau' \sigma$ therefore from Lemma 2.52 we get $(W.\theta_1, j, v_1) \in [\tau' \sigma]_V$ and $(W.\theta_2, j, v_2) \in [\tau' \sigma]_V$

And from Definition 2.33 we get $\text{ValEq}(\mathcal{A}, W, \ell' \sigma, n, \text{Lb}(v_1), \text{Lb}(v_2), \tau' \sigma)$

7. CGsub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau <: \mathbb{C} \ell'_i \ell'_o \tau'}$$

To prove: $[((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V^A \subseteq [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V^A$

IH: $[(\tau \sigma)]_V^A \subseteq [(\tau' \sigma)]_V^A$

It suffices to prove: $\forall (W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V^A. (W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V^A$

This means we are given $(W, n, e_1, e_2) \in [((\mathbb{C} \ell_i \ell_o \tau) \sigma)]_V^A$

From Definition 2.33 it means we have

$(\forall k \leq n, W_e \sqsupseteq W, H_1, H_2.(k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j.$

$(H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow_j^f (H'_2, v'_2) \wedge j < k \implies$

$\exists W' \sqsupseteq W_e.(k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau \sigma)) \wedge$

$\forall l \in \{1, 2\}. (\forall k, \theta_e \sqsupseteq W.\theta_l, H, j.(k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies$

$\exists \theta' \sqsupseteq \theta_e.(k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau \sigma]_V \wedge$

$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge$

$(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$ (Sub-CG0)

And we need to prove

$$(W, n, e_1, e_2) \in [((\mathbb{C} \ell'_i \ell'_o \tau') \sigma)]_V^A$$

Again from Definition 2.33 it means we need to prove

$$\begin{aligned} & \left(\forall k \leq n, W_e \sqsupseteq W, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \right. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \left. \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o \sigma, v'_1, v'_2, \tau' \sigma) \right) \wedge \\ & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau' \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma) \right) \end{aligned}$$

It means we need to prove:

$$\begin{aligned} \text{(a)} \quad & \forall k \leq n, W_e \sqsupseteq W. \forall H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \forall v'_1, v'_2, j. \\ & (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2) \wedge j < k \implies \\ & \exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau' \sigma): \end{aligned}$$

This means we are given $k \leq n, W_e \sqsupseteq W, H_1, H_2, v'_1, v'_2, j < k$ s.t

$$(k, H_1, H_2) \triangleright W_e, (H_1, e_1) \Downarrow_j^f (H'_1, v'_1) \wedge (H_2, e_2) \Downarrow^f (H'_2, v'_2)$$

And we need to prove

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell'_o \sigma, v'_1, v'_2, \tau' \sigma)$$

Instantiating the first conjunct of (Sub-CG0) to get

$$\exists W' \sqsupseteq W_e. (k - j, H'_1, H'_2) \triangleright W' \wedge \text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau \sigma) \quad (\text{Sub-CG1})$$

Since from (Sub-CG1) $\text{ValEq}(\mathcal{A}, W', k - j, \ell_o \sigma, v'_1, v'_2, \tau \sigma)$

Therefore from Lemma 2.55 we get $\text{ValEq}(\mathcal{A}, W', k - j, \ell'_o \sigma, v'_1, v'_2, \tau \sigma)$

$$\begin{aligned} \text{(b)} \quad & \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq \theta, H, j. (k, H) \triangleright \theta_e \wedge (H, e_l) \Downarrow_j^f (H', v'_l) \wedge j < k \implies \right. \\ & \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_l) \in [\tau' \sigma]_V \wedge \right. \\ & \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell') \wedge \right. \\ & \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma) \right): \end{aligned}$$

Case $l = 1$

Here we are given $k, \theta_e \sqsupseteq \theta, H, j < k$ s.t $(k, H) \triangleright \theta_e \wedge (H, e_1) \Downarrow_j^f (H', v'_1)$

And we need to prove

$$\text{i. } \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in [\tau' \sigma]_V:$$

Instantiating the second conjunct of (Sub-CG0) with the given k, θ_e, H, j to get

$$\exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v'_1) \in [\tau \sigma]_V$$

Since $\tau \sigma <: \tau' \sigma$ therefore from Lemma 2.52 we get $(\theta', k - j, v'_1) \in [\tau' \sigma]_V$

$$\text{ii. } (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell'):$$

Instantiating the second conjunct of (Sub-CG0) with the given v, i, k, θ_e, H, j to get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell_i \sigma \sqsubseteq \ell')$$

Since $\ell'_i \sigma \sqsubseteq \ell_i \sigma$ therefore we also get

$$(\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau'' \wedge \ell'_i \sigma \sqsubseteq \ell')$$

- iii. $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma)$:
Instantiating the second conjunct of (Sub-CG0) with the given v, i, k, θ_e, H, j to get
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell_i \sigma)$
Since $\ell'_i \sigma \sqsubseteq \ell_i \sigma$ therefore we also get
 $(\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \ell'_i \sigma)$

Case $l = 2$

Symmetric reasoning as in the previous $l = 1$ case

8. CGsub-base:

Trivial

Proof of Statement (2)

It suffice to prove that

$$\forall (W, n, e_1, e_2) \in [\tau \sigma]_E^A. (W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$$

This means given $(W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$

From Definition 2.34 it means we have

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A \quad (\text{Sub-E0})$$

And it suffices to prove $(W, n, e_1, e_2) \in [(\tau' \sigma)]_E^A$

Again from Definition 2.34 it means we need to prove

$$\forall i < n. e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2 \implies (W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$$

This means that given $i < n$ s.t $e_1 \Downarrow_i v_1 \wedge e_2 \Downarrow v_2$ we need to prove $(W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$

Instantiating (Sub-E0) with the given i we get $(W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$

From Statement (1) we get $(W, n - i, v_1, v_2) \in [\tau' \sigma]_V^A$ □

Theorem 2.57 (CG: NI). $\forall v_1, v_2, e, \tau, n.$

$$(\emptyset, n, v_1, v_2) \in [\text{Labeled } \top \mathbf{b}]_V^\perp \wedge$$

$$(\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \mathbf{b}]_E^\perp \wedge$$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge n' < n \wedge (\emptyset, e[v_2/x]) \Downarrow_{n'}^f (-, v'_2) \implies v'_1 = v'_2$$

Proof. Given some $v_1, v_2, e, \tau, H_1, H_2, W, n, j < n$ s.t

$$(\emptyset, n, v_1, v_2) \in [\text{Labeled } \top \mathbf{b}]_V^\perp \wedge$$

$$(\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \mathbf{b}]_E^\perp \wedge$$

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v'_1) \wedge n' < n \wedge (\emptyset, e[v_2/x]) \Downarrow_{n'}^f (-, v'_2)$$

We need to prove

$$v'_1 = v'_2$$

Since we are given $(\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \mathbf{b}]_E^\perp$

Therefore from Definition 2.34 we know that

$$\forall i < n. e_1[v_1/x] \Downarrow_i v_{11} \wedge e_2 \Downarrow v_{22} \implies (\emptyset, n - i, v_{11}, v_{22}) \in [\mathbb{C} \perp \perp \mathbf{b}]_V^\perp$$

From CG-val we know that $i = 0, v_{11} = e[v_1/x]$ and $v_{22} = e[v_2/x]$

Therefore we have

$$(\emptyset, n, e[v_1/x], e[v_2/x]) \in [\mathbb{C} \perp \perp \mathbf{b}]_V^\perp$$

From Definition 2.35 we have

$$\begin{aligned}
& \left(\forall k \leq n, W_e \sqsupseteq \emptyset, H_1, H_2. (k, H_1, H_2) \triangleright W_e \wedge \right. \\
& \forall v_1'', v_2'', j. (H_1, e[v_1/x]) \Downarrow_j^f (H_1', v_1'') \wedge (H_2, e[v_2/x]) \Downarrow^f (H_2', v_2'') \wedge j < k \implies \\
& \left. \exists W' \sqsupseteq W_e. (k - j, H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', k - j, \perp, v_1', v_2', \mathbf{b}) \right) \wedge \\
& \forall l \in \{1, 2\}. \left(\forall k, \theta_e \sqsupseteq W. \theta_l, H, j. (k, H) \triangleright \theta_e \wedge (H, v_l) \Downarrow_j^f (H', v_l') \wedge j < k \implies \right. \\
& \left. \exists \theta' \sqsupseteq \theta_e. (k - j, H') \triangleright \theta' \wedge (\theta', k - j, v_l') \in [\mathbf{b}]_V \wedge \right. \\
& \left. (\forall a. H(a) \neq H'(a) \implies \exists \ell'. \theta_e(a) = \text{Labeled } \ell' \tau' \wedge \perp \sqsubseteq \ell') \wedge \right. \\
& \left. (\forall a \in \text{dom}(\theta') \setminus \text{dom}(\theta_e). \theta'(a) \searrow \perp) \right)
\end{aligned}$$

Instantiating the first conjunct with $n, \emptyset, \emptyset, \emptyset$.

Since we know that

$$(\emptyset, e[v_1/x]) \Downarrow_{n'}^f (-, v_1') \wedge n' < n \wedge (\emptyset, e[v_2/x]) \Downarrow_{n'}^f (-, v_2')$$

Therefore we instantiate v_1'' with v_1' , v_2'' with v_2' , j with n' to get

$$\exists W' \sqsupseteq \emptyset. (n - n', H_1', H_2') \triangleright W' \wedge \text{ValEq}(\perp, W', k - j, \perp, v_1', v_2', \mathbf{b})$$

From Definition 2.32 and Definition 2.35 we get $v_1' = v_2'$

□

2.3 CG to FG translation

2.3.1 Type directed translation from CG to FG

CG types are translated into FG types by the following definition of $\llbracket \cdot \rrbracket$

$$\begin{array}{ll}
\llbracket \mathbf{b} \rrbracket = \mathbf{b}^\perp & \llbracket \text{ref } \ell \ \tau \rrbracket = (\text{ref } (\llbracket \tau \rrbracket) + \text{unit})^{\ell}{}^\perp \\
\llbracket \tau_1 \rightarrow \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket \xrightarrow{\top} \llbracket \tau_2 \rrbracket)^\perp & \llbracket \mathbb{C} \ \ell_i \ \ell_o \ \tau \rrbracket = (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau \rrbracket) + \text{unit})^{\ell_o}{}^\perp \\
\llbracket \tau_1 \times \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp & \llbracket c \Rightarrow \tau \rrbracket = (c \xrightarrow{\top} \llbracket \tau \rrbracket)^\perp \\
\llbracket \tau_1 + \tau_2 \rrbracket = (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp & \llbracket \forall \alpha. \tau \rrbracket = (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp \\
\llbracket \text{Labeled } \ell \ \tau \rrbracket = (\llbracket \tau \rrbracket + \text{unit})^\ell &
\end{array}$$

The translation judgment for expressions is of the form $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e_C : \tau_C \rightsquigarrow e_F}$. Its rules are shown below.

$$\begin{array}{c}
\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau \rightsquigarrow x} \text{ var} \\
\frac{\Sigma; \Psi; \Gamma, x : \tau \vdash e : \tau' \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \lambda x. e : \tau \rightarrow \tau' \rightsquigarrow \lambda x. e_F} \text{ lam} \\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau \rightarrow \tau' \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash e_1 \ e_2 : \tau' \rightsquigarrow e_{F1} \ e_{F2}} \text{ app} \\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \tau_1 \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \tau_2 \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash (e_1, e_2) : (\tau_1 \times \tau_2) \rightsquigarrow (e_{F1}, e_{F2})} \text{ prod} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \times \tau_2 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e) : \tau_1 \rightsquigarrow \text{fst}(e_F)} \text{ fst} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \times \tau_2 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{snd}(e) : \tau_1 \rightsquigarrow \text{snd}(e_F)} \text{ snd} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e) : \tau_1 + \tau_2 \rightsquigarrow \text{inl}(e_F)} \text{ inl} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{inr}(e) : \tau_1 + \tau_2 \rightsquigarrow \text{inr}(e_F)} \text{ inr} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau_1 + \tau_2 \rightsquigarrow e_F \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_1 : \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_2 : \tau \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash \text{case}(e, x. e_1, y. e_2) : \tau \rightsquigarrow \text{case}(e_F, x. e_{F1}, y. e_{F2})} \text{ case} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{Lb}_\ell(e) : (\text{Labeled } \ell \ \tau) \rightsquigarrow \text{inl}(e_F)} \text{ label} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell \ \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e) : \mathbb{C} \ \ell_i \ (\ell_i \sqcup \ell) \ \tau \rightsquigarrow \lambda_. e_F} \text{ unlabel} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \mathbb{C} \ \ell_i \ \ell_o \ \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e) : \mathbb{C} \ \ell_i \ \ell_i \ (\text{Labeled } \ell_o \ \tau) \rightsquigarrow \lambda_. \text{inl}(e_F)} \text{ toLabeled} \\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e) : \mathbb{C} \ \ell_i \ \ell_i \ \tau \rightsquigarrow \lambda_. \text{inl}(e_F)} \text{ ret}
\end{array}$$

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \mathbb{C} \ell_i \ell \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_2 : \mathbb{C} \ell \ell_o \tau' \rightsquigarrow e_{F2}}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_1, x.e_2) : \mathbb{C} \ell_i \ell_o \tau' \rightsquigarrow \lambda_.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}())} \text{bind} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{Labeled } \ell' \tau \rightsquigarrow e_F \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e : \mathbb{C} \ell \ell (\text{ref } \ell' \tau) \rightsquigarrow \lambda_.\text{inl}(\text{new } (e_F))} \text{ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \text{ref } \ell \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash !e : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \rightsquigarrow \lambda_.\text{inl}(e_F)} \text{deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e_1 : \text{ref } \ell' \tau \rightsquigarrow e_{F1} \quad \Sigma; \Psi; \Gamma \vdash e_2 : \text{Labeled } \ell' \tau \rightsquigarrow e_{F2} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_1 := e_2 : \mathbb{C} \ell \ell \text{ unit} \rightsquigarrow \lambda_.\text{inl}(e_{F1} := e_{F2})} \text{assign} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \tau' \rightsquigarrow e_F \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F} \text{sub} \\
\\
\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \Lambda e : \forall \alpha. \tau \rightsquigarrow \Lambda e_F} \text{FI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : \forall \alpha. \tau \rightsquigarrow e_F \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e [] : \tau[\ell/\alpha] \rightsquigarrow e_F[]} \text{FE} \\
\\
\frac{\Sigma; \Psi, c; \Gamma \vdash e : \tau \rightsquigarrow e_F}{\Sigma; \Psi; \Gamma \vdash \nu e : c \Rightarrow \tau \rightsquigarrow \nu e_F} \text{CI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash e : c \Rightarrow \tau \rightsquigarrow e_F \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e \bullet : \tau \rightsquigarrow e_F \bullet} \text{CE}
\end{array}$$

2.3.2 Type preservation for CG to FG translation

Assumption 2.58. $\forall e, \tau, \Sigma, \Psi, \Gamma, \ell_i, \ell_o.$

$$\Sigma; \Psi; \Gamma \vdash e : \mathbb{C} \ell_i \ell_o \tau \implies \ell_i \sqsubseteq \ell_o$$

Theorem 2.59 (CG \rightsquigarrow FG: Type preservation). $\forall \Sigma, \Psi, \Gamma, e_C, \tau.$

$$\Sigma; \Psi; \Gamma \vdash e_C : \tau \text{ is a valid typing derivation in CG} \implies$$

$$\exists e_F.$$

$$\Sigma; \Psi; \Gamma \vdash e_C : \tau \rightsquigarrow e_F \wedge$$

$$\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau] \text{ is a valid typing derivation in FG}$$

Proof. Proof by induction on the translation judgment. We show selected cases below.

1. label:

$$\frac{\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau]}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \text{inl}(e_F) : ([\tau] + \text{unit})^{\perp}} \text{FG-inl}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \text{inl}(e_F) : ([\tau] + \text{unit})^{\ell}} \text{FG-sub}$$

2. unlabel:

P1:

$$\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i \sqcup \ell \quad \overline{\Sigma; \Psi \vdash ([\tau] + \text{unit}) <: ([\tau] + \text{unit})} \text{Lemma 2.1}}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell} <: ([\tau] + \text{unit})^{\ell_i \sqcup \ell}} \text{FGsub-label}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : ([\tau] + \text{unit})^\ell} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top \quad P1}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_F : ([\tau] + \text{unit})^{\ell_i \sqcup \ell}} \text{FG-sub}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} \lambda_{-}.e_F : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^{\ell_i \sqcup \ell})^\perp} \text{FG-lam}}$$

3. toLabeled:

P2:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^{\ell_o})^\perp} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_F : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^{\ell_o})^\perp} \text{FG-sub}}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} () : \text{unit}} \text{P2} \quad \Sigma; \Psi \vdash \ell_i \sqcup \perp \sqsubseteq \ell_i \quad \Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell_o} \searrow \perp}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_F() : ([\tau] + \text{unit})^{\ell_o}} \text{FG-app}$$

Main derivation:

$$\frac{\frac{P1 \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell_i}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} \text{inl}(e_F()) : (([\tau] + \text{unit})^{\ell_o} + \text{unit})^{\ell_i}} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(e_F()) : (\text{unit} \xrightarrow{\ell_i} (([\tau] + \text{unit})^{\ell_o} + \text{unit})^{\ell_i})^\perp} \text{FG-lam}}$$

4. ret:

$$\frac{\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : [\tau]} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_F : [\tau]} \text{FG-sub} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell_i}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} \text{inl}(e_F) : ([\tau] + \text{unit})^{\ell_i}} \text{FG-sub, FG-inl}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(e_F) : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^{\ell_i})^\perp}$$

5. bind:

P1.1:

$$\frac{\frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_{F1} : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^\ell)^\perp} \text{IH1, Weakening} \quad \Sigma; \Psi \vdash \ell_i \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_{F1} : (\text{unit} \xrightarrow{\ell_i} ([\tau] + \text{unit})^\ell)^\perp} \text{FG-sub}}$$

P1:

$$\frac{P1.1 \quad \frac{}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} () : \text{unit}} \text{FG-var} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell}{\Sigma; \Psi \vdash (\ell_i \sqcup \perp) \sqsubseteq \ell_i \quad \frac{}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell \searrow \perp} \text{FG-app}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} e_{F1}() : ([\tau] + \text{unit})^\ell} \text{FG-app}$$

P2.1:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\top} e_{F2} : (\text{unit} \xrightarrow{\ell} ([\tau'] + \text{unit})^{\ell_o})^{\perp}}{\Sigma; \Psi \vdash \ell \sqsubseteq \top} \text{IH2, Weakening}}{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\ell} e_{F2} : (\text{unit} \xrightarrow{\ell} ([\tau'] + \text{unit})^{\ell_o})^{\perp}} \text{FG-sub}$$

P2:

$$\frac{\frac{\frac{P2.1 \quad \frac{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\ell} () : \text{unit}}{\Sigma; \Psi \vdash \perp \sqsubseteq \ell_o} \text{FG-var}}{\Sigma; \Psi \vdash (\ell \sqcup \perp) \sqsubseteq \ell} \quad \frac{\Sigma; \Psi \vdash ([\tau'] + \text{unit})^{\ell_o} \searrow \perp}{\Sigma; \Psi \vdash ([\tau'] + \text{unit})^{\ell_o}} \text{FG-app}}{\Sigma; \Psi; [\Gamma], - : \text{unit}, x : [\tau] \vdash_{\ell_i \sqcup \ell} e_{F2}() : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-app}}$$

P3:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit}, y : \text{unit} \vdash_{\ell} () : \text{unit}}{\Sigma; \Psi; [\Gamma], - : \text{unit}, y : \text{unit} \vdash_{\ell} \text{inr}() : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-var} \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell_o}{\Sigma; \Psi; [\Gamma], - : \text{unit}, y : \text{unit} \vdash_{\ell} \text{inr}() : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-sub, FG-inr}$$

Main derivation:

$$\frac{\frac{\frac{P1 \quad P2 \quad P3 \quad \frac{\frac{\frac{\Sigma; \Psi; \Gamma \vdash e_2 : \mathbb{C} \ell \ell_o \tau}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_o} \text{Assumption 2.58}}{\Sigma; \Psi \vdash ([\tau'] + \text{unit})^{\ell_o} \searrow \ell} \text{FG-case}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell_i} \text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : ([\tau'] + \text{unit})^{\ell_o}} \text{FG-lam, weak}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{case}(e_{F1}(), x.e_{F2}(), y.\text{inr}()) : (\text{unit} \xrightarrow{\ell_i} ([\tau'] + \text{unit})^{\ell_o})^{\perp}} \text{FG-lam, weak}}$$

6. ref:

P1:

$$\frac{\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : ([\tau] + \text{unit})^{\ell'} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_F : ([\tau] + \text{unit})^{\ell'}} \text{FG-sub}}{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_F : ([\tau] + \text{unit})^{\ell'} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell'} \searrow \ell} \text{FG-ref}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} \text{new } e_F : (\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp}} \text{FG-ref}$$

Main derivation:

$$\frac{\frac{\frac{P1 \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} \text{inl}(\text{new } e_F) : ((\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\ell}} \text{FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(\text{new } e_F) : (\text{unit} \xrightarrow{\ell} ((\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp} + \text{unit})^{\ell})^{\perp}} \text{FG-lam}}$$

7. deref:

P2:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_F : (\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp} \text{IH, Weakening} \quad \Sigma; \Psi \vdash \ell' \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} e_F : (\text{ref}([\tau] + \text{unit})^{\ell'})^{\perp}} \text{FG-sub}}$$

P1:

$$P2 \quad \frac{\frac{\Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell <: ([\tau] + \text{unit})^\ell \text{ Lemma 2.1} \quad \Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell \searrow \perp}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} !e_F : ([\tau] + \text{unit})^\ell} \text{ FG-deref}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} !e_F : ([\tau] + \text{unit})^\ell}$$

Main derivation:

$$P1 \quad \frac{\frac{\Sigma; \Psi \vdash \perp \sqsubseteq \ell' \quad \frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^\ell <: ([\tau] + \text{unit})^\ell} \text{ FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell'} \text{inl}(!e_F) : (([\tau] + \text{unit})^\ell + \text{unit})^{\ell'}} \text{ FG-lam}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(!e_F) : (\text{unit} \xrightarrow{\ell'} (([\tau] + \text{unit})^\ell + \text{unit})^{\ell'})^\perp}$$

8. assign:

P3:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_{F2} : ([\tau] + \text{unit})^{\ell'} \text{ IH2, Weakening} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F2} : ([\tau] + \text{unit})^{\ell'}} \text{ FG-sub}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F2} : ([\tau] + \text{unit})^{\ell'}}$$

P2:

$$\frac{\frac{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\top} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp \text{ IH1, Weakening} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \top}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp} \text{ FG-sub}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} : (\text{ref}([\tau] + \text{unit})^{\ell'})^\perp}$$

P1:

$$P2 \quad P3 \quad \frac{\frac{\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash ([\tau] + \text{unit})^{\ell'} \searrow (\ell \sqcup \perp)} \text{ Given}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} := e_{F2} : \text{unit}} \text{ FG-assign}}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} e_{F1} := e_{F2} : \text{unit}}$$

Main derivation:

$$\frac{\frac{P1 \quad \Sigma; \Psi \vdash \perp \sqsubseteq \ell}{\Sigma; \Psi; [\Gamma], - : \text{unit} \vdash_{\ell} \text{inl}(e_{F1} := e_{F2}) : (\text{unit} + \text{unit})^\ell} \text{ FG-inl, FG-sub}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \lambda_{-}.\text{inl}(e_{F1} := e_{F2}) : (\text{unit} \xrightarrow{\ell} (\text{unit} + \text{unit})^\ell)^\perp} \text{ FG-lam}}$$

9. sub:

$$\frac{\frac{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau'] \text{ IH} \quad \Sigma; \Psi \vdash \top \sqsubseteq \top \quad \frac{\Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi \vdash [\tau'] <: [\tau]} \text{ Lemma 2.60}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau]} \text{ FG-sub}}$$

10. FI:

$$\frac{\frac{\Sigma, \alpha; \Psi; [\Gamma] \vdash_{\top} e_F : [\tau] \text{ IH}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \Lambda e_F : (\forall \alpha. (\top, [\tau]))^\perp} \text{ FG-FI}}$$

11. FE:

$$\frac{\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : (\forall \alpha. (\top, [\tau]))^{\perp}} \text{ IH}}{\text{FV}(\ell) \in \Sigma \quad \Sigma; \Psi \vdash \top \sqcup \perp \sqsubseteq \top \quad \Sigma; \Psi \vdash [\tau[\ell/\alpha]] \searrow \perp} \text{ FG-FE}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F [] : [\tau][\ell/\alpha]} \text{ FG-FE}$$

12. CI:

$$\frac{\overline{\Sigma; \Psi, c; [\Gamma] \vdash_{\top} e_F : [\tau]} \text{ IH}}{\Sigma; \Psi; [\Gamma] \vdash_{\top} \nu e_F : (c \overset{\top}{\Rightarrow} [\tau])^{\perp}} \text{ FG-CI}$$

13. CE:

$$\frac{\overline{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F : (c \overset{\top}{\Rightarrow} [\tau])^{\perp}} \text{ IH} \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash \top \sqcup \perp \sqsubseteq \top \quad \Sigma; \Psi \vdash [\tau] \searrow \perp}{\Sigma; \Psi; [\Gamma] \vdash_{\top} e_F \bullet : [\tau]} \text{ FG-CE}$$

□

Lemma 2.60 (CG \rightsquigarrow FG: Subtyping). *For any CG types τ and τ' , Σ , and Ψ , if $\Sigma; \Psi \vdash \tau <: \tau'$, then $\Sigma; \Psi \vdash [\tau] <: [\tau']$.*

Proof. Proof by induction on CG's subtyping relation

1. CGsub-base:

$$\overline{\Sigma; \Psi \vdash [\tau] <: [\tau]} \text{ Lemma 2.1}$$

2. CGsub-arrow:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash [\tau'_1] <: [\tau_1]} \text{ IH1} \quad \overline{\Sigma; \Psi \vdash [\tau'_2] <: [\tau'_2]} \text{ IH2} \quad \Sigma; \Psi \vdash \top \sqsubseteq \top}{\Sigma; \Psi \vdash ([\tau_1] \overset{\top}{\rightarrow} [\tau_2])^{\perp} <: ([\tau'_1] \overset{\top}{\rightarrow} [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\Sigma; \Psi \vdash ((\tau_1 \overset{\ell_s}{\rightarrow} \tau_2)) <: ((\tau'_1 \overset{\ell'_s}{\rightarrow} \tau'_2))} \text{ Definition of } [\cdot]$$

3. CGsub-prod:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash [\tau_1] <: [\tau'_1]} \text{ IH1} \quad \overline{\Sigma; \Psi \vdash [\tau_2] <: [\tau'_2]} \text{ IH2}}{\Sigma; \Psi \vdash ([\tau_1] \times [\tau_2])^{\perp} <: ([\tau'_1] \times [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\Sigma; \Psi \vdash ((\tau_1 \times \tau_2)) <: ((\tau'_1 \times \tau'_2))} \text{ Definition of } [\cdot]$$

4. CGsub-sum:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash [\tau_1] <: [\tau'_1]} \text{ IH1} \quad \overline{\Sigma; \Psi \vdash [\tau_2] <: [\tau'_2]} \text{ IH2}}{\Sigma; \Psi \vdash ([\tau_1] + [\tau_2])^{\perp} <: ([\tau'_1] + [\tau'_2])^{\perp}} \text{ FGsub-arrow}}{\Sigma; \Psi \vdash ((\tau_1 + \tau_2)) <: ((\tau'_1 + \tau'_2))} \text{ Definition of } [\cdot]$$

5. CGsub-labeled:

$$\begin{array}{c}
\frac{\frac{\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket} <: \llbracket \tau'_1 \rrbracket}}{\text{IH1}} \quad \frac{\overline{\Sigma; \Psi \vdash \text{unit}} <: \text{unit}}{\text{FGsub-unit}}}{\frac{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit})} <: (\llbracket \tau'_1 \rrbracket + \text{unit})}}{\text{FGsub-sum}}}{\frac{\frac{\overline{\text{Labeled } \ell_1 \tau_1} <: \text{Labeled } \ell'_1 \tau'_1}}{\text{Given}}}{\frac{\ell_1 \sqsubseteq \ell'_1}{\text{By inversion}}}}{\frac{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit})}^{\ell_1} <: (\llbracket \tau'_1 \rrbracket + \text{unit})}^{\ell'_1}}{\text{FGsub-arrow}}}}{\frac{\overline{\Sigma; \Psi \vdash \llbracket \text{Labeled } \ell_1 \tau_1 \rrbracket} <: \llbracket \text{Labeled } \ell'_1 \tau'_1 \rrbracket}}{\text{Definition of } \llbracket \cdot \rrbracket}}
\end{array}$$

6. CGsub-monad:

P3:

$$\frac{\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket} <: \llbracket \tau'_1 \rrbracket}}{\text{IH}} \quad \frac{\overline{\Sigma; \Psi \vdash \text{unit}} <: \text{unit}}{\text{FGsub-unit}}}{\frac{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit})} <: (\llbracket \tau'_1 \rrbracket + \text{unit})}}{\text{FGsub-sum}}$$

P2:

$$\frac{P3 \quad \frac{\frac{\overline{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau_1} <: \mathbb{C} \ell'_i \ell'_o \tau'_1}}{\text{Given}}}{\frac{\overline{\Sigma; \Psi \vdash \ell_o \sqsubseteq \ell'_o}}{\text{By inversion}}}}{\frac{\overline{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket + \text{unit})}^{\ell_o} <: (\llbracket \tau'_1 \rrbracket + \text{unit})}^{\ell'_o}}{\text{FGsub-label}}$$

P1:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \text{unit}} <: \text{unit}}{\text{P2}} \quad \frac{\frac{\overline{\Sigma; \Psi \vdash \mathbb{C} \ell_i \ell_o \tau_1} <: \mathbb{C} \ell'_i \ell'_o \tau'_1}}{\text{Given}}}{\frac{\overline{\Sigma; \Psi \vdash \ell'_i \sqsubseteq \ell_i}}{\text{FGsub-arrow}}}}{\frac{\overline{\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})}^{\ell_o})} <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})}^{\ell'_o})}}$$

Main derivation:

$$\frac{P1 \quad \frac{\overline{\Sigma; \Psi \vdash \perp \sqsubseteq \perp}}{\text{FGsub-label}}}{\frac{\overline{\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_1 \rrbracket + \text{unit})}^{\ell_o})^\perp <: (\text{unit} \xrightarrow{\ell'_i} (\llbracket \tau'_1 \rrbracket + \text{unit})}^{\ell'_o})^\perp}}{\frac{\overline{\Sigma; \Psi \vdash \llbracket \mathbb{C} \ell_i \ell_o \tau_1 \rrbracket} <: \llbracket \mathbb{C} \ell'_i \ell'_o \tau'_1 \rrbracket}}{\text{Definition of } \llbracket \cdot \rrbracket}}$$

7. CGsub-forall:

P1:

$$\frac{\frac{\overline{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket} <: \llbracket \tau' \rrbracket}}{\text{IH, Weakening}} \quad \frac{\overline{\Sigma, \alpha; \Psi \vdash \top \sqsubseteq \top}}{\text{FGsub-forall}}}{\frac{\overline{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket))} <: (\forall \alpha. (\top, \llbracket \tau' \rrbracket))}}$$

Main derivation:

$$\frac{P1 \quad \frac{\overline{\Sigma, \alpha; \Psi \vdash \perp \sqsubseteq \perp}}{\text{FGsub-label}}}{\frac{\overline{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp <: (\forall \alpha. (\top, \llbracket \tau' \rrbracket))^\perp}}{\frac{\overline{\Sigma; \Psi \vdash \llbracket \forall \alpha. \tau \rrbracket} <: \llbracket \forall \alpha. \tau' \rrbracket}}$$

8. CGsub-constraint:

P1:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau \rrbracket} <: \llbracket \tau' \rrbracket} \text{IH} \quad \frac{\overline{\Sigma; \Psi \vdash \top \sqsubseteq \top} \quad \frac{\overline{\Sigma; \Psi \vdash c \Rightarrow \tau <: c' \Rightarrow \tau'} \text{Given}}{\Sigma; \Psi \vdash c' \Longrightarrow c} \text{By inversion}}{\Sigma; \Psi \vdash (c \overset{\top}{\Rightarrow} \llbracket \tau \rrbracket) <: (c' \overset{\top}{\Rightarrow} \llbracket \tau' \rrbracket)} \text{FGsub-constra}}$$

Main derivation:

$$\frac{P1 \quad \frac{\overline{\Sigma, \alpha; \Psi \vdash \perp \sqsubseteq \perp}}{\Sigma; \Psi \vdash (c \overset{\top}{\Rightarrow} \llbracket \tau \rrbracket)^\perp <: (c' \overset{\top}{\Rightarrow} \llbracket \tau' \rrbracket)^\perp} \text{FGsub-label}}{\Sigma; \Psi \vdash \llbracket c \Rightarrow \tau \rrbracket <: \llbracket c' \Rightarrow \tau' \rrbracket}}$$

□

Lemma 2.61 (CG \rightsquigarrow FG: Preservation of well-formedness). $\forall \Sigma, \Psi, \tau$.

$$\Sigma; \Psi \vdash \tau \text{ WF} \Longrightarrow \Sigma; \Psi \vdash \llbracket \tau \rrbracket \text{ WF}$$

Proof. Proof by induction on the τ WF relation.

1. CG-wff-base:

$$\frac{\overline{\Sigma; \Psi \vdash \mathbf{b} \text{ WF}} \text{FG-wff-base}}{\Sigma; \Psi \vdash \mathbf{b}^\perp \text{ WF}} \text{FG-wff-label}$$

2. CG-wff-unit:

$$\overline{\Sigma; \Psi \vdash \text{unit} \text{ WF}} \text{FG-wff-unit}$$

3. CG-wff-arrow:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket} \text{ WF}} \text{IH1} \quad \frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket} \text{ WF}} \text{IH2}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket) \text{ WF}} \text{FG-wff-arrow}}{\Sigma; \Psi \vdash (\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket)^\perp \text{ WF}} \text{FG-wff-label}$$

4. CG-wff-prod:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket} \text{ WF}} \text{IH1} \quad \frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket} \text{ WF}} \text{IH2}}{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket) \rrbracket \text{ WF}} \text{FG-wff-prod}}{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket) \rrbracket^\perp \text{ WF}} \text{FG-wff-label}$$

5. CG-wff-sum:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket} \text{ WF}} \text{IH1} \quad \frac{\overline{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket} \text{ WF}} \text{IH2}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket \text{ WF}} \text{FG-wff-prod}}{\Sigma; \Psi \vdash \llbracket (\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket) \rrbracket^\perp \text{ WF}} \text{FG-wff-label}$$

6. CG-wff-ref:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \text{ref } \ell \ \tau \ WF} \text{ Given} \\
\frac{}{\text{FV}(\tau) = \emptyset} \text{ By inversion} \\
\hline
\text{FV}(\llbracket \tau \rrbracket) = \emptyset \quad \text{Lemma 2.62} \\
\hline
\frac{}{\text{FV}(\text{unit}) = \emptyset} \text{ Given} \\
\frac{}{\Sigma; \Psi \vdash \text{ref } \ell \ \tau \ WF} \text{ By inversion} \\
\hline
\text{FV}(\llbracket \tau \rrbracket) = \emptyset \\
\hline
\Sigma; \Psi \vdash \text{FV}((\llbracket \tau \rrbracket + \text{unit})^\ell) = \emptyset \\
\hline
\Sigma; \Psi \vdash \text{ref } (\llbracket \tau \rrbracket + \text{unit})^\ell \ WF \quad \text{FG-wff-ref} \\
\hline
\Sigma; \Psi \vdash (\text{ref } (\llbracket \tau \rrbracket + \text{unit})^\ell)^\perp \ WF \quad \text{FG-wff-label}
\end{array}$$

7. CG-wff-forall:

$$\begin{array}{c}
\frac{}{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \\
\frac{}{\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket)) \ WF} \text{ FG-wff-forall} \\
\hline
\Sigma; \Psi \vdash (\forall \alpha. (\top, \llbracket \tau \rrbracket))^\perp \ WF \quad \text{CG-wff-label}
\end{array}$$

8. CG-wff-constraint:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi, c \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \\
\frac{}{\Sigma; \Psi \vdash (c \xrightarrow{\top} \llbracket \tau \rrbracket) \ WF} \text{ FG-wff-constraint} \\
\hline
\Sigma; \Psi \vdash (c \xrightarrow{\top} \llbracket \tau \rrbracket)^\perp \ WF \quad \text{CG-wff-label}
\end{array}$$

9. CG-wff-labeled:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \quad \frac{}{\Sigma; \Psi \vdash \text{unit} \ WF} \text{ FG-wff-unit} \\
\hline
\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit}) \ WF \quad \text{FG-wff-sum} \\
\hline
\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit})^\ell \ WF \quad \text{CG-wff-label}
\end{array}$$

10. CG-wff-monad:

P1:

$$\frac{}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket \ WF} \text{ IH} \quad \frac{}{\Sigma; \Psi \vdash \text{unit} \ WF} \text{ FG-wff-unit} \\
\hline
\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit}) \ WF \quad \text{FG-wff-sum}$$

Main derivation:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi \vdash \text{unit} \ WF} \text{ FG-wff-unit} \quad \frac{}{\Sigma; \Psi \vdash (\llbracket \tau \rrbracket + \text{unit})^{\ell_o} \ WF} \text{ P1} \\
\hline
\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau \rrbracket + \text{unit})^{\ell_o}) \ WF \quad \text{FG-wff-sum} \\
\hline
\Sigma; \Psi \vdash (\text{unit} \xrightarrow{\ell_i} (\llbracket \tau \rrbracket + \text{unit})^{\ell_o})^\perp \ WF \quad \text{CG-wff-label}
\end{array}$$

□

Lemma 2.62 (CG \rightsquigarrow FG: Free variable lemma). $\forall \tau. FV(\llbracket \tau \rrbracket) \subseteq FV(\tau)$

Proof. Proof by induction on the CG types, τ

1. $\tau = \mathbf{b}$:

$$\begin{aligned} & FV(\llbracket \mathbf{b} \rrbracket) \\ &= FV(\mathbf{b}^\perp) \quad \text{Definition of } \llbracket \cdot \rrbracket \\ &= \emptyset \\ &= FV(\mathbf{b}) \end{aligned}$$

2. $\tau = \mathbf{unit}$:

$$\begin{aligned} & FV(\llbracket \mathbf{b} \rrbracket) \\ &= FV(\mathbf{unit}^\perp) \quad \text{Definition of } \llbracket \cdot \rrbracket \\ &= \emptyset \\ &= FV(\mathbf{unit}) \end{aligned}$$

3. $\tau = \tau_1 \rightarrow \tau_2$:

$$\begin{aligned} & FV(\llbracket \tau_1 \rightarrow \tau_2 \rrbracket) \\ &= FV(\llbracket \tau_1 \rrbracket \overset{\top}{\rightarrow} \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ &= FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\ &\subseteq FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\ &= FV(\tau_1 \rightarrow \tau_2) \end{aligned}$$

4. $\tau = \tau_1 \times \tau_2$:

$$\begin{aligned} & FV(\llbracket \tau_1 \times \tau_2 \rrbracket) \\ &= FV(\llbracket \tau_1 \rrbracket \times \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ &= FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\ &\subseteq FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\ &= FV(\tau_1 \times \tau_2) \end{aligned}$$

5. $\tau = \tau_1 + \tau_2$:

$$\begin{aligned} & FV(\llbracket \tau_1 + \tau_2 \rrbracket) \\ &= FV(\llbracket \tau_1 \rrbracket + \llbracket \tau_2 \rrbracket)^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ &= FV(\llbracket \tau_1 \rrbracket) \cup FV(\llbracket \tau_2 \rrbracket) \\ &\subseteq FV(\tau_1) \cup FV(\tau_2) \quad \text{IH on } \tau_1 \text{ and } \tau_2 \\ &= FV(\tau_1 + \tau_2) \end{aligned}$$

6. $\tau = \text{ref } \ell_i \tau_i$:

$$\begin{aligned} & FV(\llbracket \text{ref } \ell_i \tau_i \rrbracket) \\ &= FV(\text{ref } (\llbracket \tau_i \rrbracket + \mathbf{unit})^{\ell_i})^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ &= FV(\llbracket \tau_i \rrbracket) \cup FV(\ell_i) \\ &\subseteq FV(\tau_i) \cup FV(\ell_i) \quad \text{IH} \\ &= FV(\text{ref } \ell_i \tau_i) \end{aligned}$$

7. $\tau = \forall \alpha. \tau_i$:

$$\begin{aligned} & FV(\llbracket \forall \alpha. \tau_i \rrbracket) \\ &= FV(\forall \alpha. (\top, \llbracket \tau_i \rrbracket))^\perp \quad \text{Definition of } \llbracket \cdot \rrbracket \\ &= FV(\llbracket \tau_i \rrbracket) - \{\alpha\} \\ &\subseteq FV(\tau_i) - \{\alpha\} \quad \text{IH} \\ &= FV(\forall \alpha. \tau_i) \end{aligned}$$

8. $\tau = c \Rightarrow \tau_i$:

$$\begin{aligned}
& \text{FV}(\llbracket c \Rightarrow \tau_i \rrbracket) \\
= & \text{FV}(c \xrightarrow{\top} \llbracket \tau_i \rrbracket)^\perp && \text{Definition of } \llbracket \cdot \rrbracket \\
= & \text{FV}(\llbracket c \rrbracket) \cup \text{FV}(\llbracket \tau_i \rrbracket) \\
\subseteq & \text{FV}(\llbracket c \rrbracket) \cup \text{FV}(\tau_i) && \text{IH} \\
= & \text{FV}(c \Rightarrow \tau_i)
\end{aligned}$$

9. $\tau = \text{Labeled } \ell_i \tau_i$:

$$\begin{aligned}
& \text{FV}(\llbracket \text{Labeled } \ell_i \tau_i \rrbracket) \\
= & \text{FV}(\llbracket \tau_i \rrbracket + \text{unit})^{\ell_i} && \text{Definition of } \llbracket \cdot \rrbracket \\
= & \text{FV}(\llbracket \tau_i \rrbracket) \cup \text{FV}(\ell_i) \\
\subseteq & \text{FV}(\tau_i) \cup \text{FV}(\ell_i) && \text{IH} \\
= & \text{FV}(\text{Labeled } \ell_i \tau_i)
\end{aligned}$$

10. $\tau = \mathbb{C} \ell_i \ell_o \tau_i$:

$$\begin{aligned}
& \text{FV}(\llbracket \mathbb{C} \ell_i \ell_o \tau_i \rrbracket) \\
= & \text{FV}(\text{unit} \xrightarrow{\ell_i} (\llbracket \tau_i \rrbracket + \text{unit})^{\ell_o})^\perp && \text{Definition of } \llbracket \cdot \rrbracket \\
= & \text{FV}(\llbracket \tau_i \rrbracket) \cup \text{FV}(\ell_i) \cup \text{FV}(\ell_o) \\
\subseteq & \text{FV}(\tau_i) \cup \text{FV}(\ell_i) \cup \text{FV}(\ell_o) && \text{IH} \\
= & \text{FV}(\mathbb{C} \ell_i \ell_o \tau_i)
\end{aligned}$$

□

2.3.3 Logical relation for CG to FG translation

$W : ((\text{Loc} \mapsto \text{Type}) \times (\text{Loc} \mapsto \text{Type})) \times (\text{Loc} \leftrightarrow \text{Loc})$

Definition 2.63 (CG \rightsquigarrow FG: ${}^s\theta_2$ extends ${}^s\theta_1$). ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq \forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$

Definition 2.64 (CG \rightsquigarrow FG: $\hat{\beta}_2$ extends $\hat{\beta}_1$). $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq \forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$

Definition 2.65 (CG \rightsquigarrow FG: Unary value relation).

$$\begin{aligned}
[\mathbf{b}]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{b} \rrbracket \wedge {}^t v \in \llbracket \mathbf{b} \rrbracket \wedge {}^s v = {}^t v\} \\
[\mathbf{unit}]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid {}^s v \in \llbracket \mathbf{unit} \rrbracket \wedge {}^t v \in \llbracket \mathbf{unit} \rrbracket\} \\
[\tau_1 \times \tau_2]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \mid \\
&\quad (s\theta, m, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}} \wedge (s\theta, m, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}\} \\
[\tau_1 + \tau_2]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \text{inl } {}^s v, \text{inl } {}^t v) \mid (s\theta, m, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}}\} \cup \\
&\quad \{(s\theta, m, \text{inr } {}^s v, \text{inr } {}^t v) \mid (s\theta, m, {}^s v, {}^t v) \in [\tau_2]_V^{\hat{\beta}}\} \\
[\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \lambda x.e_s, \lambda x.e_t) \mid \forall {}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}'} \\
&\quad \implies ({}^s \theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2]_V^{\hat{\beta}'}\} \\
[\forall \alpha. \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \Lambda e_s, \Lambda e_t) \mid \forall {}^s \theta' \sqsupseteq {}^s \theta, j < m, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s, e_t) \in [\tau[\ell'/\alpha]]_V^{\hat{\beta}'}\} \\
[c \Rightarrow \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, \nu e_s, \nu e_t) \mid \mathcal{L} \models c \implies \forall {}^s \theta' \sqsupseteq {}^s \theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s, e_t) \in [\tau]_V^{\hat{\beta}'}\} \\
[\text{ref } \ell \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s a, {}^t a) \mid {}^s \theta({}^s a) = \text{Labeled } \ell \tau \wedge ({}^s a, {}^t a) \in \hat{\beta}\} \\
[\text{Labeled } \ell \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid \\
&\quad \exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge (s\theta, m, {}^s v', {}^t v') \in [\tau]_V^{\hat{\beta}}\} \\
[\mathbb{C} \ell_1 \ell_2 \tau]_V^{\hat{\beta}} &\triangleq \{(s\theta, m, {}^s v, {}^t v) \mid \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_s, H_t, i, {}^s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'. \\
&\quad (k, H_s, H_t) \hat{\triangleright} ({}^s \theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \\
&\quad \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}''. (k - i, H'_s, H'_t) \hat{\triangleright} {}^s \theta' \wedge \\
&\quad \exists {}^t v''. {}^t v = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}''}\}
\end{aligned}$$

Definition 2.66 (CG \rightsquigarrow FG: Unary expression relation).

$$\begin{aligned}
[\tau]_E^{\hat{\beta}} &\triangleq \{(s\theta, n, e_s, e_t) \mid \\
&\quad \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s \theta \wedge \forall i < n, {}^s v.e_s \Downarrow_i {}^s v \implies \\
&\quad \exists H'_t, {}^t v. (H_t, e_t) \Downarrow (H'_t, {}^t v) \wedge (s\theta, n - i, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s \theta\}
\end{aligned}$$

Definition 2.67 (CG \rightsquigarrow FG: Unary heap well formedness).

$$\begin{aligned}
(n, H_s, H_t) \hat{\triangleright} {}^s \theta &\triangleq \text{dom}({}^s \theta) \subseteq \text{dom}(H_s) \wedge \\
&\quad \hat{\beta} \subseteq (\text{dom}({}^s \theta) \times \text{dom}(H_t)) \wedge \\
&\quad \forall (a_1, a_2) \in \hat{\beta}. ({}^s \theta, n - 1, H_s(a_1), H_t(a_2)) \in [{}^s \theta(a)]_V^{\hat{\beta}}
\end{aligned}$$

Definition 2.68 (CG \rightsquigarrow FG: Label substitution). $\sigma : \text{Lvar} \mapsto \text{Label}$

Definition 2.69 (CG \rightsquigarrow FG: Value substitution to values). $\delta^s : \text{Var} \mapsto \text{Val}, \delta^t : \text{Var} \mapsto \text{Val}$

Definition 2.70 (CG \rightsquigarrow FG: Unary interpretation of Γ).

$$\begin{aligned}
[\Gamma]_V^{\hat{\beta}} &\triangleq \{(s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\
&\quad \forall x \in \text{dom}(\Gamma). ({}^s \theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}\}
\end{aligned}$$

2.3.4 Soundness proof for CG to FG translation

Lemma 2.71 (CG \rightsquigarrow FG: Monotonicity). $\forall {}^s \theta, {}^s \theta', n, {}^s v, {}^t v, n', \beta, \beta'$.

$$({}^s \theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge {}^s \theta \sqsubseteq {}^s \theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s \theta', n', {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$$

Proof. Proof by induction on τ

1. Case **b**:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{b}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbf{b}]_V^{\hat{\beta}'}$$

Since $({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{b}]_V^{\hat{\beta}}$ therefore from Definition 2.65 we know that ${}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket$

Therefore from Definition 2.65 ${}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket$ we get the desired

2. Case **unit**:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{unit}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbf{unit}]_V^{\hat{\beta}'}$$

Since $({}^s\theta, n, {}^sv, {}^tv) \in [\mathbf{unit}]_V^{\hat{\beta}}$ therefore from Definition 2.65 we know that ${}^sv \in \llbracket \mathbf{unit} \rrbracket \wedge {}^tv \in \llbracket \mathbf{unit} \rrbracket$

Therefore from Definition 2.65 ${}^sv \in \llbracket \mathbf{unit} \rrbracket \wedge {}^tv \in \llbracket \mathbf{unit} \rrbracket$ we get the desired

3. Case $\tau_1 \times \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.65 we know that ${}^sv = ({}^sv_1, {}^sv_2)$ and ${}^tv = ({}^tv_1, {}^tv_2)$.

