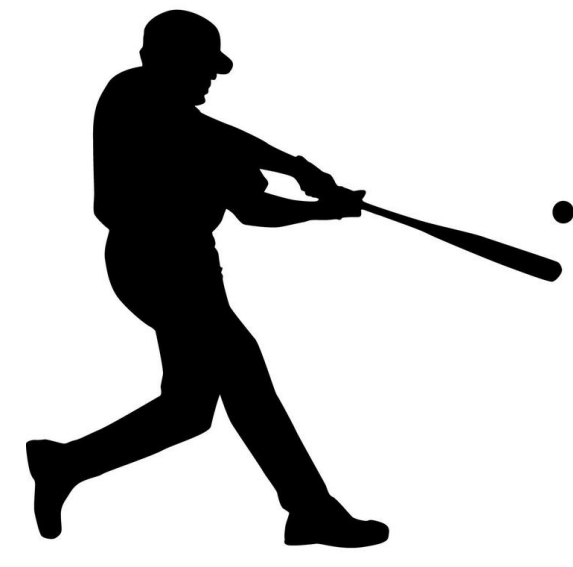


SMASH: Co-Designing Software Compression and Hardware-Accelerated Indexing for Efficient Sparse Matrix Operations



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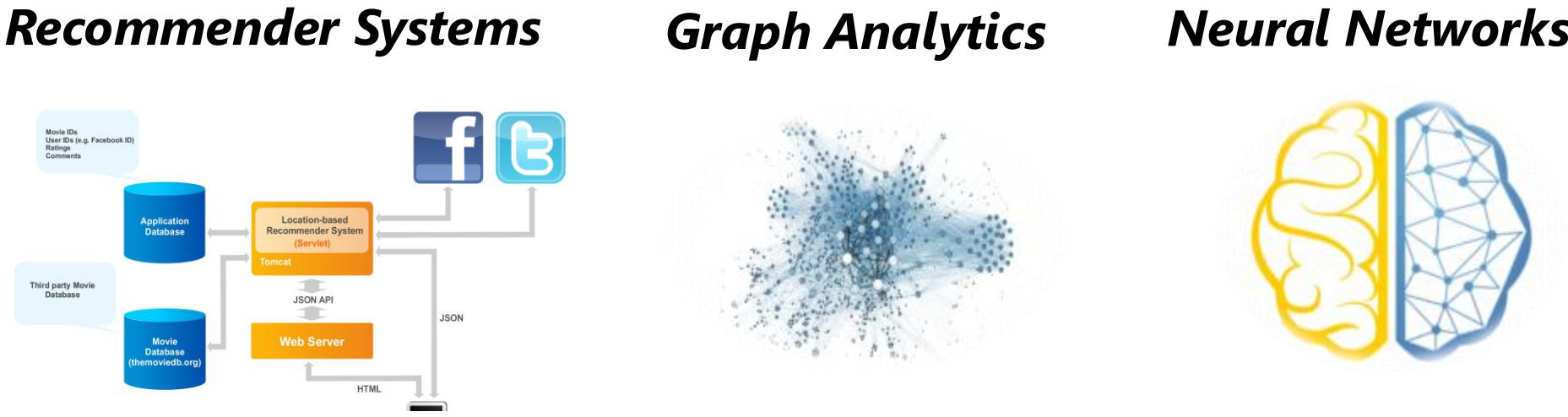
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Mellon
University

1: Summary

- Many important workloads heavily involve **Sparse Matrix Operations**
- Extremely sparse matrices require **compression** to avoid storage and computational overheads
- Shortcomings of existing compression formats
 - expensive discovery of the positions of non-zero elements** or
 - narrow applicability**
- SMASH**: hardware/software cooperative mechanism for efficient sparse matrix storage and computation
 - Software**: Efficient compression scheme using a Hierarchy of Bitmaps
 - Hardware**: Hardware unit interprets the Bitmap Hierarchy and accelerates indexing
- Performance improvement: **38% and 44%** for **SpMV** and **SpMM** over the widely used CSR format
- SMASH** is **highly efficient, low-cost** and **widely applicable**

