

Securely Compiling Partial Programs Using the Pointers-as-Capabilities Principle

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could be **buggy** or **malicious**















#include "networking.h"

void init_secret(char* s); void process(char* b, char* s);

static char secret[256];

char iobuffer[1024];

int main(void) {

init_secret(secret);
receive(iobuffer);
process(iobuffer, secret);

return 0;





could be **buggy** or **malicious**





could be **buggy** or **malicious**



Let be confidentiality



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could be **buggy** or malicious

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Isolate the memory of the different parts of the program from each other (with low performance overhead) while allowing pointer passing.



We have two requirements for



Isolate the memory of the different parts of the program from each other (with low performance overhead) while allowing pointer passing.

Want a proof technique that allows us to reuse a whole-program compiler correctness theorem.





Isolate the men

of the program from performance over **pointer passing**.

Compiler correctness is a more standard verification criterion.

Goal: avoid repeating **years**-**worth of proof effort**.





Want a proof technique that allows us to **reuse a whole-program compiler correctness theorem**.



Isolate the memory of the different parts

of the pr perform **pointer**

Hardware capabilities

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Novel proof technique (called TrICL "/ˈtrɪk(ə)l/")











Prior work on Compiler security





Prior work on Compiler security







First compiler security proof that achieves reuse of the compiler correctness proof while allowing memory sharing through pointer passing

Prior work on Compiler security

> Achieved isolation by preventing memory sharing altogether.





Novel proof technique (called TrICL "/ˈtrɪk(ə)l/")

Achieved

isolation by preventing memory sharing altogether.





First compiler security proof that achieves reuse of the compiler correctness proof while allowing memory sharing through pointer passing

Novel proof technique (called TrICL "/'trɪk(ə)l/")

security



altogether.

C-to-C
sourceexter
struct
att
att
int fi
that adds
CHERIannotations
automatically-att
att
att
att
int fi
sandb
{

extern struct cheri_object lib1; struct cheri_object lib2;

__attribute__((cheri_ccallee))
__attribute__((cheri_method_class(lib1)))
int f1(void);

__attribute__((cheri_ccall))
__attribute__((cheri_method_class(lib2)))
int f2(void);

__attribute__ ((constructor)) static void sandboxes_init(void)

lib2 = fetch_object("lib2");

int f1(void)

f2();

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altogether.

Reproved CapablePtrs

First compiler security proof that achieves reuse of the compiler correctness proof while allowing memory sharing through pointer passing



We have **two requirements** for **compiler security**

Isolate the memory of the different parts

of the pr performa pointer













Virtual memory

Every memory access instruction expects a capability as an argument.

The compiler implements pointer passing as capability passing.

No program part can forge capabilities.

capability









could be **buggy** or **malicious**





iobuffer[102

4



iobuffer[102

4







	Hardwa	218CHAPTER 7. THE CHERI-MIPS INSTRUCTION-SET REFERENCELoad Integer via Capability Register	
	R	<pre>if not (cb_val.tag) then raise_c2_exception(CapEx_TagViolation, cb)</pre>	
F.	could be	<pre>etse if cb_val.sealed then raise_c2_exception(CapEx_SealViolation, cb) else if not (cb_val.permit_load) then raise_c2_exception(CapEx_PermitLoadViolation, cb)</pre>	iobuffer
		<pre>else { let 'size = wordWidthBytes(width); let curses = cetCenCurses(ch.well);</pre>	[1024]
char	*secret_p (char*)4	<pre>let cursor = getCapCursor(cb_val); let vAddr = (cursor + unsigned(rGPR(rt)) + size*s let vAddr64 = to_bits(64, vAddr); if (vAddr + size) > getCapTop(cb_val) then</pre>	
leak	(*secret_p	<pre>raise_c2_exception(CapEx_LengthViolation, cb)</pre>	



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Hardware capabilities



Want a proof technique that allows us to

reuse a correcti Novel proof technique (called TrICL "/ˈtrɪk(ə)l/")







A partial program is **secure**



when **NO library** can distinguish two runs (with **two different secrets**) from each other.



A partial program is **secure**



The same definition for the **target language** too

when **NO library** can distinguish two runs (with **two different secrets**) from each other.





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Novel proof technique (called TrICL "/ˈtrɪk(ə)l/")



compileduntrustedprogramthird-partypartlibrary













Novel proof technique (called TrICL "/'trɪk(ə)l/")



Trace-directed Back-translation

compiled program part **untrusted** third-party library






























































More in the paper

Novel proof technique (called TrICL "/ˈtrɪk(ə)l/")



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Trace-directed Back-translation

example

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Trace-directed Back-translation

example

Summary: In CapablePtrs, we present a proof of compiler full abstraction that achieves reuse of the compiler correctness lemmas while allowing memory sharing through pointer passing.