We also know that $({}^s\theta, n, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}}$ and $({}^s\theta, n, {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}}$

$$\mathbf{IH1:} ({}^s\theta', n', {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}'}$$

$$\mathbf{IH2:} ({}^s\theta', n', {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}'}$$

Therefore from Definition 2.65, IH1 and IH2 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case $\tau_1 + \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.65 two cases arise

(a) ${}^sv = \text{inl}({}^sv')$ and ${}^tv = \text{inl}({}^tv')$:

$$\text{IH: } ({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$$

Therefore from Definition 2.65 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b) ${}^sv = \text{inr}({}^sv')$ and ${}^tv = \text{inr}({}^tv')$:

Symmetric reasoning as in the previous case

5. Case $\tau_1 \rightarrow \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \rightarrow \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.65 we know that

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta, {}^sv_1, {}^tv_1, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta'', j, {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}} \implies ({}^s\theta'', j, e_s[{}^sv_1/x], e_t[{}^tv_1/x]) \in [\tau_2]_E^{\hat{\beta}'} \quad (\text{A0})$$

Similarly from Definition 2.65 we are required to prove

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}'} \implies ({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

This means we are given some ${}^s\theta'_1 \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''$ s.t $({}^s\theta'_1, j, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}'}$ and we are required to prove

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}'}$$

Instantiating (A0) with ${}^s\theta'_1, {}^sv_2, {}^tv_2, j, \hat{\beta}''$ since ${}^s\theta'_1 \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, j < n' < n$ and $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$ therefore we get

$$({}^s\theta'_1, j, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

6. Case $\forall\alpha.\tau$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\forall\alpha.\tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\forall \alpha. \tau]_{\hat{\beta}'}^{\hat{\beta}'}$$

From Definition 2.65 we know that ${}^s v = \Lambda e'_s$ and ${}^t v = \Lambda e'_t$. And

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'' . ({}^s\theta'', j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}''} \quad (\text{F0})$$

Similarly from Definition 2.65 we are required to prove

$$\forall {}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \ell' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1 . ({}^s\theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}''_1}$$

This means we are given some ${}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \ell'' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1$

and we are required to prove

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}''_1}$$

Instantiating (F0) with ${}^s\theta''_1, j, \hat{\beta}''_1$ since ${}^s\theta''_1 \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, j < n' < n$ and $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''_1$ therefore we get

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau[\ell'/\alpha]]_{\hat{E}}^{\hat{\beta}''_1}$$

7. Case $c \Rightarrow \tau$:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [c \Rightarrow \tau]_{\hat{\beta}}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [c \Rightarrow \tau]_{\hat{\beta}'}^{\hat{\beta}'}$$

From Definition 2.65 we know that ${}^s v = \nu(e'_s)$ and ${}^t v = \nu(e'_t)$. And

$$\mathcal{L} \models c \implies \forall {}^s\theta'' \sqsupseteq {}^s\theta, j < n, \hat{\beta}' \sqsubseteq \hat{\beta}''_1 . ({}^s\theta'', j, e'_s, e'_t) \in [\tau]_{\hat{E}}^{\hat{\beta}'_1} \quad (\text{C0})$$

Similarly from Definition 2.65 we are required to prove

$$\mathcal{L} \models c \implies \forall {}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''_1 . ({}^s\theta''_1, j, e'_s, e'_t) \in [\tau]_{\hat{E}}^{\hat{\beta}''_1}$$

This means we are given some $\mathcal{L} \models c, {}^s\theta''_1 \sqsupseteq {}^s\theta', j < n', \hat{\beta}' \sqsubseteq \hat{\beta}''_1$

and we are required to prove

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau]_{\hat{E}}^{\hat{\beta}''_1}$$

Since $\mathcal{L} \models c$ and instantiating (C0) with ${}^s\theta''_1, j, \hat{\beta}''_1$ since ${}^s\theta''_1 \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, j < n' < n$ and $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''_1$ therefore we get

$$({}^s\theta''_1, j, e'_s, e'_t) \in [\tau]_{\hat{E}}^{\hat{\beta}''_1}$$

8. Case ref $\ell \tau$:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{ref } \ell \tau]_{\hat{\beta}}^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \ell \ \tau]_V^{\hat{\beta}'}$$

From Definition 2.65 we know that ${}^sv = {}^sa$ and ${}^tv = {}^ta$. We also know that ${}^s\theta({}^sa) = \text{Labeled } \ell \ \tau \wedge ({}^sa, {}^ta) \in \hat{\beta}$

From Definition 2.65, Definition 2.63 and Definition 2.64 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{ref } \ell \ \tau]_V^{\hat{\beta}'}$$

9. Case Labeled $\ell \ \tau$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}'}$$

From Definition 2.65 it means

$$\exists {}^sv', {}^tv'. {}^sv = \text{Lb}_\ell({}^sv') \wedge {}^tv = \text{inl } {}^tv' \wedge ({}^s\theta, n, {}^sv', {}^tv') \in [\tau]_V^{\hat{\beta}}$$

$$\underline{\text{IH:}} \ ({}^s\theta', n', {}^sv', {}^tv') \in [\tau]_V^{\hat{\beta}}$$

Similarly from Definition 2.65 we need to prove that

$$\exists {}^sv'', {}^tv''. {}^sv = \text{Lb}_\ell({}^sv'') \wedge {}^tv = \text{inl } {}^tv'' \wedge ({}^s\theta', n', {}^sv'', {}^tv'') \in [\tau]_V^{\hat{\beta}'}$$

We choose ${}^sv''$ as ${}^sv'$ and ${}^tv''$ as ${}^tv'$ and since from IH we know that $({}^s\theta', n', {}^sv', {}^tv') \in [\tau]_V^{\hat{\beta}'}$

Therefore from Definition 2.65 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\text{Labeled } \ell \ \tau]_V^{\hat{\beta}'}$$

10. Case $\mathbb{C} \ \ell_1 \ \ell_2 \ \tau$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\mathbb{C} \ \ell_1 \ \ell_2 \ \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\mathbb{C} \ \ell_1 \ \ell_2 \ \tau]_V^{\hat{\beta}'}$$

This means from Definition 2.65 we know that

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^sv', {}^tv', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}_1.$$

$$(k, H_s, H_t) \hat{\beta}_1 \triangleright ({}^s\theta_e) \wedge (H_s, {}^sv) \Downarrow_i^f (H'_s, {}^sv') \wedge i < k \implies$$

$$\exists {}^tv'. (H_t, {}^tv()) \Downarrow (H'_t, {}^tv') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}_1 \sqsubseteq \hat{\beta}_2. (k - i, H'_s, H'_t) \hat{\beta}_2 \triangleright {}^s\theta' \wedge$$

$$\exists {}^tv''. {}^tv' = \text{inl } {}^tv'' \wedge ({}^s\theta', {}^t\theta', k - i, {}^sv', {}^tv'') \in [\tau]_V^{\hat{\beta}_2} \wedge$$

$$(\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s\theta_e(a) = \text{Labeled } \ell' \ \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge$$

$$(\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e). {}^s\theta'(a) \searrow \ell_1) \quad (\text{CG0})$$

Similarly from Definition 2.65 we need to prove

$$\begin{aligned}
& \forall {}^s\theta'_e \sqsupseteq {}^s\theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1. \\
& (k', H'_s, H'_t) \triangleright^{\hat{\beta}'_1} ({}^s\theta'_e) \wedge (H'_s, {}^s v) \Downarrow_i^f (H''_s, {}^s v'') \wedge (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge i' < k' \implies \\
& \exists {}^t v''. (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge \exists {}^s\theta'' \sqsupseteq {}^s\theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} {}^s\theta'' \wedge \\
& \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k' - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\
& (\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e). {}^s\theta'(a) \searrow \ell_1)
\end{aligned}$$

This means we are given some ${}^s\theta'_e \sqsupseteq {}^s\theta', H'_s, H'_t, i', {}^s v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$ s.t. $(k', H'_s, H'_t) \triangleright ({}^s\theta'_e) \wedge (H'_s, {}^s v) \Downarrow_i^f (H''_s, {}^s v'') \wedge i' < k'$

And we need to prove

$$\begin{aligned}
& \exists {}^t v''. (H'_t, {}^t v()) \Downarrow (H''_t, {}^t v'') \wedge \exists {}^s\theta'' \sqsupseteq {}^s\theta'_e, \hat{\beta}'_1 \sqsubseteq \hat{\beta}'_2. (k' - i', H''_s, H''_t) \triangleright^{\hat{\beta}'_2} {}^s\theta'' \wedge \\
& \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta'', k' - i, {}^s v', {}^t v'') \in [\tau]_V^{\hat{\beta}'_2} \wedge \\
& (\forall a. H_s(a) \neq H'_s(a) \implies \exists \ell'. {}^s\theta_e(a) = \text{Labeled } \ell' \tau' \wedge \ell_1 \sqsubseteq \ell') \wedge \\
& (\forall a \in \text{dom}({}^s\theta') / \text{dom}({}^s\theta_e). {}^s\theta'(a) \searrow \ell_1)
\end{aligned}$$

Instantiating (CG0) with ${}^s\theta'_e \sqsupseteq {}^s\theta', H'_s, H'_t, i', {}^s v'', {}^t v'', k' \leq n', \hat{\beta}' \sqsubseteq \hat{\beta}'_1$ we get the desired

□

Lemma 2.72 (CG \rightsquigarrow FG: Unary monotonicity for Γ). $\forall {}^s\theta, {}^s\theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$.

$$({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies ({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

Proof. Given: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove: $({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$

From Definition 2.70 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$$

And again from Definition 2.70 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

- $\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t)$:

Given

- $\forall x \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$:

Since we know that $\forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}}$ (given)

Therefore from Lemma 2.71 we get

$$\forall x \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x), \delta^t(x)) \in [\Gamma(x)]_V^{\hat{\beta}'}$$

□

Lemma 2.73 (CG \rightsquigarrow FG: Unary monotonicity for H). $\forall {}^s\theta, H_s, H_t, n, n', \hat{\beta}, \hat{\beta}'$.

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta$$

Proof. Given: $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge n' < n$

To prove: $(n', H_s, H_t) \triangleright^{\hat{\beta}'} {}^s\theta$

From Definition 2.67 it is given that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}.({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$$

And again from Definition 2.67 we are required to prove that

$$\text{dom}({}^s\theta) \subseteq \text{dom}(H_S) \wedge \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}.({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}'}$$

- $\text{dom}({}^s\theta) \subseteq \text{dom}(H_S)$:

Given

- $\hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t))$:

Given

- $\forall (a_1, a_2) \in \hat{\beta}.({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}'}$:

Since we know that $\forall (a_1, a_2) \in \hat{\beta}.({}^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}}$ (given)

Therefore from Lemma 2.71 we get

$$\forall (a_1, a_2) \in \hat{\beta}.({}^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [{}^s\theta(a)]_V^{\hat{\beta}'}$$

□

Theorem 2.74 (CG \rightsquigarrow FG: Fundamental theorem). $\forall \Gamma, \tau, e, \delta^s, \delta^t, \sigma, {}^s\theta, n.$

$$\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t \wedge$$

$$\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$$

\implies

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

Proof. Proof by induction on the \rightsquigarrow relation

1. CF-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash x : \tau \rightsquigarrow x} \text{CF-var}$$

$$\text{Also given is: } \mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau\} \sigma]_V^{\hat{\beta}}$$

$$\text{To prove: } ({}^s\theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 it suffices to prove that

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. x \delta^s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

This means given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $x \delta^s \Downarrow_i {}^s v$

From cg-val we know that $i = 0$, ${}^s v = x \delta^s$.

And we are required to prove

$$\exists H'_t, {}^t v. (H_t, x \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n, {}^s v, {}^t v) \in [\tau \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-V0})$$

From fg-val we know that ${}^t v = x \delta^t$ and $H'_t = H_t$. So we are left with proving

$$({}^s \theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we are given $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \cup \{x \mapsto \tau \sigma\} \sigma]_{V}^{\hat{\beta}}$, therefore from Definition 2.70 we get

$$({}^s \theta, n, x \delta^s, x \delta^t) \in [\tau \sigma]_{V}^{\hat{\beta}}. \text{ And we have } (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \text{ in the context. So we are done.}$$

2. CF-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_s : \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \lambda x. e_s : \tau_1 \rightarrow \tau_2 \rightsquigarrow \lambda x. e_t} \text{ lam}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove: $({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

From Definition 2.66 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (\lambda x. e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda x. e_t) \delta^t) \Downarrow (H'_t, {}^t v) ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ and given some $i < n, {}^s v$ s.t $(\lambda x. e_s) \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that ${}^s v = (\lambda x. e_s) \delta^s$, ${}^t v = (\lambda x. e_t) \delta^t$, $H'_t = H_t$ and $i = 0$

It suffices to prove that

$$({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

We know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ from the context. So, we are only left to prove

$$({}^s \theta, n, (\lambda x. e_s) \delta^s, (\lambda x. e_t) \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_{V}^{\hat{\beta}}$$

From Definition 2.65 it suffices to prove

$$\begin{aligned} & \forall {}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'} \\ & \implies ({}^s \theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'} \end{aligned}$$

This means that we are given ${}^s \theta' \sqsupseteq {}^s \theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t $({}^s \theta', j, {}^s v, {}^t v) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'}$

And we need to prove

$$({}^s \theta', j, e_s[{}^s v/x] \delta^s, e_t[{}^t v/x] \delta^t) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'} \quad (\text{F-L0})$$

Since $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$ therefore from Lemma 2.72 we also have

$$({}^s\theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}'}$$

IH:

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_1\}, e_t \cup \{x \mapsto {}^t v_1\}) \in [\tau_2 \sigma]_E^{\hat{\beta}'}$$
 s.t

$$({}^s\theta', j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'}$$

We get (F-L0) directly from IH

3. CF-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : (\tau_1 \rightarrow \tau_2) \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \tau_1 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash e_{s1} e_{s2} : \tau_2 \rightsquigarrow e_{t1} e_{t2}} \text{ app}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, (e_{t1} e_{t2}) \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}}$

This means from Definition 2.66 it suffices to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

This further means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n, {}^s v$ s.t $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, (e_{t1} e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

(F-A0)

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \rightarrow \tau_2) \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.66 we have

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies$$

$$\exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

Instantiating with H_s, H_t and since we know that $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n$ s.t $e_{s1} \delta^s \Downarrow_j {}^s v_1$.

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

(F-A1)

IH2:

$$({}^s\theta, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.66 it suffices to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k < n - j, {}^s v_2.e_{s2} \Downarrow_i {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta'_2 \end{aligned}$$

Instantiating with H_s, H'_{t1} and since we know that $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$ therefore $\exists k < i - j < n - j$ s.t $e_{s2} \delta^s \Downarrow_k {}^s v_2$.

And we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A2})$$

Since from (F-A1) we know that $({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \rightarrow \tau_2) \sigma]_V^{\hat{\beta}}$ where ${}^s v_1 = \lambda x.e'_s$ and ${}^t v_1 = \lambda x.e'_t$

From Definition 2.65 we have

$$\begin{aligned} & \forall {}^s\theta'_3 \sqsupseteq {}^s\theta, {}^s v, {}^t v, l < n - j, \hat{\beta}_3 \sqsupseteq \hat{\beta}. ({}^s\theta'_3, l, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}_3} \\ & \implies ({}^s\theta'_3, l, e'_s[{}^s v/x], e'_t[{}^t v/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}_3} \end{aligned}$$

Instantiating with ${}^s\theta, {}^s v_2, {}^t v_2, n - j - k, \hat{\beta}$ we get

$$({}^s\theta, n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s4}, H_{t4}. (n - j - k, H_{s4}, H_{t4}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k' < n - j - k, {}^s v_4.e'_s[{}^s v_2/x] \Downarrow_{k'} {}^s v_4 \implies \\ & \exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s\theta, n - j - k - k', {}^s v_4, {}^t v_4) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge \\ & (n - j - k - k', H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H'_{t2} , from (F-A2) we know that $(n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta$. Instantiating ${}^s v_4$ with ${}^s v$ and since we know that $(e_{s1} e_{s2}) \delta^s \Downarrow_i {}^s v$ therefore $\exists k' < i - j - k < n - j - k$ s.t $e'_s[{}^s v_2/x] \delta^s \Downarrow_{k'} {}^s v$. therefore we have

$$\exists H'_{t4}, {}^t v_4. (H_{t4}, e'_t[{}^t v_2/x]) \Downarrow (H'_{t4}, {}^t v_4) \wedge ({}^s\theta, n - j - k - k', {}^s v, {}^t v_4) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k - k', H_{s4}, H'_{t4}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-A3})$$

Since from cg-app we know that $i = j + k + k'$ and $H'_t = H'_{t4}$, ${}^t v = {}^t v_4$ therefore we get (F-A0) from (F-A3) and Lemma 2.71 and Lemma 2.73

4. CF-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2) \rightsquigarrow (e_{t1}, e_{t2})} \text{prod}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, (e_{s1}, e_{s2}) \delta^s, (e_{t1}, e_{t2}) \delta^t) \in [(\tau_1 \times \tau_2) \sigma]_E^{\hat{\beta}}$

From Definition 2.66 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t, \hat{\beta}. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some $H_s, H_t, \hat{\beta}$ s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n$ s.t $(e_{s1}, e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \quad (\text{F-P0})$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n. e_{s1} \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$ therefore $\exists j < i < n$ s.t $e_{s1} \delta^s \Downarrow_j {}^s v_1$.

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-P1})$$

IH2:

$$({}^s\theta, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall k < n - j. e_{s2} \delta^s \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with $H_s, H'_{t1}, \hat{\beta}'_1$ and since we know that $(e_{s1}, e_{s2}) \delta^s \Downarrow_i ({}^s v_1, {}^s v_2)$ therefore $\exists k < i - j < n - j$ s.t $e_{s2} \delta^s \Downarrow_k {}^s v_2$.

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}} \wedge (n - j - k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-P2})$$

From cg-prod we know that $i = j + k + 1$, $H'_t = H'_{t2}$ and ${}^t v = ({}^t v_1, {}^t v_2)$ therefore from Definition 2.65 and Lemma 2.71 we get $({}^s\theta, n - i, {}^s v, {}^t v) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}}$

And since we have $(n - j - k, H_s, H'_{t2}) \hat{\triangleright}^s \theta$ therefore from Lemma 2.73 we also get

$$(n - i, H_s, H'_{t2}) \hat{\triangleright}^s \theta$$

5. CF-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 \times \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{fst}(e_t)} \text{fst}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{fst}(e_t) \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$ (F-F0)

This means from Definition 2.66 we need to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{fst}(e_s) \delta^s \Downarrow_i {}^s v &\implies \\ \exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^s) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that we are given some $H_s, H_t, \hat{\beta}$ s.t $(n, H_s, H_t) \hat{\triangleright}^s \theta$ and given some $i < n, {}^s v$ s.t $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{fst}(e_t) \delta^s) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \quad (\text{F-F0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2) \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j ({}^s v_1, -) &\implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v_1, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge \\ (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with H_s, H_t and ${}^s v_1$ with ${}^s v$ since we know that $\text{fst}(e_s) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n$ s.t $e_s \delta^s \Downarrow_j ({}^s v, -)$.

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_{t1}, (e_{t1}, e_{t2}) \delta^t) \Downarrow (H'_{t1}, ({}^t v_1, -)) \wedge ({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}} \wedge \\ (n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta \quad (\text{F-F1}) \end{aligned}$$

From cg-fst we know that $i = j + 1$, $H'_t = H'_{t1}$ and ${}^t v = {}^t v_1$. Since we know $({}^s\theta, n - j, ({}^s v, -), ({}^t v_1, -)) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}}$ therefore from Definition 2.65 and Lemma 2.71 we get $({}^s\theta, n - i, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}}$

And since from (F-F1) we have $(n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta$ therefore from Lemma 2.73 we get

$$(n - i, H_s, H'_{t1}) \hat{\triangleright}^s \theta$$

6. CF-snd:

Symmetric reasoning as in the CF-fst case

7. CF-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{inl}(e_s) : (\tau_1 + \tau_2) \rightsquigarrow \text{inl}(e_t)} \text{prod}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{inl}(e_s), \delta^s, \text{inl}(e_t), \delta^t) \in [(\tau_1 + \tau_2) \sigma]_{\hat{E}}^{\hat{\beta}}$

From Definition 2.66 it suffices to prove

$$\begin{aligned} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some $H_s, H_t, \hat{\beta}$ s.t. $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n, {}^s v$ s.t. $\text{inl}(e_s) \delta^s \Downarrow_i \text{inl}({}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v)) \in [(\tau_1 + \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - \\ i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-IL0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1 \sigma]_{\hat{E}}^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge \\ ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $\text{inl}(e_s) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n$ s.t. $e_s \delta^s \Downarrow_j {}^s v$.

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-IL1})$$

From cg-inl we know that $i = j + 1$ and $H'_t = H'_{t1}, {}^t v = {}^t v_1$. Since we know $({}^s\theta, n - j, {}^s v, {}^t v_1) \in [\tau_1 \sigma]_{\hat{V}}^{\hat{\beta}}$ therefore from Definition 2.65 and Lemma 2.71 we get

$$({}^s\theta, n - i, \text{inl}({}^s v), \text{inl}({}^t v_1)) \in [(\tau_1 + \tau_2) \sigma]_{\hat{V}}^{\hat{\beta}}$$

And since from (F-IL1) we have $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$ therefore from Lemma 2.73 we get

$$(n - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$$

8. CF-inr:

Symmetric reasoning as in the CF-inl case

9. CF-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau_1 + \tau_2 \rightsquigarrow e_t \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash e_{s2} : \tau \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash \mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \mathbf{case}(e_t, x.e_{t1}, y.e_{t2})} \text{ case}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \mathbf{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$

This means from Definition 2.66 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \mathbf{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that we are given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n$ s.t $\mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \mathbf{case}(e_t, x.e_{t1}, y.e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n-i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \text{ (F-C0)}$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2) \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n-j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $\mathbf{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n$ s.t $e_s \delta^s \Downarrow_j {}^s v_1$.

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n-j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_V^{\hat{\beta}} \wedge (n-j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \text{ (F-C1)}$$

Two cases arise:

(a) ${}^s v_1 = \mathbf{inl}({}^s v'_1)$ and ${}^t v_1 = \mathbf{inl}({}^t v'_1)$:

IH2:

$$({}^s\theta, n-j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \hat{\triangleright}^s \theta \wedge \forall k < n - j, {}^s v_2.e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v_2 \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - \\ & j - k, H_{s2}, H'_{t2}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with H_s, H'_{t1} and since we know that $\text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \Downarrow_i {}^s v$ therefore $\exists k < i - j < n - j$ s.t $e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\} \Downarrow_k {}^s v$.

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_1\}) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - k, {}^s v, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - \\ & j - k, H_s, H'_{t2}) \hat{\triangleright}^s \theta \end{aligned}$$

From cg-case1 we know that $i = j + k + 1$ and $H'_t = H'_{t2}, {}^t v = {}^t v_2$. Since we know $({}^s \theta, n - j - k, {}^s v, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}}$ therefore from Definition 2.65 and Lemma 2.71 we get $({}^s \theta, n - i, {}^s v, {}^t v_2) \in [\tau \sigma]_V^{\hat{\beta}}$

And since from (F-C2) we have $(n - j - k, H_s, H'_{t2}) \hat{\triangleright}^s \theta$ therefore from Lemma 2.73 we get $(n - i, H_s, H'_{t2}) \hat{\triangleright}^s \theta$

(b) ${}^s v_1 = \text{inr}({}^s v'_1)$ and ${}^t v_1 = \text{inr}({}^t v'_1)$:

Symmetric reasoning as in the previous case

10. CF-FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \Lambda e_s : \forall \alpha. \tau \rightsquigarrow \Lambda e_t} \text{FI}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s \theta, n, \Lambda e_s \delta^s, \Lambda e_t \delta^t) \in [(\forall \alpha. \tau) \sigma]_E^{\hat{\beta}}$

This means from Definition 2.66 we know that

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \Lambda e_s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \Lambda e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright}^s \theta$ and given some $i < n$ s.t $(\Lambda e_s) \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that ${}^s v = (\Lambda e_s) \delta^s, {}^t v = (\Lambda e_t) \delta^t, i = 0$ and $H'_t = H_t$

It suffices to prove that

$$({}^s \theta, n, (\Lambda e_s) \delta^s, (\Lambda e_t) \delta^t) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^s \theta$$

We know $(n, H_s, H_t) \hat{\triangleright}^s \theta$ from the context. So, we are only left to prove

$$({}^s \theta, n, (\Lambda e_s) \delta^s, (\Lambda e_t) \delta^t) \in [(\forall \alpha. \tau) \sigma]_V^{\hat{\beta}}$$

From Definition 2.65 it suffices to prove

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha]]_{E}^{\hat{\beta}'}$$

This means that we are given ${}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha]]_{E}^{\hat{\beta}'} \quad (\text{F-FI0})$$

Since $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$ therefore from Lemma 2.72 we also have

$$({}^s\theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}'}$$

IH:

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma \cup \{\alpha \mapsto \ell'\}]_{E}^{\hat{\beta}'}$$

We get (F-FI0) directly from IH

11. CF-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \forall \alpha. \tau \rightsquigarrow e_t \quad \text{FV}(\ell) \in \Sigma}{\Sigma; \Psi; \Gamma \vdash e_s [] : \tau[\ell/\alpha] \rightsquigarrow e_t []} \text{FE}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove: $({}^s\theta, n, e_s [] \delta^s, e_t [] \delta^t) \in [\tau[\ell/\alpha] \sigma]_{E}^{\hat{\beta}}$

From Definition 2.66 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright {}^s\theta \wedge \forall i < n, {}^s v. e_s [] \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, e_t []) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau[\ell/\alpha] \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright {}^s\theta \end{aligned}$$

This further means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright {}^s\theta$ and given some $i < n, {}^s v$ s.t $e_s [] \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, e_t []) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\tau[\ell/\alpha] \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright {}^s\theta \quad (\text{F-FE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\forall \alpha. \tau) \sigma]_{E}^{\hat{\beta}}$$

This means from Definition 2.66 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha. \tau) \sigma]_{V}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright {}^s\theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(e_s []) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n, {}^s v_1$ s.t $e_s \delta^s \Downarrow_j {}^s v_1$.

And we have

$\exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha. \tau) \sigma]_{V}^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$
(F-FE1)

From cg-FE we know that ${}^s v_1 = \Lambda e'_s$ and ${}^t v_1 = \Lambda e'_t$

Therefore we have

$({}^s \theta, n - j, \Lambda e'_s, \Lambda e'_t) \in [(\forall \alpha. \tau) \sigma]_{V}^{\hat{\beta}}$

This means from Definition 2.65 we have

$\forall {}^s \theta' \sqsubseteq {}^s \theta, k < n - j, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}_2. ({}^s \theta', k, e'_s, e'_t) \in [\tau[\ell'/\alpha] \sigma]_{E}^{\hat{\beta}_2}$

Instantiating ${}^s \theta'$ with ${}^s \theta$, k with $n - j - 1$, ℓ' with ℓ σ and $\hat{\beta}_2$ with $\hat{\beta}$ and we get

$({}^s \theta, n - j - 1, e'_s, e'_t) \in [\tau[\ell/\alpha] \sigma]_{E}^{\hat{\beta}}$

From Definition 2.66 we get

$\forall H_{s2}, H_{t2}. (n - j - 1, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}_2} {}^s \theta'_1 \wedge \forall k < n - j - 1, {}^s v_2. e'_s \Downarrow_k {}^s v_2 \implies$
 $\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell/\alpha] \sigma]_{V}^{\hat{\beta}} \wedge (n - j - 1 -$
 $k, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$

Instantiating with H_s, H'_{t1} . Since from (F-FE1) we know that $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$ therefore from Lemma 2.73 we get $(n - j - 1, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s \theta$

Since we know that $e_s \Downarrow_i \delta^s$ and from cg-FE we know that $i = j + k + 1$ (for some k) and $i < n$ therefore we have $k < n - j - 1$ s.t $e'_s \delta^s \Downarrow_k {}^s v_2$.

Therefore we have

$\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s \theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell/\alpha] \sigma]_{V}^{\hat{\beta}} \wedge (n - j - 1 -$
 $k, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta$ (F-FE2)

Since $H'_t = H_{t2'}$, ${}^s v = {}^s v_2$ and ${}^t v = {}^t v_2$ therefore we get (F-FE0) directly from (F-FE2)

12. CF-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \nu e_s : c \Rightarrow \tau \rightsquigarrow \nu e_t} \text{ CI}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove: $({}^s \theta, n, \nu e_s \delta^s, \nu e_t \delta^t) \in [(c \Rightarrow \tau) \sigma]_{E}^{\hat{\beta}}$

This means from Definition 2.66 we know that

$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n. \nu e_s \Downarrow_i {}^s v \implies$
 $\exists H'_t, {}^t v. (H_t, \nu e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(c \Rightarrow \tau) \hat{\beta} \sigma]_{V}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$

This means that given some $H_s, H_t, \hat{\beta}$ s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ and given some $i < n$ s.t

$$(\nu e_s) \delta^s \Downarrow_i {}^s v$$

From cg-val and fg-val we know that ${}^s v = (\nu e_s) \delta^s$, ${}^t v = (\nu e_t) \delta^t$, $i = 0$ and $H'_t = H_t$

It suffices to prove that

$$({}^s \theta, n, (\nu e_s) \delta^s, (\nu e_t) \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

We know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ from the context. So, we are only left to prove

$$({}^s \theta, n, (\nu e_s) \delta^s, (\nu e_t) \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{\beta}}^{\hat{\beta}}$$

From Definition 2.65 it suffices to prove

$$\mathcal{L} \models c \sigma \implies \forall {}^s \theta' \sqsupseteq {}^s \theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{\beta}'}^{\hat{\beta}'}$$

This means that we are given $\mathcal{L} \models c \sigma$ and ${}^s \theta' \sqsupseteq {}^s \theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{\beta}'}^{\hat{\beta}'} \quad (\text{F-CI0})$$

Since $({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}}^{\hat{\beta}}$ therefore from Lemma 2.72 we also have

$$({}^s \theta', j, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}'}^{\hat{\beta}'}$$

And since we know that $\mathcal{L} \models c \sigma$ therefore

$$\underline{\text{IH}}: ({}^s \theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{\hat{\beta}'}^{\hat{\beta}'}$$

We get (F-CI0) directly from IH

13. CF-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : c \Rightarrow \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash c}{\Sigma; \Psi; \Gamma \vdash e_s \bullet : \tau \rightsquigarrow e_t \bullet} \text{CE}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{\beta}}^{\hat{\beta}}$

To prove: $({}^s \theta, n, e_s \bullet \delta^s, e_t \bullet \delta^t) \in [\tau \sigma]_{\hat{\beta}}^{\hat{\beta}}$

From Definition 2.66 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. e_s \bullet \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, e_t \bullet) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This further means that given some $H_s, H_t, \hat{\beta}$ s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ and given some $i < n$ s.t $e_s \bullet \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, e_t \bullet) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-CE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(c \Rightarrow \tau) \sigma]_{\hat{\beta}_E}^{\hat{\beta}}$$

This means from Definition 2.66 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(c \Rightarrow \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(e_s \bullet) \delta^s \Downarrow_i {}^s v$ therefore $\exists j < i < n$ s.t $e_s \delta^s \Downarrow_j {}^s v_1$.

And we have

$$\exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s\theta, n - j, {}^s v_1, {}^t v_1) \in [(c \Rightarrow \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-CE1})$$

From cg-CE we know that ${}^s v_1 = \nu e'_s$ and ${}^t v_1 = \nu e'_t$

Therefore we have

$$({}^s\theta, n - j, \nu e'_s, \nu e'_t) \in [(c \Rightarrow \tau) \sigma]_{\hat{V}}^{\hat{\beta}}$$

This means from Definition 2.65 we have

$$\forall {}^s\theta' \sqsupseteq {}^s\theta'_1, k < n - j, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}_2. ({}^s\theta', k, e'_s, e'_t) \in [\tau \sigma]_{\hat{E}}^{\hat{\beta}_2}$$

Instantiating ${}^s\theta'$ with ${}^s\theta$, k with $n - j - 1$, ℓ' with $\ell \sigma$ and $\hat{\beta}_2$ with $\hat{\beta}$ and we get

$$({}^s\theta, n - j - 1, e'_s, e'_t) \in [\tau \sigma]_{\hat{E}}^{\hat{\beta}}$$

From Definition 2.66 we get

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n - j - 1, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}_2} {}^s\theta'_1 \wedge \forall k < n - j - 1. e'_s \Downarrow_k {}^s v_2 \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

Instantiating with H_s, H'_{t1} . Since from (F-CE1) we know that $(n - j, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$ therefore from Lemma 2.73 we get $(n - j - 1, H_s, H'_{t1}) \triangleright^{\hat{\beta}} {}^s\theta$

Since we know that $e_s \bullet \delta^s \Downarrow_i {}^s v$ and from cg-CE we know that $i = j + k + 1$ (for some k) and $i < n$ therefore we have $k < n - j - 1$ s.t $e'_s \delta^s \Downarrow_k {}^s v_2$.

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow (H'_{t2}, {}^t v_2) \wedge ({}^s\theta, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - i, H_s, H'_{t2}) \triangleright^{\hat{\beta}} {}^s\theta \quad (\text{F-CE2})$$

Since $H'_t = H_{t2'}$, ${}^s v = {}^s v_2$ and ${}^t v = {}^t v_2$ therefore we get (F-CE0) directly from (F-CE2)

14. CF-ret:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{ret}(e_s) : \mathbb{C} \ell_i \ell_i \tau \rightsquigarrow \lambda_. \text{inl}(e_t)} \text{ret}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{ret}(e_s) \delta^s, \lambda_{\cdot}\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_E^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. \text{ret}(e_s) \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{\cdot}\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta \end{aligned}$$

This means that given some $H_s, H_t, \hat{\beta}$ s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n$ s.t $\text{ret}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{\cdot}\text{inl}(e_t)) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$$

From CG-ret and FG-lam we know that $i = 0$, ${}^s v = \text{ret}(e_s) \delta^s$, ${}^t v = \lambda_{\cdot}\text{inl}(e_t) \delta^t$ and $H'_t = H_t$.

So we need to prove

$$({}^s\theta, n, \text{ret}(e_s) \delta^s, \lambda_{\cdot}\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$$

Since we already know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ from the context so we are left with proving

$$({}^s\theta, n, \text{ret}(e_s) \delta^s, \lambda_{\cdot}\text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell_i \ell_i \tau \sigma]_V^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\begin{aligned} & \forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}' \\ & (k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies \exists H'_t, {}^t v'. (H_t, (\lambda_{\cdot}\text{inl}(e_t) ())) \delta^t \Downarrow \\ & (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k. \text{ Also from cg-ret we know that } H'_s = H_s$$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v'. (H_t, (\lambda_{\cdot}\text{inl}(e_t) ())) \delta^t \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{F-R0}) \end{aligned}$$

IIH:

$$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k. e_s \delta^s \Downarrow_f {}^s v \implies \\ & \exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s\theta_e, k - f, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating H_{s1} with H_s and H_{t1} with H_t . And since we know that $(H_s, \text{ret}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists f < i < k \leq n$ s.t $e_s \delta^s \Downarrow_f {}^s v_h$. Therefore we have

$$\exists H'_{t1}, {}^t v. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v) \wedge ({}^s \theta_e, k-f, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \wedge (k-f, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-R1})$$

In order to prove (F-R0) we choose H'_t as H'_{t1} , ${}^t v'$ as $\text{inl}({}^t v)$, ${}^s \theta'$ as ${}^s \theta_e$, $\hat{\beta}''$ as $\hat{\beta}'$. Since from cg-ret we know that $i = f + 1$ therefore from (F-R1) and Lemma 2.73 we know that $(k - i, H_s, H'_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e$

Next we choose ${}^t v''$ as ${}^t v$ (from F-R1) and from Lemma 2.71 we get $({}^s \theta_e, k - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'}$ (we know from cg-ret that ${}^s v' = {}^s v$)

15. CF-bind:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \mathbb{C} \ell_i \ell \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, x : \tau \vdash e_{s2} : \mathbb{C} \ell \ell_o \tau' \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash \text{bind}(e_{s1}, x.e_{s2}) : \mathbb{C} \ell_i \ell_o \tau' \rightsquigarrow \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}())} \text{bind}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_E^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ and given some $i < n, {}^s v$ s.t $\text{bind}(e_{s1}, x.e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge \\ & ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta \end{aligned}$$

From cg-val and fg-val we know that $i = 0, {}^s v = \text{bind}(e_{s1}, x.e_{s2}) \delta^s,$

$${}^t v = \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t, H'_t = H_t$$

And we need to prove

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ from the context so we are left with proving

$$({}^s \theta, n, \text{bind}(e_{s1}, x.e_{s2}) \delta^s, \lambda_.\text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()) \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau') \sigma]_V^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsupseteq \hat{\beta}'.$$

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies$$

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \dots \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \\ & \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau' \sigma]_V^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda \dots \text{case}(e_{t1}(), x.e_{t2}(), y.\text{inr}()))() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - \\ & i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau' \sigma]_V^{\hat{\beta}''} \quad (\text{F-B0}) \end{aligned}$$

IH1:

$$({}^s \theta, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_E^{\hat{\beta}}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_j {}^s v_{h1} \implies \\ & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_V^{\hat{\beta}} \wedge (k - j, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \end{aligned}$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists j < i < k \leq n$ s.t $e_{s1} \delta^s \Downarrow_j {}^s v_{h1}$.

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_i \ell \tau) \sigma]_V^{\hat{\beta}} \wedge (k - j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \quad (\text{F-B1.1}) \end{aligned}$$

From Definition 2.65 we know have

$$\begin{aligned} & \forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k - j, \hat{\beta} \sqsubseteq \hat{\beta}' . \\ & (m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies \\ & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating ${}^s \theta_e$ with ${}^s \theta, H_{s3}$ with H_{s1}, H_{t3} with H'_{t2}, m with $k - j$ and $\hat{\beta}'$ with $\hat{\beta}$. Since we know that $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists b < i - j < k - j$ s.t $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$.

Therefore we have

$$\begin{aligned} & \exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1}()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - j - b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \\ & \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k - j - b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{F-B1}) \end{aligned}$$

IH2:

$$({}^s \theta'', k - j - b, e_{s2} \delta^s \cup \{x \mapsto {}^s v'_{h1}\}, e_{t2} \delta^t \cup \{x \mapsto {}^t v''_{h1}\}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_E^{\hat{\beta}''}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_{s4}, H_{t4}. (k, H_{s4}, H_{t4}) \hat{\triangleright}^{\hat{\beta}''} s\theta \wedge \forall c < (k - j - b), {}^s v_{h2}. e_{s2} \delta^s \Downarrow_j {}^s v_{h2} \implies \\ & \exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k - j - b - c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_{\hat{V}}^{\hat{\beta}''} \wedge (k - \\ & j - b - c, H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}''} s\theta'' \end{aligned}$$

Instantiating H_{s4} with H'_{s3} and H_{t4} with H'_{t3} . And since we know that $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists c < i - j - b < k - j - b$ s.t. $e_{s2} \delta^s \Downarrow_c {}^s v_{h2}$.

Therefore we have

$$\exists H'_{t4}, {}^t v_{h2}. (H_{t4}, e_{t2} \delta^t) \Downarrow (H'_{t4}, {}^t v_{h2}) \wedge ({}^s \theta'', k - j - b - c, {}^s v_{h2}, {}^t v_{h2}) \in [(\mathbb{C} \ell \ell_o \tau') \sigma]_{\hat{V}}^{\hat{\beta}''} \wedge (k - j - b - c, H_{s4}, H'_{t4}) \hat{\triangleright}^{\hat{\beta}''} s\theta'' \quad (\text{F-B2.1})$$

From Definition 2.65 we know have

$$\begin{aligned} & \forall {}^s \theta_e \sqsubseteq {}^s \theta'', H_{s5}, H_{t5}, d, {}^s v'_{h2}, {}^t v'_{h2}, m \leq k - j - b - c, \hat{\beta}'' \sqsubseteq \hat{\beta}''_1. \\ & (m, H_{s5}, H_{t5}) \hat{\triangleright}^{\hat{\beta}''_1} ({}^s \theta_e) \wedge (H_{s5}, {}^s v_{h2}) \Downarrow_d^f (H'_{s5}, {}^s v'_{h2}) \wedge d < m \implies \\ & \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s \theta''' \sqsubseteq {}^s \theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}''_2. (m - d, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} s\theta''' \wedge \\ & \exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', m - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2} \end{aligned}$$

Instantiating ${}^s \theta_e$ with ${}^s \theta''$, H_{s5} with H'_{s3} , H_{t5} with H'_{t3} , m with $k - j - b - c$ and $\hat{\beta}''$ with $\hat{\beta}''$. Since we know that $(H_{s1}, \text{bind}(e_{s1}, x.e_{s2}) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists d < i - j - b - c < k - j - b - c$ s.t. $(H'_{s3}, {}^s v_{h2}) \delta^s \Downarrow_d (H'_{s5}, {}^s v'_{h2})$.

Therefore we have

$$\begin{aligned} & \exists H'_{t5}, {}^t v'_{h2}. (H_{t5}, {}^t v_{h2}()) \Downarrow (H'_{t5}, {}^t v'_{h2}) \wedge \exists {}^s \theta''' \sqsubseteq {}^s \theta_e, \hat{\beta}'' \sqsubseteq \hat{\beta}''_2. (k - j - b - c - d, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} s\theta''' \wedge \\ & \exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2} \quad (\text{F-B2}) \end{aligned}$$

In order to prove (F-B0) we choose H'_{t1} as H'_{t5} and ${}^t v'$ as ${}^t v'_{h2}$. Next we choose ${}^s \theta'$ as ${}^s \theta'''$ and $\hat{\beta}''$ as $\hat{\beta}''_2$ (both chosen from (F-B2)). Also from cg-bind we know that in (F-B0) H'_{s1} will be H'_{s5} .

Since $(k - j - b - c - d, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} s\theta'''$ therefore Lemma 2.71 we get $(k - i, H'_{s5}, H'_{t5}) \hat{\triangleright}^{\hat{\beta}''_2} s\theta'''$

Also since from (F-B2) we have $\exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', k - j - b - c - d, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2}$

Since $i = j + b + c + d + 1$ therefore from Lemma 2.71 we get

$$\exists {}^t v''_{h2}. {}^t v'_{h2} = \text{inl } {}^t v''_{h2} \wedge ({}^s \theta''', k - i, {}^s v'_{h2}, {}^t v''_{h2}) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}''_2}$$

16. CF-label:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{Lb}_\ell(e_s) : (\text{Labeled } \ell \tau) \rightsquigarrow \text{inl}(e_t)} \text{label}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{Lb}_\ell(e_s) \delta^s, \text{inl}(e_t) \delta^t) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{E}}^{\hat{\beta}}$

From Definition 2.66 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^s \theta \wedge \forall i < n, {}^s v. \text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \end{aligned}$$

This means that we are given some $H_s, H_t, \hat{\beta}$ s.t. $(n, H_s, H_t) \hat{\triangleright}^s \theta$ and given some $i < n$ s.t. $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$.

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{inl}(e_t) \delta^t) \Downarrow (H'_t, \text{inl}({}^t v)) \wedge ({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \hat{\triangleright}^s \theta \quad (\text{F-LB0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

From Definition 2.66 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright}^s \theta \wedge \forall j < n, {}^s v_1. e_s \delta^s \Downarrow_j {}^s v_1 \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \hat{\triangleright}^s \theta \end{aligned}$$

Instantiating with H_s, H_t and since we know that $\text{Lb}_\ell(e_s) \delta^s \Downarrow_i \text{Lb}_\ell({}^s v)$ therefore $\exists j < i < n$ s.t. $e_s \delta^s \Downarrow_j {}^s v$.

Therefore we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v, {}^t v) \in [(\tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n - j, H_s, H'_{t1}) \hat{\triangleright}^s \theta \quad (\text{F-LB1})$$

Since from (F-LB0) we are required to prove $({}^s \theta, n - i, \text{Lb}_\ell({}^s v), \text{inl}({}^t v)) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}}$. Since from cg-label we know that $i = j + 1$, ${}^s v = {}^s v_1$ and ${}^t v = {}^t v_1$. Therefore we get this from Definition 2.65, (F-LB1) and Lemma 2.71.

From Lemma 2.71 we get $(n - i, H_s, H'_{t1}) \hat{\triangleright}^s \theta$

17. CF-toLabeled:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \mathbb{C} \ell_i \ell_o \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{toLabeled}(e_s) : \mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau) \rightsquigarrow \lambda_. \text{inl}(e_t ())} \text{toLabeled}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_. \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{E}}^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. \text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, (\lambda_ \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ and given some $i < n$ s.t $\text{toLabeled}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\lambda_ \text{inl } e_t()) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge \\ & (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} s\theta \end{aligned}$$

From cg-val and fg-val we know that $i = 0, {}^s v = \text{toLabeled}(e_s) \delta^s,$
 ${}^t v = (\lambda_ \text{inl } e_t()) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_ \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$$

Since we already know $(n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta$ from the context so we are left with proving

$$({}^s \theta, n, \text{toLabeled}(e_s) \delta^s, (\lambda_ \text{inl } e_t()) \delta^t) \in [(\mathbb{C} \ell_i \ell_i (\text{Labeled } \ell_o \tau)) \sigma]_{\hat{V}}^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies$$

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_ \text{inl } e_t()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell_o \tau) \sigma]_{\hat{V}}^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_ \text{inl } e_t()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ & \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell_o \tau) \sigma]_{\hat{V}}^{\hat{\beta}''} \quad (\text{F-TL0}) \end{aligned}$$

IH:

$$({}^s \theta, k, e_s \delta^s, e_t \delta^t) \in [(\mathbb{C} \ell_i \ell_o \tau) \sigma]_{\hat{E}}^{\hat{\beta}}$$

It means from Definition 2.66 that we need to prove

$$\forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_{h1}. e_s \delta^s \Downarrow_j {}^s v_{h1} \implies$$

$$\begin{aligned} & \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k - j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_i \ell_o \tau) \sigma]_{\hat{V}}^{\hat{\beta}} \wedge (k - j, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \\ & {}^s \theta \end{aligned}$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists j < i < k \leq n$ s.t $e_s \delta^s \Downarrow_j {}^s v_{h1}$.

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta, k-j, {}^s v_{h1}, {}^t v_{h1}) \in [(\mathbb{C} \ell_i \ell_o \tau) \sigma]_{V}^{\hat{\beta}} \wedge (k-j, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}} {}^s \theta \quad (\text{F-TL1.1})$$

From Definition 2.65 we know have

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s3}, H_{t3}, b, {}^s v'_{h1}, {}^t v'_{h1}, m \leq k-j, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(m, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, {}^s v_{h1}) \Downarrow_b^f (H'_{s3}, {}^s v'_{h1}) \wedge b < m \implies$$

$$\exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (m-b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', m-b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_{V}^{\hat{\beta}''}$$

Instantiating ${}^s \theta_e$ with ${}^s \theta$, H_{s3} with H_{s1} , H_{t3} with H'_{t2} , m with $k-j$ and $\hat{\beta}'$ with $\hat{\beta}$. Since we know that $(H_{s1}, \text{toLabeled}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists b < i-j < k-j$ s.t $(H_{s1}, {}^s v_{h1}) \delta^s \Downarrow_b (H'_{s3}, {}^s v'_{h1})$.