2: Sparse Matrix Operations and Compression Formats

Sparse Matrix Operations are widespread today



Limitations of Existing Compression Formats

- General formats optimize for storage → **Expensive** discovery of the positions of non-zero elements
- Specialized formats assume **specific matrix structures** and patterns (e.g., diagonals) → **Narrow applicability**

Compressed Sparse Row Indexing overhead

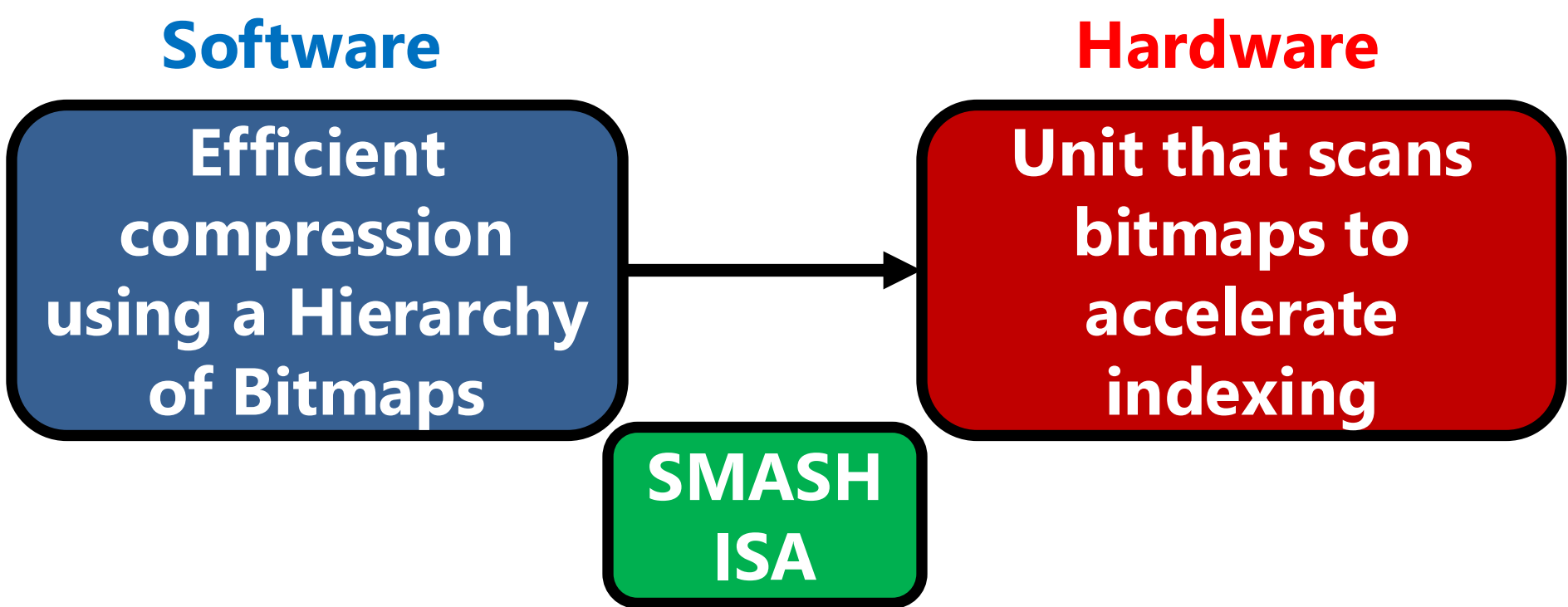
MATRIX	CSR
$A = \begin{bmatrix} 5 & 0 & 0 & 0 \\ 0 & 0 & 3 & 4 \\ 0 & 8 & 0 & 0 \\ 0 & 0 & 0 & 3 \end{bmatrix}$	$\text{row_ptr} \begin{bmatrix} 0 & 1 & 3 & 4 & 5 \end{bmatrix}$ $\text{col_ind} \begin{bmatrix} 0 & 2 & 3 & 1 & 3 \end{bmatrix}$ $\text{values} \begin{bmatrix} 5 & 3 & 4 & 8 & 3 \end{bmatrix}$