Therefore we have

$$\exists H'_{t3}, {}^t v'_{h1}. (H_{t3}, {}^t v_{h1} ()) \Downarrow (H'_{t3}, {}^t v'_{h1}) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k-j-b, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge \exists {}^t v''_{h1}. {}^t v'_{h1} = \text{inl } {}^t v''_{h1} \wedge ({}^s \theta'', k-j-b, {}^s v'_{h1}, {}^t v''_{h1}) \in [\tau \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-TL1})$$

In order to prove (F-TL0) we choose ${}^s \theta'$ as ${}^s \theta''$ and $\hat{\beta}'$ as $\hat{\beta}''$ (both chosen from (F-TL2))

Also from cg-toLabeled and fg-inl, fg-app we know that $H'_s = H'_{s3}$ and $H'_t = H'_{t3}$, and ${}^s v' = {}^s v'_{h1}$, ${}^t v' = {}^t v'_{h1}$

Therefore we get the desired from (F-TL1) and Lemma 2.71

18. CF-unlabel:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{Labeled } \ell \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash \text{unlabel}(e_s) : \mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \rightsquigarrow \lambda_{\cdot} e_t} \text{ unlabel}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{\cdot} e_t \delta^t) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{E}^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. \text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, \lambda_{\cdot} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n-i, {}^s v, {}^t v) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{V}^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ and given some $i < n, {}^s v$ s.t $\text{unlabel}(e_s) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{\cdot} e_t \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n-i, {}^s v, {}^t v) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{V}^{\hat{\beta}} \wedge (n-i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

From cg-val and fg-val we know that $i = 0$, ${}^s v = \text{unlabel}(e_s) \delta^s$, ${}^t v = \lambda_{\cdot} e_t \delta^t$, $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, {}^s v, {}^t v) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{\hat{\beta}}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ from the context so we are left with proving

$$({}^s \theta, n, \text{unlabel}(e_s) \delta^s, \lambda_{\cdot} e_t \delta^t) \in [\mathbb{C} \ell_i (\ell_i \sqcup \ell) \tau \sigma]_{\hat{\beta}}^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\cdot} e_t)() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_{\hat{\beta}''}^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{\cdot} e_t)() \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \\ \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_{\hat{\beta}''}^{\hat{\beta}''} \quad (\text{F-U0}) \end{aligned}$$

III:

$$({}^s \theta_e, k, e_s \delta^s, e_t \delta^t) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{E}}^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_h . e_s \delta^s \Downarrow_f {}^s v_h \implies \\ \exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}'} \wedge (k - \\ f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \end{aligned}$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, \text{unlabel}(e_s) \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists f < i < k \leq n$ s.t $e_s \delta^s \Downarrow_f {}^s v_h$.

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau) \sigma]_{\hat{V}}^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-U1})$$

In order to prove (F-U0) we choose H'_{t1} as H'_{t2} , ${}^t v'$ as ${}^t v_h$, ${}^s \theta'$ as ${}^s \theta_e$ and $\hat{\beta}''$ as $\hat{\beta}'$

From cg-unlabel and fg-app we also know that $H'_{s1} = H_{s1}$ and $H'_{t1} = H'_{t2}$

We need to prove

(a) $(k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$:

Since from (F-U1) we know that $(k - f, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$

Therefore from Lemma 2.73 we also get $(k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$

(b) $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_e, k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}'}$:

Since from (F-U1) we have

$({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}'}$

This means from Definition 2.65 we know that

$\exists {}^s v_i, {}^t v_i . {}^s v_h = \text{Lb}_\ell({}^s v_i) \wedge {}^t v_h = \text{inl } {}^t v_i \wedge ({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau \sigma]_V^{\hat{\beta}'}$ (F-U2)

Since we know that ${}^t v' = {}^t v_h$ and since from (F-U2) we have ${}^t v_h = \text{inl } {}^t v_i$. Therefore from we choose ${}^t v''$ as ${}^t v_i$ to get the first conjunct

From cg-unlabel we know that ${}^s v = {}^s v_i$ and since we know that $({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau \sigma]_V^{\hat{\beta}'}$

Therefore from Lemma 2.71 we also get $({}^s \theta_e, k - i, {}^s v_i, {}^t v_i) \in [\tau \sigma]_V^{\hat{\beta}'}$

19. CF-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{Labeled } \ell' \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash \text{new } e_s : \mathbb{C} \ell \ell (\text{ref } \ell' \tau) \rightsquigarrow \lambda_{-}.\text{inl}(\text{new } (e_t))} \text{ref}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{new } e_s \delta^s, \lambda_{-}.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^s v. \text{new } e_s \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(\text{new } (e_t)) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright}^{s\theta} \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$ and given some $i < n, {}^s v$ s.t $\text{new } e_s \delta^s \Downarrow_i {}^s v$

From cg-val and fg-val we know that $i = 0, {}^s v = \text{new } e_s \delta^s, {}^t v = \lambda_{-}.\text{inl}(\text{new } (e_t)) \delta^t, H'_t = H_t$

And we need to prove

$$({}^s \theta, n, \text{new } e_s \delta^s, \lambda_{-}.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}} \wedge (n, H_s, H_t) \hat{\triangleright}^{s\theta}$$

Since we already know $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$ from the context so we are left with proving

$$({}^s \theta, n, \text{new } e_s \delta^s, \lambda_{-}.\text{inl}(\text{new } (e_t)) \delta^t) \in [\mathbb{C} \ell \ell (\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^sv', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$

$(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_{s1}, \text{new } e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^sv') \wedge i < k \implies$
 $\exists H'_{t1}, {}^tv'. (H_{t1}, (\lambda_.\text{inl}(\text{new } e_t))(\delta^t)) \Downarrow (H'_{t1}, {}^tv') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}''.(k-i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge$
 $\exists {}^tv''. {}^tv' = \text{inl } {}^tv'' \wedge ({}^s\theta', k-i, {}^sv', {}^tv'') \in \llbracket (\text{ref } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}''}$

This means we are given some ${}^s\theta_e \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i, {}^sv', {}^tv', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \wedge (H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_{s1}, {}^sv') \wedge i < k.$

And we need to prove

$\exists H'_{t1}, {}^tv'. (H_{t1}, (\lambda_.\text{inl}(\text{new } e_t))(\delta^t)) \Downarrow (H'_{t1}, {}^tv') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}''.(k-i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s\theta' \wedge$
 $\exists {}^tv''. {}^tv' = \text{inl } {}^tv'' \wedge ({}^s\theta', k-i, {}^sv', {}^tv'') \in \llbracket (\text{ref } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}''} \quad (\text{F-N0})$

From cg-ref we know that ${}^sv' = a_s$ and from fg-ref, fg-inl we know that ${}^tv' = \text{inl } a_t.$

IH:

$({}^s\theta_e, k, e_s \delta^s, e_t \delta^t) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_E^{\hat{\beta}'}$

It means from Definition 2.66 that we need to prove

$\forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \wedge \forall f < k, {}^sv_h.e_s \delta^s \Downarrow_f {}^sv_h \implies$
 $\exists H'_{t2}, {}^tv_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^tv_h) \wedge ({}^s\theta_e, k-f, {}^sv_h, {}^tv_h) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k-f,$
 $H_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, \text{new } (e_s) \delta^s) \Downarrow_i^f (H'_s, {}^sv')$ therefore $\exists f < i < k \leq n$ s.t $e_s \delta^s \Downarrow_f {}^sv_h.$

Therefore we have

$\exists H'_{t2}, {}^tv_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^tv_h) \wedge ({}^s\theta_e, k-f, {}^sv_h, {}^tv_h) \in \llbracket (\text{Labeled } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k-f,$
 $H_{s1}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-N1})$

In order to prove (F-N0) we choose H'_{t1} as $H'_{t2} \cup \{a_t \mapsto {}^tv_h\}$, tv as a_t , ${}^s\theta'$ as ${}^s\theta_n$ where ${}^s\theta_n = {}^s\theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau) \sigma\}$

And we choose $\hat{\beta}''$ as $\hat{\beta}_n$ where $\hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$

From cg-ref and fg-ref we also know that $H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^sv_h\}$

We need to prove

(a) $(k-i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}_n} {}^s\theta_n:$

From Definition 2.67 it suffices to prove that

- $\text{dom}({}^s\theta_n) \subseteq \text{dom}(H'_{s1}):$

Since $\text{dom}({}^s\theta_e) \subseteq \text{dom}(H_{s1})$ (given that we have $(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$)

And since we know that

${}^s\theta_n = {}^s\theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau) \sigma\}$ and $H'_{s1} = H_{s1} \cup \{a_s \mapsto {}^sv_h\}$

Therefore we get $\text{dom}({}^s\theta_n) \subseteq \text{dom}(H'_{s1})$

- $\hat{\beta}_n \subseteq (\text{dom}({}^s\theta_n) \times \text{dom}(H'_{t1}))$:

Since $\hat{\beta}' \subseteq (\text{dom}({}^s\theta_e) \times \text{dom}(H_{t1}))$ (given that we have $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$)

And since we know that

${}^s\theta_n = {}^s\theta_e \cup \{a_s \mapsto (\text{Labeled } \ell' \tau) \sigma\}$, $H'_{t1} = H_{t1} \cup \{a_t \mapsto {}^t v_h\}$ and $\hat{\beta}_n = \hat{\beta}' \cup \{(a_s, a_t)\}$

Therefore we get $\hat{\beta}_n \subseteq (\text{dom}({}^s\theta_n) \times \text{dom}(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}_n. ({}^s\theta_n, k - i - 1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s\theta_n(a)]_V^{\hat{\beta}_n}$:

$\forall (a_1, a_2) \in \hat{\beta}_n$

– $(a_1, a_2) = (a_s, a_t)$:

Since from (F-N1) we know that $({}^s\theta_e, k - f, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_V^{\hat{\beta}'}$

From Lemma 2.71 we get $({}^s\theta_n, k - i - 1, {}^s v_h, {}^t v_h) \in [(\text{Labeled } \ell' \tau) \sigma]_V^{\hat{\beta}_n}$

– $(a_1, a_2) \neq (a_s, a_t)$:

Since we have $(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s\theta_e$ therefore
from Definition 2.67 we get

$({}^s\theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_V^{\hat{\beta}'}$

From Lemma 2.71 we get

$({}^s\theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_n(a_1)]_V^{\hat{\beta}'}$

- (b) $\exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_n, k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}_n}$:

We choose ${}^t v''$ as ${}^t v_h$ from (F-N1), fg-inl and fg-ref we know that ${}^t v' = \text{inl } {}^t v_h$

In order to prove $({}^s\theta_n, k - i, {}^s v', {}^t v'') \in [(\text{ref } \ell' \tau) \sigma]_V^{\hat{\beta}_n}$, from Definition 2.65 it suffices to prove that

${}^s\theta_n(a_s) = (\text{Labeled } \ell' \tau) \sigma \wedge (a_s, a_t) \in \hat{\beta}_n$

We get this by construction of ${}^s\theta_n$ and $\hat{\beta}_n$

20. CF-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_s : \text{ref } \ell \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash !e_s : \mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \rightsquigarrow \lambda _ . \text{inl}(e_t)} \text{deref}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, !e_s \delta^s, \lambda _ . \text{inl}(e_t) \delta^t) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_E^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. !e_s \delta^s \Downarrow_i {}^s v \implies$

$\exists H'_t. {}^t v. (H_t, \lambda _ . \text{inl}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s\theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell' \ell' (\text{Labeled } \ell \tau) \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s\theta$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n$ s.t

$!e_s \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \lambda_{-} \text{inl}(e_t) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in \llbracket \mathbb{C} \ell' \ell' \text{ (Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \triangleright^{\hat{\beta}} {}^s \theta$$

From cg-val and fg-val we know that $i = 0$, ${}^s v = !e_s \delta^s$, ${}^t v = \lambda_{-} \text{inl}(e_t) \delta^t$, $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, {}^s v, {}^t v) \in \llbracket \mathbb{C} \ell' \ell' \text{ (Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ from the context so we are left with proving

$$({}^s \theta, n, !e_s \delta^s, \lambda_{-} \text{inl}(e_t) \delta^t) \in \llbracket \mathbb{C} \ell' \ell' \text{ (Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies$$

$$\exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} \text{inl}(e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket \text{(Labeled } \ell' \tau) \sigma \rrbracket_V^{\hat{\beta}''}$$

This means we are given some ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', {}^t v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-} \text{inl}(e_t)) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in \llbracket \text{(Labeled } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}''} \quad (\text{F-D0})$$

IH:

$$({}^s \theta_e, k, e_s \delta^s, e_t \delta^t) \in \llbracket \text{(ref } \ell \tau) \sigma \rrbracket_E^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_h . e_s \delta^s \Downarrow_f {}^s v_h \implies$$

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket \text{(ref } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, !e_s \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists f < i < k \leq n$ s.t $e_s \delta^s \Downarrow_f {}^s v_h$.

Therefore we have

$$\exists H'_{t2}, {}^t v_h. (H_{t2}, e_t \delta^t) \Downarrow (H'_{t2}, {}^t v_h) \wedge ({}^s \theta_e, k - f, {}^s v_h, {}^t v_h) \in \llbracket \text{(ref } \ell \tau) \sigma \rrbracket_V^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-D1})$$

In order to prove (F-D0) we choose H'_{t1} as H'_{t2} , ${}^t v'_1$ as $H'_{t2}(a)$ (where ${}^t v_h = a_t$ from fg-deref), ${}^s \theta'$ as ${}^s \theta_e$ and we choose $\hat{\beta}''$ as $\hat{\beta}'$.

From cg-deref we also know that $H'_{s1} = H_{s1}$

We need to prove

(a) $(k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$:

Since from (F-D1) we have $(k - f, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$ and since $f < i$ therefore from Lemma 2.73 we get $(k - i, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$

(b) $\exists^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta_e, k - i, {}^s v', {}^t v'') \in [(\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}'}$:

Since from cg-deref and fg-deref we know that ${}^s v_h = a_s$ and ${}^t v_h = a_t$.

Therefore from (F-D1) and from Definition 2.65 we know that

$${}^s \theta_e(a_s) = (\text{Labeled } \ell \tau) \sigma \wedge (a_s, a_t) \in \hat{\beta}'$$

Since from (F-D1) we know that $(k - f, H_{s1}, H'_{t2}) \hat{\triangleright}^{s\theta_e}$ which means from Definition 2.67 we know that

$$({}^s \theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-D2})$$

This means from Definition 2.65 we know that

$$\exists^s v_i, {}^t v_i. H_{s1}(a_s) = \text{Lb}_{\ell}({}^s v_i) \wedge H'_{t2}(a_t) = \text{inl } {}^t v_i \wedge ({}^s \theta_e, k - f - 1, {}^s v_i, {}^t v_i) \in [\tau \sigma]_{V}^{\hat{\beta}'}$$

We choose ${}^t v''$ as ${}^t v_i$ and we know that ${}^t v' = H'_{t2}(a_t) = \text{inl } {}^t v_i$. This proves the first conjunct.

Since from (F-D2) we have $({}^s \theta, k - f - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}'}$ therefore from Lemma 2.71 we get

$$({}^s \theta, k - i - 1, H_{s1}(a_s), H'_{t2}(a_t)) \in [(\text{Labeled } \ell \tau) \sigma]_{V}^{\hat{\beta}'}$$

This proves the second conjunct.

21. CF-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash e_{s1} : \text{ref } \ell' \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash e_{s2} : \text{Labeled } \ell' \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi; \Gamma \vdash e_{s1} := e_{s2} : \mathbb{C} \ell \ell \text{ unit} \rightsquigarrow \lambda_{-}.\text{inl}(e_{t1} := e_{t2})} \text{ assign}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove: $({}^s \theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{E}^{\hat{\beta}}$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^s v. (e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v \implies \\ & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{V}^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright}^{s\theta} \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$ and given some $i < n, {}^s v$ s.t $(e_{s1} := e_{s2}) \delta^s \Downarrow_i {}^s v$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{V}^{\hat{\beta}} \wedge (n - \\ & i, H_s, H'_t) \hat{\triangleright}^{s\theta} \end{aligned}$$

From cg-val and fg-val we know that $i = 0$, ${}^s v = (e_{s1} := e_{s2}) \delta^s$, ${}^t v = \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t$, $H'_t = H_t$

And we need to prove

$$({}^s \theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{V}^{\hat{\beta}} \wedge (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$$

Since we already know $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$ from the context so we are left with proving

$$({}^s \theta, n, (e_{s1} := e_{s2}) \delta^s, \lambda_{-}.\text{inl}(e_{t1} := e_{t2}) \delta^t) \in [\mathbb{C} \ell \ell \text{ unit } \sigma]_{V}^{\hat{\beta}}$$

From Definition 2.65 it means we need to prove

$$\forall {}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$\begin{aligned} (k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k \implies \\ \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_{t1} := e_{t2})()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} \\ {}^s \theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_{V}^{\hat{\beta}''} \end{aligned}$$

This means we are given some ${}^s \theta_e \sqsupseteq {}^s \theta, H_{s1}, H_{t1}, i, {}^s v', k \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t

$$(k, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge (H_{s1}, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i^f (H'_{s1}, {}^s v') \wedge i < k.$$

And we need to prove

$$\begin{aligned} \exists H'_{t1}, {}^t v'. (H_{t1}, (\lambda_{-}.\text{inl}(e_{t1} := e_{t2})()) \delta^t) \Downarrow (H'_{t1}, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . \\ (k - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\text{unit}]_{V}^{\hat{\beta}''} \quad (\text{F-S0}) \end{aligned}$$

IH1:

$$({}^s \theta_e, k, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\text{ref } \ell' \tau) \sigma]_E^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} \forall H_{s2}, H_{t2}. (k, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \wedge \forall f < k, {}^s v_{h1}. e_{s1} \delta^s \Downarrow_f {}^s v_{h1} \implies \\ \exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta_e, k - f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k - \\ f, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \end{aligned}$$

Instantiating H_{s2} with H_{s1} and H_{t2} with H_{t1} . And since we know that $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists f < i < k \leq n$ s.t $e_s \delta^s \Downarrow_f {}^s v_{h1}$.

Therefore we have

$$\exists H'_{t2}, {}^t v_{h1}. (H_{t2}, e_{t1} \delta^t) \Downarrow (H'_{t2}, {}^t v_{h1}) \wedge ({}^s \theta_e, k - f, {}^s v_{h1}, {}^t v_{h1}) \in [(\text{ref } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k - f, H_{s1}, H'_{t2}) \triangleright^{\hat{\beta}'} {}^s \theta_e \quad (\text{F-S1})$$

IH2:

$$({}^s \theta_e, k - f, e_{s2} \delta^s, e_{t2} \delta^t) \in [(\text{Labeled } \ell' \tau) \sigma]_E^{\hat{\beta}'}$$

It means from Definition 2.66 that we need to prove

$$\begin{aligned} & \forall H_{s3}, H_{t3}. (k, H_{s3}, H_{t3}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \wedge \forall l < k - f, {}^s v_{h2}. e_{s2} \delta^s \Downarrow_l {}^s v_{h2} \implies \\ & \exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k - \\ & f - l, H_{s3}, H'_{t3}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \end{aligned}$$

Instantiating H_{s3} with H_{s1} and H_{t3} with H'_{t2} . And since we know that $(H_{s1}, e_{s1} := e_{s2} \delta^s) \Downarrow_i^f (H'_s, {}^s v')$ therefore $\exists l < i - f < k - f$ s.t. $e_{s2} \delta^s \Downarrow_l {}^s v_{h2}$.

Therefore we have

$$\exists H'_{t3}, {}^t v_{h2}. (H_{t3}, e_{t2} \delta^t) \Downarrow (H'_{t3}, {}^t v_{h2}) \wedge ({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'} \wedge (k - f - l, H_{s1}, H'_{t3}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e \quad (\text{F-S2})$$

In order to prove (F-S0) we choose H'_{t1} as $H'_{t3}[a_t \mapsto {}^t v_{h3}]$, ${}^t v'$ as $()$, ${}^s\theta'$ as ${}^s\theta_e$ and $\hat{\beta}''$ as $\hat{\beta}'$. From cg-assign and fg-assign we also know that ${}^s v_{h2} = a_s$, ${}^t v_{h2} = a_t$, $H'_{s1} = H_{s1}[a_s \mapsto {}^s v_{h3}]$ and $H'_{t1} = H'_{t3}[a_t \mapsto {}^t v_{h3}]$

We need to prove

$$(a) (k - i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e:$$

From Definition 2.67 it suffices to prove that

- $dom({}^s\theta_e) \subseteq dom(H'_{s1})$:

Since $dom({}^s\theta_e) \subseteq dom(H_{s1})$ (given that we have $(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$)

And since $dom(H_{s1}) = dom(H'_{s1})$ therefore we also get

$$dom({}^s\theta_e) \subseteq dom(H'_{s1})$$

- $\hat{\beta}' \subseteq (dom({}^s\theta_e) \times dom(H'_{t1}))$:

Since $\hat{\beta}' \subseteq (dom({}^s\theta_e) \times dom(H_{t1}))$ (given that we have $(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$)

And since $dom(H_{t1}) \subseteq dom(H'_{t1})$ therefore we also have $\hat{\beta}' \subseteq (dom({}^s\theta_e) \times dom(H'_{t1}))$

- $\forall (a_1, a_2) \in \hat{\beta}'. ({}^s\theta_e, k - i - 1, H'_{s1}(a_1), H'_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_{V}^{\hat{\beta}'}$:

$$\forall (a_1, a_2) \in \hat{\beta}'_n$$

- $(a_1, a_2) = (a_s, a_t)$:

Since from (F-S2) we know that $({}^s\theta_e, k - f - l, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'}$

From Lemma 2.71 we get $({}^s\theta_e, k - i - 1, {}^s v_{h2}, {}^t v_{h2}) \in [(\text{Labeled } \ell' \tau) \sigma]_{V}^{\hat{\beta}'}$

- $(a_1, a_2) \neq (a_s, a_t)$:

Since we have $(k, H_{s1}, H_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta_e$ therefore

from Definition 2.67 we get

$$({}^s\theta_e, k - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_{V}^{\hat{\beta}'}$$

From Lemma 2.71 we get

$$({}^s\theta_n, k - i - 1, H_{s1}(a_1), H_{t1}(a_2)) \in [{}^s\theta_e(a_1)]_{V}^{\hat{\beta}'}$$

$$(b) \exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta_e, k - i, {}^s v', {}^t v'') \in [\text{unit}]_{V}^{\hat{\beta}'_n}:$$

We choose ${}^t v''$ as $()$ from (F-S1), fg-inl and fg-assign we know that ${}^t v' = \text{inl } ()$

To prove: $({}^s\theta_n, k - i, (), ()) \in [\text{unit}]_{V}^{\hat{\beta}'_n}$,

We get this directly from Definition 2.65

□

Lemma 2.75 (CG \rightsquigarrow FG: Subtyping). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \tau, \tau'$.

1. $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^{\hat{\beta}} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}}$
2. $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^{\hat{\beta}} \subseteq [(\tau' \sigma)]_E^{\hat{\beta}}$

Proof. Proof of Statement (1)

Proof by induction on $\tau <: \tau'$

1. CGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau'_2}$$

To prove: $[((\tau_1 \rightarrow \tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^{\hat{\beta}}$

It suffices to prove: $\forall ({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^{\hat{\beta}}$

This means that given some ${}^s\theta, n$ and $\lambda x.e_i$ s.t $({}^s\theta, n, \lambda x.e_i) \in [((\tau_1 \rightarrow \tau_2) \sigma)]_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$\forall {}^s\theta' \sqsupseteq {}^s\theta, {}^s v, {}^t v, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$.

$({}^s\theta', j, {}^s v, {}^t v) \in [\tau_1]_V^{\hat{\beta}'} \implies ({}^s\theta', j, e_s[{}^s v/x], e_t[{}^t v/x]) \in [\tau_2]_E^{\hat{\beta}'}$ (S-A0)

And it suffices to prove: $({}^s\theta, n, \lambda x.e_i) \in [((\tau'_1 \rightarrow \tau'_2) \sigma)]_V^{\hat{\beta}}$

Again from Definition 2.65 it suffices to prove:

$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$.

$({}^s\theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}'_1} \implies ({}^s\theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau'_2]_E^{\hat{\beta}'_1}$

This means that given some ${}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$ s.t $({}^s\theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau'_1]_V^{\hat{\beta}'_1}$

And we are required to prove: $({}^s\theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau'_2]_E^{\hat{\beta}'_1}$

IH: $[(\tau'_1 \sigma)]_V^{\hat{\beta}'_1} \subseteq [(\tau_1 \sigma)]_V^{\hat{\beta}'_1}$ (Statement (1))

$[(\tau_2 \sigma)]_E^{\hat{\beta}'_1} \subseteq [(\tau'_2 \sigma)]_E^{\hat{\beta}'_1}$ (Sub-A0, From Statement (2))

Instantiating (S-A0) with ${}^s\theta'_1, {}^s v_1, {}^t v_1, k, \hat{\beta}'_1$

Since $({}^s\theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_1}$ therefore from IH1 we know that $({}^s\theta'_1, k, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1}$

As a result we get

$({}^s\theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1}$

From (Sub-A0), we know that

$({}^s\theta'_1, k, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_1}$

2. CGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2}$$

To prove: $\llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

IH1: $\llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}}$ (Statement (1))

IH2: $\llbracket (\tau_2 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V^{\hat{\beta}}$ (Statement (1))

It suffices to prove:

$\forall ({}^s\theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

This means that given $({}^s\theta, n, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$$({}^s\theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-P0})$$

And it suffices to prove: $({}^s\theta, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

Again from Definition 2.65, it suffices to prove:

$$({}^s\theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}}$$

Since from (S-P0) we know that $({}^s\theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}}$ therefore from IH1 we have $({}^s\theta, n, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \sigma \rrbracket_V^{\hat{\beta}}$

Similarly since from (S-P0) we have $({}^s\theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}}$ therefore from IH2 we get $({}^s\theta, n, {}^s v_2, {}^t v_2) \in \llbracket \tau'_2 \sigma \rrbracket_V^{\hat{\beta}}$

3. CGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2}$$

To prove: $\llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

IH1: $\llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}}$ (Statement (1))

IH2: $\llbracket (\tau_2 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V^{\hat{\beta}}$ (Statement (1))

It suffices to prove: $\forall ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

This means that given: $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

And it suffices to prove: $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

2 cases arise

(a) ${}^s v = \text{inl } {}^s v_i$ and ${}^t v = \text{inl } {}^t v_i$:

From Definition 2.65 we are given:

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau_1 \sigma]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

From (S-S0) and IH1 we get

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau'_1 \sigma]_V^{\hat{\beta}}$$

(b) ${}^s v = \text{inr } {}^s v_i$ and ${}^t v = \text{inr } {}^t v_i$:

Symmetric reasoning

4. CGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2}$$

To prove: $[(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}} \subseteq [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

It suffices to prove: $\forall ({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}}. ({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

This means that given: $({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_1) \sigma]_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$$\forall {}^s \theta' \sqsupseteq {}^s \theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s \theta', j, e_s, e_t) \in [\tau_1[\ell'/\alpha] \sigma]_E^{\hat{\beta}'} \quad (\text{S-F0})$$

And it suffices to prove: $({}^s \theta, n, \Lambda e_s, \Lambda e_t) \in [(\forall \alpha. \tau_2) \sigma]_V^{\hat{\beta}}$

Again from Definition 2.65, it suffices to prove:

$$\forall {}^s \theta'_1 \sqsupseteq {}^s \theta, k < n, \ell'_1 \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

This means that given ${}^s \theta_1 \sqsupseteq {}^s \theta, k < n, \ell'_1 \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we are required to prove: $({}^s \theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$

Instantiating (S-F0) with ${}^s \theta_1, k, \ell'_1, \hat{\beta}'_1$ we get

$$({}^s \theta'_1, k, e_s, e_t) \in [\tau_1[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

$$[(\tau_1 (\sigma \cup [\alpha \mapsto \ell'_1]))]_E^{\hat{\beta}'_1} \subseteq [(\tau_2 (\sigma \cup [\alpha \mapsto \ell'_1]))]_E^{\hat{\beta}'_1} \quad (\text{Sub-F0, Statement (2)})$$

From (Sub-F0), we know that

$$({}^s \theta'_1, k, e_s, e_t) \in [\tau_2[\ell'_1/\alpha] \sigma]_E^{\hat{\beta}'_1}$$

5. CGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2}$$

To prove: $\llbracket ((c_1 \Rightarrow \tau_1) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((c_2 \Rightarrow \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

It suffices to prove: $\forall ({}^s\theta, n, \nu e_s, \nu e_t) \in \llbracket ((c_1 \Rightarrow \tau_1) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, \nu e_s, \nu e_t) \in \llbracket ((c_2 \Rightarrow \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

This means that given: $({}^s\theta, n, \nu e_s, \nu e_t) \in \llbracket ((c_1 \Rightarrow \tau_1) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$$\mathcal{L} \models c_1 \sigma \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'} \quad (\text{S-C0})$$

And it suffices to prove: $({}^s\theta, n, \nu e_s, \nu e_t) \in \llbracket ((c_2 \Rightarrow \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

Again from Definition 2.65, it suffices to prove:

$$\mathcal{L} \models c_2 \sigma \implies \forall {}^s\theta'_1 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

This means that given $\mathcal{L} \models c_2, {}^s\theta'_1 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we are required to prove:

$$({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

since we know that $c_2 \implies c_1$ and since $\mathcal{L} \models c_2 \sigma$ therefore $\mathcal{L} \models c_1 \sigma$. Next we instantiate (S-C0) with ${}^s\theta'_1, k, \hat{\beta}'_1$ to get

$$({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

$$\llbracket (\tau_1 \sigma) \rrbracket_E^{\hat{\beta}'_1} \subseteq \llbracket (\tau_2 \sigma) \rrbracket_E^{\hat{\beta}'_1} \quad (\text{Sub-C0, Statement (2)})$$

Therefore from (Sub-C0), we get

$$({}^s\theta'_1, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

6. CGsub-label:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell'}{\Sigma; \Psi \vdash \text{Labeled } \ell \tau <: \text{Labeled } \ell' \tau'}$$

To prove: $\llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V^{\hat{\beta}}$

IH: $\llbracket (\tau \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau' \sigma) \rrbracket_V^{\hat{\beta}}$ (Statement (1))

It suffices to prove:

$$\forall ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\text{Labeled } \ell' \tau') \sigma) \rrbracket_V^{\hat{\beta}}$$

This means that given some $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\text{Labeled } \ell \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 2.65 we are given:

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, m, {}^s v', {}^t v') \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-L0})$$

And we are required to prove that

$$({}^s\theta, n, {}^s v, {}^t v) \in [((\text{Labeled } \ell' \tau') \sigma)]_V^{\hat{\beta}}$$

From Definition 2.65 it suffices to prove

$$\exists {}^s v', {}^t v'. {}^s v = \text{Lb}_\ell({}^s v') \wedge {}^t v = \text{inl } {}^t v' \wedge ({}^s\theta, m, {}^s v', {}^t v') \in [\tau' \sigma]_V^{\hat{\beta}}$$

We get this directly from (S-L0) and IH

7. CGsub-CG:

$$\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \ell'_1 \sqsubseteq \ell_1 \quad \Sigma; \Psi \vdash \ell_2 \sqsubseteq \ell'_2}{\Sigma; \Psi \vdash \mathbb{C} \ell_1 \ell_2 \tau <: \mathbb{C} \ell'_1 \ell'_2 \tau'}$$

$$\text{To prove: } [((\mathbb{C} \ell_i \ell_2 \tau) \sigma)]_V^{\hat{\beta}} \subseteq [((\mathbb{C} \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$$

It suffices to prove:

$$\forall ({}^s\theta, n, {}^s v, {}^t v) \in [((\mathbb{C} \ell_1 \ell_2 \tau) \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in [((\mathbb{C} \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$$

$$\text{This means that given } ({}^s\theta, n, {}^s v, {}^t v) \in [((\mathbb{C} \ell_1 \ell_2 \tau) \sigma)]_V^{\hat{\beta}}$$

Therefore from Definition 2.65 we are given:

$$\forall {}^s\theta_e \sqsupseteq {}^s\theta, H_s, H_t, i, {}^s v', k \leq m, \hat{\beta} \sqsubseteq \hat{\beta}'.$$

$$(k, H_s, H_t) \triangleright^{\hat{\beta}'} ({}^s\theta_e) \wedge (H_s, {}^s v) \Downarrow_i^f (H'_s, {}^s v') \wedge i < k \implies$$

$$\exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \triangleright^{\hat{\beta}''} {}^s\theta' \wedge$$

$$\exists {}^t v''. {}^t v' = \text{inl } {}^t v'' \wedge ({}^s\theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \quad (\text{S-M0})$$

And we are required to prove

$$({}^s\theta, n, {}^s v, {}^t v) \in [((\mathbb{C} \ell'_1 \ell'_2 \tau') \sigma)]_V^{\hat{\beta}}$$

So again from Definition 2.65 we need to prove

$$\forall {}^s\theta_{e1} \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1.$$

$$(k_1, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'_1} ({}^s\theta_{e1}) \wedge (H_{s1}, {}^s v) \Downarrow_{i_1}^f (H'_{s1}, {}^s v'_1) \wedge i_1 < k_1 \implies$$

$$\exists H'_{t1}, {}^t v'_1. (H_{t1}, {}^t v()) \Downarrow (H'_{t1}, {}^t v'_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''_1} {}^s\theta' \wedge$$

$$\exists {}^t v''_1. {}^t v'_1 = \text{inl } {}^t v''_1 \wedge ({}^s\theta', k_1 - i_1, {}^s v'_1, {}^t v''_1) \in [\tau' \sigma]_V^{\hat{\beta}''_1}$$

$$\text{This means we are given some } {}^s\theta_{e1} \sqsupseteq {}^s\theta, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1 \leq n, \hat{\beta} \sqsubseteq \hat{\beta}'_1 \text{ s.t. } (k_1, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}'_1}$$

$$({}^s\theta_{e1}) \wedge (H_{s1}, {}^s v) \Downarrow_{i_1}^f (H'_{s1}, {}^s v'_1) \wedge i_1 < k_1$$

And we need to prove

$$\exists H'_{t1}, {}^t v'_1. (H_{t1}, {}^t v()) \Downarrow (H'_{t1}, {}^t v'_1) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta_{e1}, \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1 . (k_1 - i_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''_1} {}^s\theta' \wedge$$

$$\exists {}^t v''_1. {}^t v'_1 = \text{inl } {}^t v''_1 \wedge ({}^s\theta', k_1 - i_1, {}^s v'_1, {}^t v''_1) \in [\tau' \sigma]_V^{\hat{\beta}''_1}$$

We instantiate (S-M0) with ${}^s\theta_{e1}, H_{s1}, H_{t1}, i_1, {}^s v'_1, k_1, \hat{\beta}'_1$ we get

$$\begin{aligned} & \exists H'_t, {}^t v'. (H_t, {}^t v()) \Downarrow (H'_t, {}^t v') \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_{e_1}, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_s, H'_t) \hat{\triangleright} {}^s \theta' \wedge \\ & \exists {}^t v'' . {}^t v' = \text{inl } {}^t v'' \wedge ({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''} \end{aligned}$$

$$\text{IH: } [(\tau \sigma)]_V^{\hat{\beta}''} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}} \hat{\beta}'' \text{ (Statement (1))}$$

Since we have $({}^s \theta', k - i, {}^s v', {}^t v'') \in [\tau \sigma]_V^{\hat{\beta}''}$ therefore from IH we get $({}^s \theta', k - i, {}^s v', {}^t v'') \in [(\tau' \sigma)]_V^{\hat{\beta}}$

8. CGsub-base:

Trivial

Proof of Statement(2)

It suffice to prove that

$$\forall ({}^s \theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}} . ({}^s \theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

This means that we are given $({}^s \theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}$

From Definition 2.66 it means we have

$$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} {}^s \theta \wedge \forall i < n, {}^s v. e_s \Downarrow_i {}^s v \implies$$

$$\exists H'_t, {}^t v. (H_t, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - i, H_s, H'_t) \hat{\triangleright} {}^s \theta \quad (\text{Sub-E0})$$

And we need to prove

$$({}^s \theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

From Definition 2.66 we need to prove

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \hat{\triangleright} {}^s \theta \wedge \forall j < n, {}^s v_1. e_s \Downarrow_j {}^s v_1 \implies$$

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau' \sigma)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \hat{\triangleright} {}^s \theta$$

This further means that given H_{s1}, H_{t1} s.t $(n, H_{s1}, H_{t1}) \hat{\triangleright} {}^s \theta$. Also given some $j < n, {}^s v_1$ s.t $e_s \Downarrow_j {}^s v_1$

And it suffices to prove that

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow (H'_{t1}, {}^t v_1) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v_1) \in [(\tau' \sigma)]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_{t1}) \hat{\triangleright} {}^s \theta$$

Instantiating (Sub-E0) with the given H_{s1}, H_{t1} and $j < n, {}^s v_1$. We get

$$\exists H'_t, {}^t v. (H_{t1}, e_t) \Downarrow (H'_t, {}^t v) \wedge ({}^s \theta, n - j, {}^s v_1, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}} \wedge (n - j, H_{s1}, H'_t) \hat{\triangleright} {}^s \theta$$

Since we have $({}^s \theta, n - j, {}^s v_1, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}}$ therefore from Statement(1) we get $({}^s \theta, n - j, {}^s v_1, {}^t v) \in [(\tau' \sigma)]_V^{\hat{\beta}}$

□

Theorem 2.76 (CG \rightsquigarrow FG: Deriving CG NI via compilation). $\forall e_s, {}^s v_1, {}^s v_2, {}^s v'_1, {}^s v'_2, n_1, n_2, H'_{s1}, H'_{s2}$.

let $\text{bool} = (\text{unit} + \text{unit})$.

$\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \wedge$

$\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \wedge \emptyset, \emptyset, \emptyset \vdash {}^s v_2 : \text{Labeled } \top \text{ bool} \wedge$

$(\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1}^f (H'_{s1}, {}^s v'_1) \wedge$

$(\emptyset, e_s[{}^s v_2/x]) \Downarrow_{n_2}^f (H'_{s2}, {}^s v'_2)$

\implies

${}^s v'_1 = {}^s v'_2$

Proof. From the CG to FG translation we know that $\exists e_t$ s.t

$$\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \rightsquigarrow e_t$$

Similarly we also know that $\exists {}^t v_1, {}^t v_2$ s.t

$$\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow {}^t v_1 \text{ and } \emptyset, \emptyset, \emptyset \vdash {}^s v_2 : \text{Labeled } \top \text{ bool} \rightsquigarrow {}^t v_2 \quad (\text{NI-0})$$

From type preservation theorem we know that

$$\begin{aligned} \emptyset, \emptyset, x : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \vdash_\top e_t : (\text{unit} \xrightarrow{\perp} ((\text{unit} + \text{unit})^\perp + \text{unit})^\perp)^\perp \\ \emptyset, \emptyset, \emptyset \vdash_\top {}^t v_1 : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \\ \emptyset, \emptyset, \emptyset \vdash_\top {}^t v_2 : ((\text{unit} + \text{unit})^\perp + \text{unit})^\top \quad (\text{NI-1}) \end{aligned}$$

Since we have $\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{Labeled } \top \text{ bool} \rightsquigarrow {}^t v_1$

And since ${}^s v_1$ and ${}^t v_1$ are closed terms (from given and NI-1)

Therefore from Theorem 2.74 we have (we choose n s.t $n > n_1$ and $n > n_2$)

$$(\emptyset, n, {}^s v_1, {}^t v_1) \in [\text{Labeled } \top \text{ bool}]_E^\emptyset \quad (\text{NI-2})$$

And therefore from Definition 2.70 and (NI-2) we have

$$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto {}^t v_1)) \in [x \mapsto \text{Labeled } \top \text{ bool}]_V^\emptyset$$

From (NI-0) we know that $\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_s : \mathbb{C} \perp \perp \text{ bool} \rightsquigarrow e_t$

Therefore we can apply Theorem 2.74 to get

$$(\emptyset, n, e_s[{}^s v_1/x], e_t[{}^t v_1/x]) \in [\mathbb{C} \perp \perp \text{ bool}]_E^\emptyset \quad (\text{NI-3.1})$$

Applying Definition 2.66 on (NI-3.1) we get

$$\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \hat{\triangleright}^\beta \emptyset \wedge \forall i < n. e_s[{}^s v_1/x] \Downarrow_i {}^s v \implies$$

$$\exists H'_{t2}, {}^t v. (H_{t2}, e_t[{}^t v_1/x]) \Downarrow (H'_{t2}, {}^t v) \wedge (\emptyset, n - i, {}^s v, {}^t v) \in [\mathbb{C} \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \hat{\triangleright}^\beta \emptyset$$

Instantiating with \emptyset, \emptyset . From cg-val we know that $i = 0$ and ${}^s v = e_s[{}^s v_1/x]$.

Therefore we have

$$\exists H'_{t2}, {}^t v. (H_{t2}, e_t[{}^t v_1/x]) \Downarrow (H'_{t2}, {}^t v) \wedge (\emptyset, n, {}^s v, {}^t v) \in [\mathbb{C} \perp \perp \text{ bool}]_V^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \hat{\triangleright}^\beta \emptyset$$

From translation and from (NI-1) we know that ${}^t v = e_t[{}^t v_1/x] = \lambda_. e_{b1}$ and therefore from fg-val we have $H'_{t2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[{}^s v_1/x], \lambda_. e_{b1}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\emptyset$$

Expanding $(\emptyset, n, e_s[{}^s v_1/x], \lambda_. e_{b1}) \in [\mathbb{C} \perp \perp \text{ bool}]_V^\emptyset$ using Definition 2.65 we get

$$\forall {}^s \theta_e \sqsupseteq \emptyset, H_{s3}, H_{t3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s3}, H_{t3}) \hat{\triangleright}^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, e_s[{}^s v_1/x]) \Downarrow_i^f (H'_{s1}, {}^s v''_1) \wedge i < k \implies$$

$$\begin{aligned} \exists H''_{t1}, {}^t v'', (H_{t3}, (\lambda_. e_{b1})) \Downarrow (H''_{t1}, {}^t v''_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s1}, H''_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_1. {}^t v''_1 = \\ \text{inl } {}^t v'''_1 \wedge ({}^s \theta', k - i, {}^s v''_1, {}^t v'''_1) \in [\text{bool}]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating with $\emptyset, \emptyset, \emptyset, n_1, {}^s v'_1, n, \emptyset$ we get

$$\begin{aligned} \exists H''_{t1}, {}^t v'' . (\emptyset, (\lambda_. e_{b1})) \Downarrow (H''_{t1}, {}^t v''_1) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s1}, H''_{t1}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_1. {}^t v''_1 = \\ \text{inl } {}^t v'''_1 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_1) \in [\text{bool}]_V^{\hat{\beta}''} \quad (\text{NI-3.2}) \end{aligned}$$

Since we have $\exists {}^t v'''_1. {}^t v''_1 = \text{inl } {}^t v'''_1 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_1) \in [(\text{unit} + \text{unit})]_V^{\hat{\beta}''}$, therefore from Definition 2.65 we know that 2 cases arise

- ${}^s v'_1 = \text{inl}^s v'_{i1}$ and ${}^t v'''_1 = \text{inl}^t v'_{i1}$:

And from Definition 2.65 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i1}, {}^t v'_{i1}) \in [\text{unit}]_V^{\hat{\beta}''}$$

which means ${}^s v'_{i1} = {}^t v'_{i1} = ()$

- ${}^s v'_1 = \text{inr}^s v'_{i1}$ and ${}^t v'''_1 = \text{inr}^t v'_{i1}$:

Same reasoning as in the previous case

Thus no matter which case occurs we have ${}^s v'_1 = {}^t v'''_1$ (NI-3.3)

Similarly we can apply Theorem 2.74 with the other substitution to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\mathbb{C} \perp \perp \text{bool}]_E^{\emptyset} \quad (\text{NI-4.1})$$

Applying Definition 2.66 on (NI-4.1) we get

$$\forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} \emptyset \wedge \forall i < n, {}^s v_s.e_s[{}^s v_2/x] \Downarrow_i {}^s v_s \implies \exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n - i, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \text{bool}]_V^{\hat{\beta}} \wedge (n - i, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset$$

Instantiating with \emptyset, \emptyset . From cg-val we know that $i = 0$ and ${}^s v_s = e_s[{}^s v_2/x]$.

Therefore we have

$$\exists H'_{t2}, {}^t v_s. (H_{t2}, e_t[{}^t v_2/x]) \Downarrow (H'_{t2}, {}^t v_s) \wedge (\emptyset, n, {}^s v_s, {}^t v_s) \in [\mathbb{C} \perp \perp \text{bool}]_V^{\hat{\beta}} \wedge (n, H_{s2}, H'_{t2}) \triangleright^{\hat{\beta}} \emptyset$$

Also from (NI-1) and from translation we know that ${}^t v = e_t[{}^t v_2/x] = \lambda_.e_{b2}$ and therefore from fg-val we know that $H'_{t2} = \emptyset$

Therefore we have

$$(\emptyset, n, e_s[{}^s v_2/x], \lambda_.e_{b2}) \in [\mathbb{C} \perp \perp \text{bool}]_V^{\emptyset}$$

Expanding $(\emptyset, n, e_s[{}^s v_2/x], \lambda x.e_{b2}) \in [\mathbb{C} \perp \perp \text{bool}]_V^{\emptyset}$ using Definition 2.65 we get

$$\forall {}^s \theta_e \sqsupseteq \emptyset, H_{s3}, H_{t3}, i, {}^s v'', k \leq n, \emptyset \sqsubseteq \hat{\beta}'.$$

$$(k, H_{s3}, H_{t3}) \triangleright^{\hat{\beta}'} ({}^s \theta_e) \wedge (H_{s3}, e_s[{}^s v_2/x]) \Downarrow_i^f (H'_{s2}, {}^s v''_2) \wedge i < k \implies$$

$$\exists H''_{t2}, {}^t v'', (H_{t3}, (\lambda_.e_{b2})) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta_e, \hat{\beta}' \sqsubseteq \hat{\beta}'' . (k - i, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', k - i, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_V^{\hat{\beta}''}$$

Instantiating with $\emptyset, \emptyset, \emptyset, n_2, {}^s v'_2, n, \emptyset$ we get

$$\exists H''_{t2}, {}^t v'' . (\emptyset, (\lambda_.e_{b2})) \Downarrow (H''_{t2}, {}^t v''_2) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \emptyset \sqsubseteq \hat{\beta}'' . (n - n_1, H'_{s2}, H''_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta' \wedge \exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_V^{\hat{\beta}''} \quad (\text{NI-4.2})$$

Since we have $\exists {}^t v'''_2. {}^t v''_2 = \text{inl}^t v'''_2 \wedge ({}^s \theta', n - n_1, {}^s v'_1, {}^t v'''_2) \in [\text{bool}]_V^{\hat{\beta}''}$, therefore from Definition 2.65 2 cases arise

- ${}^s v'_2 = \text{inl}^s v'_{i2}$ and ${}^t v'''_2 = \text{inl}^t v'_{i2}$:

And from Definition 2.65 we know that

$$({}^s \theta', n - n_1, {}^s v'_{i2}, {}^t v'_{i2}) \in [\text{unit}]_V^{\hat{\beta}''}$$

which means ${}^s v'_{i2} = {}^t v'_{i2} = ()$

- ${}^s v'_2 = \text{inr}^s v'_{i2}$ and ${}^t v'''_2 = \text{inr}^t v'_{i2}$:

Same reasoning as in the previous case

Thus no matter which case occurs we have ${}^s v_2' = {}^t v_2'''$ (NI-4.3)

From CG to FG translation we know that $\exists {}^t v_{i1}. {}^t v_{i1} = \text{inl } {}^t v_{i1}$ and similarly $\exists {}^t v_{i2}. {}^t v_{i2} = \text{inl } {}^t v_{i2}$

From (NI-1) since $\emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_{i1} : (\text{bool}^{\perp} + \text{unit})^{\top}$ therefore from CG-inl we know that $\emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\perp}$

And from CGsub-sum we know that $\emptyset, \emptyset, \emptyset \vdash_{\top} {}^t v_{i1} : \text{bool}^{\top}$

Therefore we also have $\emptyset, \emptyset, \emptyset \vdash_{\perp} {}^t v_{i1} : \text{bool}^{\top}$ (NI-5.1)

Similarly we also have $\emptyset, \emptyset, \emptyset \vdash_{\perp} {}^t v_{i2} : \text{bool}^{\top}$ (NI-5.2)

Next, let $e_T = (\lambda x : (\text{bool}^{\perp} + \text{unit})^{\top}. \text{case}(e_t(), y.y, z. {}^t v_b)) (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : \text{bool}^{\perp}$

where $\text{true} = \text{inl } ()$ and $\text{false} = \text{inr } ()$

We claim $\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} e_T : \text{bool}^{\perp}$

To show this we give its typing derivation

P2.3:

$$\frac{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{false} : \text{bool}^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{false} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.2:

$$\frac{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{true} : \text{bool}^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\perp}}{\text{FG-inl}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, - \vdash_{\perp} \text{inl } \text{true} : (\text{bool}^{\perp} + \text{unit})^{\top}} \text{FGSub-base}$$

P2.1:

$$\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} u : \text{bool}^{\top}}$$

P2:

$$\frac{\frac{P2.1 \quad P2.2 \quad P2.3 \quad \frac{}{\emptyset, \emptyset \models (\text{bool}^{\perp} + \text{unit})^{\top} \searrow \perp}}{\emptyset, \emptyset, u : \text{bool}^{\top} \vdash_{\perp} (\text{case}(u, -. \text{inl } \text{true}, -. \text{inl } \text{false})) : (\text{bool}^{\perp} + \text{unit})^{\top}}}$$

P1.2:

$$\frac{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t : (\text{unit} \xrightarrow{\perp} (\text{bool}^{\perp} + \text{unit})^{\perp})^{\perp}}{\text{NI-1}}}{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} () : \text{unit}}{\text{FG-unit}}}{\frac{\frac{}{\emptyset, \emptyset \models \perp \sqcup \perp \sqsubseteq \perp} \quad \frac{}{\emptyset, \emptyset \models (\text{bool}^{\perp} + \text{unit})^{\perp} \searrow \perp}}{\text{FG-app}}}}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} e_t() : (\text{bool}^{\perp} + \text{unit})^{\perp}} \text{FG-app}$$

P1.1:

$$\frac{\frac{P1.2 \quad \frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, y : \text{bool}^{\perp} \vdash_{\perp} y : \text{bool}^{\perp}}{\text{FG-var}}}{\frac{\frac{}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top}, z : \text{unit} \vdash_{\perp} \text{false} : \text{bool}^{\perp}}{\text{FG-var}} \quad \frac{}{\emptyset, \emptyset \models \text{bool}^{\perp} \searrow \perp}}{\emptyset, \emptyset, u : \text{bool}^{\top}, x : (\text{bool}^{\perp} + \text{unit})^{\top} \vdash_{\perp} \text{case}(e_t(), y.y, z. {}^t v_b) : \text{bool}^{\perp}} \text{FG-case}$$

P1:

$$\frac{\frac{P1.1}{\emptyset, \emptyset, u : \text{bool}^\top, x : (\text{bool}^\perp + \text{unit})^\top \vdash_\perp \text{case}(e_t(), y.y, z.^t v_b) : \text{bool}^\perp}}{\emptyset, \emptyset, u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top. \text{case}(e_t(), y.y, z.^t v_b)) : ((\text{bool}^\perp + \text{unit})^\top \xrightarrow{\perp} \text{bool}^\perp)^\perp}}$$

Main derivation:

$$\frac{\frac{P1 \quad P2 \quad \frac{}{\emptyset, \emptyset \models \perp \sqcup \perp \sqsubseteq \perp} \quad \frac{}{\emptyset, \emptyset \models \text{bool}^\perp \searrow \perp}}{\emptyset, \emptyset, u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top. \text{case}(e_t(), y.y, z.^t v_b)) (\text{case}(u, -.inl \text{true}, -.inl \text{false})) : \text{bool}^\perp} \text{FG-app}}{\emptyset, \emptyset, u : \text{bool}^\top \vdash_\perp (\lambda x : (\text{bool}^\perp + \text{unit})^\top. \text{case}(e_t(), y.y, z.^t v_b)) (\text{case}(u, -.inl \text{true}, -.inl \text{false})) : \text{bool}^\perp} \text{FG-app}}$$

Assuming $e_{b1}()$ reduces in n_{t1} steps in (NI-3.2) and $e_{b2}()$ reduces in n_{t2} steps in (NI-4.2).