Many instructions needed to discover the position of a non-zero element

Requires multiple data-dependent memory accesses

3: SMASH

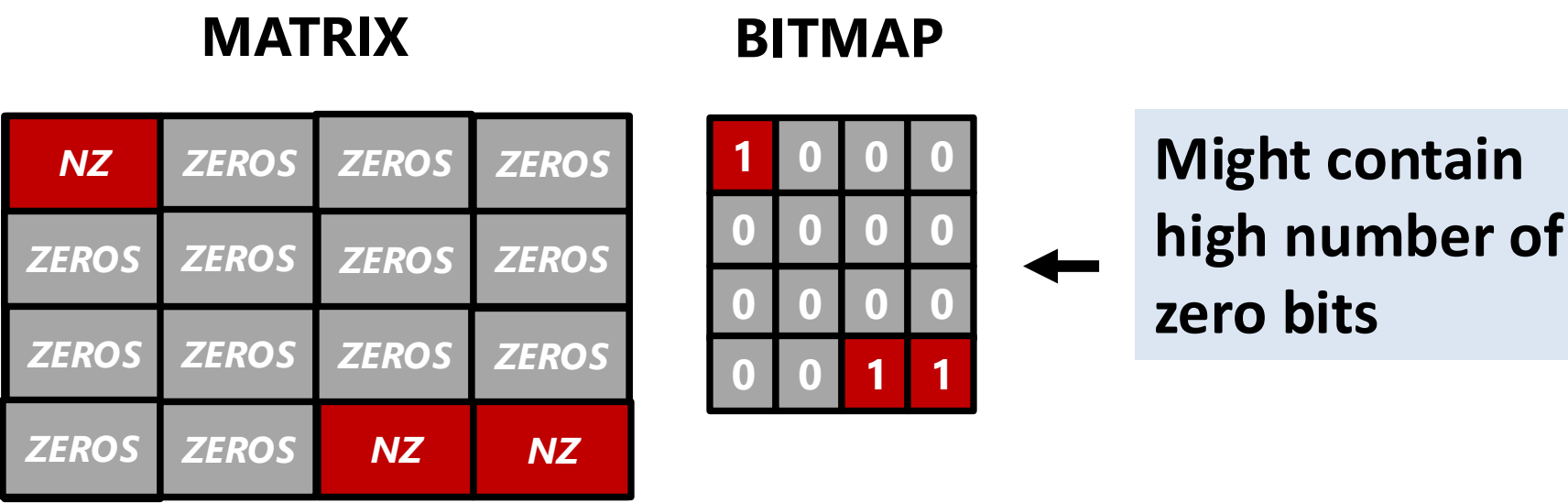
- Hardware/Software cooperative mechanism: Explain SMASH
- Enables **highly-efficient** sparse matrix compression and computation
 - General** across a diverse set of sparse matrices and sparse matrix operations



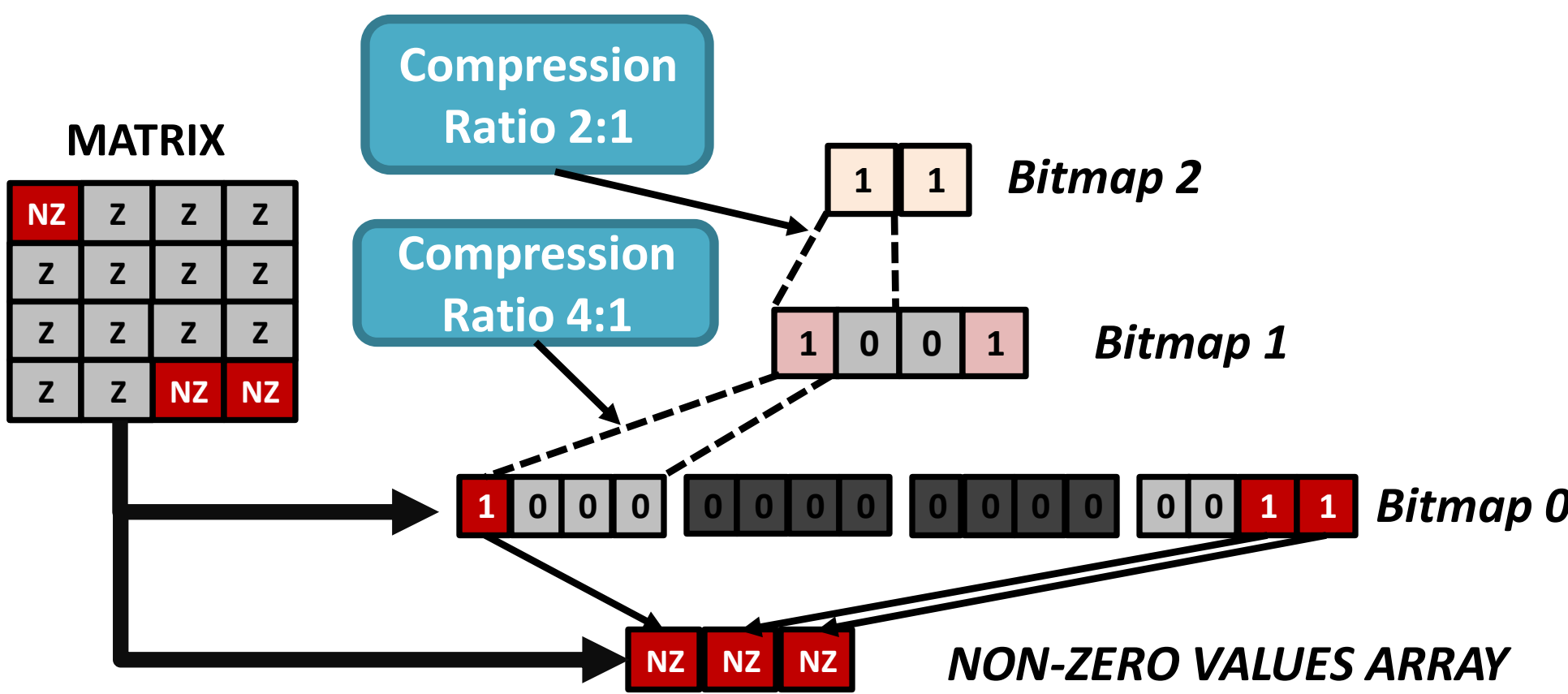
4: Software Compression Scheme

Bitmap

Encodes the presence/absence of a non-zero element in a block of the original matrix with a single bit