We instantiate Theorem 2.29 with $e_T, ^t v_{i1}, ^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H''_{t1}, H''_{t2}$ and \perp and therefore from (NI-3.3) and (NI-4.3) we get $^t v_1''' = ^t v_2'''$ and thus $^s v_1' = ^s v_2'$

□

2.4 Translation from FG to FG⁻

2.4.1 FG⁻ typesystem

Lemma 2.77 (FG⁻: Reflexivity of subtyping). *The following hold:*

1. For all $\Sigma, \Psi, \tau: \Sigma; \Psi \vdash \tau <: \tau$
2. For all $\Sigma, \Psi, A: \Sigma; \Psi \vdash A <: A$

Proof. Proof by simultaneous induction on τ and A .

Proof of statement (1)

Let $\tau = A^\ell$. Then, we have:

$$\frac{\frac{}{\Sigma; \Psi \vdash A <: A} \text{IH(2)} \quad \Sigma; \Psi \vdash \ell \sqsubseteq \ell}{\Sigma; \Psi \vdash A^\ell <: A^\ell} \text{FGsub-label}$$

Proof of statement (2)

We proceed by cases on A .

1. $A = \mathbf{b}$:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

2. $A = \text{ref } \tau$:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

3. $A = \tau_1 \times \tau_2$:

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau_1 \times \tau_2}}$$

4. $A = \tau_1 + \tau_2$:

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(1) on } \tau_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau_1 + \tau_2}}$$

5. $A = \tau_1 \xrightarrow{\ell_e} \tau_2$:

$$\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 <: \tau_1} \text{IH(1) on } \tau_1 \quad \frac{}{\Sigma; \Psi \vdash \tau_2 <: \tau_2} \text{IH(2) on } \tau_2 \quad \frac{}{\Sigma; \Psi \vdash \ell_e \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau_1 \xrightarrow{\ell_e} \tau_2}}$$

6. $A = \text{unit}$:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}}$$

Type system: $\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau}$

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow pc}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau'} \text{FG}^- \text{-var} \quad \frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^{pc}} \text{FG}^- \text{-lam} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2} \text{FG}^- \text{-app} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^{pc}} \text{FG}^- \text{-prod} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1} \text{FG}^- \text{-fst} \quad \frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^{pc}} \text{FG}^- \text{-inl} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \quad \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x. e_1, y. e_2) : \tau} \text{FG}^- \text{-case} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc'} e : \tau' \quad \Sigma; \Psi \vdash pc \sqsubseteq pc' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau} \text{FG}^- \text{-sub} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^{pc}} \text{FG}^- \text{-ref} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau'} \text{FG}^- \text{-deref} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}^{pc}} \text{FG}^- \text{-assign} \\
\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} () : \text{unit}^{pc}} \text{FG}^- \text{-unitI} \quad \frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha. (\ell_e, \tau))^{pc}} \text{FG}^- \text{-FI} \\
\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\ell \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e [] : \tau[\ell'/\alpha]} \text{FG}^- \text{-FE} \\
\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^{pc}} \text{FG}^- \text{-CI} \\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau} \text{FG}^- \text{-CE}
\end{array}$$

Figure 15: Type system for FG^-

$$\begin{array}{c}
\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell' \quad \Sigma; \Psi \vdash A <: A'}{\Sigma; \Psi \vdash A^\ell <: A'^{\ell'}} \text{FG}^- \text{-sub-label} \qquad \frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FG}^- \text{-sub-base} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FG}^- \text{-sub-ref} \qquad \frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FG}^- \text{-sub-prod} \\
\\
\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FG}^- \text{-sub-sum} \\
\\
\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FG}^- \text{-sub-arrow} \\
\\
\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FG}^- \text{-sub-unit} \qquad \frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash \forall \alpha. \tau_1 <: \forall \alpha. \tau_2} \text{FG}^- \text{-sub-forall} \\
\\
\frac{\Sigma; \Psi \vdash c_2 \implies c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2} \text{FG}^- \text{-sub-constraint}
\end{array}$$

Figure 16: FG^- subtyping

7. $A = \forall \alpha. \tau_i$:

$$\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash \forall \alpha. \tau_i <: \forall \alpha. \tau_i}$$

8. $A = c \Rightarrow \tau_i$:

$$\frac{\frac{}{\Sigma; \Psi \vdash c \implies c} \quad \frac{}{\Sigma; \Psi, c \vdash \tau_i <: \tau_i} \text{IH(1) on } \tau_i}{\Sigma; \Psi \vdash c \Rightarrow \tau <: c \Rightarrow \tau_i}$$

□

2.4.2 Type translation

We define a translation of types, indexed by a label ℓ (which represents a pc joined with all outer labels) below. This is written $\llbracket \tau \rrbracket_\ell$.

Definition 2.78 ($FG \rightsquigarrow FG^-$: Type translation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_\ell &= \mathbf{b} \\
\llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_\ell &= \forall \alpha. \alpha, (\forall \beta. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha \\
\llbracket \tau_1 \times \tau_2 \rrbracket_\ell &= \llbracket \tau_1 \rrbracket_\ell \times \llbracket \tau_2 \rrbracket_\ell \\
\llbracket \tau_1 + \tau_2 \rrbracket_\ell &= \llbracket \tau_1 \rrbracket_\ell + \llbracket \tau_2 \rrbracket_\ell \\
\llbracket \text{ref } \tau \rrbracket_\ell &= \text{ref } \llbracket \tau \rrbracket_\perp \\
\llbracket \text{unit} \rrbracket_\ell &= \text{unit} \\
\llbracket \forall \gamma. (\ell_e, \tau) \rrbracket_\ell &= \forall \alpha. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha \\
\llbracket c \xrightarrow{\ell_e} \tau \rrbracket_\ell &= \forall \alpha. \alpha, (((c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha \\
\llbracket \mathbf{A}' \rrbracket_\ell &= (\llbracket \mathbf{A} \rrbracket_{\ell \sqcup \ell'})^{\ell \sqcup \ell'}
\end{aligned}$$

Translation judgement:

$$\boxed{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_M : \llbracket \tau \rrbracket_{pc'}} \text{ where}$$

$pc' \sqsubseteq pc$ and $\forall i \in 1 \dots n. \ell_i \sqsubseteq pc'$
 $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$
 $\Gamma' = x_1 : \llbracket \tau_1 \rrbracket_{\ell_1}, \dots, x_n : \llbracket \tau_n \rrbracket_{\ell_n}$

2.4.3 Type preservation: FG to FG^-

Theorem 2.79 ($FG \rightsquigarrow FG^-$: Type preservation). *Suppose (1) $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$ and (2) $\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau$ in FG . Suppose ℓ_1, \dots, ℓ_n and pc' are arbitrary labels with free variables in Σ such that (3) $\Sigma; \Psi \vdash pc' \sqsubseteq pc$ and (4) For each $i \in [1, n]$, $\Sigma; \Psi \vdash \ell_i \sqsubseteq pc'$.*

Let Γ' be the FG^- context $x_1 : \llbracket \tau_1 \rrbracket_{\ell_1}, \dots, x_n : \llbracket \tau_n \rrbracket_{\ell_n}$. Then, $\Sigma; \Psi; \Gamma' \vdash_{pc'} e : \llbracket \tau \rrbracket_{pc'}$ in FG^- .

Proof. Proof by induction on the \rightsquigarrow relation

1. var:

$$\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} x : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} x : \llbracket \tau \rrbracket_{pc'}} \text{var}$$

$$\frac{\llbracket \tau \rrbracket_{\ell_n} <: \llbracket \tau \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} x : \llbracket \tau \rrbracket_{pc'}}$$

2. lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma, x : \llbracket \tau_1 \rrbracket_{\ell_{n+1}} \vdash_{\ell'_e} e_m : \llbracket \tau_2 \rrbracket_{\ell'_e} \quad \ell_{n+1} \sqsubseteq \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_M : T_1} \text{lam}$$

$$T_1 = (\forall \alpha. \alpha, (\forall \beta. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha)^{pc'}$$

$$T_{1.1} = (\forall \beta. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha$$

$$T_{1.2} = ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.3} = (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha$$

$$c_1 = (pc' \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha)$$

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2} \text{Given}}{\Sigma, \alpha, \beta; \Psi, c_1; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2} \text{Weakening}}{\Sigma, \alpha, \beta; \Psi, c_1; \Gamma', x : \llbracket \tau_1 \rrbracket_{\beta} \vdash_{\alpha} e_m : \llbracket \tau_2 \rrbracket_{\alpha}} \text{IH}}$$

Main derivation:

$$\frac{\frac{\frac{\frac{}{P1}}{\Sigma, \alpha, \beta; \Psi, c_1; \Gamma' \vdash_{\alpha} \lambda x. e_m : T_{1.3}} \text{FG}^- \text{-lam}}{\Sigma, \alpha, \beta; \Psi; \Gamma' \vdash_{\alpha} \nu(\lambda x. e_m) : T_{1.2}} \text{FG}^- \text{-CI}}{\Sigma, \alpha; \Psi; \Gamma' \vdash_{\alpha} \Lambda(\nu(\lambda x. e_m)) : T_{1.1}} \text{FG}^- \text{-FI}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda(\Lambda(\nu(\lambda x. e_m))) : T_1} \text{FG}^- \text{-FI}$$

3. app:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^{\ell} \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : T_1 \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_1 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_M : \llbracket \tau_2 \rrbracket_{pc'}} \text{app}}$$

$$T_1 = (\forall \alpha. \alpha, (\forall \beta. \alpha, (((pc' \sqcup \ell) \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_{\beta} \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_{\alpha})^{\alpha})^{\alpha})^{pc' \sqcup \ell}$$

$$T_{1.1} = (\forall \beta. (pc' \sqcup \ell), (((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e \wedge \beta \sqsubseteq (pc' \sqcup \ell)) \xrightarrow{(pc' \sqcup \ell)} (\llbracket \tau_1 \rrbracket_{\beta} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)}))^{(pc' \sqcup \ell)}$$

$$T_{1.2} = (((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell)) \xrightarrow{(pc' \sqcup \ell)} (\llbracket \tau_1 \rrbracket_{(pc' \sqcup \ell)} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)}))^{(pc' \sqcup \ell)}$$

$$c_1 = ((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell))$$

$$T_{1.3} = (\llbracket \tau_1 \rrbracket_{(pc' \sqcup \ell)} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$T_{1.4} = (\llbracket \tau_1 \rrbracket_{(pc')} \xrightarrow{(pc' \sqcup \ell)} \llbracket \tau_2 \rrbracket_{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

P7:

$$\overline{pc' \sqcup \ell \sqsubseteq pc' \sqcup \ell}$$

P6:

$$\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_1 \rrbracket_{pc'}} \text{IH2}$$

P5:

$$\overline{\Sigma; \Psi \vdash T_{1.3} \searrow_{pc' \sqcup \ell}} \text{Definition of } \llbracket \cdot \rrbracket$$

P4:

$$\overline{\Sigma; \Psi \vdash T_{1.2} \searrow_{pc' \sqcup \ell}} \text{Definition of } \llbracket \cdot \rrbracket$$

P3:

$$\overline{\Sigma; \Psi \vdash T_{1.1} \searrow_{pc' \sqcup \ell}} \text{Definition of } \llbracket \cdot \rrbracket$$

P2:

$$\overline{pc' \sqcup pc' \sqcup \ell \sqsubseteq pc' \sqcup \ell}$$

P1:

$$\frac{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : T_1} \text{ IH1} \quad P2 \quad P3}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} \square : T_{1.1}} \text{ FG}^- \text{-FE} \quad P2 \quad P4}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} \square \square : T_{1.2}} \text{ FG}^- \text{-FE}}$$

Main derivation:

$$\frac{\frac{\frac{P1 \quad \overline{\Sigma; \Psi \vdash c_1} \quad P2 \quad P5}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) : T_{1.3}} \text{ FG}^- \text{-CE} \quad P6 \quad P7}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) : T_{1.4}} \text{ FG}^- \text{-sub} \quad P6 \quad P7}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) e_{m2} : \llbracket \tau_2 \rrbracket_{pc' \sqcup \ell}} \text{ FG}^- \text{-app}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1} \square \square \bullet) e_{m2} : \llbracket \tau_2 \rrbracket_{pc'}} \text{ Lemma 2.82}$$

4. prod:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : \llbracket \tau_1 \rrbracket_{pc'} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_2 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1}, e_{m2}) : (\llbracket \tau_1 \rrbracket_{pc'} \times \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{ prod}}{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : \llbracket \tau_1 \rrbracket_{pc'}} \text{ IH1} \quad \overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau_2 \rrbracket_{pc'}} \text{ IH2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} (e_{m1}, e_{m2}) : (\llbracket \tau_1 \rrbracket_{pc'} \times \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{ FG}^- \text{-prod}}$$

5. fst:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'} \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{fst}(e_m) : \llbracket \tau_1 \rrbracket_{pc'}} \text{ fst}}{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'}} \text{ IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{fst}(e_m) : \llbracket \tau_1 \rrbracket_{\ell \sqcup pc'}} \text{ FG}^- \text{-fst}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{fst}(e_m) : \llbracket \tau_1 \rrbracket_{pc'}} \text{ Lemma 2.82}}$$

6. snd:

$$\frac{\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'} \quad \Sigma; \Psi \vdash \tau_2 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{snd}(e_m) : \llbracket \tau_2 \rrbracket_{pc'}} \text{ snd}}{\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{\ell \sqcup pc'} \times \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'})^{\ell \sqcup pc'}} \text{ IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{snd}(e_m) : \llbracket \tau_2 \rrbracket_{\ell \sqcup pc'}} \text{ FG}^- \text{-snd}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{snd}(e_m) : \llbracket \tau_2 \rrbracket_{pc'}} \text{ Lemma 2.82}}$$

7. inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_1 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inl}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{inl}$$

$$\frac{\frac{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_1 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inl}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inl}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{FG}^{-}\text{-inl}$$

8. inr:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_2 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{inr}$$

$$\frac{\frac{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau_2 \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inr}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{inr}(e_m) : (\llbracket \tau_1 \rrbracket_{pc'} + \llbracket \tau_2 \rrbracket_{pc'})^{pc'}} \text{FG}^{-}\text{-inr}$$

9. case:

$$\frac{\begin{array}{l} \Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc' \sqcup \ell} + \llbracket \tau_1 \rrbracket_{pc' \sqcup \ell})^{pc' \sqcup \ell} \\ \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow \Sigma; \Psi; \Gamma', x : \llbracket \tau_1 \rrbracket_{\ell_{n+1}} \vdash_{pc' \sqcup \ell} e_{m1} : \llbracket \tau \rrbracket_{pc' \sqcup \ell} \\ \Sigma; \Psi; \Gamma, y : \tau_2 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow \Sigma; \Psi; \Gamma', y : \llbracket \tau_2 \rrbracket_{\ell_{n+2}} \vdash_{pc' \sqcup \ell} e_{m2} : \llbracket \tau \rrbracket_{pc' \sqcup \ell} \end{array}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, x.e_{m2}) : \llbracket \tau \rrbracket_{pc'}} \text{case}$$

P2:

$$\frac{}{\Sigma; \Psi; \Gamma', y : \llbracket \tau_2 \rrbracket_{pc' \sqcup \ell} \vdash_{pc' \sqcup \ell} e_{m2} : \llbracket \tau \rrbracket_{pc' \sqcup \ell}} \text{IH3 @ } pc' \sqcup \ell$$

P1:

$$\frac{}{\Sigma; \Psi; \Gamma', x : \llbracket \tau_1 \rrbracket_{pc' \sqcup \ell} \vdash_{pc' \sqcup \ell} e_{m1} : \llbracket \tau \rrbracket_{pc' \sqcup \ell}} \text{IH2 @ } pc' \sqcup \ell$$

Main derivation:

$$\frac{\frac{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\llbracket \tau_1 \rrbracket_{pc' \sqcup \ell} + \llbracket \tau_1 \rrbracket_{pc' \sqcup \ell})^{pc' \sqcup \ell}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, x.e_{m2}) : \llbracket \tau \rrbracket_{pc' \sqcup \ell}} \text{IH1} \quad \text{P1} \quad \text{P2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{case}(e_m, x.e_{m1}, x.e_{m2}) : \llbracket \tau \rrbracket_{pc'}} \text{FG}^{-}\text{-case} \quad \text{Lemma 2.82}$$

10. sub:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc''} e : \tau' \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau' \rrbracket_{pc'} \quad \Sigma; \Psi \vdash_{pc} pc \sqsubseteq pc'' \quad \Sigma; \Psi \vdash \tau' <: \tau}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'}} \text{sub}$$

$$\frac{\frac{\frac{pc' \sqsubseteq pc \sqsubseteq pc''}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau' \rrbracket_{pc'}} \text{IH} \quad \frac{\tau' <: \tau}{\llbracket \tau' \rrbracket_{pc'} <: \llbracket \tau \rrbracket_{pc'}} \text{Lemma 2.80}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'}} \text{IH}$$

11. ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'} \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } e : (\text{ref } \tau)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \text{new } e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)_{pc'}} \text{ref}$$

P1:

$$\frac{\frac{\text{Given} \quad \Sigma; \Psi \vdash pc' \sqsubseteq pc}{\Sigma; \Psi \vdash \tau \searrow pc}}{\Sigma; \Psi \vdash \tau \searrow pc'} \text{Lemma 2.85}}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket_\perp \searrow pc'}$$

Main derivation:

$$\frac{\frac{\frac{\text{IH}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : \llbracket \tau \rrbracket_\perp} \text{Lemma 2.82} \quad P1}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \text{new } e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)_{pc'}} \text{FG}^- \text{-new}}$$

12. deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc' \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau' \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} !e_m : \llbracket \tau' \rrbracket_{pc'}} \text{deref}$$

$$\frac{\frac{\frac{\tau <: \tau'}{\Sigma; \Psi \vdash \llbracket \tau \rrbracket_\perp <: \llbracket \tau' \rrbracket_{pc' \sqcup \ell}} \text{Lemma 2.80} \quad \frac{\text{IH1}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc'} \text{Definition of } \searrow}{\Sigma; \Psi; \Gamma' \vdash_{pc'} !e_m : \llbracket \tau' \rrbracket_{pc' \sqcup \ell}}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} !e_m : \llbracket \tau' \rrbracket_{pc'}} \text{Lemma 2.82}$$

13. assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} : (\text{ref } \llbracket \tau \rrbracket_\perp)^\ell \sqcup pc' \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_{pc'} \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit}^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} := e_{m2} : \text{unit}^{pc'}} \text{assign}$$

P1:

$$\frac{\frac{\text{IH2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_{pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_\perp} \text{Lemma 2.82} \quad \frac{\text{Given} \quad \tau \searrow pc}{\tau \searrow pc'}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m2} : \llbracket \tau \rrbracket_\perp}$$

Main derivation:

$$\frac{\frac{\text{IH1} \quad P1 \quad \frac{\tau \searrow (\ell \sqcup pc)}{\llbracket \tau \rrbracket_\perp \searrow \ell \sqcup pc'} \text{Lemma 2.85}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_{m1} := e_{m2} : \text{unit}^{pc'}} \text{FG}^- \text{-assign}}$$

14. unitI:

$$\frac{}{\Sigma; \Psi; \Gamma \vdash_{pc} () : \text{unit}^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} () : \text{unit}^{pc'}} \text{unitI}$$

$$\frac{}{\Sigma; \Psi; \Gamma' \vdash_{pc'} () : \text{unit}^{pc'}} \text{FG}^- \text{-unitI}$$

15. FI:

$$\frac{\Sigma, \alpha; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow \Sigma, \alpha; \Psi; \Gamma' \vdash_{\ell'_e} e_m : \llbracket \tau \rrbracket_{\ell'_e} \quad \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha. (\ell_e, \tau))^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda((\nu(\Lambda e_m))) : T_1} \text{FI}$$

$$T_1 = (\forall \alpha. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^{pc'}$$

$$T_{1.1} = ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.2} = (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha$$

$$c_1 = (pc' \sqsubseteq \alpha \sqsubseteq \ell_e)$$

$$T_{1.3} = \llbracket \tau \rrbracket_\alpha$$

P1:

$$\frac{\frac{}{\Sigma, \alpha, \gamma; \Psi, c_1; \Gamma' \vdash_\alpha e_m : T_{1.3}} \text{IH with } \ell'_e \text{ as } \alpha}{\Sigma, \alpha, \gamma; \Psi, c_1; \Gamma' \vdash_\alpha \Lambda e_m : T_{1.2}} \text{FG}^- \text{-FI}$$

Main derivation:

$$\frac{\frac{P1}{\Sigma, \alpha; \Psi; \Gamma' \vdash_\alpha \nu(\Lambda e_m) : T_{1.1}} \text{FG}^- \text{-CI}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda((\nu(\Lambda e_m))) : T_1} \text{FG}^- \text{-FI}$$

16. CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow \Sigma; \Psi, c; \Gamma' \vdash_{\ell'_e} e_m : \llbracket \tau \rrbracket_{\ell'_e} \quad \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda(\nu e_m) : T_1} \text{CI}$$

$$T_1 = (\forall \alpha. \alpha, ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha)^{pc'}$$

$$T_{1.1} = ((pc' \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha$$

$$T_{1.2} = \llbracket \tau \rrbracket_\alpha$$

$$c_1 = (pc' \sqsubseteq \alpha \sqsubseteq \ell_e)$$

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi, c_1; \Gamma' \vdash_\alpha e_m : T_{1.2}} \text{IH with } \ell'_e \text{ as } \alpha}{\Sigma, \alpha; \Psi; \Gamma' \vdash_\alpha \nu e_m : T_{1.1}} \text{FG}^- \text{-CI}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} \Lambda(\nu e_m) : T_1} \text{FG}^- \text{-FI}$$

17. FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \gamma. (\ell_e, \tau))^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1 \quad \text{FV}(\ell') \in \Sigma \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\gamma] \quad \Sigma; \Psi \vdash \tau[\ell'/\gamma] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau[\ell'/\gamma] \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : \llbracket \tau[\ell'/\gamma] \rrbracket_{pc'}} \text{FE}$$

$$T_1 = (\forall \alpha. \alpha, (((pc' \sqcup \ell) \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha) \alpha)^{pc' \sqcup \ell}$$

$$T_{1.1} = (((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e) \stackrel{(pc' \sqcup \ell)}{\Rightarrow} (\forall \gamma. (pc' \sqcup \ell), \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}^{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$T_{1.2} = (\forall \gamma. (pc' \sqcup \ell), \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}^{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$c_1 = ((pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e)$$

$$T_{1.3} = \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}[\ell'/\gamma]$$

$$T_{1.31} = \llbracket \tau[\ell'/\gamma] \rrbracket_{(pc' \sqcup \ell)}$$

$$T_{1.4} = \llbracket \tau[\ell'/\gamma] \rrbracket_{pc'}$$

P5:

$$\frac{}{T_{1.2} \searrow (pc' \sqcup \ell)} \text{Definition of } \llbracket \cdot \rrbracket$$

P4:

$$\frac{}{T_{1.1} \searrow (pc' \sqcup \ell)} \text{Definition of } \llbracket \cdot \rrbracket$$

P3:

$$\frac{}{(pc' \sqcup \ell) \sqsubseteq (pc \sqcup \ell) \sqsubseteq \ell_e} \text{Given}$$

P2:

$$\frac{\frac{}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1} \text{IH} \quad P4}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.1}} \text{FG}^- \text{-FE}$$

P1:

$$\frac{P2 \quad P3 \quad P5}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.2}} \text{FG}^- \text{-CE}$$

P0:

$$\frac{P1 \quad \frac{\frac{}{\Sigma; \Psi \vdash T_{1.3} \searrow (pc' \sqcup \ell)} \text{Definition of } \llbracket \cdot \rrbracket \quad P2}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.3}} \text{FG}^- \text{-FE}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.31}} \text{Lemma 2.84}$$

Main derivation:

$$\frac{P0 \quad \frac{}{\Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell} \text{Lemma 2.82}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \square \bullet \square : T_{1.4}}$$

18. CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \stackrel{\ell_e}{\rightsquigarrow} \tau)^\ell \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1 \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow \Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : \llbracket \tau \rrbracket_{pc'}} \text{CE}$$

$$T_1 = (\forall \alpha. \alpha, ((c \wedge (pc' \sqcup \ell) \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha)^{pc' \sqcup \ell}$$

$$T_{1.1} = ((c \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e) \stackrel{(pc' \sqcup \ell)}{\Rightarrow} \llbracket \tau \rrbracket_{(pc' \sqcup \ell)})^{(pc' \sqcup \ell)}$$

$$T_{1.2} = \llbracket \tau \rrbracket_{(pc' \sqcup \ell)}$$

$$T_{1.3} = \llbracket \tau \rrbracket_{pc'}$$

$$c_1 = (c \wedge (pc' \sqcup \ell) \sqsubseteq (pc' \sqcup \ell) \sqsubseteq \ell_e)$$

P3:

$$\frac{\overline{\Sigma; \Psi \vdash (pc \sqcup \ell) \sqsubseteq \ell_e} \text{ Given}}{\Sigma; \Psi \vdash (pc' \sqcup \ell) \sqsubseteq \ell_e}$$

P2:

$$\frac{\overline{\Sigma; \Psi \vdash T_{1.2} \searrow (pc' \sqcup \ell)} \text{ Definition of } \llbracket \cdot \rrbracket}{\Sigma; \Psi \vdash T_{1.2} \searrow (pc' \sqcup \ell)}$$

P1:

$$\frac{\overline{\Sigma; \Psi \vdash T_{1.1} \searrow (pc' \sqcup \ell)} \text{ Definition of } \llbracket \cdot \rrbracket}{\Sigma; \Psi \vdash T_{1.1} \searrow (pc' \sqcup \ell)}$$

P0:

$$\frac{\frac{\overline{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m : T_1} \text{ IH} \quad P1}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.1}} \text{ FG}^- \text{-FE} \quad \frac{\overline{\Sigma; \Psi \vdash c} \text{ Given, Weakening} \quad P3}{\Sigma; \Psi \vdash c_1} \text{ P2}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.2}} \text{ FG}^- \text{-CE}$$

Main derivation:

$$\frac{P0.1 \quad \frac{\overline{\tau \searrow \ell} \text{ Given}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.3}} \text{ Lemma 2.82}}{\Sigma; \Psi; \Gamma' \vdash_{pc'} e_m \bullet : T_{1.3}}$$

□

Lemma 2.80 (FG \rightsquigarrow FG⁻: Subtyping). $\forall \Sigma, \Psi, \ell, \ell'. \Sigma; \Psi \vdash \ell \sqsubseteq \ell'$ and the following holds:

1. $\forall \tau, \tau'.$

$$\Sigma; \Psi \vdash \tau <: \tau' \implies \llbracket \tau \rrbracket_\ell <: \llbracket \tau' \rrbracket_{\ell'}$$

2. $\forall A, A'.$

$$\Sigma; \Psi \vdash A <: A' \implies \Sigma; \Psi \vdash \llbracket A \rrbracket_\ell <: \llbracket A' \rrbracket_{\ell'}$$

Proof. Proof by simultaneous induction on $\tau <: \tau$ and $A <: A$

Proof of statement (1)

Let $\tau = A_1^{\ell_1}$ and $\tau' = A_2^{\ell_2}$

P2:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\Sigma; \Psi \vdash A_1 <: A_2} \text{ By inversion } P1}{\Sigma; \Psi \vdash ([A_1]_{\ell \sqcup \ell_1}) <: ([A_2]_{\ell' \sqcup \ell_2})} \text{ IH(2) on } A_1 <: A_2$$

P1:

$$\frac{\frac{\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \text{ Given}}{\Sigma; \Psi \vdash \ell_1 \sqsubseteq \ell_2} \text{ By inversion } \quad \overline{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'} \text{ Given}}{\Sigma; \Psi \vdash \ell \sqcup \ell_1 \sqsubseteq \ell' \sqcup \ell_2}$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash ([A_1]_{\ell \sqcup \ell_1})^{\ell \sqcup \ell_1} <: ([A_2]_{\ell' \sqcup \ell_2})^{\ell' \sqcup \ell_2}}{\Sigma; \Psi \vdash [A_1^{\ell_1}]_{\ell} <: [A_2^{\ell_2}]_{\ell'}}$$

Proof of statement (2)

We proceed by cases on $A <: A$

1. FGsub-base:

$$\frac{\overline{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{ FG}^- \text{ sub-base}}{\Sigma; \Psi \vdash [\mathbf{b}]_{\ell} <: [\mathbf{b}]_{\ell'}} \text{ Definition of } [\cdot]$$

2. FGsub-ref:

$$\frac{\overline{\Sigma; \Psi \vdash \text{ref } [\tau_i]_{\perp} <: \text{ref } [\tau_i]_{\perp}} \text{ FG}^- \text{ sub-ref}}{\Sigma; \Psi \vdash [\text{ref } \tau_i]_{\ell} <: [\text{ref } \tau_i]_{\ell'}} \text{ Definition of } [\cdot]$$

3. FGsub-prod:

P1:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\Sigma; \Psi \vdash [\tau_1]_{\ell} <: [\tau'_1]_{\ell'}} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash [\tau_2]_{\ell} <: [\tau'_2]_{\ell'}} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash [\tau_1]_{\ell} \times [\tau_2]_{\ell} <: [\tau'_1]_{\ell} \times [\tau'_2]_{\ell'}}{\Sigma; \Psi \vdash [\tau_1 \times \tau_2]_{\ell} <: [\tau'_1 \times \tau'_2]_{\ell'}} \text{ FG}^- \text{ sub-prod}}{\Sigma; \Psi \vdash [\tau_1 \times \tau_2]_{\ell} <: [\tau'_1 \times \tau'_2]_{\ell'}} \text{ Definition of } [\cdot]$$

4. FGsub-sum:

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_\ell <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{ IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell + \llbracket \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 \rrbracket_{\ell'} + \llbracket \tau'_2 \rrbracket_{\ell'}} \text{ FG}^- \text{ sub-prod}}{\Sigma; \Psi \vdash \llbracket \tau_1 + \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 + \tau'_2 \rrbracket_{\ell'}} \text{ Definition of } \llbracket \cdot \rrbracket$$

5. FGsub-arrow:

$$T_1 = \forall \alpha. \alpha, (\forall \beta. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha$$

$$T_{1.0} = \forall \beta. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{1.1} = ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha)$$

$$T_{1.2} = (\llbracket \tau_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau_2 \rrbracket_\alpha)^\alpha$$

$$c_1 = (\ell \sqsubseteq \alpha \sqsubseteq \ell_e \wedge \beta \sqsubseteq \alpha)$$

$$T_2 = \forall \alpha. \alpha, (\forall \beta. \alpha, ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha)^\alpha)^\alpha$$

$$T_{2.0} = \forall \beta. \alpha, ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha)^\alpha$$

$$T_{2.1} = ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha) \xrightarrow{\alpha} (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha)$$

$$T_{2.2} = (\llbracket \tau'_1 \rrbracket_\beta \xrightarrow{\alpha} \llbracket \tau'_2 \rrbracket_\alpha)^\alpha$$

$$c_2 = (\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e \wedge \beta \sqsubseteq \alpha)$$

P3:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell'_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_\alpha <: \llbracket \tau'_2 \rrbracket_\alpha} \text{ IH(1) with } \ell = \ell' = \alpha$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell_e} \tau'_2} \text{ Given}}{\Sigma; \Psi \vdash \tau'_1 <: \tau_1} \text{ By inversion}}{\Sigma; \Psi \vdash \llbracket \tau'_1 \rrbracket_\beta <: \llbracket \tau_1 \rrbracket_\beta} \text{ IH(1) with } \ell = \ell' = \beta$$

P1:

$$\frac{P2 \quad P3}{\Sigma, \alpha, \beta; \Psi \vdash T_{1.3} <: T_{2.3}} \text{ FG}^- \text{ sub-arrow}$$

P0:

$$\begin{array}{c}
\frac{\frac{\frac{}{\Sigma, \alpha, \beta; \Psi \vdash \ell \sqsubseteq \ell'}{\text{Given, Weakening}}}{\Sigma, \alpha, \beta; \Psi \vdash \ell' \sqsubseteq \alpha} \implies \ell \sqsubseteq \alpha}{\Sigma, \alpha, \beta; \Psi \vdash c_2 \implies c_1} \quad \frac{\frac{\frac{}{\Sigma, \alpha, \beta; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\text{Given, Weakening}}}{\Sigma, \alpha, \beta; \Psi \vdash \alpha \sqsubseteq \ell'_e} \implies \alpha \sqsubseteq \ell_e}{\Sigma, \alpha, \beta; \Psi \vdash c_2 \implies c_1} \\
\frac{P1}{\Sigma, \alpha, \beta; \Psi \vdash T_{1.2} <: T_{2.2}} \text{ Weakening, FG}^- \text{-sub-label} \\
\hline
\Sigma, \alpha; \Psi \vdash T_{1.1} <: T_{2.1} \quad \text{FG}^- \text{-sub-constraint}
\end{array}$$

P0.1:

$$\frac{P0}{\Sigma, \alpha; \Psi \vdash T_{1.0} <: T_{2.0}} \text{ FG}^- \text{-sub-forall}$$

Main derivation:

$$\frac{\frac{P0.1}{\Sigma; \Psi \vdash T_1 <: T_2} \text{ FG}^- \text{-sub-label}}{\Sigma; \Psi \vdash \left[\left[\tau_1 \xrightarrow{\ell_\xi} \tau_2 \right]_\ell <: \left[\left[\tau'_1 \xrightarrow{\ell'_\xi} \tau'_2 \right]_{\ell'} \right]_{\ell'}} \text{ Definition of } \llbracket \cdot \rrbracket$$

6. FGsub-unit:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{ FG}^- \text{-sub-unit}}{\Sigma; \Psi \vdash \llbracket \text{unit} \rrbracket_\ell <: \llbracket \text{unit} \rrbracket_{\ell'}} \text{ Definition of } \llbracket \cdot \rrbracket$$

7. FGsub-forall:

$$\begin{aligned}
T_1 &= \forall \alpha. \alpha, ((\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha)^\alpha \\
T_{1.0} &= (\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha \\
T_{1.1} &= (\forall \gamma. \alpha, \llbracket \tau \rrbracket_\alpha)^\alpha \\
c_1 &= (\ell \sqsubseteq \alpha \sqsubseteq \ell_e) \\
T_{1.2} &= \llbracket \tau \rrbracket_\alpha \\
T_2 &= \forall \alpha. \alpha, ((\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau' \rrbracket_\alpha)^\alpha)^\alpha \\
T_{2.0} &= (\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \xrightarrow{\alpha} (\forall \gamma. \alpha, \llbracket \tau' \rrbracket_\alpha)^\alpha \\
T_{2.1} &= (\forall \gamma. \alpha, \llbracket \tau' \rrbracket_\alpha)^\alpha \\
c_2 &= (\ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \\
T_{2.2} &= \llbracket \tau' \rrbracket_\alpha
\end{aligned}$$

P0:

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell \sqsubseteq \ell'}{\text{Given, Weakening}}}{\Sigma, \alpha; \Psi \vdash \ell' \sqsubseteq \alpha} \implies \ell \sqsubseteq \alpha}{\Sigma, \alpha; \Psi \vdash c_2 \implies c_1} \quad \frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\text{Given, Weakening}}}{\Sigma, \alpha; \Psi \vdash \alpha \sqsubseteq \ell'_e} \implies \alpha \sqsubseteq \ell_e}{\Sigma, \alpha; \Psi \vdash c_2 \implies c_1}$$

P1:

$$\frac{\frac{\frac{}{\Sigma, \alpha, \gamma; \Psi \vdash T_{1.2} <: T_{2.2}}{\text{IH}} \quad \text{FG}^- \text{-sub-forall} \quad \frac{}{\Sigma; \Psi \vdash c_2 \implies c_1} \text{P0}}{\Sigma, \alpha; \Psi \vdash T_{1.1} <: T_{2.1}} \quad \text{FG}^- \text{-sub-constraint}}{\Sigma, \alpha; \Psi \vdash T_{1.0} <: T_{2.0}} \quad \text{FG}^- \text{-sub-forall}}{\Sigma; \Psi \vdash T_1 <: T_2} \quad \text{FG}^- \text{-sub-forall}$$

Main derivation:

$$\frac{}{\Sigma; \Psi \vdash \llbracket \forall \gamma. \tau_1 \rrbracket_\ell <: \llbracket \forall \gamma. \tau_2 \rrbracket_{\ell'}} \text{P0.1} \quad \text{Definition of } \llbracket \cdot \rrbracket$$

8. FGsub-constraint:

$$T_1 = \forall \alpha. \alpha, ((c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha$$

$$T_{1.1} = ((c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau \rrbracket_\alpha)^\alpha$$

$$T_{1.2} = \llbracket \tau \rrbracket_\alpha$$

$$c_1 = (c \wedge \ell \sqsubseteq \alpha \sqsubseteq \ell_e)$$

$$T_2 = \forall \alpha. \alpha, ((c' \wedge \ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau' \rrbracket_\alpha)^\alpha$$

$$T_{2.1} = ((c' \wedge \ell' \sqsubseteq \alpha \sqsubseteq \ell'_e) \stackrel{\alpha}{\Rightarrow} \llbracket \tau' \rrbracket_\alpha)^\alpha$$

$$T_{2.2} = \llbracket \tau' \rrbracket_\alpha$$

$$c_2 = (c' \wedge \ell' \sqsubseteq \alpha \sqsubseteq \ell'_e)$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash c_1 \Rightarrow \tau_1 <: c_2 \Rightarrow \tau_2} \text{Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau_2} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_\ell <: \llbracket \tau_2 \rrbracket_{\ell'}} \text{IH(1) on } \tau_1 <: \tau_2$$

P1:

$$\frac{\frac{}{\Sigma, \alpha; \Psi \vdash c \Rightarrow \tau <: c' \Rightarrow \tau'} \text{Given, Weakening}}{\Sigma, \alpha; \Psi \vdash c' \implies c} \text{By inversion}$$

P0:

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell \sqsubseteq \ell'} \text{Given, Weakening}}{\Sigma, \alpha; \Psi \vdash \ell' \sqsubseteq \alpha \implies \ell \sqsubseteq \alpha} \quad \frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e} \text{Given, Weakening}}{\Sigma, \alpha; \Psi \vdash \alpha \sqsubseteq \ell'_e \implies \alpha \sqsubseteq \ell_e} \text{P1}}{\Sigma, \alpha; \Psi \vdash c_2 \implies c_1}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \llbracket \tau \rrbracket_\alpha <: \llbracket \tau' \rrbracket_\alpha} \text{IH}}{\Sigma, \alpha; \Psi \vdash T_{1.1} <: T_{2.1}} \text{FG}^- \text{-sub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{Definition of } \llbracket \cdot \rrbracket_\ell}{\Sigma; \Psi \vdash \llbracket c_1 \implies \tau_1 \rrbracket_\ell <: \llbracket c_2 \implies \tau_2 \rrbracket_{\ell'}}$$

□

Lemma 2.81 (FG \rightsquigarrow FG⁻: Subtyping with label). *If $\Sigma; \Psi \vdash \ell \sqsubseteq \ell'$, then $\Sigma; \Psi \vdash \llbracket \tau \rrbracket_{\ell} <: \llbracket \tau \rrbracket_{\ell'}$ in FG⁻.*

Proof. From Lemma 2.80 with $\tau = \tau'$ and from Lemma 2.77 □

Lemma 2.82 (FG \rightsquigarrow FG⁻: Subtyping for $\tau \searrow \ell$). *If $\Sigma; \Psi \vdash \tau \searrow \ell$, then $\Sigma; \Psi \vdash \llbracket \tau \rrbracket_{\ell \sqcup \ell'} <: \llbracket \tau \rrbracket_{\ell'}$ in FG⁻.*

Proof. Since $\Sigma; \Psi \vdash \tau \searrow \ell$, there exists ℓ'' such that $\tau = \mathbf{A}^{\ell''}$ and $\Sigma; \Psi \vdash \ell \sqsubseteq \ell''$. Now we have:

$$\begin{aligned} & \Sigma; \Psi \vdash \llbracket \tau \rrbracket_{\ell \sqcup \ell'} <: \llbracket \tau \rrbracket_{\ell'} \\ = & \Sigma; \Psi \vdash \llbracket \mathbf{A}^{\ell''} \rrbracket_{\ell \sqcup \ell'} <: \llbracket \mathbf{A}^{\ell''} \rrbracket_{\ell'} && (\tau = \mathbf{A}^{\ell''}) \\ = & \Sigma; \Psi \vdash (\llbracket \mathbf{A} \rrbracket_{\ell \sqcup \ell' \sqcup \ell''})^{\ell \sqcup \ell' \sqcup \ell''} <: (\llbracket \mathbf{A} \rrbracket_{\ell' \sqcup \ell''})^{\ell' \sqcup \ell''} && (\text{Definition of } \llbracket \cdot \rrbracket) \\ = & \Sigma; \Psi \vdash \llbracket \mathbf{A}^{\ell'} \rrbracket_{\ell \sqcup \ell''} <: \llbracket \mathbf{A}^{\ell'} \rrbracket_{\ell''} && (\text{Definition of } \llbracket \cdot \rrbracket) \end{aligned}$$

The last statement holds by Lemma 2.81, since $\Sigma; \Psi \vdash \ell \sqcup \ell'' \sqsubseteq \ell''$ follows from our earlier assertion that $\Sigma; \Psi \vdash \ell \sqsubseteq \ell''$. □

Lemma 2.83 (FG \rightsquigarrow FG⁻: Lemma for protection relation). $\forall \Sigma, \Psi, \alpha, \tau, \ell, \ell'$.

$$\Sigma, \alpha; \Psi \vdash \tau \searrow \ell \implies \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell[\ell'/\alpha], \text{ where } FV(\ell') \in \Sigma$$

Proof. Say $\tau = \mathbf{A}^{\ell_g}$

$$\frac{\frac{\frac{}{\Sigma, \alpha; \Psi \vdash \ell \sqsubseteq \ell_g} \text{By inversion on } \Sigma, \alpha; \Psi \vdash \tau \searrow \ell}}{\Sigma; \Psi \vdash \ell[\ell'/\alpha] \sqsubseteq \ell_g[\ell'/\alpha]} \text{Substitution over constraints}}{\Sigma; \Psi \vdash \mathbf{A}^{\ell_g}[\ell'/\alpha] \searrow \ell[\ell'/\alpha]} \text{Definition of } \searrow$$

□

Lemma 2.84 (FG \rightsquigarrow FG⁻: Substitution lemma). *For all ℓ, ℓ' the following hold:*

1. $\forall \tau. \llbracket \tau \rrbracket_{\ell}[\ell'/\alpha] = \llbracket \tau[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}$
2. $\forall \mathbf{A}. \llbracket \mathbf{A} \rrbracket_{\ell}[\ell'/\alpha] = \llbracket \mathbf{A}[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}$

Proof. Proof by simultaneous induction on τ and \mathbf{A}

Proof of statement (1)

$$\begin{aligned} & \text{Let } \tau = \mathbf{A}^{\ell_i} \\ & \llbracket \mathbf{A}^{\ell_i} \rrbracket_{\ell}[\ell'/\alpha] \\ = & (\llbracket \mathbf{A} \rrbracket_{\ell_i \sqcup \ell})^{\ell_i \sqcup \ell}[\ell'/\alpha] && \text{Definition of } \llbracket \cdot \rrbracket \\ = & (\llbracket \mathbf{A} \rrbracket_{\ell_i \sqcup \ell}[\ell'/\alpha])^{\ell_i[\ell'/\alpha] \sqcup \ell[\ell'/\alpha]} \\ = & (\llbracket \mathbf{A}[\ell'/\alpha] \rrbracket_{\ell_i[\ell'/\alpha] \sqcup \ell[\ell'/\alpha]})^{\ell_i[\ell'/\alpha] \sqcup \ell[\ell'/\alpha]} && \text{IH(2) on } \mathbf{A} \\ = & \llbracket (\mathbf{A}[\ell'/\alpha])^{\ell_i[\ell'/\alpha]} \rrbracket_{\ell[\ell'/\alpha]} \\ = & \llbracket \mathbf{A}^{\ell_i}[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} \end{aligned}$$

Proof of statement (2)

We consider cases of \mathbf{A}

1. $\mathbf{A} = \mathbf{b}$:

$$\begin{aligned} & \llbracket \mathbf{b} \rrbracket_{\ell}[\ell'/\alpha] \\ = & \mathbf{b}[\ell'/\alpha] && (\text{Definition of } \llbracket \cdot \rrbracket) \\ = & \mathbf{b} && \alpha \notin FV(\mathbf{b}) \\ = & \llbracket \mathbf{b} \rrbracket_{\ell} && (\text{Definition of } \llbracket \cdot \rrbracket) \\ = & \llbracket \mathbf{b}[\ell'/\alpha] \rrbracket_{\ell} \end{aligned}$$

2. $A = \text{ref } \tau_i$:

$$\begin{aligned}
& \llbracket \text{ref } \tau_i \rrbracket_{\ell}[\ell'/\alpha] \\
&= \text{ref } \llbracket \tau_i \rrbracket_{\perp}[\ell'/\alpha] && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \text{ref } (\llbracket \tau_i \rrbracket_{\perp}[\ell'/\alpha]) \\
&= \text{ref } \llbracket \tau_i[\ell'/\alpha] \rrbracket_{\perp} && \text{IH(1) on } \tau_i \\
&= \llbracket \text{ref } \tau_i[\ell'/\alpha] \rrbracket_{\ell}
\end{aligned}$$