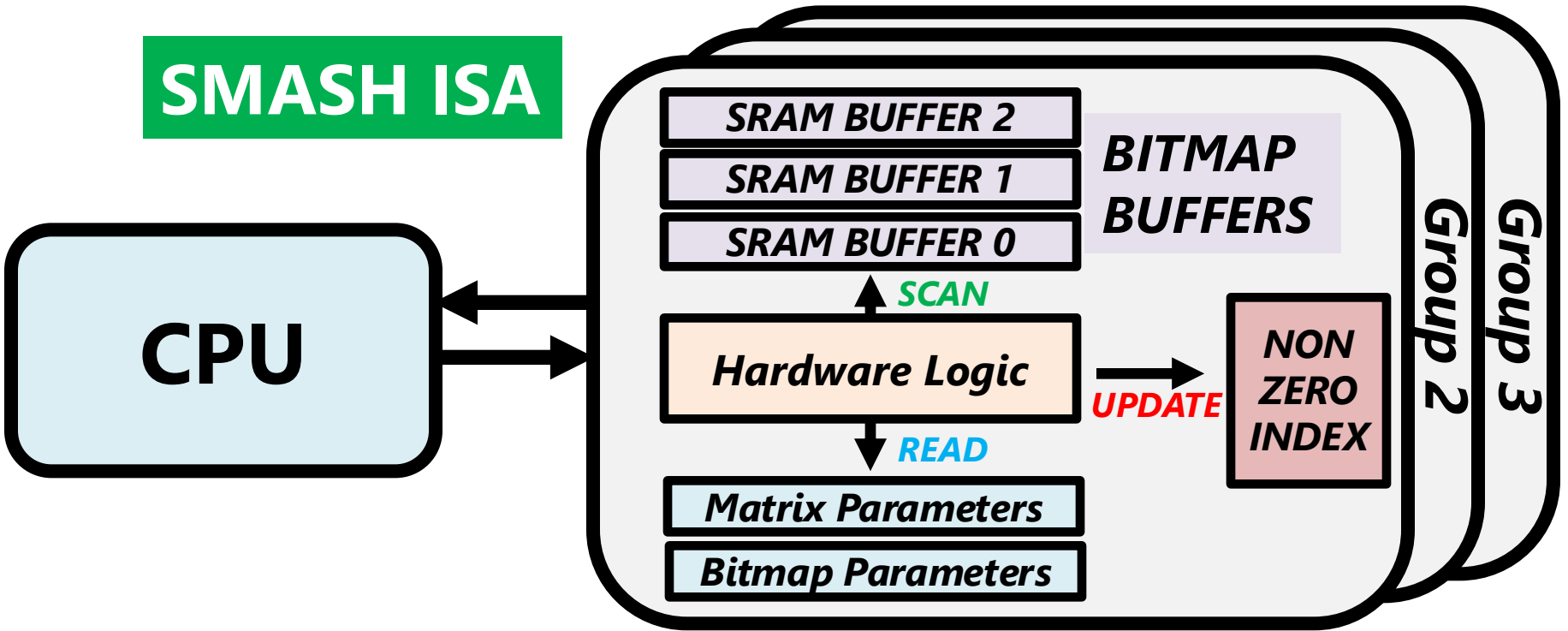


Hierarchy of Bitmaps



5: Hardware Acceleration Unit

Bitmap Management Unit (BMU)



6: Cross-Layer Interface

Need for a cross-layer interface that enables software to control the BMU

- Communicate** the **parameters** needed to calculate the index
- Query** the BMU to **retrieve the index** of the next non-zero element

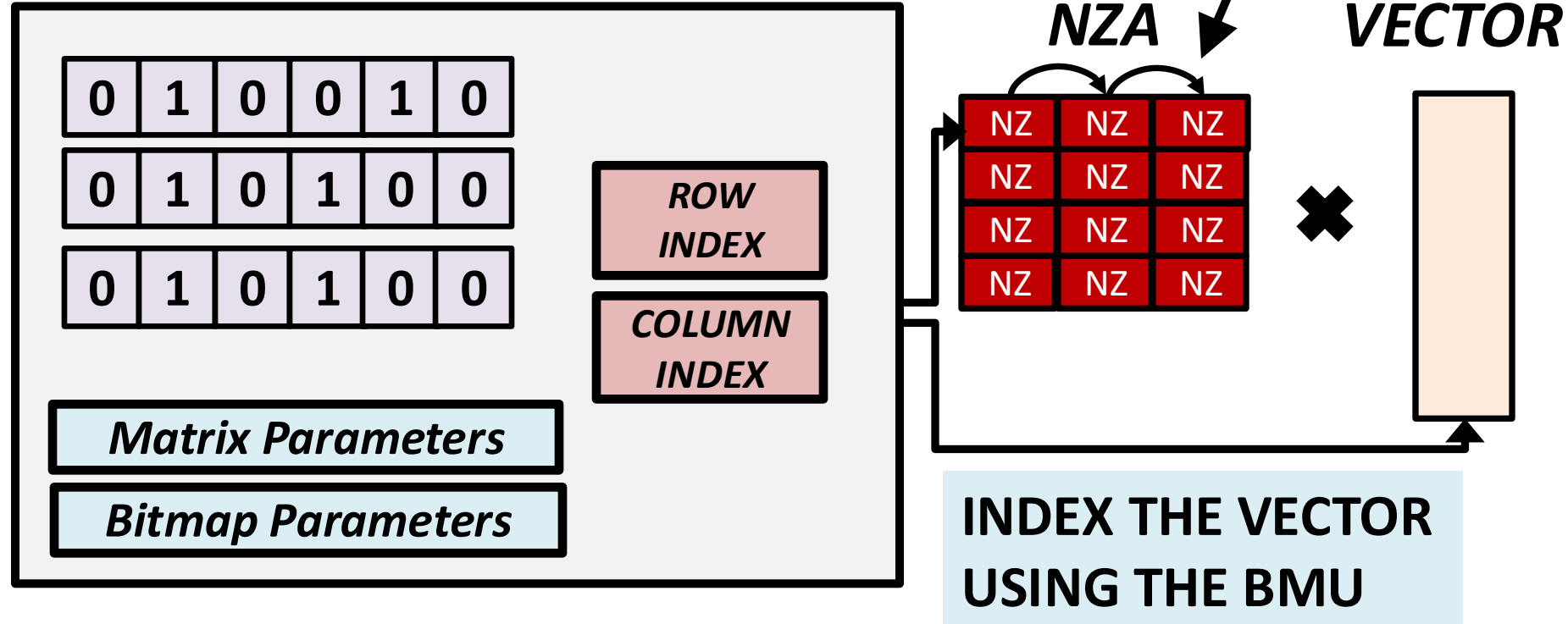
Enables SMASH to flexibly accelerate a diverse range of operations on any sparse matrix