3. $A = \tau_1 \times \tau_2$:

$$\begin{aligned}
& \llbracket \tau_1 \times \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= (\llbracket \tau_1 \rrbracket_{\ell} \times \llbracket \tau_2 \rrbracket_{\ell})[\ell'/\alpha] && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket \tau_1 \rrbracket_{\ell}[\ell'/\alpha] \times \llbracket \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= \llbracket \tau_1[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} \times \llbracket \tau_2[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \llbracket (\tau_1[\ell'/\alpha] \times \tau_2[\ell'/\alpha]) \rrbracket_{\ell[\ell'/\alpha]} && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket (\tau_1 \times \tau_2)[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

4. $A = \tau_1 + \tau_2$:

$$\begin{aligned}
& \llbracket \tau_1 + \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= (\llbracket \tau_1 \rrbracket_{\ell} + \llbracket \tau_2 \rrbracket_{\ell})[\ell'/\alpha] && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket \tau_1 \rrbracket_{\ell}[\ell'/\alpha] + \llbracket \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= \llbracket \tau_1[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} + \llbracket \tau_2[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]} && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
&= \llbracket (\tau_1[\ell'/\alpha] + \tau_2[\ell'/\alpha]) \rrbracket_{\ell[\ell'/\alpha]} && \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \llbracket (\tau_1 + \tau_2)[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

5. $A = \tau_1 \xrightarrow{\ell_e} \tau_2$:

$$\begin{aligned}
& \llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_{\ell}[\ell'/\alpha] \\
&= \forall \beta_1. \beta_1, (\forall \beta. \beta_1, ((\ell \sqsubseteq \beta_1 \sqsubseteq \ell_e \wedge \beta \sqsubseteq \beta_1) \xrightarrow{\beta_1} (\llbracket \tau_1 \rrbracket_{\beta} \xrightarrow{\beta_1} \llbracket \tau_2 \rrbracket_{\beta_1})^{\beta_1})^{\beta_1})^{\beta_1}[\ell'/\alpha] \\
& \quad \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \forall \beta_1. \beta_1, (\forall \beta. \beta_1, ((\ell[\ell'/\alpha] \sqsubseteq \beta_1 \sqsubseteq \ell_e[\ell'/\alpha] \wedge \beta \sqsubseteq \beta_1) \xrightarrow{\beta_1} (\llbracket \tau_1 \rrbracket_{\beta}[\ell'/\alpha] \xrightarrow{\beta_1} \llbracket \tau_2 \rrbracket_{\beta_1}[\ell'/\alpha])^{\beta_1})^{\beta_1})^{\beta_1} \\
&= \forall \beta_1. \beta_1, (\forall \beta. \beta_1, ((\ell[\ell'/\alpha] \sqsubseteq \beta_1 \sqsubseteq \ell_e[\ell'/\alpha] \wedge \beta \sqsubseteq \beta_1) \xrightarrow{\beta_1} (\llbracket \tau_1[\ell'/\alpha] \rrbracket_{\beta} \xrightarrow{\beta_1} \llbracket \tau_2[\ell'/\alpha] \rrbracket_{\beta_1})^{\beta_1})^{\beta_1})^{\beta_1} \\
& \quad \text{(IH1 on } \tau_1 \text{ and } \tau_2 \text{)} \\
&= \llbracket (\tau_1[\ell'/\alpha] \xrightarrow{\ell_e[\ell'/\alpha]} \tau_2[\ell'/\alpha]) \rrbracket_{\ell[\ell'/\alpha]} \\
&= \llbracket (\tau_1 \xrightarrow{\ell_e} \tau_2)[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

6. $A = \forall \gamma. \tau_i$:

$$\begin{aligned}
& \llbracket \forall \gamma. \tau_i \rrbracket_{\ell}[\ell'/\alpha] \\
&= \forall \beta. \beta, ((\ell \sqsubseteq \beta \sqsubseteq \ell_e) \xrightarrow{\beta} (\forall \gamma. \beta, \llbracket \tau_i \rrbracket_{\beta})^{\beta})^{\beta}[\ell'/\alpha] \\
& \quad \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
&= \forall \beta. \beta, ((\ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} (\forall \gamma. \beta, \llbracket \tau_i \rrbracket_{\beta}[\ell'/\alpha])^{\beta})^{\beta} \\
&= \forall \beta. \beta, ((\ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} (\forall \gamma. \beta, \llbracket \tau_i[\ell'/\alpha] \rrbracket_{\beta})^{\beta})^{\beta} \\
& \quad \text{IH1 on } \tau_i \\
&= \llbracket \forall \beta. \ell_e[\ell'/\alpha], \tau_i[\ell'/\alpha] \rrbracket_{\ell[\ell'/\alpha]}
\end{aligned}$$

7. $A = c \Rightarrow \tau_i$:

$$\begin{aligned}
& \llbracket c \Rightarrow \tau_i \rrbracket_{\ell}[\ell'/\alpha] \\
= & \forall \beta. \beta, (((c \wedge \ell \sqsubseteq \beta \sqsubseteq \ell_e) \xrightarrow{\beta} \llbracket \tau \rrbracket_{\beta})^{\beta})^{\beta}[\ell'/\alpha] \\
& \text{(Definition of } \llbracket \cdot \rrbracket \text{)} \\
= & \forall \beta. \beta, (((c[\ell'/\alpha] \wedge \ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} \llbracket \tau \rrbracket_{\beta}[\ell'/\alpha])^{\beta})^{\beta} \\
= & \forall \beta. \beta, (((c[\ell'/\alpha] \wedge \ell[\ell'/\alpha] \sqsubseteq \beta \sqsubseteq \ell_e[\ell'/\alpha]) \xrightarrow{\beta} \llbracket \tau[\ell'/\alpha] \rrbracket_{\beta})^{\beta})^{\beta} \\
& \text{IH1 on } \tau_i \\
= & \left[\left[(c[\ell'/\alpha] \xrightarrow{\ell_e[\ell'/\alpha]} \tau_i[\ell'/\alpha]) \right] \right]_{\ell[\ell'/\alpha]} \\
= & \left[(c \xrightarrow{\ell_e} \tau_i)[\ell'/\alpha] \right]_{\ell[\ell'/\alpha]}
\end{aligned}$$

□

Lemma 2.85 (FG \rightsquigarrow FG⁻: Preservation of protection relation). $\forall \tau, \ell, \ell'$.

$$\tau \searrow \ell \implies \llbracket \tau \rrbracket_{\ell'} \searrow \ell$$

Proof. Let $\tau = \mathbf{A}^{\ell_i}$

$$\begin{array}{c}
\frac{}{\tau \searrow \ell} \text{ Given} \\
\frac{}{\ell \sqsubseteq \ell_i} \text{ Given} \\
\frac{}{\ell \sqsubseteq (\ell' \sqcup \ell_i)} \text{ By inversion} \\
\frac{}{(\llbracket \mathbf{A} \rrbracket_{\ell' \sqcup \ell_i})^{\ell' \sqcup \ell_i} \searrow \ell} \text{ Definition of } \llbracket \cdot \rrbracket \\
\frac{}{\llbracket \mathbf{A}^{\ell_i} \rrbracket_{\ell'} \searrow \ell} \\
\hline
\llbracket \tau \rrbracket_{\ell'} \searrow \ell
\end{array}$$

□

2.5 FG to CG translation

2.5.1 Type directed (direct) translation from FG to CG

Definition 2.86 (FG \rightsquigarrow CG: Type translation).

$$\begin{aligned}
\langle \mathbf{b} \rangle_\ell &= \mathbf{b} \\
\langle \mathbf{unit} \rangle_\ell &= \mathbf{unit} \\
\langle \tau_1 \xrightarrow{\ell_e} \tau_2 \rangle_\ell &= \forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha \\
\langle \tau_1 \times \tau_2 \rangle_\ell &= \langle \tau_1 \rangle_\ell \times \langle \tau_2 \rangle_\ell \\
\langle \tau_1 + \tau_2 \rangle_\ell &= \langle \tau_1 \rangle_\ell + \langle \tau_2 \rangle_\ell \\
\langle \mathbf{ref} \ A^{\ell'} \rangle_\ell &= \mathbf{ref} \ \ell' \ \langle A \rangle_{\ell'} \\
\langle \forall \alpha. (\ell_e, \tau) \rangle_\ell &= \forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha'} \\
\langle c \xrightarrow{\ell_e} \tau \rangle_\ell &= \forall \alpha, \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_\alpha \\
\langle A^{\ell'} \rangle_\ell &= \mathbf{Labeled} \ (\ell \sqcup \ell') \ \langle A \rangle_{\ell \sqcup \ell'}
\end{aligned}$$

For $\Gamma = x_1 : \tau_1, \dots, x_n : \tau_n$ and $\bar{\ell} = \ell_1, \dots, \ell_n$, define $\langle \Gamma \rangle_{\bar{\ell}} = x_1 : \langle \tau_1 \rangle_{\ell_1}, \dots, x_n : \langle \tau_n \rangle_{\ell_n}$.

We use a coercion function defined as follows:

$$\begin{aligned}
\mathbf{coerce_taint} &: \mathbb{C} \gamma \alpha_c \tau' \rightarrow \mathbb{C} \gamma \gamma \tau' \quad \text{when } \tau' = \mathbf{Labeled} \ \alpha'_c \ \tau \text{ and } \Sigma, \Psi \models \alpha_c \sqsubseteq \alpha'_c \\
\mathbf{coerce_taint} &\triangleq \lambda x. \mathbf{toLabeled}(\mathbf{bind}(x, y. \mathbf{unlabel}(y)))
\end{aligned}$$

$$\begin{aligned}
&\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \mathbf{ret} \ x} \text{FC-var} \\
&\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \mathbf{ret}(\mathbf{Lb}(\mathbf{\Lambda}\mathbf{\Lambda}\mathbf{\Lambda}(\nu(\lambda x. e_{c1}))))} \text{FC-lam} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 \ e_2 : \tau_2 \rightsquigarrow \mathbf{coerce_taint}(\mathbf{bind}(e_{c1}, a. \mathbf{bind}(e_{c2}, b. \mathbf{bind}(\mathbf{unlabel} \ a, c. (c \bullet b)))))} \text{FC-app} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_{c1}, a. \mathbf{bind}(e_{c2}, b. \mathbf{ret}(\mathbf{Lb}(a, b))))} \text{FC-prod} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \mathbf{fst}(e) : \tau_1 \rightsquigarrow \mathbf{coerce_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ (a), b. \mathbf{ret}(\mathbf{fst}(b)))))} \text{FC-fst} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \mathbf{snd}(e) : \tau_2 \rightsquigarrow \mathbf{coerce_taint}(\mathbf{bind}(e_c, a. \mathbf{bind}(\mathbf{unlabel} \ (a), b. \mathbf{ret}(\mathbf{snd}(b)))))} \text{FC-snd} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \mathbf{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{ret}(\mathbf{Lbinl}(a)))} \text{FC-inl} \\
&\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \mathbf{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \mathbf{bind}(e_c, a. \mathbf{ret}(\mathbf{Lbinr}(a)))} \text{FC-inr}
\end{aligned}$$

$$\begin{array}{c}
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))))} \text{FC-case} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } (e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb}b)))} \text{FC-ref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{FC-deref} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{FC-assign} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha_g. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda (\nu(e_c))))} \text{FC-FI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha_g. (\ell_e, \tau))^\ell \rightsquigarrow e_c \quad \text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e[] : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])))} \text{FC-FE} \\
\\
\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \stackrel{\ell_e}{\Rightarrow} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda (\nu(e_c))))} \text{FC-CI} \\
\\
\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \stackrel{\ell_e}{\Rightarrow} \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])))} \text{FC-CE}
\end{array}$$

2.5.2 Type preservation for FG to CG translation

Lemma 2.87 (Coercion lemma - typing). $\forall \Sigma, \Psi, \Gamma, \alpha_c, \alpha'_c, \tau.$

$$\Sigma, \Psi \models \alpha_c \sqsubseteq \alpha'_c \implies$$

$$\Sigma; \Psi; \Gamma \vdash \text{coerce_taint} : \mathbb{C} \ \gamma \ \alpha_c \ \text{Labeled } \alpha'_c \ \tau \rightarrow \mathbb{C} \ \gamma \ \gamma \ \text{Labeled } \alpha'_c \ \tau$$

Proof. $T_{c4} = \text{Labeled } \alpha'_c \ \tau$

$$T_{c3} = \mathbb{C} \ \alpha_c \ \alpha'_c \ \tau$$

$$T_{c2} = \mathbb{C} \ \gamma \ \alpha'_c \ \tau$$

$$T_{c1} = \mathbb{C} \ \gamma \ \gamma \ \text{Labeled } \alpha'_c \ \tau$$

$$T_{c0} = \mathbb{C} \ \gamma \ \alpha_c \ \text{Labeled } \alpha'_c \ \tau$$

$$T_c = T_{c0} \rightarrow T_{c1}$$

Pc2:

$$\frac{\frac{\Sigma; \Psi; \Gamma, x : T_{c0}, y : T_{c4} \vdash y : T_{c4}}{\Sigma; \Psi; \Gamma, x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-unlabel} \quad \frac{}{\Sigma, \Psi \models \alpha_c \sqsubseteq \alpha'_c} \text{Given}}{\Sigma; \Psi; \Gamma, x : T_{c0}, y : T_{c4} \vdash \text{unlabel}(y) : T_{c3}} \text{CG-var}$$

Pc1:

$$\frac{}{\Sigma; \Psi; \Gamma, x : T_{c0} \vdash x : T_{c0}} \text{CG-var}$$

Pc0:

$$\frac{\frac{Pc1 \quad Pc2}{\Sigma; \Psi; \Gamma, x : T_{c0} \vdash \text{bind}(x, y.\text{unlabel}(y)) : T_{c2}} \text{CG-bind}}{\Sigma; \Psi; \Gamma, x : T_{c0} \vdash \text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_{c1}} \text{CG-tolabeled}$$

Pc:

$$\frac{\frac{Pc0}{\Sigma; \Psi; \Gamma \vdash \lambda x.\text{toLabeled}(\text{bind}(x, y.\text{unlabel}(y))) : T_c} \text{CG-lam}}{\Sigma; \Psi; \Gamma \vdash \text{coerce_taint} : T_c} \text{From Definition of coerce_taint}$$

□

Theorem 2.88 (FG \rightsquigarrow CG: Type preservation). *Suppose $\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau$ in FG. Then, there exists e' such that $\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e'$ and for any α', β', γ' with $\beta' \sqcup \gamma' \sqsubseteq pc \sqcap \alpha'$, there is a derivation of $\Sigma; \Psi; (\Gamma)_{\beta'} \vdash e' : \mathbb{C} \gamma' \gamma' (\tau)_{\alpha'}$ in CG.*

Proof. Proof by induction on the \rightsquigarrow relation

1. FC-var:

$$\frac{\frac{\frac{}{(\Gamma)_{\beta'_o}(x) = (\tau)_{\beta'_o}} \text{Given}}{\Sigma; \Psi; (\Gamma)_{\beta'_o} \vdash x : (\tau)_{\beta'_o}} \text{CG-var} \quad \frac{\frac{}{\Sigma; \Psi \vdash \beta'_o \sqcup \gamma'_o \sqsubseteq \alpha'_o \sqcap pc} \text{Given}}{\Sigma; \Psi \vdash \beta'_o \sqsubseteq \alpha'_o} \text{Lemma 2.89, CG-sub}}{\Sigma; \Psi; (\Gamma)_{\beta'_o} \vdash x : (\tau)_{\alpha'_o}} \text{CG-ret}}{\Sigma; \Psi; (\Gamma)_{\beta'_o} \vdash \text{ret } x : \mathbb{C} \gamma'_o \gamma'_o (\tau)_{\alpha'_o}} \text{CG-ret}$$

2. FC-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e : \tau_2 \rightsquigarrow e_{c1}}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x.e : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x.e_{c1}))))} \text{FC-lam}$$

$$\begin{aligned} T_0 &= \mathbb{C} \gamma'_j \gamma'_j ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp)_{\alpha'_j} = \mathbb{C} \gamma'_j \gamma'_j \text{Labeled } \alpha'_j ((\tau_1 \xrightarrow{\ell_e} \tau_2))_{\alpha'_j} \\ T_1 &= \mathbb{C} \gamma'_j \gamma'_j \text{Labeled } \alpha'_j \forall \alpha_t, \beta_t, \gamma_t. (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.0} &= \text{Labeled } \alpha'_j \forall \alpha_t, \beta_t, \gamma_t. (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.1} &= \forall \alpha_t, \beta_t, \gamma_t. (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.2} &= (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.3} &= (\tau_1)_{\beta_t} \rightarrow \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \\ T_{1.4} &= \mathbb{C} \gamma_t \gamma_t (\tau_2)_{\alpha_t} \end{aligned}$$

P3:

$$\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \overline{\beta'_j} \sqcup \gamma_j \sqsubseteq \alpha'_j \sqcap pc \quad \text{Given, Weakening}}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \overline{\beta'_j} \sqsubseteq \alpha'_j}$$

P2:

$$\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e \quad P3}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e) \vdash \overline{\beta'_j} \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e}$$

P1:

$$\frac{\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e); (\Gamma)_{\overline{\beta'_j}}, x : (\tau_1)_{\beta_t} \vdash e_{c1} : T_{1.4} \quad \text{IH}}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e); (\Gamma)_{\overline{\beta'_j}} \vdash \lambda x.e_{c1} : T_{1.3}} \quad P2}{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi, (\alpha'_j \sqcup \beta_t \sqcup \gamma_t \sqsubseteq \alpha_t \sqcap \ell_e); (\Gamma)_{\overline{\beta'_j}} \vdash \lambda x.e_{c1} : T_{1.3}} \quad \text{CG-lam}$$

P0:

$$\frac{\Sigma; \Psi \vdash \overline{\beta'_j} \sqcup \gamma'_j \sqsubseteq \alpha'_j \quad \text{Given}}{\Sigma; \Psi \vdash \gamma_j \sqsubseteq \alpha_j}$$

Main derivation:

$$\frac{\frac{\frac{\Sigma, \alpha_t, \beta_t, \gamma_t; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \nu(\lambda x.e_{c1}) : T_{1.2} \quad P1}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \Lambda\Lambda\Lambda(\nu(\lambda x.e_{c1})) : T_{1.1}} \quad \text{CG-CI}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x.e_{c1}))) : T_{1.0}} \quad \text{3 applications CG-FI} \quad P0}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'_j}} \vdash \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x.e_{c1})))) : T_1} \quad \text{CG-label} \quad \text{CG-ret}$$

3. FC-app:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_1 \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 e_2 : \tau_2 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c \square \square \square \bullet) b))))} \quad \text{FC-app}$$

$$\beta' = \bigcup_{\beta_i \in \overline{\beta'}} \beta_i$$

$$T_0 = \mathbb{C} \gamma' \gamma' ((\tau_1 \xrightarrow{\ell_e} \tau_2)^\ell)_{\beta' \sqcup \gamma'} = \mathbb{C} \gamma' \gamma' \text{Labeled } (\beta' \sqcup \gamma' \sqcup \ell) ((\tau_1 \xrightarrow{\ell_e} \tau_2))_{\beta' \sqcup \gamma' \sqcup \ell}$$

$$T_1 = \mathbb{C} \gamma' \gamma' \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.1} = \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.2} = \mathbb{C} \gamma' (\gamma' \sqcup (\beta' \sqcup \gamma') \sqcup \ell) \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.3} = \forall \alpha, \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.4} = \forall \beta, \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup \beta \sqcup \gamma \sqsubseteq ((\beta' \sqcup \gamma') \sqcup \ell) \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.5} = \forall \gamma. (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup (\beta' \sqcup \gamma') \sqcup \gamma \sqsubseteq ((\beta' \sqcup \gamma') \sqcup \ell) \sqcap \ell_e) \Rightarrow (\tau_1)_{(\beta' \sqcup \gamma')} \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.6} = (((\beta' \sqcup \gamma') \sqcup \ell) \sqcup (\beta' \sqcup \gamma') \sqcup (\beta' \sqcup \gamma' \sqcup \ell) \sqsubseteq ((\beta' \sqcup \gamma') \sqcup \ell) \sqcap \ell_e) \Rightarrow T_{1.7}$$

$$T_{1.7} = (\tau_1)_{(\beta' \sqcup \gamma')} \rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.8} = \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.9} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\tau_2)_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.10} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A}^{\ell_i})_{((\beta' \sqcup \gamma') \sqcup \ell)}$$

$$T_{1.11} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{Labeled} (\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A})_{(\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{1.12} = \mathbb{C} (\gamma') (\gamma') \text{Labeled} (\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell) (\mathbf{A})_{(\ell_i \sqcup \beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{1.13} = \mathbb{C} (\gamma') (\gamma') \text{Labeled} (\ell_i \sqcup \beta' \sqcup \gamma') (\mathbf{A})_{(\ell_i \sqcup \beta' \sqcup \gamma')}$$

$$T_2 = \mathbb{C} (\gamma') (\gamma') (\tau_2)_{(\beta' \sqcup \gamma')}$$

P8:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash b : (\tau_2)_{(\beta' \sqcup \gamma')}} \text{CG-var}$$

P7:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e} \text{Given}}{\Sigma; \Psi \vdash pc \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \alpha' \sqcap pc \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq \alpha' \sqcap pc \sqsubseteq \ell_e}$$

P6:

$$\frac{P7 \quad \frac{}{\Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e} \text{Given}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_e}}{\Sigma; \Psi \vdash (\ell \sqcup \beta' \sqcup \gamma') \sqsubseteq \ell_e}$$

P5:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c : T_{1.3}} \text{CG-var}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c[] : T_{1.4}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c[][] : T_{1.5}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c[][][] : T_{1.6}} \text{CG-FE} \quad P6}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash c[][][] \bullet : T_{1.7}} \text{CG-CE}$$

P4:

$$\frac{\frac{P5 \quad P8}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')}, c : T_{1.3} \vdash (c[][][] \bullet) b : T_{1.8}} \text{CG-app}}$$

P3:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')} \vdash a : T_{1.1}} \text{CG-var}$$

P2:

$$\frac{\frac{P3}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')} \vdash \text{unlabel } a : T_{1.2}} \text{CG-unlabel} \quad P4}}{\Sigma; \Psi; (\Gamma)_{\overline{\beta'}}, a : T_{1.1}, b : (\tau_2)_{(\beta' \sqcup \gamma')} \vdash \text{bind}(\text{unlabel } a, c.(c[][][] \bullet) b) : T_{1.9}} \text{CG-bind}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash e_{c2} : T_2 \quad \text{IH2, Weakening} \quad P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)) : T_{1.9}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash e_{c1} : T_1} \text{IH1 with } (\beta' \sqcup \gamma'), \bar{\beta}', \gamma' \quad P1$$

Main derivation:

$$\frac{\frac{\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash e_{c1} : T_1}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)) : T_{1.9}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)) : T_{1.10}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)) : T_{1.11}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{coerce_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) : T_{1.12}} \text{Lemma 2.87}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{coerce_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) : T_{1.13}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{1.1} \vdash \text{coerce_taint}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) : T_2} \text{CG-bind}$$

4. FC-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : \tau_1 \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau_2 \rightsquigarrow e_{c2}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_1, e_2) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))))} \text{FC-prod}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\tau_1 \times \tau_2)^\perp)_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((\tau_1 \times \tau_2))_{\alpha'}$$

$$T_3 = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}$$

$$T_{3.1} = \text{Labeled } \alpha' (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}$$

$$T_4 = \mathbb{C} \gamma' \gamma' (\tau_1)_{\alpha'}$$

$$T_5 = \mathbb{C} \gamma' \gamma' (\tau_2)_{\alpha'}$$

P4:

$$\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}} \text{CG-var}$$

P3:

$$\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash b : (\tau_2)_{\alpha'}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash b : (\tau_2)_{\alpha'}} \text{CG-var}$$

P2:

$$\frac{\frac{\frac{\frac{P3 \quad P4}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash (a, b) : (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}} \text{CG-prod}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_2)_{\alpha'} \vdash \text{Lb}(a, b) : T_{3.1}} \text{CG-label}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_2)_{\alpha'} \vdash \text{ret}(\text{Lb}(a, b)) : T_3} \text{CG-ret}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'}, b : (\tau_1)_{\alpha'} \vdash (a, b) : (\tau_1)_{\alpha'} \times (\tau_2)_{\alpha'}} \text{CG-label}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'} \vdash e_{c2} : T_5 \quad \text{IH2} \quad P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b))) : T_3} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : (\tau_1)_{\alpha'} \vdash e_{c2} : T_5} \text{IH2} \quad P2$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_{c1} : T_4} \text{IH1} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_3} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{ret}(\text{Lb}(a, b)))) : T_1} \text{Definition 2.86}}$$

5. FC-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e) : \tau_1 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))))} \text{FC-fst}$$

$$T_1 = \mathbb{C} \gamma' \gamma' (\tau_1)_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' ((\tau_1 \times \tau_2)^\ell)_{\alpha'}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup \alpha' ((\tau_1 \times \tau_2)^\ell)_{\alpha' \sqcup \ell}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup \alpha' (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup \alpha' (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{2.4} = (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_{2.5} = \mathbb{C} (\gamma') (\gamma' \sqcup \alpha' \sqcup \ell) (\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}$$

$$T_3 = \mathbb{C} (\gamma' \sqcup \alpha' \sqcup \ell) (\gamma' \sqcup \alpha' \sqcup \ell) (\tau_1)_{\alpha' \sqcup \ell}$$

$$T_{3.1} = \mathbb{C} (\gamma') (\gamma' \sqcup \alpha' \sqcup \ell) (\tau_1)_{\alpha' \sqcup \ell}$$

$$T_{3.2} = \mathbb{C} (\gamma') (\alpha' \sqcup \ell) (\tau_1)_{\alpha' \sqcup \ell}$$

$$T_{3.3} = \mathbb{C} (\gamma') (\alpha' \sqcup \ell) (\mathbf{A}^{\ell_i})_{\alpha' \sqcup \ell}$$

$$T_{3.4} = \mathbb{C} (\gamma') (\alpha' \sqcup \ell) \text{Labeled } \ell \sqcup \ell_i \sqcup \alpha' (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i}$$

$$T_{3.5} = \mathbb{C} (\gamma') (\gamma') \text{Labeled } \ell \sqcup \ell_i \sqcup \alpha' (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i}$$

$$T_{3.6} = \mathbb{C} (\gamma') (\gamma') \text{Labeled } \ell_i \sqcup \alpha' (\mathbf{A})_{\alpha' \sqcup \ell_i}$$

$$T_{3.7} = \mathbb{C} (\gamma') (\gamma') (\mathbf{A}^{\ell_i})_{\alpha'}$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}} \text{CG-var}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{fst}(b) : (\tau_1)_{\alpha' \sqcup \ell}} \text{CG-fst}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{fst}(b)) : T_3} \text{CG-ret}}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}} \text{CG-unlabel} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))) : T_{3.1}} \text{CG-bind}}$$

P0:

$$\begin{array}{c}
\frac{}{\Sigma; \Psi; (\Gamma)_{\beta'} \vdash e_c : T_{2.2}} \text{IH} \quad P1 \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.1} \quad \text{CG-bind} \\
\frac{}{\Sigma; \Psi \vdash \gamma' \sqsubseteq \alpha'} \text{Given} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.2} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.3} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))) : T_{3.4} \quad \text{Definition 2.86} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.5} \quad \text{Lemma 2.87}
\end{array}$$

Main derivation:

$$\begin{array}{c}
P0 \quad \frac{\frac{}{\Sigma; \Psi \vdash \mathbf{A}^{\ell_i} \searrow \ell} \text{By inversion}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{By inversion} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.6} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_{3.7} \quad \text{Definition 2.86} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) : T_1
\end{array}$$

6. FC-snd:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{snd}(e) : \tau_2 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b))))))} \text{FC-snd}$$

$$\begin{aligned}
T_1 &= \mathbb{C} \ \gamma' \ \gamma' \ (\tau_2)_{\alpha'} \\
T_2 &= \mathbb{C} \ \gamma' \ \gamma' \ ((\tau_1 \times \tau_2)^\ell)_{\alpha'} \\
T_{2.1} &= \mathbb{C} \ \gamma' \ \gamma' \ \text{Labeled } \ell \sqcup \alpha' \ ((\tau_1 \times \tau_2)_{\alpha' \sqcup \ell}) \\
T_{2.2} &= \mathbb{C} \ \gamma' \ \gamma' \ \text{Labeled } \ell \sqcup \alpha' \ ((\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}) \\
T_{2.3} &= \text{Labeled } \ell \sqcup \alpha' \ ((\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}) \\
T_{2.4} &= ((\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}) \\
T_{2.5} &= \mathbb{C} \ (\gamma') \ (\gamma' \sqcup \alpha' \sqcup \ell) \ ((\tau_1)_{\alpha' \sqcup \ell} \times (\tau_2)_{\alpha' \sqcup \ell}) \\
T_3 &= \mathbb{C} \ (\gamma' \sqcup \alpha' \sqcup \ell) \ (\gamma' \sqcup \alpha' \sqcup \ell) \ ((\tau_2)_{\alpha' \sqcup \ell}) \\
T_{3.1} &= \mathbb{C} \ (\gamma') \ (\gamma' \sqcup \alpha' \sqcup \ell) \ ((\tau_2)_{\alpha' \sqcup \ell}) \\
T_{3.2} &= \mathbb{C} \ (\gamma') \ (\alpha' \sqcup \ell) \ ((\tau_2)_{\alpha' \sqcup \ell}) \\
T_{3.3} &= \mathbb{C} \ (\gamma') \ (\alpha' \sqcup \ell) \ (\mathbf{A}^{\ell_i})_{\alpha' \sqcup \ell} \\
T_{3.4} &= \mathbb{C} \ (\gamma') \ (\alpha' \sqcup \ell) \ \text{Labeled } \ell \sqcup \ell_i \sqcup \alpha' \ (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i} \\
T_{3.5} &= \mathbb{C} \ (\gamma') \ (\gamma') \ \text{Labeled } \ell \sqcup \ell_i \sqcup \alpha' \ (\mathbf{A})_{\alpha' \sqcup \ell \sqcup \ell_i} \\
T_{3.6} &= \mathbb{C} \ (\gamma') \ (\gamma') \ \text{Labeled } \ell_i \sqcup \alpha' \ (\mathbf{A})_{\alpha' \sqcup \ell_i} \\
T_{3.7} &= \mathbb{C} \ (\gamma') \ (\gamma') \ (\mathbf{A}^{\ell_i})_{\alpha'}
\end{aligned}$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash b : T_{2.4}}{\text{CG-var}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{snd}(b) : (\tau_2)_{\alpha' \sqcup \ell}} \text{CG-snd}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.4} \vdash \text{ret}(\text{snd}(b)) : T_3} \text{CG-ret}$$

P1:

$$\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{unlabel}(a) : T_{2.5}}{\text{CG-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b))) : T_{3.1}} \text{CG-bind}$$

P0:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_{2.2}}{\text{IH}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.1}} \text{CG-bind}}{\Sigma; \Psi \vdash \gamma' \sqsubseteq \alpha'} \text{Given}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.2}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.3}} \text{Definition 2.86}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))) : T_{3.4}} \text{Lemma 2.87}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_{3.5}} \text{Lemma 2.87}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \mathbf{A}^{\ell_i} \searrow \ell} \text{By inversion}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{By inversion}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_{3.6}} \text{Definition 2.86}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_{3.7}} \text{Definition 2.86}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{snd}(b)))))) : T_1}$$

7. FC-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a)))} \text{FC-inl}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\tau_1 + \tau_2)^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((\tau_1 + \tau_2))_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_{1.3} = \text{Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' (\tau_1)_{\alpha'}$$

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}}{\text{CG-var}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{inl}(a) : (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}} \text{CG-inl}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{Lbinl}(a) : T_{1.3}} \text{CG-label}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{ret}(\text{Lbinl}(a)) : T_{1.2}} \text{CG-ret}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_2} \text{IH} \quad P1}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a))) : T_{1.2}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinl}(a))) : T_1} \text{Definition 2.86}$$

8. FC-inr:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_2 \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inr}(e) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a)))} \text{FC-inr}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\tau_1 + \tau_2)^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((\tau_1 + \tau_2))_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_{1.3} = \text{Labeled } \alpha' (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' (\tau_2)_{\alpha'}$$

P1:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash a : (\tau_1)_{\alpha'}} \text{CG-var}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{inr}(a) : (\tau_1)_{\alpha'} + (\tau_2)_{\alpha'}} \text{CG-inr}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{Lbinr}(a) : T_{1.3}} \text{CG-label}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : (\tau_1)_{\alpha'} \vdash \text{ret}(\text{Lbinr}(a)) : T_{1.2}} \text{CG-ret}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_c : T_2} \text{IH} \quad P1}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a))) : T_{1.2}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_c, a.\text{ret}(\text{Lbinr}(a))) : T_1} \text{Definition 2.86}$$

9. FC-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_1 : \tau \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_2 : \tau \rightsquigarrow e_{c2} \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e, x.e_1, y.e_2) : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))))} \text{FC-case}$$

$$\beta' = \bigcup_{\beta_i \in \vec{\beta}'} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' (\tau)_{(\alpha')}$$

$$T_2 = \mathbb{C} \gamma' \gamma' ((\tau_1 + \tau_2)^\ell)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) (\tau_1 + \tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) ((\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} + (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell})$$

$$\begin{aligned}
T_{2.3} &= \text{Labeled } ((\beta' \sqcup \gamma') \sqcup \ell) \ ((\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} + (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell}) \\
T_{2.4} &= \mathbb{C} \ \gamma' \ (\gamma' \sqcup (\beta' \sqcup \gamma') \sqcup \ell) \ ((\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} + (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell}) \\
T_{2.5} &= ((\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} + (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell}) \\
T_3 &= \mathbb{C} \ (\beta' \sqcup \gamma' \sqcup \ell) \ (\beta' \sqcup \gamma' \sqcup \ell) \ (\tau)_{(\beta' \sqcup \gamma' \sqcup \ell)} \\
T_4 &= \mathbb{C} \ (\gamma') \ (\beta' \sqcup \gamma' \sqcup \ell) \ (\tau)_{(\beta' \sqcup \gamma' \sqcup \ell)} \\
T_5 &= \mathbb{C} \ (\gamma') \ (\beta' \sqcup \gamma' \sqcup \ell) \ (\mathbf{A}^{\ell_i})_{(\beta' \sqcup \gamma' \sqcup \ell)} \\
T_{5.1} &= \mathbb{C} \ (\gamma') \ (\beta' \sqcup \gamma' \sqcup \ell) \ \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma' \sqcup \ell) \ (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma' \sqcup \ell)} \\
T_{5.2} &= \mathbb{C} \ (\gamma') \ (\gamma') \ \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma' \sqcup \ell) \ (\mathbf{A})_{\ell_i \sqcup (\beta' \sqcup \gamma' \sqcup \ell)} \\
T_{5.3} &= \mathbb{C} \ (\gamma') \ (\gamma') \ \text{Labeled } \ell_i \sqcup \beta' \sqcup \gamma' \ (\mathbf{A})_{\ell_i \sqcup \beta' \sqcup \gamma'} \\
T_{5.4} &= \mathbb{C} \ (\gamma') \ (\gamma') \ (\mathbf{A}^{\ell_i})_{\beta' \sqcup \gamma'} \\
T_{5.5} &= \mathbb{C} \ (\gamma') \ (\gamma') \ (\tau)_{\beta' \sqcup \gamma'} \\
\text{P2:} &
\end{aligned}$$

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{CG-var}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5}, x : (\tau_1)_{(\beta' \sqcup \gamma') \sqcup \ell} \vdash e_{c1} : T_3} \text{IH2, Weakening}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5}, y : (\tau_2)_{(\beta' \sqcup \gamma') \sqcup \ell} \vdash e_{c2} : T_3} \text{IH3, Weakening}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{case}(b, x.e_{c1}, y.e_{c2}) : T_3} \text{CG-case}}$$

$$\text{P1:} \quad \frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\text{CG-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})) : T_4} \text{CG-bind}}$$

$$\text{P0:} \quad \frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2}}{\text{IH1}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_4} \text{CG-bind}}$$

$$\text{P0.1:} \quad \frac{\frac{\frac{}{P0}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_5} \text{Definition 2.86}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2}))) : T_{5.1}} \text{Lemma 2.87}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.2}}$$

Main derivation:

$$\text{P0.1} \quad \frac{\frac{\frac{}{\Sigma; \Psi \vdash \mathbf{A}^{\ell_i} \searrow \ell} \text{Given}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{By inversion}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.3}} \text{Definition 2.86}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.4}} \text{Definition 2.86}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_{5.5}} \text{Given}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, e_c : T_{2.2} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{c1}, y.e_{c2})))) : T_1} \text{CG-}$$

10. FC-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau \searrow_{pc}}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } (e) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})))} \text{FC-ref}$$

$$\beta' = \bigcup_{\beta_i \in \bar{\beta}'} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\text{ref } \tau)^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' ((\text{ref } A^{\ell_i})^\perp)_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((\text{ref } A^{\ell_i}))_{\alpha'}$$

$$T_{1.3} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \text{ref } \ell_i (A)_{\ell_i}$$

$$T_2 = \mathbb{C} \gamma' \gamma' (\tau)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' (A^{\ell_i})_{(\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.3} = \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.4} = \mathbb{C} \gamma' \gamma' \text{ref } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.5} = \text{ref } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.51} = \text{Labeled } \alpha' \text{ref } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.6} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \text{ref } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.7} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \text{ref } \ell_i (A)_{\ell_i}$$

P3:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash A^{\ell_i} \searrow_{pc}}{\text{Given}} \quad \text{By inversion} \quad \frac{}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq pc} \text{Given}}{\Sigma; \Psi \vdash pc \sqsubseteq \ell_i}}{\Sigma; \Psi \vdash (\beta' \sqcup \gamma') \sqsubseteq \ell_i}$$

P2:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{CG-var}} \quad \text{CG-label} \quad \frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{Lbb} : T_{2.51}}{\text{CG-ret}} \quad P3}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{2.6}}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash \text{ret}(\text{Lbb}) : T_{1.3}}$$

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{new } (a) : T_{2.4}}{\text{CG-new}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb})) : T_{1.3}} \text{CG-bind}}$$

Main derivation:

$$\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash e_c : T_{2.2}}{\text{IH}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb}))) : T_{1.3}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lbb}))) : T_1} \text{Definition 2.86}$$

11. FC-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (\text{ref } \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{FC-deref}$$

$$\beta' = \bigcup_{\beta_i \in \bar{\beta}'} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' (\tau')_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' (\mathbf{A}^{\ell'_i})_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell'_i \sqcup \alpha' (\mathbf{A}')_{\ell'_i \sqcup \alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' ((\text{ref } \tau)^\ell)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) ((\text{ref } \tau))_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) ((\text{ref } \mathbf{A}^{\ell_i}))_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.3} = \mathbb{C} \gamma' \gamma' \text{Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) (\text{ref } \ell_i (\mathbf{A}))_{\ell_i}$$

$$T_{2.4} = \text{Labeled } (\ell \sqcup (\beta' \sqcup \gamma')) (\text{ref } \ell_i (\mathbf{A}))_{\ell_i}$$

$$T_{2.5} = \mathbb{C} \gamma' \beta' \sqcup \gamma' \sqcup \ell (\text{ref } \ell_i (\mathbf{A}))_{\ell_i}$$

$$T_{2.6} = (\text{ref } \ell_i (\mathbf{A}))_{\ell_i}$$

$$T_{2.7} = \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) (\text{Labeled } \ell_i (\mathbf{A}))_{\ell_i}$$

$$T_{2.8} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\text{Labeled } \ell_i (\mathbf{A}))_{\ell_i}$$

$$T_{2.9} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) (\text{Labeled } \ell'_i (\mathbf{A}'))_{\ell'_i}$$

$$T_{2.10} = \mathbb{C} (\gamma') (\gamma') (\text{Labeled } \beta' \sqcup \gamma' \sqcup \ell \sqcup \ell'_i (\mathbf{A}'))_{\ell'_i}$$

$$T_{2.11} = \mathbb{C} (\gamma') (\gamma') (\text{Labeled } \alpha \sqcup \ell'_i (\mathbf{A}'))_{\ell'_i}$$

P2:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.4}, b : T_{2.6} \vdash b : T_{2.6}}{\text{CG-var}}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.4}, b : T_{2.6} \vdash !b : T_{2.7}} \text{CG-deref}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.4} \vdash \text{unlabel } a : T_{2.5}}{\text{CG-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.4} \vdash \text{bind}(\text{unlabel } a, b.!b) : T_{2.8}} \text{CG-bind}$$

P0:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash e_c : T_{2.3}}{P1}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.8}} \text{CG-bind}$$

Main derivation:

$$\frac{\frac{\frac{\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.9}}{P0}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b)) : T_{2.10}} \text{Lemma 2.87}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{2.10}} \text{Lemma 2.87}}{\frac{\frac{\Sigma; \Psi \vdash \mathbf{A}^{\ell_i} \searrow \ell}{\text{Given}}}{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_i} \text{By inversion}}{\frac{\frac{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq \alpha'}{\text{Given}}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a, b.!b))) : T_{1.1}} \text{CG-sub}}$$

12. FC-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 : (\text{ref } \tau)^\ell \rightsquigarrow e_{c1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_2 : \tau \rightsquigarrow e_{c2} \quad \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_1 := e_2 : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{FC-assign}$$

$$\beta' = \bigcup_{\beta_i \in \bar{\beta}'} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' (\text{unit})_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{unit}$$

$$T_2 = \mathbb{C} \gamma' \gamma' ((\text{ref } \tau)^\ell)_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') ((\text{ref } \tau))_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') ((\text{ref } A^{\ell_i}))_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.3} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \text{ref } \ell_i (A)_{\ell_i}$$

$$T_{2.4} = \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \text{ref } \ell_i (A)_{\ell_i}$$

$$T_{2.5} = \mathbb{C} \gamma' \ell \sqcup (\beta' \sqcup \gamma') \text{ref } \ell_i (A)_{\ell_i}$$

$$T_{2.6} = \text{ref } \ell_i (A)_{\ell_i}$$

$$T_{2.7} = \mathbb{C} \ell \sqcup (\beta' \sqcup \gamma') \ell \sqcup (\beta' \sqcup \gamma') \text{unit}$$

$$T_{2.8} = \mathbb{C} \gamma' \ell \sqcup (\beta' \sqcup \gamma') \text{unit}$$

$$T_{2.9} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \text{unit}$$

$$T_3 = \mathbb{C} \gamma' \gamma' (\tau)_{(\beta' \sqcup \gamma')}$$

$$T_{3.1} = \mathbb{C} \gamma' \gamma' (A^{\ell_i})_{(\beta' \sqcup \gamma')}$$

$$T_{3.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{3.3} = \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') (A)_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{3.4} = \text{Labeled } \ell_i (A)_{\ell_i}$$

P4:

$$\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c : T_{2.6}} \text{CG-var}$$

P5:

$$\frac{\frac{\frac{\frac{}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash b : T_{3.3}}{\text{CG-var}}}{\Sigma; \Psi \vdash \tau = A^{\ell_i} \searrow (pc \sqcup \ell)} \text{Given}}{\Sigma; \Psi \vdash (pc \sqcup \ell) \sqsubseteq \ell_i} \text{By inversion}}{\Sigma; \Psi \vdash \beta' \sqcup \gamma' \sqsubseteq \ell_i} \text{Given}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash b : T_{3.4}} \text{CG-var}$$

P3:

$$\frac{P4 \quad P5}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.4}, b : T_{3.3}, c : T_{2.6} \vdash c := b : T_{2.7}} \text{CG-assign}$$

P2:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3} \vdash \text{unlabel } a : T_{2.5}}{\text{CG-unlabel}} \quad P3}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4}, b : T_{3.3} \vdash \text{bind}(\text{unlabel } a, c.c := b) : T_{2.8}} \text{CG-bind}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash e_{c2} : T_{3.2}}{\text{IH2}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'}, a : T_{2.4} \vdash \text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)) : T_{2.8}} \text{CG-bind}$$

P0:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash e_{c1} : T_{2.3}}{\text{IH1}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))) : T_{2.8}} \text{CG-bind}$$

P0.1:

$$\frac{P0}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b)))) : T_{2.9}} \text{CG-toLabeled}$$

Main derivation:

$$\frac{P0.1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}()) : T_{1.1}} \text{CG-bind}$$

13. FC-FI:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e : (\forall \alpha_g. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_c))))} \text{FC-FI}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((\forall \alpha. (\ell_e, \tau))^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((\forall \alpha. (\ell_e, \tau))_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \forall \alpha. \forall \alpha_i, \gamma_i. (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_2 = \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.1} = (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.2} = \forall \alpha, \alpha_i, \gamma_i. (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{2.3} = \text{Labeled } \alpha' \forall \alpha, \alpha_i, \gamma_i. (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

Main derivation:

$$\frac{\frac{\frac{\frac{\Sigma, \alpha, \alpha_i, \gamma_i; \Psi, (\alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e); (\Gamma)_{\vec{\beta}'} \vdash e_c : T_2}{\text{IH, Weakening}}}{\Sigma, \alpha, \alpha_i, \gamma_i; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \nu(e_c) : T_{2.1}} \text{CG-CI}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \Lambda \Lambda \Lambda(\nu(e_c)) : T_{2.2}} \text{CG-FI}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{Lb}(\Lambda \Lambda \Lambda(\nu(e_c))) : T_{2.3}} \text{CG-label}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}'} \vdash \text{ret}(\text{Lb}(\Lambda \Lambda \Lambda(\nu(e_c)))) : T_{1.2}} \text{CG-ret}$$

14. FC-FE:

$$\frac{\text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi; \Gamma \vdash_{pc} e : (\forall \alpha_g. (\ell_e, \tau))^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e[] : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.\text{b}[]\bullet)))} \text{FC-FE}$$

$$\beta' = \bigcup_{\beta_i \in \overline{\beta'}} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' \langle \tau[\ell''/\alpha] \rangle_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' \langle (\forall \alpha. (\ell_e, \tau))^\ell \rangle_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \langle (\forall \alpha. (\ell_e, \tau)) \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}$$

$$T_{2.4} = \mathbb{C} \gamma' (\gamma' \sqcup \ell \sqcup (\beta' \sqcup \gamma')) \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}$$

$$T_{2.5} = \forall \alpha. \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}$$

$$T_{2.6} = \forall \alpha_i, \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\alpha_i}[\ell''/\alpha]$$

$$T_{2.7} = \forall \gamma_i. ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup \gamma_i \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.8} = ((\ell \sqcup (\beta' \sqcup \gamma')) \sqcup (\beta' \sqcup \gamma' \sqcup \ell) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.81} = ((\ell \sqcup (\beta' \sqcup \gamma')) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha]) \Rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.9} = \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.10} = \mathbb{C} (\gamma') \langle \tau \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}[\ell''/\alpha]$$

$$T_{2.11} = \mathbb{C} (\gamma') \langle \tau[\ell''/\alpha] \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{2.12} = \mathbb{C} (\gamma') \langle \mathbf{A}^{\ell_i}[\ell''/\alpha] \rangle_{(\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{2.13} = \mathbb{C} (\gamma') \langle \mathbf{A}[\ell''/\alpha] \rangle_{\ell_i[\ell''/\alpha] \sqcup (\beta' \sqcup \gamma' \sqcup \ell)}$$

$$T_{2.14} = \mathbb{C} (\gamma') \langle \mathbf{A}[\ell''/\alpha] \rangle_{\ell_i[\ell''/\alpha] \sqcup \beta' \sqcup \gamma'}$$

$$T_{2.15} = \mathbb{C} (\gamma') \langle \mathbf{A}[\ell''/\alpha] \rangle_{\ell_i[\ell''/\alpha] \sqcup \beta' \sqcup \gamma'}$$

$$T_{2.16} = \mathbb{C} (\gamma') \langle \mathbf{A}[\ell''/\alpha]^{\ell_i[\ell''/\alpha]} \rangle_{(\beta' \sqcup \gamma')}$$

P3:

$$\frac{\frac{\frac{\Sigma; \Psi \vdash \ell \sqsubseteq \ell_e[\ell''/\alpha]}{\text{Given}} \quad \frac{\Sigma; \Psi \vdash (\beta' \sqcup \gamma') \sqsubseteq pc \sqsubseteq \ell_e[\ell''/\alpha]}{\text{Given}}}{\Sigma; \Psi \vdash ((\ell \sqcup (\beta' \sqcup \gamma')) \sqsubseteq \ell_e[\ell''/\alpha])}}{\Sigma; \Psi \vdash ((\ell \sqcup (\beta' \sqcup \gamma')) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e[\ell''/\alpha])}$$

P2:

$$\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\text{CG-var}}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[] : T_{2.6}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[][] : T_{2.7}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[][]\bullet : T_{2.81}} \text{CG-FE} \quad P3}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}}, a : T_{2.3}, b : T_{2.5} \vdash b[][]\bullet : T_{2.9}} \text{CG-CE}$$

P1:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\text{CG-unlabel}} \quad P2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.10}} \text{CG-bind}$$

P0:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash e_c : T_{2.2}}{\text{IH}} \quad P1}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_{2.10}} \text{CG-bind}$$

P0.1:

$$\frac{\frac{\Sigma; \Psi \vdash \mathbf{A}[\ell''/\alpha]^{l_i[\ell''/\alpha]} \searrow \ell}{\text{Given}}}{\Sigma; \Psi \vdash \ell \sqsubseteq l_i[\ell''/\alpha]} \text{By inversion}$$

P0.2:

$$\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.11}}{\text{Definition 2.86}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.12}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.13}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.14}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_{2.15}} \text{Lemma 2.87}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.11}} \text{P0}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.12}} \text{Definition 2.86}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.13}} \text{Definition 2.86}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet]) : T_{2.14}} \text{Definition 2.86}}}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_{2.15}} \text{Lemma 2.87}} \text{P0.1}$$

Main derivation:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_1}{\text{Definition 2.86}} \quad P0.2}{\Sigma; \Psi; (\Gamma)_{\vec{\beta}} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[\square\square\square\bullet])) : T_1}}$$

15. FC-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c))))} \text{FC-CI}$$

$$T_1 = \mathbb{C} \gamma' \gamma' ((c \xrightarrow{\ell_e} \tau)^\perp)_{\alpha'}$$

$$T_{1.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' ((c \xrightarrow{\ell_e} \tau))_{\alpha'}$$

$$T_{1.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \alpha' \forall \alpha_i, \gamma_i. (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{1.3} = \text{Labeled } \alpha' \forall \alpha_i, \gamma_i. (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{1.4} = \forall \alpha_i, \gamma_i. (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_{1.5} = (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

$$T_2 = \mathbb{C} \gamma_i \gamma_i (\tau)_{\alpha_i}$$

Main derivation:

$$\frac{\frac{\frac{\frac{\frac{\frac{\Sigma, \alpha_i, \gamma_i; \Psi, (c \wedge \alpha' \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e); (\Gamma) \vdash e_c : T_2}{\text{IH, Weakening}}}{\Sigma; \Psi; \Gamma \vdash \nu(e_c) : T_{1.5}} \text{CG-CI}}}{\Sigma; \Psi; \Gamma \vdash \Lambda\Lambda(\nu(e_c)) : T_{1.4}} \text{CG-FI}}}{\Sigma; \Psi; \Gamma \vdash \text{Lb}(\Lambda\Lambda(\nu(e_c))) : T_{1.3}} \text{CG-label}}}{\Sigma; \Psi; \Gamma \vdash \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) : T_{1.2}} \text{CG-ret}$$

16. FC-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : (c \xrightarrow{\ell_e} \tau)^\ell \rightsquigarrow e_c \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e \bullet : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[]\bullet)))} \text{FC-CE}$$

$$\beta' = \bigcup_{\beta_i \in \bar{\beta}'} \beta_i$$

$$T_1 = \mathbb{C} \gamma' \gamma' \langle \tau \rangle_{\alpha'}$$

$$T_2 = \mathbb{C} \gamma' \gamma' \langle (c \xrightarrow{\ell_e} \tau)^\ell \rangle_{(\beta' \sqcup \gamma')}$$

$$T_{2.1} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \langle (c \xrightarrow{\ell_e} \tau) \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.2} = \mathbb{C} \gamma' \gamma' \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup \alpha_i}$$

$$T_{2.3} = \text{Labeled } \ell \sqcup (\beta' \sqcup \gamma') \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup \alpha_i}$$

$$T_{2.4} = \mathbb{C} \gamma' (\gamma' \sqcup \ell \sqcup (\beta' \sqcup \gamma')) \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup \alpha_i}$$

$$T_{2.5} = \forall \alpha_i, \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq \alpha_i \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup \alpha_i}$$

$$T_{2.6} = \forall \gamma_i. (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup \gamma_i \sqsubseteq (\beta' \sqcup \gamma') \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma_i \gamma_i \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.7} = (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqcup (\beta' \sqcup \gamma')) \sqsubseteq (\beta' \sqcup \gamma' \sqcup \ell) \sqcap \ell_e \Rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.71} = (c \wedge (\beta' \sqcup \gamma' \sqcup \ell) \sqsubseteq (\beta' \sqcup \gamma') \sqcap \ell_e) \Rightarrow \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.8} = \mathbb{C} (\beta' \sqcup \gamma' \sqcup \ell) (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.9} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \langle \tau \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.10} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \langle \mathbf{A}^{\ell_i} \rangle_{\ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.11} = \mathbb{C} (\gamma') (\beta' \sqcup \gamma' \sqcup \ell) \text{Labeled } \ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma') \langle \mathbf{A} \rangle_{\ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.12} = \mathbb{C} (\gamma') (\gamma') \text{Labeled } \ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma') \langle \mathbf{A} \rangle_{\ell_i \sqcup \ell \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.13} = \mathbb{C} (\gamma') (\gamma') \text{Labeled } \ell_i \sqcup (\beta' \sqcup \gamma') \langle \mathbf{A} \rangle_{\ell_i \sqcup (\beta' \sqcup \gamma')}$$

$$T_{2.14} = \mathbb{C} (\gamma') (\gamma') \langle \mathbf{A}^{\ell_i} \rangle_{(\beta' \sqcup \gamma')}$$

P2:

$$\frac{\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b : T_{2.5}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b[] : T_{2.6}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b[][] : T_{2.71}} \text{CG-FE}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3}, b : T_{2.5} \vdash b[][]\bullet : T_{2.8}} \text{CG-CE}} \text{CG-var}$$

P1:

$$\frac{\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a.b.b[]\bullet) : T_{2.9}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{unlabel } a : T_{2.4}} \text{CG-unlabel} \quad P2}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'}, a : T_{2.3} \vdash \text{bind}(\text{unlabel } a.b.b[]\bullet) : T_{2.9}} \text{CG-bind}$$

P0:

$$\frac{\frac{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash e_c : T_{2.2}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[]\bullet)) : T_{2.9}} \text{CG-bind}}{\Sigma; \Psi; (\Gamma)_{\bar{\beta}'} \vdash e_c : T_{2.2}} \text{IH} \quad P1$$

Main derivation:

$$\begin{array}{c}
P0 \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])) : T_{2.10} \quad \text{CG-bind} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet])) : T_{2.11} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet]))) : T_{2.12} \quad \text{Lemma 2.87} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet]))) : T_{2.13} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet]))) : T_{2.14} \\
\hline
\Sigma; \Psi; (\Gamma)_{\beta'} \vdash \text{coerce_taint}(\text{bind}(e_c, a.\text{bind}(\text{unlabel } a.b.b[][\bullet]))) : T_1
\end{array}$$

□

Lemma 2.89 (FG \rightsquigarrow CG: Subtyping preservation). $\forall \Sigma, \Psi, \ell, \ell'. \Sigma; \Psi \vdash \ell \sqsubseteq \ell'$ and the following holds:

1. $\forall \tau, \tau'$.

$$\Sigma; \Psi \vdash \tau <: \tau' \implies \llbracket \tau \rrbracket_{\ell} <: \llbracket \tau' \rrbracket_{\ell'}$$

2. $\forall A, A'$.

$$\Sigma; \Psi \vdash A <: A' \implies \Sigma; \Psi \vdash \llbracket A \rrbracket_{\ell} <: \llbracket A' \rrbracket_{\ell'}$$

Proof. Proof by simultaneous induction on $\tau <: \tau$ and $A <: A$

Proof of statement (1)

Let $\tau = A_1^{\ell_1}$ and $\tau' = A_2^{\ell_2}$

P2:

$$\begin{array}{c}
\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \quad \text{Given} \\
\hline
\Sigma; \Psi \vdash A_1 <: A_2 \quad \text{By inversion} \quad P1 \\
\hline
\Sigma; \Psi \vdash (\llbracket A_1 \rrbracket_{\ell \sqcup \ell_1}) <: (\llbracket A_2 \rrbracket_{\ell' \sqcup \ell_2}) \quad \text{IH(2) on } A_1 <: A_2
\end{array}$$

P1:

$$\begin{array}{c}
\overline{A_1^{\ell_1} <: A_2^{\ell_2}} \quad \text{Given} \\
\hline
\Sigma; \Psi \vdash \ell_1 \sqsubseteq \ell_2 \quad \text{By inversion} \quad \overline{\Sigma; \Psi \vdash \ell \sqsubseteq \ell'} \quad \text{Given} \\
\hline
\Sigma; \Psi \vdash \ell \sqcup \ell_1 \sqsubseteq \ell' \sqcup \ell_2
\end{array}$$

Main derivation:

$$\begin{array}{c}
P1 \quad P2 \\
\hline
\Sigma; \Psi \vdash \text{Labeled } \ell \sqcup \ell_1 (\llbracket A_1 \rrbracket_{\ell \sqcup \ell_1}) <: \text{Labeled } \ell' \sqcup \ell_2 (\llbracket A_2 \rrbracket_{\ell' \sqcup \ell_2}) \quad \text{CGsub-labeled} \\
\hline
\Sigma; \Psi \vdash \llbracket A_1^{\ell_1} \rrbracket_{\ell} <: \llbracket A_2^{\ell_2} \rrbracket_{\ell'}
\end{array}$$

Proof of statement (2)

We proceed by cases on $A <: A$

1. FGsub-base:

$$\begin{array}{c}
\overline{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \quad \text{CG-refl} \\
\hline
\Sigma; \Psi \vdash \llbracket \mathbf{b} \rrbracket_{\ell} <: \llbracket \mathbf{b} \rrbracket_{\ell'} \quad \text{Definition 2.86}
\end{array}$$

2. FGsub-ref:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{ref } \ell_i \llbracket \mathbf{A} \rrbracket_{\ell_i} <: \text{ref } \ell_i \llbracket \mathbf{A} \rrbracket_{\ell_i}} \text{CG-ref}}{\Sigma; \Psi \vdash \llbracket \text{ref } \mathbf{A}^{\ell_i} \rrbracket_{\ell} <: \llbracket \text{ref } \mathbf{A}^{\ell_i} \rrbracket_{\ell'}} \text{Definition 2.86}$$

3. FGsub-prod:

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} \times \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell} \times \llbracket \tau'_2 \rrbracket_{\ell'}} \text{CGsub-prod}}{\Sigma; \Psi \vdash \llbracket \tau_1 \times \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 \times \tau'_2 \rrbracket_{\ell'}} \text{Definition 2.86}$$

4. FGsub-sum:

P1:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{Given}}{\Sigma; \Psi \vdash \tau_1 <: \tau'_1} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell'}} \text{IH(1) on } \tau_1 <: \tau'_1$$

P2:

$$\frac{\frac{\frac{}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{Given}}{\Sigma; \Psi \vdash \tau_2 <: \tau'_2} \text{By inversion}}{\Sigma; \Psi \vdash \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_2 \rrbracket_{\ell'}} \text{IH(1) on } \tau_2 <: \tau'_2$$

Main derivation:

$$\frac{\frac{P1 \quad P2}{\Sigma; \Psi \vdash \llbracket \tau_1 \rrbracket_{\ell} + \llbracket \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 \rrbracket_{\ell} + \llbracket \tau'_2 \rrbracket_{\ell'}} \text{CGsub-prod}}{\Sigma; \Psi \vdash \llbracket \tau_1 + \tau_2 \rrbracket_{\ell} <: \llbracket \tau'_1 + \tau'_2 \rrbracket_{\ell'}} \text{Definition 2.86}$$

5. FGsub-arrow:

$$T_1 = \forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta} \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha}$$

$$T_{1.0} = \forall \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta} \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha}$$

$$T_{1.1} = \forall \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_{\beta} \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha}$$

$$T_{1.2} = (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow (\tau_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{1.3} = (\tau_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$c_1 = (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e)$$

$$T_2 = \forall \alpha, \beta, \gamma. (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.0} = \forall \beta, \gamma. (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.1} = \forall \gamma. (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.2} = (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$T_{2.3} = (\tau'_1)_\beta \rightarrow \mathbb{C} \gamma \gamma (\tau'_2)_\alpha$$

$$c_2 = (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e)$$

P3:

$$\frac{\frac{\frac{}{\text{Given}}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \tau_2 <: \tau'_2} \text{By inversion, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \mathbb{C} \gamma \gamma (\tau_2)_\alpha <: \mathbb{C} \gamma \gamma (\tau'_2)_\alpha} \text{IH(1) with } \ell = \ell' = \alpha, \text{ CGsub-monad}$$

P2:

$$\frac{\frac{\frac{}{\text{Given}}}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_\xi} \tau_2 <: \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \tau'_1 <: \tau_1} \text{By inversion, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \llbracket \tau'_1 \rrbracket_\beta <: \llbracket \tau_1 \rrbracket_\beta} \text{IH(1) with } \ell = \ell' = \beta$$

P1:

$$\frac{P2 \quad P3}{\Sigma; \Psi \vdash T_{1.3} <: T_{2.3}} \text{CGsub-arrow}$$

P0.1:

$$\frac{\frac{\frac{\frac{}{\text{Given, Weakening}}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \ell \sqsubseteq \ell'}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \alpha) \Rightarrow (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha)}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash \ell'_e \sqsubseteq \ell_e} \text{Given, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash (\ell' \sqcup \beta \sqcup \gamma \sqsubseteq \ell'_e) \Rightarrow (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \ell_e)} \text{Given, Weakening}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash c_2 \Rightarrow c_1}$$

P0:

$$\frac{P0.1 \quad \frac{\frac{P1}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash T_{1.3} <: T_{2.3}}{\Sigma, \alpha, \beta, \gamma; \Psi \vdash T_{1.2} <: T_{2.2}} \text{CGsub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{CGsub-forall}}$$

Main derivation:

$$\frac{P0}{\Sigma; \Psi \vdash \llbracket \tau_1 \xrightarrow{\ell_\xi} \tau_2 \rrbracket_\ell <: \llbracket \tau'_1 \xrightarrow{\ell'_\xi} \tau'_2 \rrbracket_{\ell'}} \text{Definition 2.86}$$

6. FGsub-unit:

$$\frac{\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{CGsub-unit}}{\Sigma; \Psi \vdash \llbracket \text{unit} \rrbracket_{\ell} <: \llbracket \text{unit} \rrbracket_{\ell'}} \text{Definition 2.86}$$

7. FGsub-forall:

$$T_1 = \forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.0} = \forall \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.1} = \forall \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.2} = (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$T_{1.3} = \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'}$$

$$c_1 = (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e)$$

$$T_2 = \forall \alpha, \alpha', \gamma. (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.0} = \forall \alpha', \gamma. (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.1} = \forall \gamma. (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.2} = (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$T_{2.3} = \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}$$

$$c_2 = (\ell' \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell'_e)$$

P3:

$$\frac{\frac{\frac{}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash \tau_1 <: \tau_2} \text{Given, Weakening}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\tau_1)_{\alpha'} <: \tau_{2\alpha'}} \text{IH(1) with } \ell = \ell' = \alpha'}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash \mathbb{C} \gamma \gamma (\tau_1)_{\alpha'} <: \mathbb{C} \gamma \gamma (\tau_2)_{\alpha'}} \text{IH(1) with } \ell = \ell' = \alpha'$$

P2:

$$\frac{\frac{}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell'_e \sqsubseteq \ell_e)} \text{Given}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell' \sqcup \gamma \sqsubseteq \ell'_e) \Rightarrow (\ell \sqcup \gamma \sqsubseteq \ell_e)} \text{IH(1) with } \ell = \ell' = \alpha'$$

P1:

$$\frac{\frac{}{(\ell \sqsubseteq \ell')} \text{Given}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell' \sqcup \gamma \sqsubseteq \alpha') \Rightarrow (\ell \sqcup \gamma \sqsubseteq \alpha')} \text{IH(1) with } \ell = \ell' = \alpha'$$

P0:

$$\frac{\frac{}{P1} \quad \frac{}{P2}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash c_2 \Rightarrow c_1} \text{IH(1) with } \ell = \ell' = \alpha'$$

Main derivation:

$$\frac{\frac{\frac{}{P0} \quad \frac{}{P3}}{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash T_{1.2} <: T_{2.2}} \text{CGsub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{CGsub-forall}}{\Sigma; \Psi \vdash \llbracket \forall \alpha. \tau_1 \rrbracket_{\ell} <: \llbracket \forall \alpha. \tau_2 \rrbracket_{\ell'}} \text{Definition 2.86}$$

8. FGsub-constraint:

$$T_1 = \forall \alpha, \gamma. (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_\alpha$$

$$T_{1.0} = \forall \gamma. (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_\alpha$$

$$T_{1.1} = (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_1)_\alpha$$

$$T_{1.2} = \mathbb{C} \gamma \gamma (\tau_1)_\alpha$$

$$C_1 = (c_1 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e)$$

$$T_2 = \forall \alpha, \gamma. (c_2 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{2.0} = \forall \gamma. (c_2 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{2.1} = (c_2 \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$T_{2.2} = \mathbb{C} \gamma \gamma (\tau_2)_\alpha$$

$$C_2 = (c_2 \wedge \ell' \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e)$$

P1:

$$\frac{\frac{}{\Sigma, \alpha, \gamma; \Psi \vdash \tau_1 <: \tau_2} \text{Given, Weakening}}{\Sigma, \alpha, \gamma; \Psi \vdash (\tau_1)_\alpha <: \tau_{2\alpha}} \text{IH(1) with } \ell = \ell' = \alpha}{\Sigma, \alpha, \gamma; \Psi \vdash \mathbb{C} \gamma \gamma (\tau_1)_\alpha <: \mathbb{C} \gamma \gamma (\tau_2)_\alpha}$$

P0:

$$\frac{\frac{}{\Sigma; \Psi \vdash c_2 \Rightarrow c_1} \text{Given}}{\Sigma, \alpha, \gamma; \Psi \vdash c_2 \wedge (\ell' \sqcup \gamma \sqsubseteq \alpha \sqcap \ell'_e) \Rightarrow c_1 \wedge (\ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e)} \text{Weakening, } \ell \sqsubseteq \ell', \ell'_e \sqsubseteq \ell_e}{\Sigma, \alpha, \gamma; \Psi \vdash C_2 \Rightarrow C_1}$$

Main derivation:

$$\frac{\frac{\frac{P0 \quad P1}{\Sigma, \alpha, \gamma; \Psi \vdash T_{1.1} <: T_{2.1}} \text{CGsub-constraint}}{\Sigma; \Psi \vdash T_1 <: T_2} \text{CGsub-forall}}{\Sigma; \Psi \vdash \left[\left[c_1 \xrightarrow{\ell_e} \tau_1 \right]_\ell <: \left[\left[c_2 \xrightarrow{\ell'_e} \tau_2 \right]_{\ell'} \right]} \text{Definition 2.86}$$

□

Lemma 2.90 (FG \rightsquigarrow CG: Preservation of well-formedness). *For all Σ, Ψ and ℓ s.t. $FV(\ell) \in \Sigma$ the following hold:*

$$1. \forall \tau. \Sigma; \Psi \vdash \tau \text{ WF} \Rightarrow \Sigma; \Psi \vdash (\tau)_\ell \text{ WF}$$

$$2. \forall \mathbf{A}. \Sigma; \Psi \vdash \mathbf{A} \text{ WF} \Rightarrow \Sigma; \Psi \vdash (\mathbf{A})_\ell \text{ WF}$$

Proof. Proof by simultaneous induction on the WF relation of FG

Proof of statement (1)

Let $\tau = \mathbf{A}^{\ell'}$

$$\frac{\frac{\overline{\text{FV}(\ell') \in \Sigma} \text{ By inversion}}{\text{FV}(\ell' \sqcup \ell) \in \Sigma} \text{ IH(2) on A}}{\Sigma; \Psi \vdash \langle \mathbf{A} \rangle_{\ell' \sqcup \ell} WF} \text{ CG-wff-labeled}}{\Sigma; \Psi \vdash \text{Labeled } \ell' \sqcup \ell \langle \mathbf{A} \rangle_{\ell' \sqcup \ell} WF} \text{ CG-wff-labeled}$$

Proof of statement (2)

We proceed by case analyzing the last rule of given WF judgment.

1. FG-wff-base:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} WF} \text{ CG-wff-base}$$

2. FG-wff-unit:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} WF} \text{ CG-wff-unit}$$

3. FG-wff-arrow:

P1:

$$\frac{\frac{\overline{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \langle \tau_2 \rangle_\alpha WF} \text{ IH(1) on } \tau_2}}{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha WF} \text{ CG-wff-monad}}$$

P0:

$$\frac{\frac{\overline{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \langle \tau_1 \rangle_\beta WF} \text{ IH(1) on } \tau_1 \quad P1}}{\Sigma, \alpha, \beta, \gamma; \Psi, (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash (\langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha) WF} \text{ CG-wff-arrow}}$$

Main derivation:

$$\frac{\frac{\overline{\Sigma, \alpha, \beta, \gamma; \Psi \vdash ((\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha) WF} \text{ CG-wff-constraint}}{\Sigma; \Psi \vdash (\forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha) WF} \text{ CG-wff-forall}}{P0}$$

4. FG-wff-prod:

$$\frac{\frac{\overline{\Sigma; \Psi \vdash \langle \tau_1 \rangle_\ell WF} \text{ IH(1) on } \tau_1 \quad \overline{\Sigma; \Psi \vdash \langle \tau_2 \rangle_\ell WF} \text{ IH(1) on } \tau_2}}{\Sigma; \Psi \vdash \langle \tau_1 \rangle_\ell \times \langle \tau_2 \rangle_\ell WF} \text{ CG-wff-prod}}$$

5. FG-wff-sum:

$$\frac{\overline{\Sigma; \Psi \vdash \langle \tau_1 \rangle_\ell WF} \text{ IH(1) on } \tau_1 \quad \overline{\Sigma; \Psi \vdash \langle \tau_2 \rangle_\ell WF} \text{ IH(1) on } \tau_2}{\Sigma; \Psi \vdash \langle \tau_1 \rangle_\ell + \langle \tau_2 \rangle_\ell WF} \text{ CG-wff-prod}}$$

6. FG-wff-ref:

Let $\tau = A^{\ell'}$

$$\frac{\frac{\overline{FV(A) = \emptyset} \text{ By inversion} \quad \overline{FV(\ell') = \emptyset} \text{ By inversion}}{\overline{FV(\langle A \rangle_{\ell'}) = \emptyset}} \text{ Lemma 2.91}}{\Sigma; \Psi \vdash \text{ref } \ell' \langle A \rangle_{\ell'} WF} \text{ CG-wff-ref}$$

7. FG-wff-forall:

$$\frac{\frac{\frac{\overline{\Sigma, \alpha, \alpha', \gamma; \Psi, (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \vdash \langle \tau \rangle_{\alpha'} WF} \text{ IH(1) on } \tau}{\overline{\Sigma, \alpha, \alpha', \gamma; \Psi, (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \vdash \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha'} WF}} \text{ CG-wff-monad}}{\overline{\Sigma, \alpha, \alpha', \gamma; \Psi \vdash (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha'} WF}} \text{ CG-wff-constraint}}{\overline{\Sigma; \Psi \vdash (\forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha'}) WF}} \text{ CG-wff-forall}$$

8. FG-wff-constraint:

$$\frac{\frac{\frac{\overline{\Sigma, \alpha, \gamma; \Psi, (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \langle \tau \rangle_{\alpha} WF} \text{ IH(1) on } \tau}{\overline{\Sigma, \alpha, \gamma; \Psi, (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \vdash \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha} WF}} \text{ CG-wff-monad}}{\overline{\Sigma, \alpha, \gamma; \Psi \vdash (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha} WF}} \text{ CG-wff-constraint}}{\overline{\Sigma; \Psi \vdash (\forall \alpha, \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_{\alpha}) WF}} \text{ CG-wff-forall}$$

□

Lemma 2.91 (FG \rightsquigarrow CG: Free variable lemma). $\forall \Sigma, \ell. FV(\ell) \in \Sigma$, the following hold

1. $\forall \tau. FV(\langle \tau \rangle_{\ell}) \subseteq FV(\tau) \cup FV(\ell)$
2. $\forall A. FV(\langle A \rangle_{\ell}) \subseteq FV(A) \cup FV(\ell)$

Proof. Proof by simultaneous induction on τ and A

Proof for (1)

Let $\tau = A^{\ell_i}$

$$\begin{aligned} & FV(\langle A^{\ell_i} \rangle) \\ &= FV(\text{Labeled } \ell_i \sqcup \ell \langle A \rangle_{\ell_i \sqcup \ell}) && \text{Definition 2.86} \\ &= FV(\ell_i) \cup FV(\ell) \cup FV(\langle A \rangle_{\ell_i \sqcup \ell}) \\ &\subseteq FV(\ell_i) \cup FV(\ell) \cup FV(A) && \text{IH(2) on } A \\ &= FV(A^{\ell_i}) \cup FV(\ell) \end{aligned}$$

Proof for (2)

1. $A = \mathbf{b}$:

$$\begin{aligned} & FV(\langle \mathbf{b} \rangle_{\ell}) \\ &= FV(\mathbf{b}) && \text{Definition 2.86} \\ &\subseteq FV(\mathbf{b}) \cup FV(\ell) \end{aligned}$$

2. $A = \text{unit}$:

$$\begin{aligned} & FV(\langle \text{unit} \rangle_{\ell}) \\ &= FV(\text{unit}) && \text{Definition 2.86} \\ &\subseteq FV(\text{unit}) \cup FV(\ell) \end{aligned}$$

3. $A = \tau_1 \xrightarrow{\ell_e} \tau_2$:
- $$\begin{aligned}
& \text{FV}(\langle \tau_1 \xrightarrow{\ell_e} \tau_2 \rangle_\ell) \\
= & \text{FV}(\forall \alpha, \beta, \gamma. (\ell \sqcup \beta \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \langle \tau_1 \rangle_\beta \rightarrow \mathbb{C} \gamma \gamma \langle \tau_2 \rangle_\alpha) && \text{Definition 2.86} \\
= & \text{FV}(\ell) \cup \text{FV}(\langle \tau_1 \rangle_\beta) \cup \text{FV}(\ell_e) \cup \text{FV}(\langle \tau_2 \rangle_\alpha) \\
\subseteq & \text{FV}(\tau_1) \cup \text{FV}(\ell_e) \cup \text{FV}(\tau_2) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
= & \text{FV}(\tau_1 \xrightarrow{\ell_e} \tau_2) \cup \text{FV}(\ell)
\end{aligned}$$
4. $A = \tau_1 \times \tau_2$:
- $$\begin{aligned}
& \text{FV}(\langle \tau_1 \times \tau_2 \rangle_\ell) \\
= & \text{FV}(\langle \tau_1 \rangle_\ell \times \langle \tau_2 \rangle_\ell) && \text{Definition 2.86} \\
= & \text{FV}(\langle \tau_1 \rangle_\ell) \cup \text{FV}(\langle \tau_2 \rangle_\ell) \cup \text{FV}(\ell) \\
\subseteq & \text{FV}(\tau_1) \cup \text{FV}(\tau_2) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
= & \text{FV}(\tau_1 \times \tau_2) \cup \text{FV}(\ell)
\end{aligned}$$
5. $A = \tau_1 + \tau_2$:
- $$\begin{aligned}
& \text{FV}(\langle \tau_1 + \tau_2 \rangle_\ell) \\
= & \text{FV}(\langle \tau_1 \rangle_\ell + \langle \tau_2 \rangle_\ell) && \text{Definition 2.86} \\
= & \text{FV}(\langle \tau_1 \rangle_\ell) \cup \text{FV}(\langle \tau_2 \rangle_\ell) \cup \text{FV}(\ell) \\
\subseteq & \text{FV}(\tau_1) \cup \text{FV}(\tau_2) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_1 \text{ and } \tau_2 \\
= & \text{FV}(\tau_1 + \tau_2) \cup \text{FV}(\ell)
\end{aligned}$$
6. $A = \text{ref } \tau_i$:
- Let $\tau_i = A_i^{\ell_i}$
- $$\begin{aligned}
& \text{FV}(\langle \text{ref } \tau_i \rangle_\ell) \\
= & \text{FV}(\text{ref } \ell_i \langle A_i \rangle) && \text{Definition 2.86} \\
= & \text{FV}(\ell_i) \cup \text{FV}(\langle A_i \rangle) \\
\subseteq & \text{FV}(\ell_i) \cup \text{FV}(A_i) \cup \text{FV}(\ell) && \text{IH(2) on } A_i \\
= & \text{FV}(\text{ref } A_i^{\ell_i}) \cup \text{FV}(\ell) \\
= & \text{FV}(\text{ref } \tau_i) \cup \text{FV}(\ell)
\end{aligned}$$
7. $A = \forall \alpha. (\ell_e, \tau_i)$:
- $$\begin{aligned}
& \text{FV}(\langle \forall \alpha. (\ell_e, \tau_i) \rangle) \\
= & \text{FV}(\forall \alpha, \alpha', \gamma. (\ell \sqcup \gamma \sqsubseteq \alpha' \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau_i \rangle_{\alpha'}) && \text{Definition 2.86} \\
= & \text{FV}(\ell) \cup \text{FV}(\ell_e) \cup \text{FV}(\langle \tau_i \rangle) \\
\subseteq & \text{FV}(\ell) \cup \text{FV}(\ell_e) \cup \text{FV}(\tau_i) && \text{IH(1) on } \tau_i \\
= & \text{FV}(\ell) \cup \text{FV}(\forall \alpha. (\ell_e, \tau_i))
\end{aligned}$$
8. $A = c \xrightarrow{\ell_e} \tau_i$:
- $$\begin{aligned}
& \text{FV}(\langle c \xrightarrow{\ell_e} \tau_i \rangle) \\
= & \text{FV}(\forall \alpha, \gamma. (c \wedge \ell \sqcup \gamma \sqsubseteq \alpha \sqcap \ell_e) \Rightarrow \mathbb{C} \gamma \gamma \langle \tau \rangle_\alpha) && \text{Definition 2.86} \\
= & \text{FV}(\ell_e) \cup \text{FV}(c) \cup \text{FV}(\langle \tau_i \rangle) \cup \text{FV}(\ell) \\
\subseteq & \text{FV}(\ell_e) \cup \text{FV}(c) \cup \text{FV}(\tau_i) \cup \text{FV}(\ell) && \text{IH(1) on } \tau_i \\
= & \text{FV}(c \xrightarrow{\ell_e} \tau_i) \cup \text{FV}(\ell)
\end{aligned}$$

□

2.5.3 Logical relation for FG to CG translation

Definition 2.92 (FG \rightsquigarrow CG: ${}^s\theta_2$ extends ${}^s\theta_1$). ${}^s\theta_1 \sqsubseteq {}^s\theta_2 \triangleq$
 $\forall a \in {}^s\theta_1. {}^s\theta_1(a) = \tau \implies {}^s\theta_2(a) = \tau$

Definition 2.93 (FG \rightsquigarrow CG: $\hat{\beta}_2$ extends $\hat{\beta}_1$). $\hat{\beta}_1 \sqsubseteq \hat{\beta}_2 \triangleq$
 $\forall (a_1, a_2) \in \hat{\beta}_1. (a_1, a_2) \in \hat{\beta}_2$

Definition 2.94 (FG \rightsquigarrow CG: Unary value relation).

$$\begin{aligned}
\llbracket \mathbf{b} \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in \llbracket \mathbf{b} \rrbracket \wedge {}^tv \in \llbracket \mathbf{b} \rrbracket \wedge {}^sv = {}^tv\} \\
\llbracket \mathbf{unit} \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, {}^tv) \mid {}^sv \in \llbracket \mathbf{unit} \rrbracket \wedge {}^tv \in \llbracket \mathbf{unit} \rrbracket\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, ({}^sv_1, {}^sv_2), ({}^tv_1, {}^tv_2)) \mid \\
&\quad ({}^s\theta, m, {}^sv_1, {}^tv_1) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^sv_2, {}^tv_2) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 + \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \text{inl } {}^sv, \text{inl } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}}\} \cup \\
&\quad \{({}^s\theta, m, \text{inr } {}^sv, \text{inr } {}^tv) \mid ({}^s\theta, m, {}^sv, {}^tv) \in \llbracket \tau_2 \rrbracket_V^{\hat{\beta}}\} \\
\llbracket \tau_1 \xrightarrow{\ell_e} \tau_2 \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \lambda x. e_s, \Lambda \Lambda(\nu(\lambda x. e_t))) \mid \\
&\quad \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv, {}^tv, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^sv, {}^tv) \in \llbracket \tau_1 \rrbracket_V^{\hat{\beta}'} \implies \\
&\quad ({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in \llbracket \tau_2 \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \forall \alpha. (\ell_e, \tau) \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \Lambda e_s, \Lambda \Lambda(\nu(e_t))) \mid \\
&\quad \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau[\ell'/\alpha] \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket c \xrightarrow{\ell_e} \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, \nu e_s, \Lambda(\nu(e_t))) \mid \\
&\quad \mathcal{L} \models c \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau \rrbracket_E^{\hat{\beta}'}\} \\
\llbracket \text{ref } \tau \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, a_s, a_t) \mid {}^s\theta(a_s) = \tau \wedge ({}^sa, {}^ta) \in \hat{\beta}\} \\
\llbracket \mathbf{A}' \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, m, {}^sv, \text{Lb}({}^tv)) \mid ({}^s\theta, m, {}^sv, {}^tv) \in \llbracket \mathbf{A} \rrbracket_V^{\hat{\beta}}\}
\end{aligned}$$

Definition 2.95 (FG \rightsquigarrow CG: Unary expression relation).

$$\begin{aligned}
\llbracket \tau \rrbracket_E^{\hat{\beta}} &\triangleq \{({}^s\theta, n, e_s, e_t) \mid \\
&\quad \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^sv. (H_s, e_s) \Downarrow_i (H'_s, {}^sv) \implies \\
&\quad \exists H'_t, {}^tv. (H_t, e_t) \Downarrow^f (H'_t, {}^tv) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \\
&\quad \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in \llbracket \tau \rrbracket_V^{\hat{\beta}'}\}
\end{aligned}$$

Definition 2.96 (FG \rightsquigarrow CG: Unary heap well formedness).

$$\begin{aligned}
(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta &\triangleq \text{dom}({}^s\theta) \subseteq \text{dom}(H_s) \wedge \\
&\quad \hat{\beta} \subseteq (\text{dom}({}^s\theta) \times \text{dom}(H_t)) \wedge \\
&\quad \forall (a_1, a_2) \in \hat{\beta}. ({}^s\theta, n - 1, H_s(a_1), H_t(a_2)) \in \llbracket {}^s\theta(a_1) \rrbracket_V^{\hat{\beta}}
\end{aligned}$$

Definition 2.97 (FG \rightsquigarrow CG: Label substitution). $\sigma : \text{Lvar} \mapsto \text{Label}$

Definition 2.98 (FG \rightsquigarrow CG: Value substitution to values). $\delta^s : \text{Var} \mapsto \text{Val}, \delta^t : \text{Var} \mapsto \text{Val}$

Definition 2.99 (FG \rightsquigarrow CG: Unary interpretation of Γ).

$$\begin{aligned}
\llbracket \Gamma \rrbracket_V^{\hat{\beta}} &\triangleq \{({}^s\theta, n, \delta^s, \delta^t) \mid \text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \\
&\quad \forall x \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x), \delta^t(x)) \in \llbracket \Gamma(x) \rrbracket_V^{\hat{\beta}}\}
\end{aligned}$$

2.5.4 Soundness proof for FG to CG translation

Lemma 2.100 (FG \rightsquigarrow CG: Monotonicity). $\forall {}^s\theta, {}^s\theta', n, {}^s v, {}^t v, n', \beta, \beta'$.

1. $\forall A. ({}^s\theta, n, {}^s v, {}^t v) \in [A]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^s v, {}^t v) \in [A]_V^{\hat{\beta}'}$
2. $\forall \tau. ({}^s\theta, n, {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n \implies ({}^s\theta', n', {}^s v, {}^t v) \in [\tau]_V^{\hat{\beta}'}$

Proof. Proof by simultaneous induction on A and τ

Proof of statement (1)

We case analyze A in the last step

1. Case **b**:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [b]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [b]_V^{\hat{\beta}'}$$

Since $({}^s\theta, n, {}^s v, {}^t v) \in [b]_V^{\hat{\beta}}$ therefore from Definition 2.94 we know that ${}^s v \in \llbracket b \rrbracket \wedge {}^t v \in \llbracket b \rrbracket$ and ${}^s v = {}^t v$

Therefore from Definition 2.94 we get the desired

2. Case **unit**:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}'}$$

Since $({}^s\theta, n, {}^s v, {}^t v) \in [\text{unit}]_V^{\hat{\beta}}$ therefore from Definition 2.94 we know that ${}^s v \in \llbracket \text{unit} \rrbracket \wedge {}^t v \in \llbracket \text{unit} \rrbracket$

Therefore from Definition 2.94 we get the desired

3. Case $\tau_1 \times \tau_2$:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that ${}^s v = ({}^s v_1, {}^s v_2)$ and ${}^t v = ({}^t v_1, {}^t v_2)$.

We also know that $({}^s\theta, n, {}^s v_1, {}^t v_1) \in [\tau_1]_V^{\hat{\beta}}$ and $({}^s\theta, n, {}^s v_2, {}^t v_2) \in [\tau_2]_V^{\hat{\beta}}$

IH1: $({}^s\theta', n', {}^sv_1, {}^tv_1) \in [\tau_1]_V^{\hat{\beta}'}$ (From Statement (2))

IH2: $({}^s\theta', n', {}^sv_2, {}^tv_2) \in [\tau_2]_V^{\hat{\beta}'}$ (From Statement (2))

Therefore from Definition 2.94, IH1 and IH2 we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \times \tau_2]_V^{\hat{\beta}'}$$

4. Case $\tau_1 + \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.94 two cases arise

(a) ${}^sv = \text{inl}({}^sv')$ and ${}^tv = \text{inl}({}^tv')$:

IH: $({}^s\theta', n', {}^sv', {}^tv') \in [\tau_1]_V^{\hat{\beta}'}$ (From Statement (2))

Therefore from Definition 2.94 and IH we get

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 + \tau_2]_V^{\hat{\beta}'}$$

(b) ${}^sv = \text{inr}({}^sv')$ and ${}^tv = \text{inr}({}^tv')$:

Symmetric reasoning as in the previous case

5. Case $\tau_1 \xrightarrow{\ell_e} \tau_2$:

Given:

$$({}^s\theta, n, {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^sv, {}^tv) \in [\tau_1 \xrightarrow{\ell_e} \tau_2]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that

sv is of the form $\lambda x.e_s$ (for some e_s) and tv is of the form $\Lambda\Lambda\Lambda(\nu(\lambda x.e_t))$ (for some e_t) s.t

$$({}^s\theta', j, e_s[{}^sv/x], e_t[{}^tv/x]) \in [\tau_2]_E^{\hat{\beta}'_1} \quad (\text{A0})$$

Similarly from Definition 2.94 we are required to prove

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' . ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}''} \implies ({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

This means we are given some

$${}^s\theta'' \sqsupseteq {}^s\theta', {}^sv_2, {}^tv_2, k < n', \hat{\beta}' \sqsubseteq \hat{\beta}'' \text{ s.t } ({}^s\theta'', k, {}^sv_2, {}^tv_2) \in [\tau_1]_V^{\hat{\beta}''}$$

and we are required to prove

$$({}^s\theta'', k, e_s[{}^sv_2/x], e_t[{}^tv_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$$

Instantiating (A0) with ${}^s\theta'', {}^s v_2, {}^t v_2, k, \hat{\beta}''$ since
 ${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n$ and $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$ therefore we get
 $({}^s\theta'', k, e_s[{}^s v_2/x], e_t[{}^t v_2/x]) \in [\tau_2]_E^{\hat{\beta}''}$

6. Case $\forall\alpha.\tau$:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\forall\alpha.(\ell_e, \tau)]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\forall\alpha.(\ell_e, \tau)]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that ${}^s v = \Lambda e'_s$ and ${}^t v = \Lambda\Lambda(\nu(e_t))$ s.t

$$\forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta', j, e_s, e_t) \in [\tau[\ell'/\alpha]]_E^{\hat{\beta}'_1} \quad (\text{F0})$$

Similarly from Definition 2.94 we are required to prove

$$\forall {}^s\theta'' \sqsupseteq {}^s\theta', k < n', \ell'' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''_1. ({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''_1}$$

This means we are given ${}^s\theta''_1 \sqsupseteq {}^s\theta', k < n', \ell'' \in \mathcal{L}, \hat{\beta}' \sqsubseteq \hat{\beta}''$

and we are required to prove

$$({}^s\theta'', k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

Instantiating (F0) with ${}^s\theta''_1, k, \hat{\beta}''$ since ${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta, k < n' < n$ and $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$
therefore we get

$$({}^s\theta''_1, k, e_s, e_t) \in [\tau[\ell''/\alpha]]_E^{\hat{\beta}''}$$

7. Case $c \stackrel{\ell_e}{\Rightarrow} \tau$:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [c \stackrel{\ell_e}{\Rightarrow} \tau]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that ${}^s v = \nu(e'_s)$ and ${}^t v = \Lambda\Lambda(\nu(e_t))$. And

$$\mathcal{L} \models c \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta', j, e_s, e_t) \in [\tau]_E^{\hat{\beta}'_1} \quad (\text{C0})$$

Similarly from Definition 2.94 we are required to prove

$$\mathcal{L} \models c \implies \forall {}^s\theta'' \sqsupseteq {}^s\theta', k < n', \hat{\beta}' \sqsubseteq \hat{\beta}''_1. ({}^s\theta'', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''_1}$$

This means we are given $\mathcal{L} \models c, {}^s\theta'' \sqsupseteq {}^s\theta', k < n', \hat{\beta}' \sqsubseteq \hat{\beta}''$

and we are required to prove

$$({}^s\theta'', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$$

Since $\mathcal{L} \models c$ and instantiating (C0) with ${}^s\theta''_1, k, \hat{\beta}''$ since ${}^s\theta'' \sqsupseteq {}^s\theta' \sqsupseteq {}^s\theta$, $k < n' < n$ and $\hat{\beta} \sqsubseteq \hat{\beta}' \sqsubseteq \hat{\beta}''$ therefore we get

$$({}^s\theta', k, e_s, e_t) \in [\tau]_E^{\hat{\beta}''}$$

8. Case ref τ :

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\text{ref } \tau]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \tau]_V^{\hat{\beta}'}$$

From Definition 2.94 we know that ${}^s v = a_s$ and ${}^t v = a_t$. We also know that ${}^s\theta(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}$

From Definition 2.94, Definition 2.92 and Definition 2.93 we get

$$({}^s\theta', n', {}^s v, {}^t v) \in [\text{ref } \tau]_V^{\hat{\beta}'}$$

Proof of Statement (2)

Let $\tau = \mathbf{A}^{\ell''}$:

Given:

$$({}^s\theta, n, {}^s v, {}^t v) \in [\mathbf{A}^{\ell''}]_V^{\hat{\beta}} \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \wedge n' < n$$

From Definition 2.94 we know that

$$\exists {}^t v_i. {}^t v = \text{Lb}({}^t v_i) \text{ and } ({}^s\theta, n, {}^s v, {}^t v_i) \in [\mathbf{A}]_V^{\hat{\beta}}$$

To prove:

$$({}^s\theta', n', {}^s v, {}^t v) \in [\mathbf{A}^{\ell''}]_V^{\hat{\beta}'}$$

This means from Definition 2.94 we need to prove

$$({}^s\theta', n', {}^s v, {}^t v_i) \in [\mathbf{A}]_V^{\hat{\beta}'}$$

$$\text{IH: } ({}^s\theta', n', {}^s v, {}^t v_i) \in [\mathbf{A}]_V^{\hat{\beta}'} \quad (\text{From Statement (1)})$$

Therefore we get the desired directly from IH. □

Lemma 2.101 (FG \rightsquigarrow CG: Unary monotonicity for Γ). $\forall {}^s\theta, {}^s\theta', \delta, \Gamma, n, n', \hat{\beta}, \hat{\beta}'$.

$$({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}' \implies ({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$$

Proof. Given: $({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}} \wedge n' < n \wedge {}^s\theta \sqsubseteq {}^s\theta' \wedge \hat{\beta} \sqsubseteq \hat{\beta}'$

To prove: $({}^s\theta', n', \delta^s, \delta^t) \in [\Gamma]_V^{\hat{\beta}'}$

From Definition 2.99 it is given that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}}$$

And again from Definition 2.99 we are required to prove that

$$\text{dom}(\Gamma) \subseteq \text{dom}(\delta^s) \wedge \text{dom}(\Gamma) \subseteq \text{dom}(\delta^t) \wedge \forall x_i \in \text{dom}(\Gamma). ({}^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$$

- $dom(\Gamma) \subseteq dom(\delta^s) \wedge dom(\Gamma) \subseteq dom(\delta^t)$:

Given

- $\forall x_i \in dom(\Gamma). (^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$:

Since we know that $\forall x_i \in dom(\Gamma). (^s\theta, n, \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}}$ (given)

Therefore from Lemma 2.100 we get

$$\forall x_i \in dom(\Gamma). (^s\theta', n', \delta^s(x_i), \delta^t(x_i)) \in [\Gamma(x_i)]_V^{\hat{\beta}'}$$

□

Lemma 2.102 (FG \rightsquigarrow CG: Unary monotonicity for H). $\forall ^s\theta, H_s, H_t, n, n', \hat{\beta}$.

$$(n, H_s, H_t) \triangleright^{\hat{\beta}} ^s\theta \wedge n' < n \implies (n', H_s, H_t) \triangleright^{\hat{\beta}} ^s\theta$$

Proof. Given: $(n, H_s, H_t) \triangleright^{\hat{\beta}} ^s\theta \wedge n' < n$

To prove: $(n', H_s, H_t) \triangleright^{\hat{\beta}'} ^s\theta$

From Definition 2.96 it is given that

$$dom(^s\theta) \subseteq dom(H_S) \wedge \hat{\beta} \subseteq (dom(^s\theta) \times dom(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}}$$

And again from Definition 2.96 we are required to prove that

$$dom(^s\theta) \subseteq dom(H_S) \wedge \hat{\beta} \subseteq (dom(^s\theta) \times dom(H_t)) \wedge \forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}'}$$

- $dom(^s\theta) \subseteq dom(H_S)$:

Given

- $\hat{\beta} \subseteq (dom(^s\theta) \times dom(H_t))$:

Given

- $\forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}'}$:

Since we know that $\forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}}$ (given)

Therefore from Lemma 2.100 we get

$$\forall (a_1, a_2) \in \hat{\beta}. (^s\theta, n'-1, H_s(a_1), H_t(a_2)) \in [^s\theta(a)]_V^{\hat{\beta}'}$$

□

Lemma 2.103 (Coercion lemma). $\forall H, e, v$.

$$(H, e) \Downarrow_-^f (H', \text{Lb } v) \implies (H, \text{coerce_taint } e) \Downarrow_-^f (H', \text{Lb } v)$$

Proof. Given: $(H, e) \Downarrow_-^f (H', \text{Lb } v)$

To prove: $(H, \text{coerce_taint } e) \Downarrow_-^f (H', \text{Lb } v)$

From Definition of `coerce_taint` and cg-app it suffices to prove that

$$(H, \text{toLabeled}(\text{bind}(e, y.\text{unlabel}(y)))) \Downarrow_-^f (H', \text{Lb } v)$$

From cg-tolabeled it suffices to prove that
 $(H, \text{bind}(e, y.\text{unlabel}(y))) \Downarrow_{-}^f (H', v)$

From cg-bind it suffices to prove that

1. $(H, e) \Downarrow_{-}^f (H'_1, v_1)$:

We are given that $(H, e) \Downarrow_{-}^f (H', v)$ therefore we have $H'_1 = H'$ and $v'_1 = \text{Lb } v$

2. $(H'_1, \text{unlabel}(y)[v_1/y]) \Downarrow_{-}^f (H', v)$:

It suffices to prove that

$(H', \text{unlabel}(\text{Lb } v)) \Downarrow_{-}^f (H', v)$:

We get this directly from cg-unlabel

□

Theorem 2.104 (FG \rightsquigarrow CG: Fundamental theorem). $\forall \Sigma, \Psi, \Gamma, \tau, e_s, e_t, pc, \mathcal{L}, \delta^s, \delta^t, \sigma, {}^s\theta, n, \hat{\beta}$.