SMASH ISA

MATINFO
BMAPINFO
RDBMAP
PBMAP
RDIND

7: Use Case: SpMV

BMU



8: Evaluation

Methodology

Simulator: ZSim Simulator

Workloads:

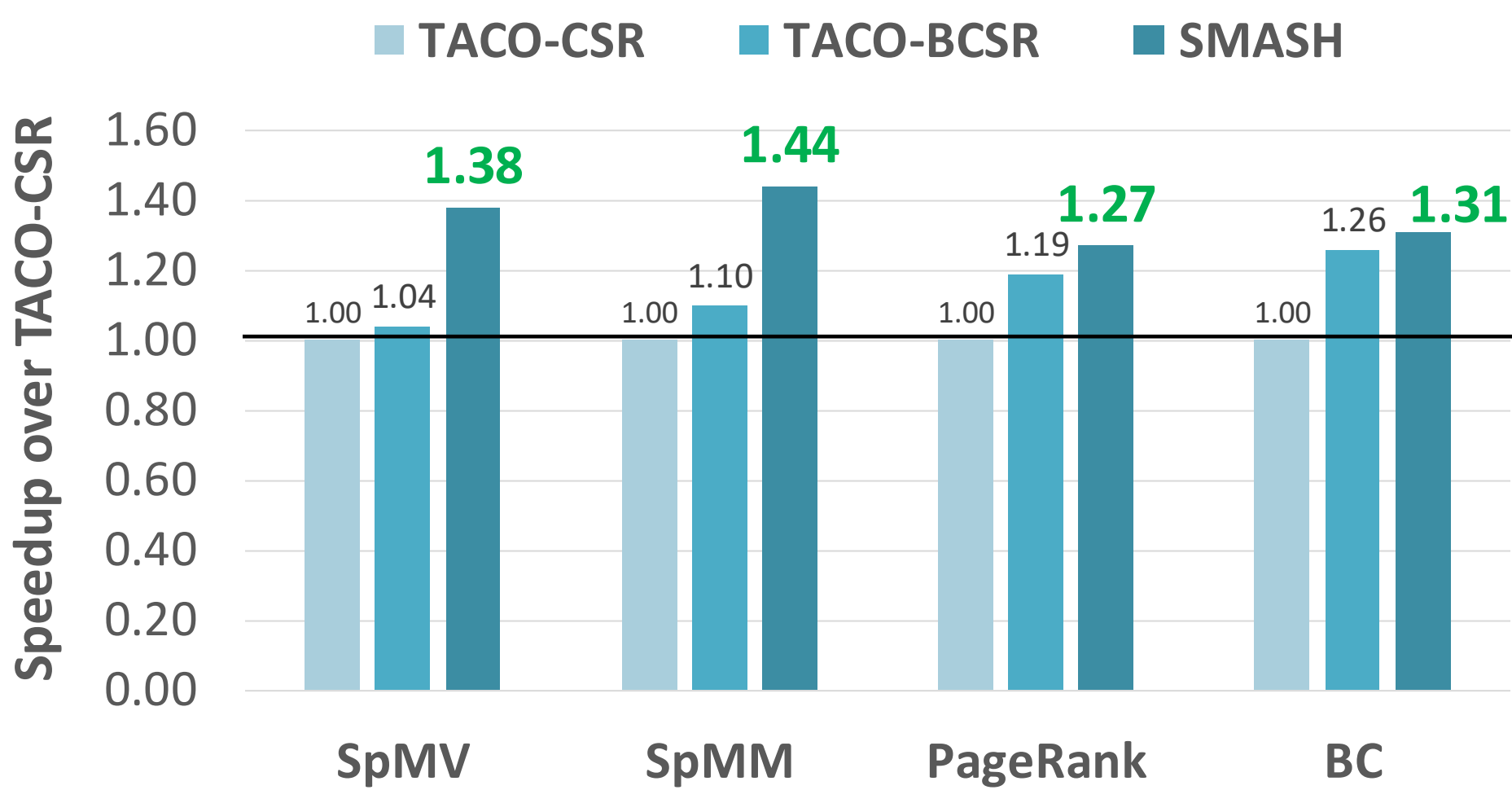
- Sparse Matrix Kernels**
SpMV & SpMM from TACO
- Graph Applications**
PageRank & Betweenness Centrality from Ligra