$\Sigma; \Psi; \Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \wedge$

$\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

\implies

$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

Proof. Proof by induction on the \rightsquigarrow relation

1. FC-var:

$$\frac{}{\Sigma; \Psi; \Gamma, x : \tau \vdash_{pc} x : \tau \rightsquigarrow \text{ret } x} \text{FC-var}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\}) \sigma]_{V}^{\hat{\beta}}$

To prove: $({}^s\theta, n, x \delta^s, \text{ret}(x) \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

From Definition 2.95 it suffices to prove that

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(x) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge \\ & ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau]_{V}^{\hat{\beta}'} \end{aligned}$$

This means given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, x \delta^s) \Downarrow_i (H'_s, {}^s v)$

From fg-val we know that $i = 0, {}^s v = x \delta^s$. Also from cg-ret we know that ${}^t v = x \delta^t$ and $H'_t = H_t$

And we are required to prove

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n, {}^s v, {}^t v) \in [\tau]_{V}^{\hat{\beta}'} \quad (\text{F-V0})$$

We choose ${}^s\theta'$ as ${}^s\theta$ and $\hat{\beta}'$ as $\hat{\beta}$

(a) $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$: Given

(b) $({}^s\theta, n, {}^sv, {}^tv) \in [\tau]_V^{\hat{\beta}}$:

Since we are given $({}^s\theta, n, \delta^s, \delta^t) \in [(\Gamma \cup \{x \mapsto \tau\}) \sigma]_V^{\hat{\beta}}$, therefore from Definition 2.99 we get $({}^s\theta, n, {}^sv, {}^tv) \in [\tau \sigma]_V^{\hat{\beta}}$

2. FC-lam:

$$\frac{\Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{\ell_e} e_s : \tau_2 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \lambda x. e_s : (\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t))))} \text{FC-lam}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, (\lambda x. e_s) \delta^s, \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_E^{\hat{\beta}}$

From Definition 2.95 it suffices to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^sv. (H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^sv) \implies \\ & \exists H'_t, {}^tv. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t) \Downarrow^f (H'_t, {}^tv) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n - i, {}^sv, {}^tv) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$ and given some $i < n, {}^sv$ s.t $(H_s, (\lambda x. e_s) \delta^s) \Downarrow_i (H'_s, {}^sv)$

From fg-val we know that ${}^sv = (\lambda x. e_s) \delta^s$, $H'_s = H_s$ and $i = 0$. Also from cg-ret, cg-label and cg-FI we know that $H'_t = H_t$ and ${}^tv = (\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t$

It suffices to prove that

$$\exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n, H_s, H_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n, {}^sv, {}^tv) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}'}$$

We choose ${}^s\theta'$ as ${}^s\theta$ and $\hat{\beta}'$ as $\hat{\beta}$

(a) $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$: Given

(b) $({}^s\theta, n, \lambda x. e_s \delta^s, (\text{Lb}(\Lambda\Lambda\Lambda(\nu(\lambda x. e_t)))) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2)^\perp \sigma]_V^{\hat{\beta}}$:

From Definition 2.94 it suffices to prove that

$$({}^s\theta, n, \lambda x. e_s \delta^s, (\Lambda\Lambda\Lambda(\nu(\lambda x. e_t))) \delta^t) \in [(\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma]_V^{\hat{\beta}}$$

Again from Definition 2.94 it suffices to prove that

$$\begin{aligned} & \forall {}^s\theta' \sqsupseteq {}^s\theta, {}^sv_d, {}^tv_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, {}^sv_d, {}^tv_d) \in [\tau_1 \sigma]_V^{\hat{\beta}'} \implies \\ & ({}^s\theta', j, e_s[{}^sv_d/x] \delta^s, e_t[{}^tv_d/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'} \end{aligned}$$

This further means that given ${}^s\theta' \sqsupseteq {}^s\theta, {}^sv_d, {}^tv_d, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$ s.t $({}^s\theta', j, {}^sv_d, {}^tv_d) \in [\tau_1 \sigma]_V^{\hat{\beta}'}$

And we are required to prove

$$({}^s\theta', j, e_s[{}^sv_d/x] \delta^s, e_t[{}^tv_d/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'} \quad (\text{F-L0})$$

Since we are given $({}^s\theta', j, {}^s v_d, {}^t v_d) \in [\tau_1 \sigma]_{V'}^{\hat{\beta}'}$, therefore from Definition 2.99 and Lemma 2.101 we have

$$({}^s\theta', j, \delta^s \cup \{x \mapsto {}^s v_d\}, \delta^t \cup \{x \mapsto {}^t v_d\}) \in [(\Gamma \cup \{x \mapsto \tau_1\}) \sigma]_{V'}^{\hat{\beta}'}$$

Therefore from IH we get

$$({}^s\theta', j, e_s \delta^s \cup \{x \mapsto {}^s v_d\}, e_t \delta^t \cup \{x \mapsto {}^t v_d\}) \in [\tau_2 \sigma]_E^{\hat{\beta}'}$$

We get (F-L0) directly from IH

3. FC-app:

$$\frac{\begin{array}{c} \Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} : (\tau_1 \xrightarrow{\ell_s} \tau_2)^\ell \rightsquigarrow e_{t1} \\ \Sigma; \Psi; \Gamma \vdash_{pc} e_{s2} : \tau_1 \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \ell \sqcup pc \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell \end{array}}{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} e_{s2} : \tau_2 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b))))} \text{FC-app}}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_{s1} e_{s2}) \delta^s, \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b)))) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 it suffices to prove

$$\begin{array}{l} \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_{V'}^{\hat{\beta}'} \end{array}$$

This further means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$ and given some $i < n, {}^s v$ s.t $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{array}{l} \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[\square\square\square]\bullet) b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau_2 \sigma]_{V'}^{\hat{\beta}'} \quad (\text{F-A0}) \end{array}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in [(\tau_1 \xrightarrow{\ell_s} \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{array}{l} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge \\ ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_s} \tau_2)^\ell \sigma]_{V'}^{\hat{\beta}'_1} \end{array}$$

We instantiate with H_s, H_t . And since we know that $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t $(H_{s1}, e_{s1}) \Downarrow_j (H'_{s1}, {}^s v_1)$.

This means we have

$$\exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_{\zeta}} \tau_2)^\ell \sigma]_{V}^{\hat{\beta}'_1} \quad (\text{F-A1.0})$$

Since we know that $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \xrightarrow{\ell_{\zeta}} \tau_2)^\ell \sigma]_{V}^{\hat{\beta}'_1}$ therefore from Definition 2.94 we know that $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$ s.t

$$({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \xrightarrow{\ell_{\zeta}} \tau_2) \sigma]_{V}^{\hat{\beta}'_1} \quad (\text{F-A1.1})$$

From Definition 2.94 we know that ${}^s v_1 = \lambda x. e'_s$ and ${}^t v_i = \Lambda \Lambda \Lambda (\nu (\lambda x. e'_t))$ s.t

$$\forall {}^s \theta''_1 \sqsupseteq {}^s \theta'_1, {}^s v', {}^t v', l < (n - j), \hat{\beta}'_1 \sqsubseteq \hat{\beta}''_1.$$

$$({}^s \theta''_1, l, {}^s v', {}^t v') \in [\tau_1 \sigma]_{V}^{\hat{\beta}''_1} \implies ({}^s \theta''_1, l, e'_s[{}^s v'/x], e'_t[{}^t v'/x]) \in [\tau_2 \sigma]_{E}^{\hat{\beta}''_1} \quad (\text{F-A1})$$

IH2:

$$({}^s \theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in [\tau_1 \sigma]_{E}^{\hat{\beta}'_1}$$

This means from Definition 2.95 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n - j, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2 \delta^t) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'_2} \end{aligned}$$

We instantiate with H'_{s1}, H'_{t1} . And since we know that $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists k < i - j < n - j$ s.t $(H'_{s1}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$.

This means we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'_2} \quad (\text{F-A2})$$

We instantiate (F-A1) with θ''_1 as θ'_2 , ${}^s v'$ as ${}^s v_2$, ${}^t v'$ as ${}^t v_2$, l as $n - j - k$ and $\hat{\beta}''_1$ as $\hat{\beta}'_2$. Therefore we get

$$({}^s \theta'_2, n - j - k, e'_s[{}^s v_2/x], e'_t[{}^t v_2/x]) \in [\tau_2 \sigma]_{E}^{\hat{\beta}'_2}$$

From Definition 2.95 we have

$$\begin{aligned} \forall H_s, H_t. (n - j - k, H_s, H_t) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge \forall a < n - j - k, {}^s v. (H_s, e'_s[{}^s v_2/x]) \Downarrow_i (H'_{s3}, {}^s v_3) \implies \\ \exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s \theta'_3 \sqsupseteq {}^s \theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2. \\ (n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'_3} {}^s \theta'_3 \wedge ({}^s \theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2 \sigma]_{V}^{\hat{\beta}'_3} \end{aligned}$$

Instantiating with H'_{s2}, H'_{t2} . since we know that $(H_s, (e_{s1} e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists a < i - j - k < n - j - k$ s.t $(H'_{s2}, e'_s[{}^s v/x] \delta^s) \Downarrow_a (H'_{s3}, {}^s v_3)$

Therefore we have

$$\exists H'_{t3}, {}^t v_3. (H_t, e'_t[{}^t v_2/x]) \Downarrow^f (H'_{t3}, {}^t v_3) \wedge \exists {}^s \theta'_3 \sqsupseteq {}^s \theta'_2, \hat{\beta}'_3 \sqsupseteq \hat{\beta}'_2.$$

$$(n - j - k - a, H'_{s3}, H'_{t3}) \triangleright^{\hat{\beta}'_3} {}^s \theta'_3 \wedge ({}^s \theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in [\tau_2 \sigma]_{V}^{\hat{\beta}'_3} \quad (\text{F-A3})$$

Let $\tau_2 \sigma = A_2^{\ell_i}$, since $\tau_2 \sigma \searrow \ell \sigma$ therefore $\ell \sigma \sqsubseteq \ell_i$ and

$$({}^s\theta'_3, n - j - k - a, {}^s v_3, {}^t v_3) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3}$$

Therefore from Definition 2.94 we know that

$$({}^s\theta'_3, n - j - k - a, {}^s v_3, \text{Lb}^t v_{3i}) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3} \quad (\text{F-A3.1})$$

In order to prove (F-A0) we choose H'_t as H'_{t3} and ${}^t v$ as $\text{Lb}({}^t v_{3i})$. We need to prove:

$$(a) (H_t, \text{coerce_taint}(\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i}):$$

From Lemma 2.103 it suffices to prove that

$$(H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)))) \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i})$$

From cg-bind it further suffices to show that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1):$

We get this directly from (F-A1.0)

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i}):$

From cg-bind it suffices to prove that

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2):$

We get this directly from (F-A2)

- $(H'_{t2}, \text{bind}(\text{unlabel } a, c.(c[] [] \bullet) b) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i}):$

From cg-bind again it suffices to prove

- * $(H'_{t2}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t31}, {}^t v_{t2}):$

Since from (F-A1.1) we know that $\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel and (F-A1) we know that $H'_{t31} = H'_{t2}$ and ${}^t v_{t2} = {}^t v_i = \Lambda\Lambda\Lambda(\nu(\lambda x.e'_t))$

- * $((c[] [] \bullet) b) [{}^t v_2/b] [{}^t v_{t2}/c] \delta^t) \Downarrow {}^t v_{t21}:$

It suffices to prove that

$$(((\Lambda\Lambda\Lambda(\nu(\lambda x.e'_t))) [] [] \bullet) {}^t v_2) \delta^t) \Downarrow {}^t v_{t21}$$

From cg-FE it suffices to prove that

$$(((\Lambda\Lambda(\nu(\lambda x.e'_t))) [] \bullet) {}^t v_2) \delta^t) \Downarrow {}^t v_{t21}$$

Again from cg-FE applied two times it suffices to prove that

$$((\nu(\lambda x.e'_t) \bullet) {}^t v_2) \delta^t) \Downarrow {}^t v_{t21}$$

From cg-CE it suffices to prove that

$$(((\lambda x.e'_t) {}^t v_2) \delta^t) \Downarrow {}^t v_{t21}$$

From cg-app we know that

$${}^t v_{t21} = e'_t [{}^t v_2/x] \delta^t$$

- * $(H'_{t2}, {}^t v_{t21}) \Downarrow^f (H'_{t3}, \text{Lb}^t v_{3i}):$

We get this from (F-A3) and (F-A3.1)

$$(b) \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \lfloor \tau_2 \sigma \rfloor_V^{\hat{\beta}'_3}:$$

We choose ${}^s\theta'$ as ${}^s\theta'_3$ and $\hat{\beta}'$ as $\hat{\beta}'_3$. From fg-app we know that $i = j + k + a + 1$, ${}^s v = {}^s v_3$ and $H'_s = H'_{s3}$. Also from the termination proof (previous point) we know that $H'_t = H'_{t3}$ and ${}^t v = \text{Lb}({}^t v_3)$

We get $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_3} {}^s\theta'$ from (F-A3) and Lemma 2.102

Since ${}^t v = \text{Lb}({}^t v_3)$ therefore from Definition 2.94 it suffices to prove that

$$({}^s \theta'_3, n - j - k - a - 1, {}^s v_3, {}^t v_3) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}'_3}$$

We get this directly from (F-A3) and Lemma 2.100

4. FC-prod:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} : \tau_1 \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{s2} : \tau_2 \rightsquigarrow e_{t2}}{\Sigma; \Psi; \Gamma \vdash_{pc} (e_{s1}, e_{s2}) : (\tau_1 \times \tau_2)^\perp \rightsquigarrow \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b))))} \text{prod}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in \llbracket \Gamma \sigma \rrbracket_V^{\hat{\beta}}$

To prove: $({}^s \theta, n, (e_{s1}, e_{s2}) \delta^s, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \in \llbracket (\tau_1 \times \tau_2)^\perp \sigma \rrbracket_E^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v_1, {}^s v_2. (H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2)) \implies \\ & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket (\tau_1 \times \tau_2)^\perp \sigma \rrbracket_V^{\hat{\beta}'} \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$. Also given some $i < n, {}^s v_1, {}^s v_2$ s.t $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in \llbracket (\tau_1 \times \tau_2)^\perp \sigma \rrbracket_V^{\hat{\beta}'} \quad (\text{F-P0}) \end{aligned}$$

IH1:

$$({}^s \theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$ therefore $\exists j < i < n$ s.t $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-P1}) \end{aligned}$$

IH2:

$$({}^s \theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_1}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}} s\theta'_1 \wedge \forall k < n - j, {}^s v_1. (H_{s2}, e_{s2} \delta^s) \Downarrow_j (H'_{s2}, {}^s v_1) \implies \\ & \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_1) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} s\theta'_2 \wedge ({}^s \theta'_1, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}'_2} \end{aligned}$$

Instantiating with H'_{s1}, H'_{t1} and since we know that $(H_s, (e_{s1}, e_{s2})) \Downarrow_i (H'_s, ({}^s v_1, {}^s v_2))$ therefore $\exists k < i - j < n - j$ s.t $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_1. (H_{t2}, e_{t2}) \Downarrow^f (H'_{t2}, {}^t v_1) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \\ & (n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} s\theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau_2 \sigma]_V^{\hat{\beta}'_2} \quad (\text{F-P2}) \end{aligned}$$

In order to prove (F-P0) we choose H_t as H'_{t2} and ${}^t v$ as $\text{Lb}({}^t v_1, {}^t v_2)$

$$(a) (H_t, (\text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))))) \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2)):$$

From cg-bind it suffices to prove that

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{tb1}, {}^t v_{tb1})$:
From (F-P1) we know that $H'_{tb1} = H'_{t1}$ and ${}^t v_{tb1} = {}^t v_1$
- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{ret}(\text{Lb}(a, b)))) [{}^t v_1/a] \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$:
From cg-bind it suffices to prove that
 - $(H_t, e_{t2} \delta^t) \Downarrow^f (H'_{tb2}, {}^t v_{tb2})$:
From (F-P2) we know that $H'_{tb2} = H'_{t2}$ and ${}^t v_{tb2} = {}^t v_2$
 - $(H'_{t2}, \text{ret}(\text{Lb}(a, b))) [{}^t v_1/a] [{}^t v_2/b] \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_1, {}^t v_2))$:
From cg-ret, (F-P1) and (F-P2)

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} s\theta' \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^{\hat{\beta}'}$$

We choose ${}^s \theta'$ as ${}^s \theta'_2$ and $\hat{\beta}'$ as $\hat{\beta}'_2$ and since from fg-prod $i = j + k + 1$ and $H'_s = H'_{s2}$. Therefore from (F-P2) and Lemma 2.102 we get

$$(n - i, H'_s, H'_{t2}) \triangleright^{\hat{\beta}'_2} s\theta'_2$$

In order to prove $({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v) \in [(\tau_1 \times \tau_2)^\perp \sigma]_V^{\hat{\beta}'_2}$

From Definition 2.94 it suffices to prove

$$\exists {}^t v_i. {}^t v = \text{Lb}({}^t v_i) \wedge ({}^s \theta', n - i, ({}^s v_1, {}^s v_2), {}^t v_i) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}'_2}$$

Since ${}^t v = \text{Lb}({}^t v_1, {}^t v_2)$ therefore we get the desired from (F-P1), (F-P2), Definition 2.94 and Lemma 2.100

5. FC-fst:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\tau_1 \times \tau_2)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \tau_1 \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{fst}(e_s) : \tau_1 \rightsquigarrow \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b))))))} \text{fst}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, \text{fst}(e_s) \delta^s, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \delta^t \in [\tau_1 \sigma]_E^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$

We need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{F-F0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 \times \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, \text{fst}(e_s)) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1}) \end{aligned}$$

Since we know that $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 \times \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1}$ therefore from Definition 2.94 we know that ${}^t v_1 = \text{Lb}({}^t v_i)$ s.t

$$({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 \times \tau_2) \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1.1})$$

From Definition 2.94 we know that ${}^s v_1 = ({}^s v_{i1}, {}^s v_{i2})$ and ${}^t v_i = ({}^t v_{i1}, {}^t v_{i2})$ s.t

$$({}^s\theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1.2})$$

Let $\tau_1 \sigma = A_1^{\ell_i}$, since $\tau_1 \sigma \searrow \ell \sigma$ therefore $\ell \sigma \sqsubseteq \ell_i$ and

$$\text{Since } ({}^s\theta'_1, n - j, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1}$$

Therefore from Definition 2.94 we know that

$$({}^s\theta'_1, n - j, {}^s v_{i1}, \text{Lb}({}^t v_{i1})) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-F1.3})$$

In order to prove (F-F0) we choose H'_t as H'_{t1} and ${}^t v$ as $\text{Lb}({}^t v_{i1})$ as we need to prove

(a) $(H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_{t1}, \text{Lb}^t v_{i11})$:

From Lemma 2.103 it suffices to prove that

$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))))) \Downarrow^f (H'_{t1}, \text{Lb}^t v_{i11})$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1)$:

We get this from (F-F1)

- $(H'_{t1}, \text{bind}(\text{unlabel}(a), b.\text{ret}(\text{fst}(b)))[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}^t v_{i11})$:

Again from cg-bind it suffices to prove that

- $(H'_{t1}, \text{unlabel}(a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$:

Since ${}^t v_1 = \text{Lb}({}^t v_{i1}, {}^t v_{i2})$ from (F-F1.1) and (F-F1.2) therefore we get the desired from cg-unlabel

So, $H_{t21} = H'_{t1}$ and ${}^t v_{t21} = ({}^t v_{i1}, {}^t v_{i2})$

- $(H'_{t1}, \text{ret}(\text{fst}(b))[({}^t v_{i1}, {}^t v_{i2})/b] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}^t v_{i11})$:

We get this from cg-fst, cg-ret and (F-F1.2) and (F-F1.3)

(b) $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau_1 \sigma]_{\mathcal{V}}^{\hat{\beta}'}$:

We choose ${}^s \theta'$ as ${}^s \theta'_1$ and $\hat{\beta}'$ as $\hat{\beta}'_1$. And from fg-fst we know that $i = j + 1$ and $H'_s = H'_{s1}$ therefore from (F-F1) and Lemma 2.102 we get

$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$

Since from fg-fst we know that ${}^s v = {}^s v_{i1}$ therefore from (F-F1.2) and Lemma 2.100 we get

$({}^s \theta', n - i, {}^s v_{i1}, {}^t v_{i1}) \in [\tau_1 \sigma]_{\mathcal{V}}^{\hat{\beta}'_1}$

6. FC-snd:

Symmetric reasoning as in the FC-fst case

7. FC-inl:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e : \tau_1 \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{inl}(e_s) : (\tau_1 + \tau_2)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a)))} \text{inl}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\mathcal{V}}^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{inl}(e_s) \delta^s, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \in [(\tau_1 + \tau_2)^\perp \sigma]_{\mathcal{E}}^{\hat{\beta}}$

This means from Definition 2.95 we have

$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v) \implies$
 $\exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$

$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\mathcal{V}}^{\hat{\beta}'}$

This means that we are given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$. Also given some $i < n, {}^s v$ s.t $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-IL0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau_1 \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, \text{inl}(e_s)) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_t, e_{t1}) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-IL1}) \end{aligned}$$

In order to prove (F-IL0) we choose H'_t as H'_{t1} and ${}^t v$ as $(\text{Lb inl}({}^t v_1))$ and we need to prove:

$$(a) (H'_{t1}, \text{bind}(e_t, a.\text{ret}(\text{Lbinl}(a))) \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1))):$$

From cg-bind it suffices to prove that

$$i. (H'_{t1}, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11}):$$

From (F-IL1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t11} = {}^t v_1$

$$ii. (H'_{t1}, \text{ret}(\text{Lbinl}(a))[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, (\text{Lb inl}({}^t v_1))):$$

We get this from cg-ret, (F-IL1)

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{V}^{\hat{\beta}'}$$

We choose ${}^s \theta'$ as ${}^s \theta'_1$ and $\hat{\beta}'$ as $\hat{\beta}'_1$. Since from fg-inl we know that $i = j + 1$ and $H'_s = H'_{s1}$ therefore from (F-IL1) and Lemma 2.102 we get

$$(n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$$

Now we need to prove $({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2)^\perp \sigma]_{V}^{\hat{\beta}'}$

Since ${}^s v = \text{inl } {}^s v_1$ and ${}^t v = \text{Lb}(\text{inl}({}^t v_1))$ therefore from Definition 2.94 it suffices to prove that

$$({}^s \theta', n - i, \text{inl } {}^s v_1, \text{inl } {}^t v_1) \in [(\tau_1 + \tau_2) \sigma]_{V}^{\hat{\beta}'}$$

Since from (F-IL1) we know that $({}^s \theta', n - j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_{V}^{\hat{\beta}'}$

Therefore from Lemma 2.100 and Definition 2.94 we get

$$({}^s \theta', n - i, {}^s v, {}^t v) \in [(\tau_1 + \tau_2) \sigma]_{V}^{\hat{\beta}'}$$

8. FC-inr:

Symmetric reasoning as in the FC-inl case

9. FC-case:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\tau_1 + \tau_2)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s1} : \tau \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma, x : \tau_1 \vdash_{pc \sqcup \ell} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{case}(e_s, x.e_{s1}, y.e_{s2}) : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2}))))} \text{case}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_V^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \in [\tau \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \end{aligned}$$

This means we are given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \\ & \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{F-C0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\tau_1 + \tau_2)^\ell \sigma]_E^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1} \quad (\text{F-C1}) \end{aligned}$$

Since from (F-C1) we have $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\tau_1 + \tau_2)^\ell \sigma]_V^{\hat{\beta}'_1}$ therefore from Definition 2.94 we know that

$$\exists^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\tau_1 + \tau_2) \sigma]_{V'}^{\hat{\beta}'_1} \quad (\text{F-C1.1})$$

2 cases arise

(a) ${}^s v_1 = \text{inl}({}^s v_{i1})$ and ${}^t v_i = \text{inl}({}^t v_{i1})$:

Also from Lemma 2.101 and Definition 2.99 we know that

$$({}^s \theta'_1, n - j, \delta^s \cup \{x \mapsto {}^s v_1\}, \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [(\Gamma, \{x \mapsto {}^s v_1\}) \sigma]_{V'}^{\hat{\beta}'_1}$$

IH2:

$$({}^s \theta'_1, n - j, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \in [\tau \sigma]_E^{\hat{\beta}'_1}$$

This means from Definition 2.95 we have

$$\begin{aligned} \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s1} \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_j (H'_{s2}, {}^s v_2) \implies \\ \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \end{aligned}$$

$$(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2}$$

Instantiating with H'_{s1}, H'_{t1} and since we know that $(H_s, \text{case}(e_s, x.e_{s1}, y.e_{s2}) \delta^s \cup \{x \mapsto {}^s v_1\}) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists k < i - j < n - j$ s.t. $(H'_{s1}, e_{s1}) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t1} \delta^t \cup \{x \mapsto {}^t v_{i1}\}) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1.$$

$$(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2} \quad (\text{F-C2})$$

Let $\tau \sigma = \mathbf{A}_2^{\ell_i}$, since $\tau \sigma \searrow \ell \sigma$ therefore $\ell \sigma \sqsubseteq \ell_i$ and

$$({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2}$$

Therefore from Definition 2.94 we know that

$$({}^s \theta'_2, n - j - k, {}^s v_2, \text{Lb}({}^t v_{2i})) \in [\tau \sigma]_{V'}^{\hat{\beta}'_2} \quad (\text{F-C2.1})$$

In order to prove (F-C0) we choose H'_t as H'_{t2} and ${}^t v$ as $\text{Lb}({}^t v_{2i})$

And we need to prove:

i. $(H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$:

From Lemma 2.103 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})))) \delta^t) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$:

From (F-C1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t11} = {}^t v_1$

- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.\text{case}(b, x.e_{t1}, y.e_{t2})) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t1}, \text{Lb}({}^t v_{2i}))$:

From cg-bind it suffices to prove that

- $(H'_{t1}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21})$:

Since from (F-C1.1) we know that ${}^t v_1 = \text{Lb}({}^t v_i)$ therefore from cg-unlabel we know that

$$H'_{t21} = H'_{t1} \text{ and } {}^t v_{t21} = {}^t v_i$$

- $(\text{case}(b, x.e_{t1}, y.e_{t2}) [{}^t v_i/b] \delta^t) \Downarrow {}^t v_{t22}$:

Since we know that in this case ${}^t v_i = \text{inl}({}^t v_{i1})$

Therefore from cg-case we know that ${}^t v_{t22} = e_{t1} [{}^t v_{i1}/x] \delta^t$

– $(H'_{t1}, e_{t1}[{}^t v_{i1}/x] \delta^t) \Downarrow (H'_{t2}, \mathbf{Lb}^t v_{2i})$:
 We get this from (F-C2) and (F-C2.1)

ii. $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'}$:
 We choose ${}^s \theta'$ as ${}^s \theta'_2$ and $\hat{\beta}'$ as $\hat{\beta}'_2$. Since from fg-case we know that $i = j + k + 1$ and $H'_s = H'_{s2}$ therefore from (F-C2) and Lemma 2.102 we get

$$(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s \theta'_2$$

Now we need to prove $({}^s \theta'_2, n - i, {}^s v, {}^t v) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$

Since ${}^s v = {}^s v_2$ and ${}^t v = {}^t v_2$ and since from (F-C2) we know that

$$({}^s \theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$$

Therefore from Lemma 2.100 and Definition 2.94 we get

$$({}^s \theta'_2, n - i, {}^s v_2, {}^t v_2) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_2}$$

(b) ${}^s v_1 = \text{inr}({}^s v_{i1})$ and ${}^t v_1 = \text{inr}({}^t v_{i1})$:

Symmetric reasoning as in the previous case

10. FC-FI:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{\ell_e} e_s : \tau \rightsquigarrow e_t}{\Sigma; \Psi; \Gamma \vdash_{pc} \Lambda e_s : (\forall \alpha_g. (\ell_e, \tau))^\perp \rightsquigarrow \text{ret}(\mathbf{Lb}(\Lambda \Lambda \Lambda(\nu(e_t))))} \text{FI}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in \lfloor \Gamma \sigma \rfloor_V^{\hat{\beta}}$

To prove: $({}^s \theta, n, \Lambda e_s \delta^s, \text{ret}(\mathbf{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t) \in \lfloor \tau \sigma \rfloor_E^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \Lambda e_s \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(\mathbf{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\forall \alpha_g. (\ell_e, \tau))^\perp \sigma \rfloor_V^{\hat{\beta}'} \end{aligned}$$

This means given some H_s, H_t s.t. $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$. Also given some $i < n, {}^s v$ s.t. $(H_s, \Lambda e_s \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{ret}(\mathbf{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\forall \alpha_g. (\ell_e, \tau))^\perp \sigma \rfloor_V^{\hat{\beta}'} \end{aligned}$$

From fg-val we know that ${}^s v = (\Lambda e_s) \delta^s$, $H'_s = H_s$ and $i = 0$. Also from cg-ret we know that $H'_t = H_t$ and ${}^t v = (\mathbf{Lb}(\Lambda \Lambda \Lambda(\nu(e_t)))) \delta^t$

It suffices to prove that

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\forall \alpha_g. (\ell_e, \tau))^\perp \sigma \rfloor_V^{\hat{\beta}'}$$

We choose ${}^s \theta'$ as ${}^s \theta$ and $\hat{\beta}'$ as $\hat{\beta}$

(a) $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta$: Given

(b) $({}^s\theta, n, \Lambda e_s \delta^s, (\text{Lb}(\Lambda\Lambda\Lambda(\nu(e_t)))) \delta^t) \in [(\forall\alpha_g.(\ell_e, \tau))^\perp \sigma]_{\hat{\beta}}^{\hat{\beta}}$:

From Definition 2.94 it suffices to prove that

$$({}^s\theta, n, \Lambda e_s \delta^s, (\Lambda\Lambda\Lambda(\nu(e_t))) \delta^t) \in [(\forall\alpha_g.(\ell_e, \tau)) \sigma]_{\hat{\beta}}^{\hat{\beta}}$$

Again from Definition 2.94 it suffices to prove that

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha_g] \sigma]_{\hat{E}}^{\hat{\beta}'_1}$$

This further means that given ${}^s\theta'_1 \sqsupseteq {}^s\theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1$

And we need to prove

$$({}^s\theta'_1, j, e_s \delta^s, e_t \delta^t) \in [\tau[\ell'/\alpha_g]]_{\hat{E}}^{\hat{\beta}'_1} \quad (\text{F-FI0})$$

$$\underline{\text{IH}}: ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in [\tau \sigma \cup \{\alpha_g \mapsto \ell'\}]_{\hat{E}}^{\hat{\beta}'_1}$$

We get (F-FI0) directly from IH

11. FC-FE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\forall\alpha_g.(\ell_e, \tau))^\ell \rightsquigarrow e_t \quad \text{FV}(\ell') \subseteq \Sigma \quad \Sigma; \Psi \vdash_{pc} \sqcup \ell \sqsubseteq \ell_e[\ell'/\alpha] \quad \Sigma; \Psi \vdash \tau[\ell'/\alpha] \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_s [] : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[] [] [] \bullet)))} \text{FE}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove: $({}^s\theta, n, e_s [] \delta^s, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[] [] [] \bullet))) \delta^t) \in [\tau \sigma]_{\hat{E}}^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, e_s []) \Downarrow_i (H'_s, {}^s v) \implies \\ \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[] [] [] \bullet)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'}$$

This means given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, e_s []) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[] [] [] \bullet)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'} \quad (\text{F-FE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\forall\alpha_g.(\ell_e, \tau))^\ell \sigma]_{\hat{E}}^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}.$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall\alpha_g.(\ell_e, \tau))^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1}$$

Instantiating with H_s, H_t and since we know that $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t. $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha_g. (\ell_e, \tau))^\ell \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-FE1}) \end{aligned}$$

Since from (F-FE1) we have $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\forall \alpha_g. (\ell_e, \tau))^\ell \sigma]_{V}^{\hat{\beta}'}$ therefore from Definition 2.94 we know that

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\forall \alpha_g. (\ell_e, \tau)) \sigma]_{V}^{\hat{\beta}'} \quad (\text{F-FE1.1})$$

Therefore from Definition 2.94 we have

$$\begin{aligned} & {}^s v_1 = \Lambda e'_s \text{ and } {}^t v_i = \Lambda \Lambda \Lambda \nu e'_t \\ & \forall {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \ell'' \in \mathcal{L}, k < n - j, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_2, k, e'_s, e'_t) \in [\tau[\ell''/\alpha_g] \sigma]_{E}^{\hat{\beta}'_1} \quad (\text{F-FE1.2}) \end{aligned}$$

We instantiate with ${}^s \theta'_1, \ell', n - j - 1, \hat{\beta}'$ we get $({}^s \theta'_1, n - j - 1, e'_s, e'_t) \in [\tau[\ell'/\alpha_g] \sigma]_{E}^{\hat{\beta}'}$

From Definition 2.95 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \hat{\triangleright}^{\hat{\beta}'} {}^s \theta'_1 \wedge \forall k < (n - j - 1), {}^s v_2. (H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - j - 1 - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''} \end{aligned}$$

Instantiating with H'_{s1}, H'_{t1} and since we know that $(H_s, e_s[]) \Downarrow_i (H'_s, {}^s v)$ and from fg-FE we know that $i = j + k + 1 < n$ therefore we know that $k < n - j - 1$ s.t. $(H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2)$. Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H'_{t1}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - j - 1 - k, H'_{s2}, H'_{t2}) \hat{\triangleright}^{\hat{\beta}''} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-FE1.3}) \end{aligned}$$

Let $\tau[\ell'/\alpha] \sigma = \mathbf{A}^{\ell_i}$, since $\tau[\ell'/\alpha] \sigma \searrow \ell \sigma$ therefore $\ell \sigma \sqsubseteq \ell_i$ and

$$\begin{aligned} & ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''} \\ & \text{Therefore from Definition 2.94 we know that} \\ & ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, \text{Lb}({}^t v_{2i})) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''} \quad (\text{F-FE1.4}) \end{aligned}$$

In order to prove (F-FE0) we choose H'_t as H'_{t2} and ${}^t v$ as $\text{Lb}({}^t v_{2i})$. We need to prove

$$(a) (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[][\bullet]))) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i})):$$

From Lemma 2.103 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b[][\bullet]))) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$:
From (F-FE1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t11} = {}^t v_1$
- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.b[][\bullet]))[{}^t v_1/a] \delta^t \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_{2i}))$:
Again from cg-bind it suffices to prove that

- $(H'_{t1}, (\text{unlabel } a)[^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12})$:
From (F-FE1.1) we know that ${}^t v_1 = \text{Lb}({}^t v_i)$
Therefore from cg-unlabel we have $H'_{t12} = H'_{t1}$ and ${}^t v_{t12} = {}^t v_i$
- $(b[\square\square\square\bullet][^t v_i/b] \delta^t) \Downarrow {}^t v_{t13}$:
From (F-FE1.2) we know that ${}^s v_1 = \Lambda e'_s$ and ${}^t v_i = \Lambda\Lambda\Lambda\nu e'_t$
Therefore from cg-FE and cg-CE we know that ${}^t v_{t13} = e'_t$
- $(H'_{t1}, e'_t \Downarrow^f (H'_{t2}, \text{Lb}{}^t v_{2i}))$
From (F-FE1.3) and (F-FE1.4) we get the desired.

(b) $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}'}$:

We choose ${}^s \theta'$ as ${}^s \theta'_2$ and $\hat{\beta}'$ as $\hat{\beta}''$. From fg-FE we know that $i = j + k + 1$, ${}^s v = {}^s v'_2$, ${}^t v = {}^t v'_2$, $H'_s = H'_{s2}$ and $H'_t = H'_{t2}$.

Therefore from (F-FE1.3) we get the $(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta'_2$

To prove: $({}^s \theta'_2, n - i, {}^s v'_2, {}^t v'_2) \in [\tau[\ell'/\alpha_g] \sigma]_{V}^{\hat{\beta}''}$

We get this directly from (F-FE1.3)

12. FC-CI:

$$\frac{\Sigma; \Psi, c; \Gamma \vdash_{\ell_e} e : \tau \rightsquigarrow e_c}{\Sigma; \Psi; \Gamma \vdash_{pc} \nu e : (c \xrightarrow{\ell_e} \tau)^\perp \rightsquigarrow \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c))))} \text{ CI}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V}^{\hat{\beta}}$

To prove: $({}^s \theta, n, \nu e \delta^s, \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \in [\tau \sigma]_{E}^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \nu e_s \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_{V}^{\hat{\beta}'} \end{aligned}$$

This means given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$. Also given some $i < n, {}^s v$ s.t $(H_s, \nu e_s \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{ret}(\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_{V}^{\hat{\beta}'} \end{aligned}$$

From fg-val we know that ${}^s v = (\nu e_s) \delta^s$, $H'_s = H_s$ and $i = 0$. Also from cg-ret we know that $H'_t = H_t$ and ${}^t v = (\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t$

It suffices to prove that

$$\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(c \xrightarrow{\ell_e} \tau)^\perp \sigma]_{V}^{\hat{\beta}'}$$

We choose ${}^s \theta'$ as ${}^s \theta$ and $\hat{\beta}'$ as $\hat{\beta}$

(a) $(n, H_s, H_t) \hat{\triangleright}^{s\theta}$: Given

(b) $({}^s\theta, n, \nu e_s \delta^s, (\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\perp \sigma \rrbracket_V^{\hat{\beta}}$:

From Definition 2.94 it suffices to prove that

$$({}^s\theta, n, \Lambda e_s \delta^s, (\text{Lb}(\Lambda\Lambda(\nu(e_c)))) \delta^t) \in \llbracket (c \xrightarrow{\ell_\xi} \tau) \sigma \rrbracket_V^{\hat{\beta}}$$

Again from Definition 2.94 it suffices to prove that

$$\mathcal{L} \models c \sigma \implies \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'. ({}^s\theta', j, e_s, e_t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'}$$

This further means that given $\mathcal{L} \models c \sigma$ and ${}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'$

And we need to prove

$$({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'} \quad (\text{F-CI0})$$

$$\underline{\text{IH}}: ({}^s\theta', j, e_s \delta^s, e_t \delta^t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'}$$

We get (F-CI0) directly from IH

13. FC-CE:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (c \xrightarrow{\ell_\xi} \tau)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash c \quad \Sigma; \Psi \vdash pc \sqcup \ell \sqsubseteq \ell_e \quad \Sigma; \Psi \vdash \tau \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} e_s \bullet : \tau \rightsquigarrow \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b \square \square \bullet)))} \text{CE}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in \llbracket \Gamma \sigma \rrbracket_V^{\hat{\beta}}$

To prove: $({}^s\theta, n, e_s \bullet \delta^s, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b \square \square \bullet)))) \delta^t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright}^{s\theta} \wedge \forall i < n, {}^s v. (H_s, e_s \square) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b \square \square \bullet)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \end{aligned}$$

$$(n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'}$$

This means given some H_s, H_t s.t $(n, H_s, H_t) \hat{\triangleright}^{\gamma, \hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, e_s \square) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a.b \square \square \bullet)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}.$$

$$(n - i, H'_s, H'_t) \hat{\triangleright}^{s\theta'} \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}'} \quad (\text{F-CE0})$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\ell \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, e_s \llbracket \cdot \rrbracket) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < i < n$ s.t $(H_s, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

This means we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-CE1}) \end{aligned}$$

Since from (F-CE1) we have $({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (c \xrightarrow{\ell_\xi} \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1}$ therefore from Definition 2.94 we know that

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_i) \in \llbracket (c \xrightarrow{\ell_\xi} \tau) \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-CE1.1})$$

Therefore from Definition 2.94 we have

$$\begin{aligned} & {}^s v_1 = \Lambda e'_s \text{ and } {}^t v_i = \Lambda \Lambda \nu e'_t \\ & \mathcal{L} \models c \sigma \implies \forall {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, k < n - j, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_2, k, e'_s, e'_t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'_1} \quad (\text{F-CE1.2}) \end{aligned}$$

Since we know that $\mathcal{L} \models c \sigma$, we instantiate with ${}^s \theta'_1, n - j - 1, \hat{\beta}'$ to get

$$({}^s \theta'_1, n - j - 1, e'_s, e'_t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'_1}$$

From Definition 2.95 we have

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge \forall k < (n - j - 1), {}^s v_2. (H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'_1. \\ & (n - j - 1 - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''_1} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''_1} \end{aligned}$$

Instantiating with H'_{s1}, H'_{t1} and since we know that $(H_s, e_s \llbracket \cdot \rrbracket) \Downarrow_i (H'_s, {}^s v)$ and since from fg-CE we know that $i = j + k + 1 < n$ therefore we know that $k < n - j - 1$ s.t $(H_{s2}, e'_s) \Downarrow_k (H'_{s2}, {}^s v_2)$. Therefore we have

$$\begin{aligned} & \exists H'_{t2}, {}^t v_2. (H'_{t1}, e'_t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s \theta'_2 \sqsupseteq {}^s \theta'_1, \hat{\beta}'' \sqsupseteq \hat{\beta}'_1. \\ & (n - j - 1 - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''_1} {}^s \theta'_2 \wedge ({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''_1} \quad (\text{F-CE1.3}) \end{aligned}$$

Let $\tau \sigma = \mathbf{A}^{\ell_i}$, since $\tau \sigma \searrow \ell \sigma$ therefore $\ell \sigma \sqsubseteq \ell_i$ and

$$({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''_1}$$

Therefore from Definition 2.94 we know that

$$({}^s \theta'_2, n - j - 1 - k, {}^s v_2, \text{Lb}({}^t v_2)) \in \llbracket \tau \sigma \rrbracket_V^{\hat{\beta}''_1} \quad (\text{F-CE1.4})$$

In order to prove (F-CE0) we choose H'_t as H'_{t2} and ${}^t v$ as $\text{Lb}({}^t v_2)$. We need to prove

$$(a) (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b \llbracket \cdot \rrbracket \bullet)))) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_2)):$$

From Lemma 2.103 it suffices to prove that

$$(H_t, (\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.b \llbracket \cdot \rrbracket \bullet)))) \Downarrow^f (H'_{t2}, \text{Lb}({}^t v_2))$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11})$:
From (F-CE1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t11} = {}^t v_1$
- $(H'_{t1}, \text{bind}(\text{unlabel } a, b.b[][\bullet]) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, \text{Lb} {}^t v_{2i})$:
Again from cg-bind it suffices to prove that
 - $(H'_{t1}, (\text{unlabel } a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12})$:
From (F-CE1.1) we know that ${}^t v_1 = \text{Lb}({}^t v_i)$
Therefore from cg-unlabel we have $H'_{t12} = H'_{t1}$ and ${}^t v_{t12} = {}^t v_i$
 - $(b[][\bullet]) [{}^t v_i/b] \delta^t \Downarrow {}^t v_{t13}$:
From (F-CE1.2) we know that ${}^s v_1 = \Lambda e'_s$ and ${}^t v_i = \Lambda \Lambda \nu e'_t$
Therefore from cg-FE and cg-CE we know that ${}^t v_{t13} = e'_t$
 - $(H'_{t1}, e'_t \Downarrow^f (H'_{t2}, \text{Lb} {}^t v_{2i}))$
We get the desired from From (F-CE1.3) and (F-CE1.4)

(b) $\exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{V'}^{\hat{\beta}'}$:

We choose ${}^s \theta'$ as ${}^s \theta'_2$ and $\hat{\beta}'$ as $\hat{\beta}''$. From fg-CE we know that $i = j + k + 1$, ${}^s v = {}^s v'_2$, ${}^t v = {}^t v'_2$, $H'_s = H'_{s2}$ and $H'_t = H'_{t2}$.

Therefore from (F-CE1.3) we get the $(n - i, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}''} {}^s \theta'_2$

To prove: $({}^s \theta'_2, n - i, {}^s v'_2, {}^t v'_2) \in [\tau \sigma]_{V'}^{\hat{\beta}''}$

From (F-CE1.3) we know that $({}^s \theta'_2, n - j - 1 - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{V'}^{\hat{\beta}''}$

14. FC-ref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : \tau \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \tau \searrow pc}{\Sigma; \Psi; \Gamma \vdash_{pc} \text{new } (e_s) : (\text{ref } \tau)^\perp \rightsquigarrow \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb } b)))} \text{ref}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s \theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{V'}^{\hat{\beta}}$

To prove: $({}^s \theta, n, \text{new } (e_s) \delta^s, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb } b)) \delta^t) \delta^t) \in [(\text{ref } \tau)^\perp \sigma]_{E'}^{\hat{\beta}}$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s \theta \wedge \forall i < n, {}^s v. (H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_{V'}^{\hat{\beta}'} \end{aligned}$$

This means that given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s \theta$. Also given some $i < n, {}^s v$ s.t $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$.

And we are required to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(e_t, a.\text{bind}(\text{new } (a), b.\text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [(\text{ref } \tau)^\perp \sigma]_{V'}^{\hat{\beta}'} \quad (\text{F-R0}) \end{aligned}$$

IH:

$$({}^s \theta, n, e_s \delta^s, e_t \delta^t) \in [\tau \sigma]_{E'}^{\hat{\beta}}$$

This means from Definition 2.95 we have

$$\begin{aligned} \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, \text{new } (e_s) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore we know that $\exists j < n$ s.t $(H_s, e_s \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$.

Therefore we have

$$\begin{aligned} \exists H'_{t1}, {}^t v_1. (H_t, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in \lfloor \tau \sigma \rfloor_V^{\hat{\beta}'_1} \quad (\text{F-R1}) \end{aligned}$$

In order to prove (F-R0) we choose H'_t as $H'_1 \cup \{a_t \mapsto {}^t v_1\}$, ${}^t v = \text{Lb}(a_t)$, ${}^s \theta'$ as ${}^s \theta'_1 \cup \{a_s \mapsto \tau \sigma\}$ and $\hat{\beta}'$ as $\hat{\beta}'_1 \cup \{(a_s, a_t)\}$

And we need to prove:

$$(a) (H_t, \text{bind}(e_t, a. \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b))) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

- $(H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1})$:
From (F-R1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t1} = {}^t v_1$
- $(H'_1, \text{bind}(\text{new } (a), b. \text{ret}(\text{Lb } b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, {}^t v_{t2})$:

From cg-bind it suffices to prove that

- i. $(H'_1, \text{new } (a) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t2}, {}^t v_{t2})$:
From cg-new we know that $H'_{t2} = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$ and ${}^t v_{t2} = a_t$
- ii. $(H'_1 \cup \{a_t \mapsto {}^t v_1\}, \text{ret}(\text{Lb } b)) [{}^t v_1/a] [a_t/b] \delta^t) \Downarrow^f (H'_t, {}^t v_t)$:
From cg-ret we know that $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$ and ${}^t v_t = \text{Lb}(a_t)$

$$(b) \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_V^{\hat{\beta}'_1}:$$

From (F-R1) we know that $(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1$ and since $H'_s = H'_{s1} \cup \{a_s \mapsto {}^s v_1\}$, $H'_t = H'_{t1} \cup \{a_t \mapsto {}^t v_1\}$, ${}^s \theta' = {}^s \theta'_1 \cup \{a_s \mapsto \tau \sigma\}$

Therefore from Definition 2.96 and Lemma 2.102 we get $(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta'$

To prove: $({}^s \theta', n - i, {}^s v, {}^t v) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_V^{\hat{\beta}'_1}$

Since we know that ${}^s v = a_s$ and ${}^t v = \text{Lb } a_t$ therefore we need to prove

$$({}^s \theta', n - i, a_s, \text{Lb}(a_t)) \in \lfloor (\text{ref } \tau)^\perp \sigma \rfloor_V^{\hat{\beta}'_1}$$

From Definition 2.94 it suffices to prove that

$$({}^s \theta', n - i, a_s, a_t) \in \lfloor (\text{ref } \tau) \sigma \rfloor_V^{\hat{\beta}'_1}$$

Again from Definition 2.94 it suffices to prove that

$${}^s \theta'(a_s) = \tau \sigma \wedge (a_s, a_t) \in \hat{\beta}'_1$$

We get this by construction

15. FC-deref:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_s : (\text{ref } \tau)^\ell \rightsquigarrow e_t \quad \Sigma; \Psi \vdash \tau <: \tau' \quad \Sigma; \Psi \vdash \tau' \searrow \ell}{\Sigma; \Psi; \Gamma \vdash_{pc} !e_s : \tau' \rightsquigarrow \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))} \text{deref}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove: $({}^s\theta, n, !e \delta^s, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b))) \delta^t) \in [\tau' \sigma]_{\hat{E}}^{\hat{\beta}}$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, !e_s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}'} \end{aligned}$$

This means that we are given some H_s, H_t s.t $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}'} \quad (\text{F-DR0}) \end{aligned}$$

IH:

$$({}^s\theta, n, e_s \delta^s, e_t \delta^t) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{E}}^{\hat{\beta}}$$

This means from Definition 2.95 we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, !e_s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < n$ s.t $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsupseteq {}^s\theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-DR1}) \end{aligned}$$

From (F-DR1) we have $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in [(\text{ref } \tau)^\ell \sigma]_{\hat{V}}^{\hat{\beta}'_1}$

From Definition 2.94 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{\hat{V}}^{\hat{\beta}'_1} \quad (\text{F-DR1.1})$$

From Definition 2.94 we know that ${}^s v_1 = a_s$ and ${}^t v_i = a_t$

$${}^s\theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \quad (\text{F-DR1.2})$$

Let $\tau' \sigma = A^{\ell_i}$, since $\tau' \sigma \searrow \ell \sigma$ therefore $\ell \sigma \sqsubseteq \ell_i$ and

Let $v_g = H_t(a_t)$ therefore from Definition 1.76 we have

$$({}^s\theta, n-1, H_s(a_s), \text{Lb}v_{gi}) \in [\tau']_{\hat{V}}^{\hat{\beta}} \quad (\text{F-DR1.3})$$

In order to prove (F-DR0) we choose H'_t as H'_{t1} and ${}^t v$ as $H'_{t1}(a_t) = v_g = \text{Lb}v_{gi}$

$$(a) \ (H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \delta^t \Downarrow^f (H'_{t1}, \text{Lb}v_{gi}):$$

From Lemma 2.103 it suffices to prove that

$$(H_t, \text{coerce_taint}(\text{bind}(e_t, a.\text{bind}(\text{unlabel } a, b.!b)))) \delta^t \Downarrow^f (H'_{t1}, \text{Lb}v_{gi})$$

From cg-bind it suffices to prove

$$i. \ (H_t, e_t \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t1}):$$

From (F-DR1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t1} = {}^t v_1$

$$ii. \ (H'_{t1}, \text{bind}(\text{unlabel } a, b.!b)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t2}):$$

From cg-bind it suffices to prove that

$$A. \ (H'_{t1}, (\text{unlabel } a)[{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$$

From (F-DR1.1) we know that ${}^t v_1 = \text{Lb}({}^t v_i)$

Therefore from cg-unlabel we know that $H'_{t21} = H'_{t1}$ and ${}^t v_{t21} = {}^t v_i$

$$B. \ (H'_{t1}, (!b)[{}^t v_1/a][{}^t v_i/b] \delta^t) \Downarrow^f (H'_t, \text{Lb}v_{gi}):$$

Since from (F-DR1.2) we know that ${}^t v_i = a_t$ therefore from cg-deref we know that $H'_t = H'_{t1}$ and ${}^t v = H'_{t1}(a_t) = v_g = \text{Lb}v_{gi}$

$$(b) \ \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n-i, H'_s, H'_t) \hat{\triangleright}^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n-i, {}^s v, {}^t v) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}'}$$

We choose ${}^s\theta'$ as ${}^s\theta'_1$ and $\hat{\beta}'$ as $\hat{\beta}'_1$

Therefore from (F-DR1) we get $(n-j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$ and since $i = j+1$ therefore from Lemma 2.102 we get $(n-i, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$

Since from (F-DR1.2) we know that $(a_s, a_t) \in \hat{\beta}'_1$ and ${}^s\theta'_1(a_s) = \tau$. Also from (F-DR1) we have $(n-j, H'_{s1}, H'_{t1}) \hat{\triangleright}^{\hat{\beta}'_1} {}^s\theta'_1$. Therefore from Definition 2.95 we have $(n-j-1, H'_{s1}(a_s), H'_{t1}(a_t)) \in [{}^s\theta'_1(a_s)]_{\hat{V}}^{\hat{\beta}'_1}$

Since $i = j+1$, ${}^s\theta'_1(a_s) = \tau$, $H'_{s1}(a_s) = {}^s v$ and $H'_{t1}(a_t) = {}^t v$

Therefore we get

$$({}^s\theta', n-i, {}^s v, {}^t v) \in [\tau \sigma]_{\hat{V}}^{\hat{\beta}'}$$

Finally from Lemma 2.105 we get

$$({}^s\theta', n-i, {}^s v, {}^t v) \in [\tau' \sigma]_{\hat{V}}^{\hat{\beta}'}$$

16. FC-assign:

$$\frac{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} : (\text{ref } \tau)^\ell \rightsquigarrow e_{t1} \quad \Sigma; \Psi; \Gamma \vdash_{pc} e_{s2} : \tau \rightsquigarrow e_{t2} \quad \Sigma; \Psi \vdash \tau \searrow (pc \sqcup \ell)}{\Sigma; \Psi; \Gamma \vdash_{pc} e_{s1} := e_{s2} : \text{unit} \rightsquigarrow \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}())} \text{assign}$$

Also given is: $\mathcal{L} \models \Psi \sigma \wedge ({}^s\theta, n, \delta^s, \delta^t) \in [\Gamma \sigma]_{\hat{V}}^{\hat{\beta}}$

To prove:

$$({}^s\theta, n, (e_{s1} := e_{s2}) \delta^s, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}())) \delta^t) \in \llbracket \text{unit } \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 2.95 we are required to prove

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} {}^s\theta \wedge \forall i < n, {}^s v. (H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}())) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsubseteq {}^s\theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}'} \end{aligned}$$

This means that given some H_s, H_t s.t. $(n, H_s, H_t) \triangleright^{\gamma, \hat{\beta}} {}^s\theta$. Also given some $i < n, {}^s v$ s.t. $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$

And we need to prove

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))), d.\text{ret}())) \delta^t) \Downarrow^f \\ & (H'_t, {}^t v) \wedge \exists {}^s\theta' \sqsubseteq {}^s\theta, \hat{\beta}' \sqsubseteq \hat{\beta}. (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in \llbracket \text{unit} \rrbracket_V^{\hat{\beta}'} \quad (\text{F-AN0}) \end{aligned}$$

IH1:

$$({}^s\theta, n, e_{s1} \delta^s, e_{t1} \delta^t) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_E^{\hat{\beta}}$$

This means from Definition 2.95 we are required to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\gamma, \hat{\beta}} {}^s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsubseteq {}^s\theta, \hat{\beta}'_1 \sqsubseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \end{aligned}$$

Instantiating with H_s, H_t and since we know that $(H_s, (e_{s1} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists j < n$ s.t. $(H_{s1}, e_{s1} \delta^s) \Downarrow_j (H'_{s1}, {}^s v_1)$

Therefore we have

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_{t1} \delta^t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s\theta'_1 \sqsubseteq {}^s\theta, \hat{\beta}'_1 \sqsubseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s\theta'_1 \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-AN1}) \end{aligned}$$

Since from (F-AN1) we know that $({}^s\theta'_1, n - j, {}^s v_1, {}^t v_1) \in \llbracket (\text{ref } \tau)^\ell \sigma \rrbracket_V^{\hat{\beta}'_1}$ therefore from Definition 2.94 we have

$$\exists {}^t v_i. {}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in \llbracket (\text{ref } \tau) \sigma \rrbracket_V^{\hat{\beta}'_1} \quad (\text{F-AN1.1})$$

From Definition 2.94 this further means that

$${}^s\theta'_1(a_s) = \tau \wedge (a_s, a_t) \in \hat{\beta}'_1 \text{ where } {}^s v_1 = a_s \text{ and } {}^t v_1 = a_t \quad (\text{F-AN1.2})$$

IH2:

$$({}^s\theta'_1, n - j, e_{s2} \delta^s, e_{t2} \delta^t) \in \llbracket \tau \sigma \rrbracket_E^{\hat{\beta}'_1}$$

This means from Definition 2.95 we are required to prove

$$\begin{aligned} & \forall H_{s2}, H_{t2}. (n, H_{s2}, H_{t2}) \stackrel{\hat{\beta}'_1}{\triangleright} {}^s\theta'_1 \wedge \forall k < n - j, {}^s v_2. (H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2) \implies \\ & \exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1. \end{aligned}$$

$$(n - j - k, H'_{s2}, H'_{t2}) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\mathcal{V}}^{\hat{\beta}'_2}$$

Instantiating with H'_{s1}, H'_{t1} and since we know that $(H_s, (e_{s2} := e_{s2}) \delta^s) \Downarrow_i (H'_s, {}^s v)$ therefore $\exists k < n - j$ s.t $(H_{s2}, e_{s2} \delta^s) \Downarrow_k (H'_{s2}, {}^s v_2)$

Therefore we have

$$\exists H'_{t2}, {}^t v_2. (H_{t2}, e_{t2} \delta^t) \Downarrow^f (H'_{t2}, {}^t v_2) \wedge \exists {}^s\theta'_2 \sqsupseteq {}^s\theta'_1, \hat{\beta}'_2 \sqsupseteq \hat{\beta}'_1.$$

$$(n - j - k, H'_{s2}, H'_{t2}) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta'_2 \wedge ({}^s\theta'_2, n - j - k, {}^s v_2, {}^t v_2) \in [\tau \sigma]_{\mathcal{V}}^{\hat{\beta}'_2} \wedge$$

(F-AN2)

In order to prove (F-AN0) we choose H'_t as $H'_{t2}[a_t \mapsto {}^s v_2]$, ${}^t v$ as ()

We need to prove

$$(a) (H_t, \text{bind}(\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f (H'_t, {}^t v):$$

From cg-bind it suffices to prove that

$$- (H_t, (\text{toLabeled}(\text{bind}(e_{c1}, a.\text{bind}(e_{c2}, b.\text{bind}(\text{unlabel } a, c.c := b))))), d.\text{ret}()) \delta^t) \Downarrow^f (H'_T, {}^t v_T):$$

From cg-toLabeled it suffices to prove that

$$(H_t, \text{bind}(e_{t1}, a.\text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b))) \delta^t) \Downarrow^f (H'_T, {}^t v_{Ti})$$

where ${}^t v_T = \text{Lb}^t v_{Ti}$

From cg-bind it further suffices to prove that:

- $(H_t, e_{t1} \delta^t) \Downarrow^f (H'_{t11}, {}^t v_{t11}):$

From (F-AN1) we know that $H'_{t11} = H'_{t1}$ and ${}^t v_{t11} = {}^t v_1$

- $(H'_{t1}, \text{bind}(e_{t2}, b.\text{bind}(\text{unlabel } a, c.c := b)) [{}^t v_1/a] \delta^t) \Downarrow^f (H'_{t12}, {}^t v_{t12}):$

From cg-bind it suffices to prove

- $(H'_{t1}, e_{t2} \delta^t) \Downarrow^f (H'_{t13}, {}^t v_{t13}):$

From (F-AN2) we know that $H'_{t13} = H'_{t2}$ and ${}^t v_{t13} = {}^t v_2$

- $(H'_{t1}, \text{bind}(\text{unlabel } a, c.c := b) [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_t, {}^t v):$

From cg-bind it suffices to prove that

- * $(H'_{t1}, \text{unlabel } a [{}^t v_1/a] [{}^t v_2/b] \delta^t) \Downarrow^f (H'_{t21}, {}^t v_{t21}):$

From (F-AN1.1) we know that

$${}^t v_1 = \text{Lb}({}^t v_i) \wedge ({}^s\theta'_1, n - j, {}^s v_1, {}^t v_i) \in [(\text{ref } \tau) \sigma]_{\mathcal{V}}^{\hat{\beta}'_1}$$

Therefore from cg-unlabel we know that $H'_{t21} = H'_{t1}$ and ${}^t v_{t21} = {}^t v_i = a_t$

- * $(H'_{t1}, (c := b) [{}^t v_1/a] [{}^t v_2/b] [{}^t v_i/c] \delta^t) \Downarrow^f (H'_T, {}^t v_{Ti}):$

From cg-assign we know that $H'_T = H'_{t1}[a_t \mapsto {}^t v_2]$ and ${}^t v_{Ti} = ()$

Since ${}^t v_{t12} = {}^t v_{Ti} = ()$ therefore ${}^t v_T = \text{Lb}()$

$$- (H'_T, \text{ret}() [{}^t v_T/d]) \delta^t) \Downarrow^f (H'_t, ()):$$

From cg-ret and cg-val

$$(b) \exists {}^s\theta' \sqsupseteq {}^s\theta, \hat{\beta}' \sqsupseteq \hat{\beta}. (n - i, H'_s, H'_t) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta' \wedge ({}^s\theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_{\mathcal{V}}^{\hat{\beta}'_2}:$$

We choose ${}^s\theta'$ as ${}^s\theta'_2$ and $\hat{\beta}'$ as $\hat{\beta}'_2$

In order to prove $(n - i, H'_s, H'_t) \stackrel{\hat{\beta}'_2}{\triangleright} {}^s\theta'_2$ it suffices to prove

- $\text{dom}({}^s\theta'_2) \subseteq \text{dom}(H'_s)$:

Since from (F-AN2) we know that $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2$ therefore from Definition 2.96 we get $\text{dom}({}^s\theta'_2) \subseteq \text{dom}(H'_s)$

- $\hat{\beta}'_2 \subseteq (\text{dom}({}^s\theta'_2) \times \text{dom}(H'_t))$:

Since from (F-AN2) we know that $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2$ therefore from Definition 2.96 we get

$$\hat{\beta}'_2 \subseteq (\text{dom}({}^s\theta'_2) \times \text{dom}(H'_t))$$

- $\forall (a_1, a_2) \in \hat{\beta}'_2. ({}^s\theta'_2, n - i - 1, H'_s(a_1), H'_t(a_2)) \in [{}^s\theta'_2(a_1)]_V^{\hat{\beta}'_2}$:

$\forall (a_1, a_2) \in \hat{\beta}'_2$.

- $a_1 = a_s$ and $a_1 = a_t$:

Since from (F-AN2) we know that $({}^s\theta'_2, n - j - k, {}^sv_2, {}^tv_2) \in [\tau \sigma]_V^{\hat{\beta}'_2}$

Also from (F-AN1.2) and Definition 2.92 we know that ${}^s\theta'_2(a_1) = \tau \sigma$

Therefore from Lemma 2.100 we get

$$({}^s\theta'_2, n - i - 1, {}^sv_2, {}^tv_2) \in [\tau \sigma]_V^{\hat{\beta}'_2}$$

- $a_1 \neq a_s$ and $a_1 \neq a_t$:

From (F-AN2) since we know that $(n - j - k, H'_{s2}, H'_{t2}) \triangleright^{\hat{\beta}'_2} {}^s\theta'_2$ therefore from Definition 2.96 we get

$$({}^s\theta'_2, n - j - k - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s\theta'_2(a_1) \sigma]_V^{\hat{\beta}'_2}$$

Since $i = j + k + 1$ therefore from Lemma 2.100 we get

$$({}^s\theta'_2, n - i - 1, H'_{s2}(a_1), H'_{t2}(a_2)) \in [{}^s\theta'_2(a_1) \sigma]_V^{\hat{\beta}'_2}$$

- $a_1 = a_s$ and $a_1 \neq a_t$:

This case cannot arise

- $a_1 \neq a_s$ and $a_1 = a_t$:

This case cannot arise

And in order to prove $({}^s\theta', n - i, {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}'_2}$

Since we know that ${}^sv = ()$ and ${}^tv = ()$ therefore from Definition 2.94 we get $({}^s\theta', n - i, {}^sv, {}^tv) \in [\text{unit}]_V^{\hat{\beta}'_2}$

□

Lemma 2.105 (FG \rightsquigarrow CG: Semantic Subtyping lemma). *The following holds:*

$\forall \Sigma, \Psi, \sigma, \mathcal{L}, \hat{\beta}$.

1. $\forall A, A'$.

$$(a) \Sigma; \Psi \vdash A <: A' \wedge \mathcal{L} \models \Psi \sigma \implies [(A \sigma)]_V^{\hat{\beta}} \subseteq [(A' \sigma)]_V^{\hat{\beta}}$$

2. $\forall \tau, \tau'$.

$$(a) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_V^{\hat{\beta}} \subseteq [(\tau' \sigma)]_V^{\hat{\beta}}$$

$$(b) \Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma \implies [(\tau \sigma)]_E^{\hat{\beta}} \subseteq [(\tau' \sigma)]_E^{\hat{\beta}}$$

Proof. Proof by simultaneous induction on $A <: A'$ and $\tau <: \tau'$

Proof of statement 1(a)

We analyse the different cases of $A <: A'$ in the last step:

1. FGsub-arrow:

Given:

$$\frac{\Sigma; \Psi \vdash \tau'_1 <: \tau_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \tau_1 \xrightarrow{\ell_e} \tau_2 <: \tau'_1 \xrightarrow{\ell'_e} \tau'_2} \text{FGsub-arrow}$$

$$\text{To prove: } [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^{\hat{\beta}} \subseteq [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^{\hat{\beta}}$$

$$\text{IH1: } [(\tau'_1 \sigma)]_V^{\hat{\beta}} \subseteq [(\tau_1 \sigma)]_V^{\hat{\beta}} \text{ (Statement 2(a))}$$

$$\text{It suffices to prove: } \forall ({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^{\hat{\beta}}. \\ ({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^{\hat{\beta}}$$

This means that given some ${}^s\theta, m$ and $\lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))$ s.t

$$({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))) \in [((\tau_1 \xrightarrow{\ell_e} \tau_2) \sigma)]_V^{\hat{\beta}}$$

Therefore from Definition 2.94 we are given:

$$\forall {}^s\theta'_1 \sqsupseteq {}^s\theta, {}^s v_1, {}^t v_1, j < m, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, {}^s v_1, {}^t v_1) \in [\tau_1 \sigma]_V^{\hat{\beta}'_1} \implies \\ ({}^s\theta'_1, j, e_s[{}^s v_1/x] \delta^s, e_t[{}^t v_1/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_1} \quad (\text{S-L0})$$

$$\text{And it suffices to prove: } ({}^s\theta, m, \lambda x.e_s, \Lambda\Lambda(\nu(\lambda x.e_t))) \in [((\tau'_1 \xrightarrow{\ell'_e} \tau'_2) \sigma)]_V^{\hat{\beta}}$$

Again from Definition 2.94, it suffices to prove:

$$\forall {}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_2} \implies \\ ({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-L1})$$

This means that given ${}^s\theta'_2 \sqsupseteq {}^s\theta, {}^s v_2, {}^t v_2, k < m, \hat{\beta} \sqsubseteq \hat{\beta}'_2$ s.t $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_2}$

And we need to prove

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-L2})$$

Instantiating (S-L0) with ${}^s\theta'_2, {}^s v_2, {}^t v_2, k, \hat{\beta}'_2$. Since we have $({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau'_1 \sigma]_V^{\hat{\beta}'_2}$ therefore from IH1 we also have

$$({}^s\theta'_2, k, {}^s v_2, {}^t v_2) \in [\tau_1 \sigma]_V^{\hat{\beta}'_2}$$

Therefore we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau_2 \sigma]_E^{\hat{\beta}'_2}$$

$$\text{IH2: } [(\tau_2 \sigma)]_E^{\hat{\beta}} \subseteq [(\tau'_2 \sigma)]_E^{\hat{\beta}} \text{ (Statement 2(b))}$$

Finally using IH2 we get

$$({}^s\theta'_2, k, e_s[{}^s v_2/x] \delta^s, e_t[{}^t v_2/x] \delta^t) \in [\tau'_2 \sigma]_E^{\hat{\beta}'_2}$$

2. FGsub-prod:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 \times \tau_2 <: \tau'_1 \times \tau'_2} \text{FGsub-prod}$$

To prove: $\llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

IH1: $\llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}}$ (Statement 2(a))

IH2: $\llbracket (\tau_2 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V^{\hat{\beta}}$ (Statement 2(a))

It suffices to prove:

$$\forall ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

This means that given some ${}^s\theta, n$ and ${}^s v_1, {}^s v_2, {}^t v_1, {}^t v_2$ s.t

$$({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau_1 \times \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

Therefore from Definition 2.94 we are given:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-P0})$$

And it suffices to prove: $({}^s\theta, m, ({}^s v_1, {}^s v_2), ({}^t v_1, {}^t v_2)) \in \llbracket ((\tau'_1 \times \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

Again from Definition 2.94, it suffices to prove:

$$({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \sigma \rrbracket_V^{\hat{\beta}} \wedge ({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau'_2 \sigma \rrbracket_V^{\hat{\beta}} \quad (\text{S-P1})$$

Since from (S-P0) we know that $({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau_1 \sigma \rrbracket_V^{\hat{\beta}}$ therefore from IH1 we have $({}^s\theta, m, {}^s v_1, {}^t v_1) \in \llbracket \tau'_1 \sigma \rrbracket_V^{\hat{\beta}}$

Similarly since we have $({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau_2 \sigma \rrbracket_V^{\hat{\beta}}$ from (S-P0) therefore from IH2 we have $({}^s\theta, m, {}^s v_2, {}^t v_2) \in \llbracket \tau'_2 \sigma \rrbracket_V^{\hat{\beta}}$

3. FGsub-sum:

Given:

$$\frac{\Sigma; \Psi \vdash \tau_1 <: \tau'_1 \quad \Sigma; \Psi \vdash \tau_2 <: \tau'_2}{\Sigma; \Psi \vdash \tau_1 + \tau_2 <: \tau'_1 + \tau'_2} \text{FGsub-sum}$$

To prove: $\llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

IH1: $\llbracket (\tau_1 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_1 \sigma) \rrbracket_V^{\hat{\beta}}$ (Statement 2(a))

IH2: $\llbracket (\tau_2 \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket (\tau'_2 \sigma) \rrbracket_V^{\hat{\beta}}$ (Statement 2(a))

It suffices to prove: $\forall ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

This means that given: $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau_1 + \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

And it suffices to prove: $({}^s\theta, n, {}^s v, {}^t v) \in \llbracket ((\tau'_1 + \tau'_2) \sigma) \rrbracket_V^{\hat{\beta}}$

2 cases arise

(a) ${}^s v = \text{inl } {}^s v_i$ and ${}^t v = \text{inl } {}^t v_i$:

From Definition 2.94 we are given:

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau_1 \sigma]_V^{\hat{\beta}} \quad (\text{S-S0})$$

And we are required to prove that:

$$({}^s \theta, n, {}^s v_i, {}^t v_i) \in [\tau_1' \sigma]_V^{\hat{\beta}}$$

From (S-S0) and IH1 get this

(b) ${}^s v = \text{inr } {}^s v_i$ and ${}^t v = \text{inr } {}^t v_i$:

Symmetric reasoning as in the previous case

4. FGsub-forall:

Given:

$$\frac{\Sigma, \alpha; \Psi \vdash \tau_1 <: \tau_2 \quad \Sigma, \alpha; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash \forall \alpha. (\ell_e, \tau_1) <: \forall \alpha. (\ell'_e, \tau_2)} \text{FGsub-forall}$$

To prove: $[(\forall \alpha. (\ell_e, \tau_1)) \sigma]_V^{\hat{\beta}} \subseteq [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^{\hat{\beta}}$

It suffices to prove:

$$\forall ({}^s \theta, n, \Lambda e_s, \Lambda \Lambda \Lambda (\nu(e_t))) \in [(\forall \alpha. (\ell_e, \tau_1)) \sigma]_V^{\hat{\beta}}. ({}^s \theta, n, \Lambda e_s, \Lambda \Lambda \Lambda (\nu(e_t))) \in [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^{\hat{\beta}}$$

This means that given $({}^s \theta, n, \Lambda e_s, \Lambda \Lambda \Lambda (\nu(e_t))) \in [(\forall \alpha. (\ell_e, \tau_1)) \sigma]_V^{\hat{\beta}}$

Therefore from Definition 2.94 we have:

$$\forall {}^s \theta'_1 \sqsupseteq {}^s \theta, j < n, \ell' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s \theta'_1, j, e_s, e_t) \in [\tau_1[\ell'/\alpha] \sigma]_E^{\hat{\beta}'_1} \quad (\text{S-F0})$$

And we need to prove

$$({}^s \theta, n, \Lambda e_s, \Lambda \Lambda \Lambda (\nu(e_t))) \in [(\forall \alpha. (\ell'_e, \tau_2)) \sigma]_V^{\hat{\beta}}$$

Again from Definition 2.94 it means we need to prove

$$\forall {}^s \theta'_2 \sqsupseteq {}^s \theta, k < n, \ell'' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s \theta'_2, k, e_s, e_t) \in [\tau_2[\ell''/\alpha] \sigma]_E^{\hat{\beta}'_2}$$

This means that given ${}^s \theta'_2 \sqsupseteq {}^s \theta, k < n, \ell'' \in \mathcal{L}, \hat{\beta} \sqsubseteq \hat{\beta}'_2$

And we need to prove

$$({}^s \theta'_2, k, e_s, e_t) \in [\tau_2[\ell''/\alpha] \sigma]_E^{\hat{\beta}'_2} \quad (\text{S-F1})$$

Instantiating (S-F0) with ${}^s \theta'_2, k, \ell'', \hat{\beta}'_2$ and we get

$$({}^s \theta'_2, k, e_s, e_t) \in [\tau_1[\ell''/\alpha] \sigma]_E^{\hat{\beta}'_2}$$

IH: $[(\tau_1 \sigma \cup \{\alpha \mapsto \ell''\})]_E^{\hat{\beta}'_2} \subseteq [(\tau_2 \sigma \cup \{\alpha \mapsto \ell''\})]_E^{\hat{\beta}'_2}$ (Statement 2(b))

Therefore from IH we get the desired

5. FGsub-constraint:

Given:

$$\frac{\Sigma; \Psi \vdash c_2 \Longrightarrow c_1 \quad \Sigma; \Psi, c_2 \vdash \tau_1 <: \tau_2 \quad \Sigma; \Psi \vdash \ell'_e \sqsubseteq \ell_e}{\Sigma; \Psi \vdash c_1 \xrightarrow{\ell'_e} \tau_1 <: c_2 \xrightarrow{\ell'_e} \tau_2} \text{FGsub-constraint}$$

To prove: $\llbracket ((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$

It suffices to prove:

$$\forall ({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in \llbracket ((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in \llbracket ((c_2 \xrightarrow{\ell'_e} \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

This means that given: $({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in \llbracket ((c_1 \xrightarrow{\ell'_e} \tau_1) \sigma) \rrbracket_V^{\hat{\beta}}$

Therefore from Definition 2.94 we are given:

$$\mathcal{L} \models c_1 \sigma \Longrightarrow \forall {}^s\theta' \sqsupseteq {}^s\theta, j < n, \hat{\beta} \sqsubseteq \hat{\beta}'_1. ({}^s\theta'_1, j, e_s, e_t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'_1} \quad (\text{S-C0})$$

And it suffices to prove:

$$({}^s\theta, n, \nu e_s, \Lambda\Lambda(\nu(e_t))) \in \llbracket ((c_1 \xrightarrow{\ell'_e} \tau_2) \sigma) \rrbracket_V^{\hat{\beta}}$$

Again from Definition 2.94 it means that we need to prove:

$$\mathcal{L} \models c_2 \sigma \Longrightarrow \forall {}^s\theta'_2 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_2. ({}^s\theta'_2, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_2}$$

This means that given that $\mathcal{L} \models c_2 \sigma$ and ${}^s\theta'_2 \sqsupseteq {}^s\theta, k < n, \hat{\beta} \sqsubseteq \hat{\beta}'_2$

And we need to prove

$$({}^s\theta'_2, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_2} \quad (\text{S-C1})$$

Instantiating (S-C0) with ${}^s\theta'_2, k, \hat{\beta}'_2$ we get $({}^s\theta'_2, k, e_s, e_t) \in \llbracket \tau_1 \sigma \rrbracket_E^{\hat{\beta}'_2}$

IH: $\llbracket (\tau_1 \sigma) \rrbracket_E^{\hat{\beta}'_2} \subseteq \llbracket (\tau_2 \sigma) \rrbracket_E^{\hat{\beta}'_2}$ (Statement 2(b))

Finally from IH we get $({}^s\theta'_2, k, e_s, e_t) \in \llbracket \tau_2 \sigma \rrbracket_E^{\hat{\beta}'_2}$

6. FGsub-ref:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{ref } \tau <: \text{ref } \tau} \text{FGsub-ref}$$

To prove: $\llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}} \subseteq \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

It suffices to prove: $\forall ({}^s\theta, n, a_s, a_t) \in \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}. ({}^s\theta, n, a_s, a_t) \in \llbracket ((\text{ref } \tau) \sigma) \rrbracket_V^{\hat{\beta}}$

We get this directly from Definition 2.94

7. FGsub-base:

Given:

$$\frac{}{\Sigma; \Psi \vdash \mathbf{b} <: \mathbf{b}} \text{FGsub-base}$$

$$\text{To prove: } [((\mathbf{b}) \sigma)]_V^{\hat{\beta}} \subseteq [((\mathbf{b}) \sigma)]_V^{\hat{\beta}}$$

Directly from Definition 2.94

8. FGsub-unit:

Given:

$$\frac{}{\Sigma; \Psi \vdash \text{unit} <: \text{unit}} \text{FGsub-unit}$$

$$\text{To prove: } [((\text{unit}) \sigma)]_V^{\hat{\beta}} \subseteq [((\text{unit}) \sigma)]_V^{\hat{\beta}}$$

Directly from Definition 2.94

Proof of statement 2(a)

Given:

$$\frac{\Sigma; \Psi \vdash \ell' \sqsubseteq \ell'' \quad \Sigma; \Psi \vdash \mathbf{A} <: \mathbf{A}'}{\Sigma; \Psi \vdash \mathbf{A}^{\ell'} <: \mathbf{A}^{\ell''}} \text{FGsub-label}$$

$$\text{To prove: } [(\mathbf{A}^{\ell'}) \sigma]_V^{\hat{\beta}} \subseteq [(\mathbf{A}^{\ell''}) \sigma]_V^{\hat{\beta}}$$

This means from Definition 2.94 we need to prove

$$\forall ({}^s\theta, n, {}^sv, \text{Lb}({}^tv_i)) \in [(\mathbf{A}^{\ell'} \sigma)]_V^{\hat{\beta}}. ({}^s\theta, n, {}^sv, \text{Lb}({}^tv_i)) \in [(\mathbf{A}^{\ell''} \sigma)]_V^{\hat{\beta}}$$

This means that given $({}^s\theta, n, {}^sv, \text{Lb}({}^tv_i)) \in [(\mathbf{A}^{\ell'} \sigma)]_V^{\hat{\beta}}$

From Definition 2.94 it further means that we are given

$$({}^s\theta, n, {}^sv, {}^tv_i) \in [(\mathbf{A} \sigma)]_V^{\hat{\beta}} \quad (\text{S-LB0})$$

And we need to prove

$$({}^s\theta, n, {}^sv, \text{Lb}({}^tv_i)) \in [(\mathbf{A}^{\ell''} \sigma)]_V^{\hat{\beta}}$$

Again from Definition 2.94 it suffices to prove that

$$({}^s\theta, n, {}^sv, {}^tv_i) \in [(\mathbf{A}' \sigma)]_V^{\hat{\beta}}$$

Since $\ell' \sqsubseteq \ell''$ and $\mathbf{A}' <: \mathbf{A}''$ therefore from IH (Statement 1(a)) and (S-LB0) we get the desired

Proof of statement 2(b)

Given: $\Sigma; \Psi \vdash \tau <: \tau' \wedge \mathcal{L} \models \Psi \sigma$

$$\text{To prove: } [(\tau \sigma)]_E^{\hat{\beta}} \subseteq [(\tau' \sigma)]_E^{\hat{\beta}}$$

This means we need to prove that

$$\forall ({}^s\theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}. ({}^s\theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$$

This means given $({}^s\theta, n, e_s, e_t) \in [(\tau \sigma)]_E^{\hat{\beta}}$

This means from Definition 2.95 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}} s\theta \wedge \forall i < n, {}^s v. (H_s, e_s) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \end{aligned}$$

$$(n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v) \in [\tau \sigma]_V^{\hat{\beta}'} \quad (\text{S-E0})$$

And it suffices to prove that $({}^s \theta, n, e_s, e_t) \in [(\tau' \sigma)]_E^{\hat{\beta}}$

Again from Definition 2.95 it means we need to prove

$$\begin{aligned} & \forall H_{s1}, H_{t1}. (n, H_{s1}, H_{t1}) \triangleright^{\hat{\beta}} s\theta \wedge \forall j < n, {}^s v_1. (H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1) \implies \\ & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \end{aligned}$$

$$(n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}'_1}$$

This means that given some H_{s1}, H_{t1} s.t $(n, H_{s1}, H_{t1}) \triangleright^{\ell_2, \hat{\beta}} s\theta$. Also given some $j < n, {}^s v_1$ s.t $(H_{s1}, e_s) \Downarrow_j (H'_{s1}, {}^s v_1)$

And we need to prove

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}'_1} \quad (\text{S-E1}) \end{aligned}$$

Instantiating (S-E0) with H_{s1}, H_{t1} and with $j, {}^s v_1$. Then we get

$$\begin{aligned} & \exists H'_t, {}^t v. (H_t, e_t) \Downarrow^f (H'_t, {}^t v) \wedge \exists {}^s \theta' \sqsupseteq {}^s \theta, \hat{\beta}' \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_t) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

Since we have $\tau <: \tau'$. Therefore from IH (Statement 2(a)) we get

$$\begin{aligned} & \exists H'_{t1}, {}^t v_1. (H_{t1}, e_t) \Downarrow^f (H'_{t1}, {}^t v_1) \wedge \exists {}^s \theta'_1 \sqsupseteq {}^s \theta, \hat{\beta}'_1 \sqsupseteq \hat{\beta}. \\ & (n - j, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}'_1} {}^s \theta'_1 \wedge ({}^s \theta'_1, n - j, {}^s v_1, {}^t v_1) \in [\tau' \sigma]_V^{\hat{\beta}'_1} \end{aligned}$$

□

Theorem 2.106 (FG \rightsquigarrow CG: Deriving FG NI via compilation). $\forall e_s, {}^s v_1, {}^s v_2, n_1, n_2, H'_{s1}, H'_{s2}, pc$.

Let $\text{bool} = (\text{unit} + \text{unit})$

$$\begin{aligned} & \emptyset, \emptyset, x : \text{bool}^\top \vdash_{pc} e_s : \text{bool}^\perp \wedge \\ & \emptyset, \emptyset, \emptyset \vdash_{pc} {}^s v_1 : \text{bool}^\top \wedge \emptyset, \emptyset, \emptyset \vdash_{pc} {}^s v_2 : \text{bool}^\top \wedge \\ & (\emptyset, e_s[{}^s v_1/x]) \Downarrow_{n_1} (H'_{s1}, {}^s v'_1) \wedge \\ & (\emptyset, e_s[{}^s v_2/x]) \Downarrow_{n_2} (H'_{s2}, {}^s v'_2) \wedge \\ & \implies \\ & {}^s v'_1 = {}^s v'_2 \end{aligned}$$

Proof. From the FG to CG translation we know that $\exists e_t$ s.t

$$\emptyset, \emptyset, x : \text{bool}^\top \vdash e_s : \text{bool}^\perp \rightsquigarrow e_t$$

Similarly we also know that $\exists {}^t v_1, {}^t v_2$ s.t

$$\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{bool}^\top \rightsquigarrow {}^t v_1 \text{ and } \emptyset, \emptyset, \emptyset \vdash {}^s v_2 : \text{bool}^\top \rightsquigarrow {}^t v_2 \quad (\text{NI-0})$$

From type preservation theorem (choosing $\alpha = \gamma = \bar{\beta} = \perp$) we know that

$$\emptyset, \emptyset, x : \text{Labeled} \top \text{ bool} \vdash e_t : \mathbb{C} \perp \perp \text{Labeled} \perp \text{ bool}$$

$$\emptyset, \emptyset, \emptyset \vdash {}^t v_1 : \mathbb{C} \perp \perp \text{Labeled} \top \text{ bool}$$

$$\emptyset, \emptyset, \emptyset \vdash {}^t v_2 : \mathbb{C} \perp \perp \text{Labeled} \top \text{ bool} \quad (\text{NI-1})$$

Since we have $\emptyset, \emptyset, \emptyset \vdash {}^s v_1 : \text{bool}^\top \rightsquigarrow {}^t v_1$

And since ${}^s v_1$ and ${}^t v_1$ are closed terms (from given and NI-1)

Therefore from Theorem 2.104 we have (we choose $n > n_1$ and $n > n_2$)

$$(\emptyset, n, {}^s v_1, {}^t v_1) \in [\text{bool}^\top]_E^\emptyset \quad (\text{NI-2})$$

Therefore from Definition 2.95 we have

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_1) \Downarrow_i (H'_s, {}^s v) \implies \\ & \exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ & (n - i, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{11}) \in [\mathbf{bool}^\top \sigma]_V^{\hat{\beta}'} \end{aligned}$$

Instantiating with \emptyset, \emptyset and from fg-val we know that $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$. Therefore we have

$$\begin{aligned} & \exists H'_t, {}^t v_{11}. (H_t, {}^t v_1) \Downarrow^f (H'_t, {}^t v_{11}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ & (n, H'_s, H'_t) \triangleright^{\hat{\beta}'} {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{11}) \in [\mathbf{bool}^\top \sigma]_V^{\hat{\beta}'} \quad (\text{NI-2.1}) \end{aligned}$$

From Definition 2.94 we know that

$${}^t v_{11} = \mathbf{Lb}({}^t v_{i11}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{i11}) \in [(\mathbf{unit} + \mathbf{unit}) \sigma]_V^{\hat{\beta}'}$$

Again from Definition 2.94 we know that

$$\begin{aligned} & \text{Either a) } {}^s v_1 = \mathbf{inl}() \text{ and } {}^t v_{i11} = \mathbf{inl}() \text{ or b) } {}^s v_1 = \mathbf{inr}() \text{ and } {}^t v_{i11} = \mathbf{inr}() \\ & \text{But in either case we have that } \emptyset, \emptyset, \emptyset \vdash {}^t v_{i11} : (\mathbf{unit} + \mathbf{unit}) \quad (\text{NI-2.2}) \end{aligned}$$

$$\text{As a result we have } \emptyset, \emptyset, \emptyset \vdash {}^t v_{11} : \mathbf{Labeled} \top (\mathbf{unit} + \mathbf{unit}) \quad (\text{NI-2.3})$$

We give it typing derivation

$$\frac{\overline{\emptyset, \emptyset, \emptyset \vdash {}^t v_{i11} : (\mathbf{unit} + \mathbf{unit})} \quad (\text{NI-2.2})}{\emptyset, \emptyset, \emptyset \vdash \mathbf{Lb}({}^t v_{i11}) : \mathbf{Labeled} \top (\mathbf{unit} + \mathbf{unit})}$$

From Definition 2.99 and (NI-2.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_1), (x \mapsto {}^t v_{11})) \in [x \mapsto \mathbf{bool}^\top]_V^{\hat{\beta}'}$$

Therefore we can apply Theorem 2.104 to get

$$(\emptyset, n, e_s[{}^s v_1/x], e_t[{}^t v_{11}/x]) \in [\mathbf{bool}^\perp]_E^{\hat{\beta}'} \quad (\text{NI-2.4})$$

From Definition 2.95 we get

$$\begin{aligned} & \forall H_s, H_t. (n, H_s, H_t) \triangleright^{\hat{\beta}'} \emptyset \wedge \forall i < n, {}^s v''_1. (H_s, e_s[{}^s v_1/x]) \Downarrow_i (H'_{s1}, {}^s v''_1) \implies \\ & \exists H'_{t1}, {}^t v''_1. (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v''_1) \wedge \exists {}^s \theta'' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - i, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge ({}^s \theta'', n - i, {}^s v''_1, {}^t v''_1) \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''} \end{aligned}$$

Instantiating with $\emptyset, \emptyset, n_1, {}^s v'_1$ we get

$$\begin{aligned} & \exists H'_{t1}, {}^t v''_1. (H_t, e_t[{}^t v_{11}/x]) \Downarrow^f (H'_{t1}, {}^t v''_1) \wedge \exists {}^s \theta'' \sqsupseteq {}^s \theta, \hat{\beta}'' \sqsupseteq \hat{\beta}'. \\ & (n - n_1, H'_{s1}, H'_{t1}) \triangleright^{\hat{\beta}''} {}^s \theta'' \wedge ({}^s \theta'', n - n_1, {}^s v'_1, {}^t v''_1) \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''} \quad (\text{NI-2.5}) \end{aligned}$$

Since we have $({}^s \theta'', n - n_1, {}^s v'_1, {}^t v''_1) \in [\mathbf{bool}^\perp \sigma]_V^{\hat{\beta}''}$ therefore from Definition 2.94 we have

$$\exists {}^t v_{i1}. {}^t v'' = \mathbf{Lb}({}^t v_{i1}) \wedge ({}^s \theta'', n - n_1, {}^s v'_1, {}^t v_{i1}) \in [\mathbf{bool} \sigma]_V^{\hat{\beta}''}$$

Since $({}^s \theta'', n - n_1, {}^s v'_1, {}^t v_{i1}) \in [(\mathbf{unit} + \mathbf{unit})]_V^{\hat{\beta}''}$ therefore from Definition 2.94 two cases arise

- ${}^s v'_1 = \mathbf{inl} {}^s v_{i11}$ and ${}^t v_{i1} = \mathbf{inl} {}^t v_{i11}$:

From Definition 2.94 we have

$$({}^s \theta'', n - n_1, {}^s v_{i11}, {}^t v_{i11}) \in [\mathbf{unit}]_V^{\hat{\beta}''}$$

which means we have ${}^s v_{i11} = {}^t v_{i11}$

- ${}^s v'_1 = \text{inr } {}^s v_{i11}$ and ${}^t v_{i1} = \text{inr } {}^t v_{i11}$:

Symmetric reasoning as in the previous case

So no matter which case arise we have ${}^s v'_1 = {}^t v_{i1}$

Similarly with other substitution we have $(\emptyset, n, {}^s v_2, {}^t v_2) \in [\text{bool}^\top]_E^\emptyset$ (NI-3)

Therefore from Definition 2.95 we have

$$\forall H_s, H_t. (n, H_s, H_t) \triangleright \emptyset \wedge \forall i < n, {}^s v. (H_s, {}^s v_2) \Downarrow_i (H'_s, {}^s v) \implies \exists H'_t, {}^t v_{22}. (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset.$$

$$(n - i, H'_s, H'_t) \hat{\triangleright} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v, {}^t v_{22}) \in [\text{bool}^\top \sigma]_{V}^{\hat{\beta}'}$$

Instantiating with \emptyset, \emptyset and from fg-val we know that $H'_s = H_s = \emptyset, {}^s v = {}^s v_1$. Therefore we have

$$\begin{aligned} & \exists H'_t, {}^t v_{22}. (H_t, {}^t v_2) \Downarrow^f (H'_t, {}^t v_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}' \sqsupseteq \emptyset. \\ (n, H'_s, H'_t) \hat{\triangleright} & {}^s \theta' \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{22}) \in [\text{bool}^\top \sigma]_{V}^{\hat{\beta}'} \quad (\text{NI-3.1}) \end{aligned}$$

From Definition 2.94 we know that

$${}^t v_2 = \text{Lb}({}^t v_{i22}) \wedge ({}^s \theta', n, {}^s v_1, {}^t v_{i22}) \in [(\text{unit} + \text{unit}) \sigma]_{V}^{\hat{\beta}'}$$

Again from Definition 2.94 we know that

Either a) ${}^s v_2 = \text{inl}()$ and ${}^t v_{i22} = \text{inl}()$ or b) ${}^s v_2 = \text{inr}()$ and ${}^t v_{i22} = \text{inr}()$
But in either case we have that $\emptyset, \emptyset, \emptyset \vdash {}^t v_{i22} : (\text{unit} + \text{unit})$ (NI-3.2)

As a result we have $\emptyset, \emptyset, \emptyset \vdash {}^t v_{22} : \text{Labeled } \top (\text{unit} + \text{unit})$ (NI-3.3)

We give it typing derivation

$$\frac{\overline{\emptyset, \emptyset, \emptyset \vdash {}^t v_{i22} : (\text{unit} + \text{unit})} \quad (\text{NI-3.2})}{\emptyset, \emptyset, \emptyset \vdash \text{Lb}({}^t v_{i22}) : \text{Labeled } \top (\text{unit} + \text{unit})}$$

From Definition 2.99 and (NI-3.1) we know that

$$(\emptyset, n, (x \mapsto {}^s v_2), (x \mapsto {}^t v_{22})) \in [x \mapsto \text{bool}^\top]_{V}^{\hat{\beta}'}$$

Therefore we can apply Theorem 2.104 to get

$$(\emptyset, n, e_s[{}^s v_2/x], e_t[{}^t v_{22}/x]) \in [\text{bool}^\perp]_E^{\hat{\beta}'} \quad (\text{NI-3.4})$$

From Definition 2.95 we get

$$\forall H_s, H_t. (n, H_s, H_t) \hat{\triangleright} \emptyset \wedge \forall i < n, {}^s v''_2. (H_s, e_s[{}^s v_2/x]) \Downarrow_i (H'_{s2}, {}^s v''_2) \implies \exists H'_{t2}, {}^t v''_{22}. (H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H'_{t2}, {}^t v''_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \emptyset.$$

$$(n - i, H'_{s2}, H'_{t2}) \hat{\triangleright} {}^s \theta' \wedge ({}^s \theta', n - i, {}^s v''_2, {}^t v''_{22}) \in [\text{bool}^\perp \sigma]_{V}^{\hat{\beta}''}$$

Instantiating with $\emptyset, \emptyset, n_2, {}^s v''_2$ we get

$$\begin{aligned} & \exists H'_{t2}, {}^t v''_{22}. (H_t, e_t[{}^t v_{22}/x]) \Downarrow^f (H'_{t2}, {}^t v''_{22}) \wedge \exists {}^s \theta' \sqsupseteq \emptyset, \hat{\beta}'' \sqsupseteq \emptyset. \\ (n - n_1, H'_s, H'_{t2}) \hat{\triangleright} & {}^s \theta' \wedge ({}^s \theta', n - n_1, {}^s v''_2, {}^t v''_{22}) \in [\text{bool}^\perp \sigma]_{V}^{\hat{\beta}''} \quad (\text{NI-3.5}) \end{aligned}$$

Since we have $({}^s \theta', n - n_2, {}^s v''_2, {}^t v''_{22}) \in [\text{bool}^\perp \sigma]_{V}^{\hat{\beta}''}$ therefore from Definition 2.94 we have

$$\exists {}^t v_{i2}. {}^t v''_{22} = \text{Lb}({}^t v_{i2}) \wedge ({}^s \theta', n - n_2, {}^s v''_2, {}^t v_{i2}) \in [\text{bool} \sigma]_{V}^{\hat{\beta}''}$$

Since $({}^s \theta', n - n_2, {}^s v''_2, {}^t v_{i2}) \in [(\text{unit} + \text{unit})]_{V}^{\hat{\beta}''}$ therefore from Definition 2.94 two cases arise

- ${}^s v'_2 = \text{inl } {}^s v_{i22}$ and ${}^t v_{i2} = \text{inl } {}^t v_{i22}$:

From Definition 2.94 we have

$$({}^s \theta', n - n_2, {}^s v_{i22}, {}^t v_{i22}) \in [\text{unit}]_{V'}^{\hat{\beta}''}$$

which means we have ${}^s v_{i22} = {}^t v_{i22}$

- ${}^s v'_1 = \text{inr } {}^s v_{i22}$ and ${}^t v_{i2} = \text{inr } {}^t v_{i22}$:

Symmetric reasoning as in the previous case

So no matter which case arise we have ${}^s v'_2 = {}^t v_{i2}$

We know that $\emptyset, \emptyset, \emptyset \vdash {}^t v_{i1} : \text{Labeled } \top \text{ bool}$ (NI-2.3)

Also we have $\emptyset, \emptyset, \emptyset \vdash {}^t v_{i2} : \text{Labeled } \top \text{ bool}$ (NI-3.3)

Let $e_T = \text{bind}(e_t, y.\text{unlabel}(y))$

We show that $\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_T : \mathbb{C} \perp \perp \text{ bool}$ by giving a typing derivation P2:

$$\frac{\frac{}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash y : \text{Labeled } \perp \text{ bool}}{\text{CG-var}}}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool}, y : \text{Labeled } \perp \text{ bool} \vdash \text{unlabel}(y) : \mathbb{C} \perp \perp \text{ bool}} \text{CG-unlabel}$$

P1:

$$\frac{}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash e_t : \mathbb{C} \perp \perp \text{Labeled } \perp \text{ bool}} \text{From (NI-1)}$$

Main derivation:

$$\frac{P1 \quad P2}{\emptyset, \emptyset, x : \text{Labeled } \top \text{ bool} \vdash \text{bind}(e_t, y.\text{unlabel}(y)) : \mathbb{C} \perp \perp \text{ bool}}$$

Say $e_t[{}^t v_{i1}/x]$ reduces in n_{t1} steps in (NI-2.5) and $e_t[{}^t v_{i2}/x]$ reduces in n_{t2} steps in (NI-3.5)

We instantiate Theorem 2.57 with $e_T, {}^t v_{i1}, {}^t v_{i2}, {}^t v_{i1}, {}^t v_{i2}, n_{t1} + 2, n_{t2} + 2, H'_{t1}, H'_{t2}$ and from (NI-2.5) and (NI-3.5) we have ${}^t v_{i1} = {}^t v_{i2}$ and thus ${}^s v'_1 = {}^s v'_2$

□