Input dataset:

15 diverse sparse matrices & 4 graphs from the Sparse Suite Collection [4]

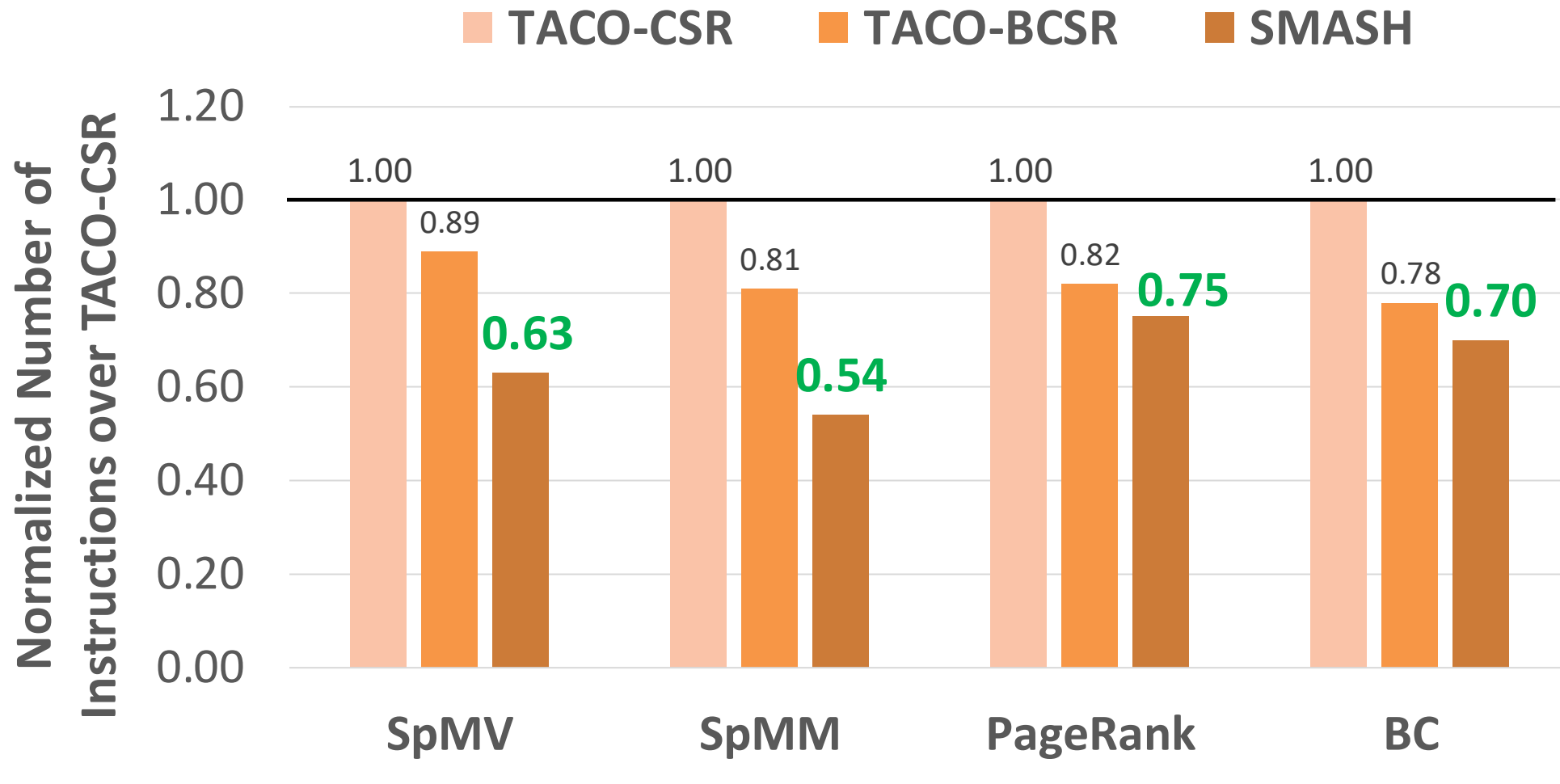
Sparsity ranges from 0.01% to 8.79%

Performance Improvement using SMASH



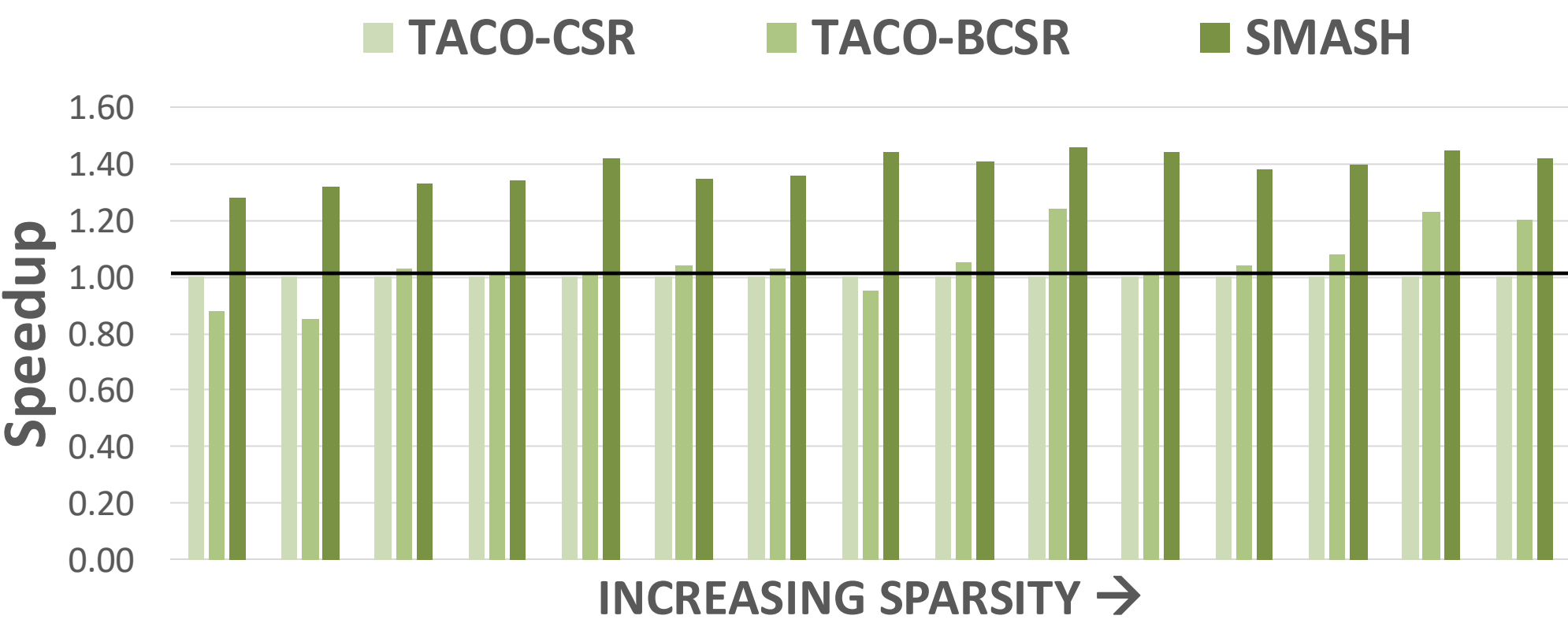
SMASH provides significant performance improvements over state-of-the-art formats

Number of Executed Instructions



SMASH significantly reduces the number of executed instructions

Sparsity sweep for SpMV



SMASH provides speedups regardless of the sparsity of the matrix

Hardware Overhead

SMASH configuration

- Support for 4 matrices in the BMU
- 256 bytes / SRAM buffer
- 140 bytes for registers & counters

0.076% area overhead over an Intel Xeon CPU

SMASH incurs negligible area overhead

Other Results in the Paper

- Compression ratio sensitivity analysis
- Distribution of non-zero elements
- Detailed results for SpMM
- Conversion from CSR to SMASH overhead
- Software-only approaches

Scan for full paper



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